
The notion of noise sensitivity was introduced by Benjamini, Kalai, and Schramm in 1998, in the context of percolation theory. Since then noise sensitivity has found applications in many fields, including concentration of measure, social choice theory, and theoretical computer science.

Say we have $n$ people voting for one of two candidates, $\mathcal{D}$ or $\mathcal{R}$. One can write their choices as $\omega = (\omega(1), \ldots, \omega(n)) \in \{\mathcal{D}, \mathcal{R}\}^n$, and the election is decided by some voting rule $f(\omega) \in \{\mathcal{D}, \mathcal{R}\}$. Errors occur during the ballot counting process, and the election is actually decided according to slightly perturbed data $\omega\epsilon$. The subject of noise sensitivity studies how susceptible the voting rule is to the noise as a function of the number of voters $n$ and the level of noise $\epsilon$. Intuitively, $f(\cdot)$ is noise sensitive if even for a small level of noise, given enough voters, $f(\omega\epsilon)$ gives us very little information on the true outcome of the election $f(\omega)$. On the other hand, $f(\cdot)$ is noise stable if for a small enough level of noise the election is decided by the unperturbed data $\omega$, regardless of the number of voters.

A Boolean function is a function from the hypercube $\Omega_n = \{-1, 1\}^n$ into $\{-1, 1\}$ (or $\{0, 1\}$). We call the elements of $\Omega_n$ configurations. Consider the hypercube $\Omega_n$ endowed with the uniform measure $\mathbf{P}_n^\omega = (\frac{1}{2} \delta_{-1} + \frac{1}{2} \delta_1)^\otimes n$. Let $\omega \in \Omega_n$ be a configuration sampled according to $\mathbf{P}_n^\omega$, and let $\omega'\epsilon$ be an independent configuration. For every $\epsilon > 0$, denote by $\omega\epsilon$ the “noisy” configuration obtained from $\omega$ by resampling each bit independently with probability $\epsilon$, i.e., independently for every $x \in [n] := \{1, 2, \ldots, n\}$,

$$\omega\epsilon(x) = \begin{cases} \omega(x) & \text{with probability } 1 - \epsilon, \\ \omega'(x) & \text{with probability } \epsilon. \end{cases}$$

We denote by $\mathbf{P}$ the joint distribution of $\omega$ and $\omega\epsilon$, and by $\mathbf{E}$ the expectation with respect to $\mathbf{P}$.

Next we introduce two main concepts regarding sensitivity to noise.

**Definition 1.** Let $m_n$ be an increasing sequence of integers, and let $\{f_n\}$ be a sequence of functions $f_n : \Omega_{m_n} \to \{-1, 1\}$. The sequence $\{f_n\}$ is called noise sensitive if for every $\epsilon > 0$,

$$\lim_{n \to \infty} \left( \mathbf{E}[f_n(\omega)f_n(\omega\epsilon)] - \mathbf{E}[f_n(\omega)]^2 \right) = 0.$$

Note that since $f_n$ are Boolean, zero covariance implies independence. Next we present a trivial but important example of a noise sensitive sequence of functions.

**Example 2.** Define the parity function, $\text{PAR}_n(\omega(1), \ldots, \omega(n)) = \prod_{i=1}^n \omega(i)$. Note that for every $i$, with probability $1 - \epsilon$, $\omega(i) = \omega\epsilon(i)$, and thus their product is 1. With probability $\epsilon$, $\omega(i)$ and $\omega\epsilon(i)$ are independent and thus their product has zero

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expectation. A direct calculation shows
\[
\mathbb{E}[\text{PAR}_n(\omega)] = \prod_{i=1}^n \mathbb{E}[\omega(i)\omega(i+1)] = ((1 - \epsilon) \cdot 1 + \epsilon \cdot 0)^n \rightarrow 0.
\]
This, together with the fact that \(\mathbb{E}[\text{PAR}_n(\omega)] = 0\), yields that \(\{\text{PAR}_n\}\) is a noise sensitive sequence.

The other extreme case to noise sensitivity is noise stability.

**Definition 3.** The sequence of functions \(\{f_n\}\) is noise stable if
\[
\lim_{\epsilon \to 0} \sup_n \mathbb{P}[f_n(\omega) \neq f_n(\omega^\epsilon)] = 0.
\]

The next example shows an important noise stable sequence of Boolean functions.

