HOW WAS $F_6$ FACTORED?

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Dedicated to the memory of my friend, D. H. Lehmer

Abstract. In 1880 at the age of 82 Fortuné Landry factored the 20 digit number $F_6 = 2^{64} + 1$. How did he do it? Landry himself never described how he factored $F_6$; however, he did leave enough clues in his work and letters to provide some indication of the ideas with which he was working. In this paper we present a likely reconstruction of Landry's technique.

1. Introduction

On July 12, 1880 Fortuné Landry [8] published the following short announcement:

"I have just factored the number

$$2^{64} + 1 = 18446744073709551617.$$ 

This number is the product of two factors 274177 which is prime, and 67280421310721. I do not currently know if this last factor is a prime."

This remarkable achievement, made even more so when one considers that Landry was 82 at the time, has unfortunately not received the attention that it deserves. There are two reasons for this: first, the factor 274177 is so small that many would tend to regard it as having been found by the simple process of trial division. This was not, however, the way Landry factored $F_6$. This brings us to the second reason that his work has been largely ignored: he never published any account of how he did factor $F_6$. Nevertheless, he did leave several clues as to how he did it, and the purpose of this paper is to use these clues to reconstruct the probable method that he used.

It should be mentioned that according to Lucas [17] both Landry and LeLasseur succeeded later in proving that the larger factor of $F_6$ is a prime. This seems to be the only contribution that LeLasseur made to the problem of factoring $F_6$.

In 1869 Landry [7] published a pamphlet containing the complete factorization of each integer of the form $2^n \pm 1$ for $n \leq 64$, with the exception of four numbers that he was unable to factor: $2^{59} - 1$, $2^{61} - 1$, $(2^{61} + 1)/3$, $2^{64} + 1$. 

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In [15, pp. 238–240] Lucas states that Landry later factored $2^{59} - 1$ and that Landry thought (probably because they were difficult to factor) that the remaining three were prime. In fact, in 1877 Lucas [13] knew that $F_6 = 2^{64} + 1$ is composite, but his technique (see [15, p. 238, p. 313]) did not reveal the factors. He [14] also proved the following

**Theorem.** The prime divisors of $2^{4q} + 1$ are of the form $16h^q + 1$. □

In 1891 Lucas [17] stated that by using this result for simplifying the search for factors of $F_6$, Landry was able to factor $F_6$ after a labor of several months. Landry, however, said very little indeed about how he did this. One of the few explicit references in print that Landry made to his technique is in a letter to Lucas, which, fortunately, Gérardin published in *Sphinx-Oedipe* [9]. Because of its importance in this investigation, we reproduce it in full here. It should be mentioned that all the work reported on in this paper was originally written in French; thus, all quotes are translations of this French into English.

Paris, 7 July 80

Sir

The number $2^{64} + 1$ or $18446744073709551617$ is the product of two factors $67280421310721$ and $274177$, the latter being prime.

Notice of the result has been sent to the Academy today.

Your advice proved to be ever the wiser, considering that the smaller factor is relatively small and would require only about 1000 operations to discover by using the standard method of successive divisions.

I had set up my calculations with the intention of taking the process to its end and thus I would be able to announce a result. There was not much left to do when I unexpectedly came upon the factor $274177$. After such a lengthy process, I could hardly believe it. The procedure I use first considers the largest divisors, less than $\sqrt{N}$ of course, of the number to be factored.

One could very well conclude from this that the other factor is a prime, but besides the possibility that I could have erred during the course of such a lengthy process, in order to facilitate my work, I purposely left out divisors of particular forms which would be easy to deal with separately.

I thank you very much for your interest in this work. This last 14-digit number will allow me to describe the methods that I use. I hope to arrive for the Congress quite early. Could you please let me know where I would be able to obtain the pamphlet where you will mention my work.

My thanks to you.

