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**L. Egri, P. Hell, B. Larose\*** (benoit.larose@concordia.ca) and **A. Rafiey**. *Space complexity of list  $H$ -coloring: a dichotomy.*

The Dichotomy Conjecture for constraint satisfaction problems (CSPs) states that every CSP is in P or is NP-complete (Feder-Vardi, 1993). It has been verified for conservative problems (also known as list homomorphism problems) by A. Bulatov (2003). We augment this result by showing that for digraph templates  $H$ , every conservative CSP, denoted  $\text{LHOM}(H)$ , is solvable in logspace or is hard for NL. More precisely, we introduce a digraph structure we call a circular  $N$ , and prove the following. if  $\mathbf{A}(H)$  denotes an algebra whose term operations are the conservative polymorphisms of the digraph  $H$ , then the following conditions are equivalent: (1) the variety generated by  $\mathbf{A}(H)$  admits only the Boolean type; (2) the variety generated by  $\mathbf{A}(H)$  is congruence  $k$ -permutable for some  $k$ ; (3) the digraph  $H$  contains no circular  $N$ . If one of these conditions holds then the problem  $\text{LHOM}(H)$  is solvable in logspace, otherwise it is NL-hard. Moreover, we show that the presence of a circular  $N$  can be decided in time polynomial in the size of  $H$ . (Joint work with L. Egri, P. Hell and A. Rafiey.) (Received August 01, 2013)