**Example 4.** Define the dictator function, \(\text{DICT}_n(\omega(1), \ldots, \omega(n)) := \omega(1)\). It is immediate that
\[
\mathbb{P}[\text{DICT}_n(\omega) \neq \text{DICT}_n(\omega^\epsilon)] = \epsilon \mathbb{P}[\omega(1) \neq \omega'(1)] = \epsilon/2,
\]
which tends to zero uniformly on \(n\) as \(\epsilon \to 0\). Thus, the dictator function is noise stable.

Note that there are examples of sequences that are neither noise sensitive nor noise stable, and a sequence can be both if and only if \(\lim_{n \to \infty} \text{Var}[f_n] = 0\).

Historically, the most important example of noise sensitivity is the crossing event in critical bond percolation. Denote by \(\mathcal{E}(\mathbb{Z}^2)\) the set of edges of the planar integer lattice. For \(p \in [0, 1]\), consider the product measure \(\mathbb{P}_p = ((1 - p)\delta_0 + p\delta_1) \otimes \mathcal{E}(\mathbb{Z}^2)\) on \(\Omega = \{0, 1\}^{\mathcal{E}(\mathbb{Z}^2)}\). One can think of a configuration \(\omega \in \{0, 1\}^{\mathcal{E}(\mathbb{Z}^2)}\) as a subgraph of \(\mathbb{Z}^2\) by prescribing which edges in \(\mathbb{Z}^2\) are included. We call two edges neighbors if they have a common vertex. We say that two edges \(e, e' \in \mathcal{E}(\mathbb{Z}^2)\) are connected in a configuration \(\omega\) if there is a path of edges \(\{e_i\}_{i=1}^m\) satisfying for all \(i \in [m]\), \(e_i\) is a neighbor of \(e_{i+1}\), \(\omega(e_i) = 1\) and \(e_1 = e, e_m = e'\). The behavior of large clusters in percolation exhibits a sharp phase transition. There is some critical number \(p_c\) such that for \(p > p_c\) there is an infinite connected component almost surely, and if \(p \leq p_c\) all connected components are finite almost surely. Harry Kesten proved in [15] that \(p_c(\mathbb{Z}^2) = 1/2\).

Given \(a, b \in \mathbb{N}\), consider bond percolation at the critical point \(p_c(\mathbb{Z}^2) = \frac{1}{2}\), in the increasing sequence of rectangles, \(B_n := [0, an] \times [0, bn] \cap \mathbb{Z}^2\).

We say that there is a left right crossing if there is an edge in \(\{0\} \times [0, bn]\) which is connected to an edge in \(\{an\} \times [0, bn]\).
Define the sequence of events \( A_n = \{ \text{there is a left right crossing} \} \), and the sequence of Boolean functions \( f_n : \{0, 1\}^{\mathcal{E}(B_n)} \to \{0, 1\} \),

\[
f_n(\omega) = \# A_n(\omega) = \begin{cases} 1 & \text{if } \omega \in A_n, \\ 0 & \text{else.} \end{cases}
\]

A fundamental result of Benjamini, Kalai, and Schramm [3, Theorem 1.2] states that \( \{ f_n \} \) is noise sensitive.

In the heart of the theory of noise sensitivity lies the notion of influence. This concept first arose in political science to measure the power of individual voters in a voting scheme. For a configuration \( \omega \in \Omega_n \), denote by \( \omega^k \) the configuration obtained from \( \omega \) by flipping the \( k \)th coordinate,

\[
\omega^k(x) = \begin{cases} \omega(x) & \text{if } x \neq k, \\ -\omega(x) & \text{if } x = k. \end{cases}
\]

Abbreviate the discrete partial derivatives \( \forall k \in [n], \nabla_k f(\omega) = f(\omega) - f(\omega^k) \).

**Definition 5.** The influence of the \( k \)th coordinate is defined as \( I_k(f) := \|\nabla_k f\|_1 = \mathbb{E}[|\nabla_k f|] \).