Landry

According to Lucas's theorem, any prime factor of $F_6$ must be of the form $256k + 1$. Since $274177 = 256 \cdot 1071 + 1$, we see what Landry was talking about in his third paragraph. Also, in [16] Lucas points out that the smallest factor for each of the Fermat numbers $F_n = 2^{2^n} + 1$ of (then) known factorization is the smallest prime of the form $k2^{n+2} + 1$. Thus, he was likely of the belief
that one factor of a composite Fermat number might tend to be small. This, undoubtedly, was the advice that he gave to Landry. It appears, however, that Landry did not take Lucas's advice because he states in his next paragraph that he wanted to set up his technique in order to test for all possibilities. Since $\sqrt{F_6}/256 \approx 1 \cdot 7 \times 10^7$ is so large, trial division just wouldn't do the job. Indeed, in 1867 he [6] had estimated that attempting to prove $F_6$ a prime by trial division (using the then-known form of the factors as $128k + 1$ instead of $256k + 1$) could take up to 3000 years. We will explain certain other features of this letter as they become relevant to our discussion.

2. LANDRY'S EARLY FACTORING IDEAS

In order to get some familiarity with how Landry approached the factoring problem, it is of some value to discuss his work prior to 1869. Apart from some simple tables of primes and factors, Landry's first significant work on factoring is his 1859 pamphlet [5] in which he proved that $2^{31} - 1$ is a prime. This had also been done in 1772 by Euler, but Landry's interest in the problem was in reducing the amount of work to be done.

Euler had realized that any prime factor $p$ of $N = 2^{31} - 1$ must have the form $1 + 31k$. He also was aware that since $(2/p) = 1$, $p$ must be congruent to $\pm1$ modulo 8. Thus, $p = 248k + 1$ or $248k + 63$ for some integer $k$. Since $\sqrt{N}/248 = 186$, this means that after 372 trial divisions (or fewer if one excludes composite values of $248k + 1$ and $248k + 63$), $N$ can be shown prime.

Landry approached this problem as follows. If $N$ is composite, then

$$N = (62x + 1)(62x' + 1),$$

where $x, x' \in \mathbb{Z}$, and with no loss of generality $62x + 1$ is a prime such that $62x + 1 < \sqrt{N}$, i.e., $x < 748$. Since

$$N = 62^2 \cdot 558658 + 62 \cdot 37 + 1,$$

we get

$$x + x' = 62h + 37,$$

$$xx' = 558658 - h$$

for some $h \in \mathbb{Z}^+$. It follows from (2.1) and (2.2) that $2|h$, $h \equiv 3$ or $4 \pmod{5}$, and $h \equiv 1$ or $6 \pmod{9}$. Hence,

$$h = 90h' + k,$$

where $k \in \{24, 28, 64, 78\}$. If we solve (2.1) and (2.2) for $x$ and $x'$, we get

$$2x, 2x' = 62h + 37 \pm \sqrt{62^2h^2 + 4592h - 2233263}.$$  

If we eliminate $x'$ from (2.1) and (2.2), we get

$$h = \frac{558658 - x(37 - x)}{62x + 1} \in \mathbb{Z}^+.$$

Substitute $90h' + k$ for $h$ and we find

$$h' = \frac{558658 - k - x(62k - 137 - x)}{90(62x + 1)} \in \mathbb{Z}_{\geq 0}.$$
Now we may assume that \( x < 748 \); but, if \( x \geq 80 \), then from (2.5) we see that \( h' < 1 \) and therefore \( h' = 0 \). If, however, \( h' = 0 \), then from (2.3)
\[
62^2k^2 + 4592k - 2233263 = t^2 \quad (t \in \mathbb{Z})
\]
for some \( k \in \{24, 28, 64, 78\} \). As this is not the case, we must have \( x < 80 \). In fact, if \( x \geq 60 \), then from (2.5) one has \( h < 2 \). Since
\[
62^2(k + 90)^2 + 4592(k + 90) - 2233263 \neq t^2 \quad (t \in \mathbb{Z})
\]
for any \( k \in \{24, 28, 64, 78\} \), we must have \( x < 60 \).

