1056-05-60Patrick Bahls* (pbahls@unca.edu), CPO #2350, One University Heights, Asheville, NC28804-8511. The L(2, 1) channel-assignment problem on trees. Preliminary report.

Let G = (V, E) be a simple graph. We say that a non-negative integer labeling ℓ of its vertices V is called an L(2, 1)labeling if for every pair $\{u, v\}$ of adjacent vertices $|\ell(u) - \ell(v)| \ge 2$, and for every pair $\{u, v\}$ satisfying $\rho(u, v) = 2$, $|\ell(u) - \ell(v)| \ge 1$, where ρ is the usual path metric on V. (Such labelings model the assignment of non-interfering "channels" to nearby radio transmitters.) The L(2, 1)-span of a graph G, $\lambda(G)$, is defined to be the minimum value, over all L(2, 1)-labelings of G, of $\max_{v \in V} \ell(v)$.

In 1992 J.R. Griggs and R.K. Yeh proved that for a tree T with maximal vertex degree Δ , $\lambda(T) \in \{\Delta + 1, \Delta + 2\}$, but conjectured that for an arbitrary tree determining which of these values obtains would prove to be NP-hard.

We describe a deterministic algorithm for computing $\lambda(T)$ in the case $\Delta = 3$ and indicate how this algorithm can be generalized to arbitrary maximal degree Δ . The algorithm has exponential time complexity, and its construction shows why no more efficient deterministic algorithm can be found. (Received July 21, 2009)