For Boolean functions (to \( \{-1, 1\} \)), \( \nabla_k f \in \{-2, 0, 2\} \), and thus

\[
I_k(f) = 2\mathbb{P}[f(\omega) \neq f(\omega^k)].
\]

This relates to the classical concept of pivotal edges from percolation theory. Say that \( k \in [n] \) is pivotal for \( f \) in the configuration \( \omega \) if \( f(\omega) \neq f(\omega^k) \). Note that the event \( \{ f(\omega) \neq f(\omega^k) \} \) is independent of \( \omega(k) \). If we denote by \( \mathcal{P} = \mathcal{P}_f(\omega) := \{ i \in [n] : i \text{ is pivotal for } f \text{ in } \omega \} \), we can write \( I_k(f) = 2\mathbb{P}[k \in \mathcal{P}] \).

**Definition 6.** The total influence, \( I(f) \), is defined by

\[
I(f) := \sum_{k=1}^{n} I_k(f) = 2\mathbb{E}[|\mathcal{P}|].
\]

We are now in a position to state one of the main theorems in the theory of noise sensitivity:

**Theorem 7** (\[3, Theorem 1.3\]). Let \( \{ f_n \} \) be a sequence of Boolean functions. If

\[
\lim_{n \to \infty} \sum_{k=1}^{n} I_k(f_n)^2 = 0,
\]

then \( \{ f_n \} \) is noise sensitive.

The converse is not true in general, as can be seen by Example [2]. However, it is true for a sequence of monotonic functions.

**Definition 8.** A function \( f : \Omega_n \to \mathbb{R} \) is monotone if \( f(x) \leq f(y) \) whenever \( \forall j \in [n], x(j) \leq y(j) \).

**Theorem 9** (\[3, Theorem 1.4\]). Let \( \{ f_n \} \) be a sequence of monotone Boolean functions. If

\[
\inf_{n} \sum_{k=1}^{n} I_k(f_n)^2 > 0,
\]

then \( \{ f_n \} \) is not noise sensitive (but not necessarily noise stable).
In recent years, Theorem 11 and other results such as the KKL theorem (Theorem 11 below), were generalized by Keller, Mossel, and Sen 13,14 to continuous distributions such as the Gaussian measure.

**Tools from discrete harmonic analysis**

The main tool in the theory we are discussing is the *Fourier–Walsh expansion*. It is useful to define the expansion for non-Boolean functions, thus we consider the space $L^2(\Omega_n)$ of real valued functions. The inner product is defined relative to the uniform measure $P^n$, $(f,g) := \sum_{\omega} 2^{-n} f(\omega)g(\omega) = \mathbf{E}[fg]$. As the complete orthonormal Fourier base we take the Walsh functions: for any subset $S \subset [n]$, define

$$\chi_S(\omega) := \prod_{i \in S} \omega(i),$$

with the convention that $\chi_\emptyset \equiv 1$. Thus any function can be decomposed according to the Fourier–Walsh series

$$f = \sum_{S \subset [n]} \hat{f}(S)\chi_S := \sum_{S \subset [n]} \langle f, \chi_S \rangle \chi_S.$$

This allows us to study concentration and noise sensitivity of functions by their Fourier–Walsh spectrum. The catch phrase one can learn from Benjamini, Kalai, and Schramm 3 is that functions of high frequency are noise sensitive and functions of low frequencies are noise stable.

In the case of a monotonic function $f : \Omega_n \to \{-1, 1\}$, we get

$$I_k(f) = \mathbf{E}[|\nabla_k f|] = \mathbf{E}[|\nabla_k f|^2] = 2\mathbf{E}[\hat{f}(S)\chi_S],$$

where the third equality is due to the monotonicity of $f$. For every function

$$\mathbf{E}[f \cdot \omega(k)\chi_k] = 0,$

we get for monotonic Boolean functions $f(\{k\}) = 1_k f/2$. The Cauchy–Schwarz inequality yields for monotonic functions that $I(f) \leq \sqrt{n}$. A calculation shows the majority function has total influence equal to $\sqrt{n}$, hinting that the majority function is maximal among all monotonic functions.

**Influence, concentration of measure, and hypercontractivity**

The study of influences is very important in many scientific disciplines, e.g., learning, information theory, and social choice theory.