Consider now \( 62x + 1 \). We know from Euler's observation that \( 62x + 1 \equiv \pm 1 \pmod{8} \). Also, \( 62x + 1 \not\equiv 0 \pmod{3, 5, 7} \). It follows that since \( x < 60 \), we can only have
\[
x \in \{5, 9, 20, 21, 24, 33, 41, 44, 45, 48, 53, 56\}.
\]
However, (2.4) is not satisfied for any of these 12 values of \( x \); hence, \( N \) must be a prime. Notice that with just 12 trial divisions (with the dividend much smaller than \( N \)) and 8 perfect square tests, Landry was able to show that \( N \) is a prime.

Later in 1867 Landry [6] announced that he had discovered a very simple principle which he had used to obtain a number of factorizations of numbers of the form \( 2^n \pm 1 \). He did not reveal this method, however; instead he presented several of his more impressive factorizations. This work was followed by the table in [7]. What was this new method? Landry finally described it in a letter [10] that he wrote to Charles Henry.

Let \( N = ab \), where \( a, b \) are odd and \( a > b \). Put \( x = (a + b)/2 \), \( y = (a - b)/2 \), and we get
\[
(2.6) \quad x^2 - N = y^2.
\]
Determine for \( x = \lceil \sqrt{N} \rceil + 1, \lceil \sqrt{N} \rceil + 2, \ldots \) a value for \( y \) such that (2.6) holds. Certain values of \( x \) can be easily eliminated by making use of moduli \( m \) such that for these values of \( x \) we cannot have
\[
x^2 - N \equiv y^2 \pmod{m}.
\]
For example, if \( m = 5 \) and \( N \equiv 2 \pmod{5} \), then \( x \not\equiv 0, 4 \pmod{5} \). Today, we call such moduli exclusion moduli.

This method of factoring, which is particularly effective if the two factors \( a \) and \( b = N/a \) are close in value, had been discovered many years earlier by Fermat [3]. Landry, however, remained unaware of Fermat's work until Henry, one of the people involved in publishing Fermat's complete works, put it in print in his *Recherches sur les Manuscrits de Fermat*. It was the similarity of Fermat's idea to that of Landry that particularly struck Landry and caused him to write his letter to Henry.

Landry also noted in this letter that if \( N = 2^n \pm 1 \), then for "certain" values of \( n \) the process can be accelerated because we know that
\[
a = 2nu + 1, \quad b = 2nv + 1 \quad (u, v, \in \mathbb{Z}).
\]
Hence, \( x = 1 + ns, \quad y = nt \pmod{s, t \in \mathbb{Z}} \). Although he did not say so, he must have been aware that from (2.6) these latter results allow us to restrict the value of \( x \) to a single residue class modulo \( n^2 \), a result first put in print by Pepin.
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[18]. He went on to say that one could (by using other moduli) accelerate the course of the calculations, but that it is then necessary to follow up each series of possible values of $x$ separately.

Is this the method that Landry used to factor $F_6$? No, for he mentioned that it was found to be insufficient for the larger numbers that he turned to after 1867. In fact, he said that he had to find another method, and it was by this new method that he had recently factored $F_6$. Also, at the time of writing this letter (July or August of 1880) he stated that he was busy writing up this new method.

3. Landry's later factoring ideas

It appears that Landry published most of this new method in the Proceedings of the Congress that he referred to at the end of his letter to Lucas. This annual conference was sponsored by the French Association for the Advancement of Science, and Landry’s paper [11] was presented at the session devoted to mathematics on August 16, 1880. Unfortunately, this discussion of his techniques gives the impression that they are almost trivial; this is certainly well illustrated by the synopsis of the paper given in Dickson [2, p. 371]. In fact, however, if one reads the paper carefully, it is clear that the new method is a further development of the idea presented in [5]. In view of this, we will rearrange the order of topics in this paper and modify them somewhat for our presentation here. Also, in order to aid us in showing what is going on, we use upper case letters to refer to quantities that are either known or easily calculated and lower case letters to represent those quantities that are difficult to evaluate.

