

Thus the values of  $p_n$  for  $2100 \leq n \leq 2500$  are still significantly low but higher than the value of  $p_n$  at  $n = 2000$ .

Note that the general size and trend of  $p_n$ , as well as its sudden deviation at  $n = 2000$ , indicate a non random character in the digits of  $e$ .

More detailed investigations are in progress and will be reported later.

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<sup>1</sup> Both  $e$  and  $1/e$  were computed somewhat beyond 2500 D and the results checked by actual multiplication.

*Notes on Numerical Analysis—2*  
*Note on the Condition of Matrices*

1. The object of this note is to establish the following theorem.

**THEOREM.** *Let  $A$  be a real  $n \times n$  non-singular matrix and  $A'$  be its transpose. Then  $AA'$  is more "ill-conditioned" than  $A$ .*

This theorem confirms an opinion expressed by Dr. L. FOX<sup>1</sup> based on his practical experience. The term "condition of a matrix" has been used rather vaguely for a long time. The most common measure of the condition of a matrix has been the size of its determinant, ill-conditioned matrices being those with a "small" determinant. With this interpretation imposed, the theorem is clearly correct. More adequate measures of the condition of a matrix have been proposed recently by JOHN VON NEUMANN & H. H. GOLDSTINE<sup>2</sup> and by A. M. TURING.<sup>3</sup> Their definitions concern all matrices, not just the ill-conditioned ones, characterized by very large condition numbers. The following two of these definitions will form a basis for the proof of the above-mentioned theorem:

The  $P$ -condition number is  $|\lambda_{\max}|/|\lambda_{\min}|$ , where  $\lambda_{\max}$  and  $\lambda_{\min}$  are the characteristic roots of largest and smallest modulus.<sup>2</sup>

The  $N$ -condition number is  $N(A)N(A^{-1})/n$ , where<sup>3</sup>

$$N(A) = \left( \sum_{i,k} a_{ik}^2 \right)^{\frac{1}{2}}.$$

2. Proof of the theorem in the  $P$  case:

Let  $\lambda_i$  be the characteristic roots of  $A$  and  $\mu_i$  those of  $AA'$  (which are in general distinct from the squares of the absolute values of  $\lambda_i$ ). E. T. BROWNE<sup>4</sup> has shown that

$$\mu_{\min} \leq \lambda_i \bar{\lambda}_i \leq \mu_{\max}.$$

From this it follows that

$$1 \leq \left| \frac{\lambda_{\max}}{\lambda_{\min}} \right| \leq \left| \frac{\lambda_{\max}}{\lambda_{\min}} \right|^2 \leq \frac{\mu_{\max}}{\mu_{\min}},$$

which implies the required result.

3. Proof of the theorem in the  $N$  case:

It is known that  $N(A)$  is the square root of the trace of  $AA'$  and therefore equal to  $(\sum \mu_i)^{\frac{1}{2}}$ . The numbers  $\mu_i$  are all positive since  $AA'$  is symmetric and positive definite. Since the characteristic roots of  $A'A$  and  $AA'$  are the

same and since the characteristic roots of the inverse of a matrix are the reciprocals of those of the original matrix, it follows that

$$N(A^{-1}) = (\text{tr} A^{-1}(A^{-1})')^{\frac{1}{2}} = (\text{tr}(A'A)^{-1})^{\frac{1}{2}} = (\sum \mu_i^{-1})^{\frac{1}{2}}.$$

The  $N$ -condition number of  $A$  is therefore

$$\frac{1}{n} (\sum \mu_i)^{\frac{1}{2}} (\sum \mu_i^{-1})^{\frac{1}{2}}.$$

In a similar way it can be shown that the  $N$ -condition number of  $AA'$  is

$$\frac{1}{n} (\sum \mu_i^2)^{\frac{1}{2}} (\sum \mu_i^{-2})^{\frac{1}{2}}.$$

The theorem follows from the inequality

$$\sum \mu_i^2 \sum \mu_i^{-2} \geq \sum \mu_i \sum \mu_i^{-1},$$

which is in fact true for all real and positive numbers. (It is, indeed, true when the first power on the right is replaced by an arbitrary power  $r$  and the second power on the left by a power  $s > r$ .) The proof of the inequality is as follows:

$$\begin{aligned} \sum \mu_i^2 \sum \mu_i^{-2} - \sum \mu_i \sum \mu_i^{-1} &= n + \sum_{i \neq j} \mu_i^2 \mu_j^{-2} - n - \sum_{i \neq j} \mu_i \mu_j^{-1} \\ &= \sum_{i < j} (\mu_i^2 \mu_j^{-2} + \mu_j^2 \mu_i^{-2}) - \sum_{i < j} (\mu_i \mu_j^{-1} + \mu_j \mu_i^{-1}) \\ &= \sum_{i < j} \{ (\mu_i \mu_j^{-1} + \mu_j \mu_i^{-1})(\mu_i \mu_j^{-1} + \mu_j \mu_i^{-1} - 1) - 2 \} \geq 0, \end{aligned}$$

since

$$\mu_i \mu_j^{-1} + \mu_j \mu_i^{-1} \geq 2, \quad \text{and} \quad \mu_i \mu_j^{-1} + \mu_j \mu_i^{-1} - 1 \geq 1.$$

There is equality if and only if

$$\mu_1 = \mu_2 = \dots = \mu_n.$$

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<sup>1</sup> In a course of lectures given at the British Admiralty by himself and D. H. SADLER in 1949.

<sup>2</sup> J. VON NEUMANN & H. H. GOLDSTINE, "Numerical inverting of matrices of high order," Amer. Math. Soc., *Bull.*, v. 53, 1947, p. 1021-1099. (These authors consider symmetric matrices only, but it is reasonable to apply the definition to the general case.)

<sup>3</sup> A. M. TURING, "Rounding-off errors in matrix processes," *Quart. Jn. Mech. Appl. Math.*, v. 1, 1948, p. 287-308.

<sup>4</sup> E. T. BROWNE, "The characteristic equation of a matrix," Amer. Math. Soc., *Bull.*, v. 34, 1928, p. 363-368.

#### BIBLIOGRAPHY Z-XI

1. E. G. ANDREWS, "The Bell Computer, Model VI," *Electrical Engineering*, v. 68, 1949, p. 751-756, 7 figs., 5 tables. 22.2 × 29.5 cm.

Controlled from remote stations, this new digital computer of the relay type reduces punched-tape instructions to a minimum. With novel control features similar to those used in recent automatic dial-telephone developments, this "upper-class" computer possesses six "intelligence levels." Sub-