For a Boolean function (to $\{-1, 1\}$), a calculation shows that

$$\nabla_k f(S) = \begin{cases} 2\hat{f}(S) & \text{if } k \in S, \\ 0 & \text{otherwise.} \end{cases}$$

Since $\nabla_k f \in \{-2, 0, 2\}$, $\|\nabla_k f\|_2^2 = \mathbf{E}[|\nabla_k f|^2] = 4\mathbf{E}[|\nabla_k f|] = 4\|\nabla_k f\|_1$, and together with Parseval’s formula we obtain that

$$I_k(f) = \sum_{S \subset [n]: k \in S} \hat{f}(s)^2$$

and

$$I(f) = \sum_{S \subset [n]} |S|\hat{f}(S)^2.$$

The variance of a function $f$ can be easily represented with Parseval’s formula:

$$\mathbf{Var}[f] = \mathbf{E}[f^2] - \mathbf{E}[f]^2 = \langle f, f \rangle - \langle f, \chi_\emptyset \rangle^2 = \sum_{S \subset [n]} \hat{f}(S)^2 - \hat{f}(\emptyset)^2.$$
This together with (1) immediately gives us the \textit{discrete Poincaré inequality} for Boolean functions.

\textbf{Theorem 10.} \textit{For any Boolean function } $f : \Omega_n \to \{-1, 1\}$, \begin{equation*} \text{Var}[f] \leq \sum_{k} I_k(f). \end{equation*} \end{center} This means that \begin{equation*} \max_{i} I_i(f) \geq \frac{\text{Var}[f]}{n}. \end{equation*} The celebrated KKL theorem of Kahn, Kalai, and Linial \[11\] gives a logarithmic correction to the above inequality. Though it might seem that a log correction is not significant, the fact that it diverges was enough for Kahn, Kalai, and Linial to show that in an idealized voting system with two candidates it is enough to bribe $o(1)$ fraction of the voters in order to control the elections with high probability.

\textbf{Theorem 11 (\[11\]).} \textit{There is a universal constant } $c > 0$ \textit{such that for any Boolean function} \begin{equation*} \max_{i} I_i(f) \geq c \frac{\text{Var}[f]}{n}. \end{equation*} Remarkably, the logarithmic correction is sharp, as shown by the tribes example of Ben-Or and Linial \[2\]: Partition $[n]$ into blocks of length $\log_2(n) - \log_2(\log_2(n))$, and connect the leftover arbitrarily. Define $f_n$ to be 1 if there is a block of all 1’s and 0 otherwise.

An inequality of Talagrand \[20\], which improves upon some results of Kahn, Kalai, and Linial \[11\], proved to be instrumental in the study of fluctuations in \textit{first passage percolation}. For $p \in [0, 1]$, consider the measure $P^n_p = ((1 - p)\delta_{-1} + p\delta_1)^\otimes n$ on $\Omega_n$.

\textbf{Theorem 12 (\[20\] Theorem 1.5).} \textit{Let } $f : \Omega_n \to \mathbb{R}$ \textit{then for any } $p \in [0, 1]$ \textit{there is a } $K = K(p) \in (0, \infty)$ \textit{such that} \begin{equation*} \text{Var}_p(f) \leq K \cdot \sum_{k=1}^{n} \frac{\|\nabla_k f\|_2^2}{1 + \log \left(\frac{\|\nabla_k f\|_2}{\|\nabla_k f\|_1}\right)}. \end{equation*} Talagrand’s inequality also proved useful in the study of percolation isoperimetry \[5, 9, 19\], showing concentration of the Cheeger constant and proving a weak version of a conjecture by Benjamini.

Talagrand’s approach for concentration estimates, such as Theorem 12, is based on \textit{hypercontractivity}. In \[7\], Sourav Chatterjee presents a modern general approach to hypercontractivity. A semigroup $P_t$ is called hypercontractive if for any $p > 1$ and $t > 0$ there exists a $q(t, p) > p$ such that \begin{equation*} \|P_t f\|_{L^q} \leq \|f\|_{L^p}. \end{equation*} This approach began with Nelson’s \[18\] proof of the hypercontractivity of the Ornstein–Uhlenbeck semigroup and Gross’s proof \[10\] that a semigroup is hypercontractive if the associated Dirichlet form satisfies a logarithmic Sobolev inequality. First passage percolation (FPP) is a model for a random metric on a graph; see \[1\] for a thorough review of FPP. Consider the edge set $\mathcal{E}$ of $\mathbb{Z}^d$. For every edge $e \in \mathcal{E}$ we associate a weight $\omega(e)$ distributed according to some measure $\mu$, such that $\{\omega(e)\}_e$ are i.i.d. For any two vertices $x, y \in \mathbb{Z}^d$ and any path connecting
them $\gamma : x \mapsto y$, $\gamma = (e_1, \ldots, e_l)$, we associate a length $\lambda(\gamma) = \sum_{i=1}^{l} \omega(e_i)$. This defines in general a semimetric $d_\omega(x, y) = \inf_{\gamma : x \mapsto y} \lambda(\gamma)$. By Kingman’s subadditive ergodic theorem [16,17], under weak assumptions on $\mu$, one obtains that balls in the metric converge in the Hausdorff sense to a ball $D$, of a deterministic norm on $\mathbb{R}^d$,

$$\lim_{n \to \infty} D_n := \lim_{n \to \infty} \frac{1}{n} \{ x \in \mathbb{Z}^d : d_\omega(0, x) < n \} = D.$$ 

This can be thought of as a geometric version of the law of large numbers. Many open questions remain about the asymptotic shape $D$; e.g., if $\mu$ is a continuous distribution, is $D$ strictly convex? What are the conditions for $\mu$ under which corners appear in $D$? The question relevant to our discussion is that of the fluctuations around the limit shape.

Benjamini, Kalai, and Schramm [4] studied the simple case in which $\mu = p\delta_a + (1-p)\delta_b$, where $0 < a < b$. This allows us to study the metric $d_\omega$ in the framework of Boolean functions. They prove that the variance of the FPP metric is sublinear and thus is also little $o$ of the expectation.

**Theorem 13** ([4, Theorem 1]). There is a constant $C = C(d, a, b)$ such that for any $v \in \mathbb{Z}^d$, $\|v\|_1 > 2$,

$$\text{Var}_p[d_\omega(0, v)] \leq C \frac{\|v\|_1}{\log \|v\|_1}.$$ 

This theorem was generalized by Damron, Hanson, and Sosoe [8] under weak assumptions on $\mu$. This is a good point to demonstrate a simple but powerful idea developed by Benjamini, Kalai, and Schramm. Since we took as our weight distribution a measure supported on two points $\{a, b\}$, any path $\gamma$ connecting $0$ and $v$ for which $\lambda(\gamma) = d_\omega(0, v)$ must satisfy

$$a|\gamma| \leq d_\omega(0, v) \leq b\|v\|_1.$$ 

Thus we get $|\gamma| \leq \frac{b}{a}\|v\|_1$. This means that we need to only look at a finite box around the origin in order to find $d_\omega(0, v)$. Moreover, only edges contained in all minimizing paths (not necessarily unique) can have nonzero influence. Changing the weight of a path from $a$ to $b$ or from $b$ to $a$ can change $d_\omega$ by no more than $|b - a|$. This leads to

$$\sum_{e} \|\nabla_e d_\omega(0, v)\|_2^2 \leq |\mathcal{G}|(b - a)^2 \leq \frac{b}{a}(b - a)^2\|v\|_1,$$ 

providing justification to the nominator in Theorem 13. This simple idea of controlling the sum of influences by some geometric a priori knowledge has proved useful in many applications, e.g., concentration of maximal cardinality matchings for general graphs [6].

Note that if we choose $\mu$ to be the exponential distribution, $D$ is the limiting shape of the *Eden model*, which is the aggregation process one gets by adding edges according to the uniform distribution on the boundary of the aggregate. Kardar, Parisi, and Zhang [12] explain with their KPZ equation that $D_n$ should have $n^{1/3}$ fluctuations around the limiting shape $D$, which is very far from the known bound given by Benjamini, Kalai, and Schramm.
ABOUT THE BOOK

The book under review is a refined version of lecture notes that have been circulating for a few years. Considerable effort was made to make the book as thorough and concise as possible but still readable and friendly. There are many interesting and important subjects covered in the book that we did not discuss in this review. We believe this book is worthy for any departmental or scholarly library. It is clear that it will turn out to be the “go to” source for studying the subject of noise sensitivity of Boolean functions.

REFERENCES


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