Suppose we wish to factor $N = f_1 f_2$, where we know in advance that

$$f_1 = Pn_1 + A, \quad f_2 = Pn_2 + A \quad (n_1, n_2 \geq 1).$$

Landry referred to factors $f_1, f_2$ such that $f_1 \equiv f_2 \pmod{P}$ as being “similar” modulo $P$. By substitution we get

$$Pn_1 n_2 + A(n_1 + n_2) = (N - A^2)/P = PQ + R;$$

hence,

$$A(n_1 + n_2) \equiv R \pmod{P}$$

and

$$n_1 + n_2 = R' + Ph,$$

where $R' \equiv A^{-1}R \pmod{P} \ (0 \leq R' < P)$. Putting (3.2) into (3.1) and using

$$Q' = Q - (AR' - R)/P,$$

we get

$$n_1 n_2 = Q' - Ah.$$

Solving (3.2) and (3.3) for $h$, we find that

$$h = \frac{Q' - n_1(R' - n_1)}{Pn_1 + A} \in \mathbb{Z}^{>0}.$$

With regard to the simultaneous equations (3.2), (3.3), Landry was aware of the following
Theorem. If \( p \) is an odd prime and \( p \nmid PN \), then there are exactly \((p+(N/p))/2\) values of \( h \) \((\text{mod } p)\) such that (3.2) and (3.3) can hold simultaneously modulo \( p \). \( \square \)

Furthermore, if \((N/p) = 1\), we can get a case of similar factors modulo \( p \), i.e., \( n_1 \equiv n_2 \) \((\text{mod } p)\). To a single value of \( h \) \((\text{mod } p)\), there correspond two possible forms \( n_1 \equiv a, n_1 \equiv b \) \((\text{mod } p)\). These forms of \( n_1 \) Landry called the "conjugate" values to that of \( h \) modulo \( p \).

Landry preferred the use of (3.4) to the usual trial division process because the numerator was less than \( N \). Also, we can assume that \( n_1 < \sqrt{N}/p \). Landry states that "when the values of \( N \) become large, it becomes necessary, in order to avoid numerous operations, to resort to forms of \( h \) and \( n \), for moduli 2, 3, 6, 5, ... and to put aside for separate treatment the similar factors of \( N \) which result when \( n_1 \equiv n_2 \) for any of these moduli." That is, we attempt to restrict \( h \) by using certain moduli and (3.2), (3.3) with respect to these moduli. With the exception of the number\(^1\) 6, Landry's moduli were intended to be the small primes.

Suppose \( G \) is a product of certain moduli and \( K \) a possible value of \( h \) modulo \( G \). We have

\[
h = Gh' + K \quad (h' \geq 0, \ 0 \leq K < G),
\]

and

\[
h' = \frac{Q' - n_1(R' - n_1)}{G(Pn_1 + A)} - \frac{K}{G}.
\]

Notice that if \( n_1 \) exceeds some bound \( B \) \((\approx Q'/(PG))\), we must get \( h' < 1 \). Since this forces \( h' = 0 \), we must have

\[
\begin{align*}
\{ & n_1 + n_2 = R' + PK, \\
& n_1n_2 = Q' - AK
\end{align*}
\]

when \( n_1 > B \); but, if (3.5) has no solution in integers, then this possible value of \( h \) \((\text{mod } G)\) can be eliminated. On the other hand, if (3.5) has a solution \((n_1, n_2)\), then \( N = f_1f_2 \), where \( f_1 = Pn_1 + A \), \( f_2 = Pn_2 + A \).

Notice that there are two problems here: (1) find a convenient value of \( G \) in order to get \( B \) small, (2) determine a fast method of resolving (3.5). With respect to the first of these problems, Landry suggested the use of several moduli, but pointed out that "it becomes necessary to adopt for the work a particular setup which allows us to conveniently group together the residues of the forms of \( h \) and then those of the \( n_1 \) values which are conjugate to them." He hoped to show how this could be done "soon", but never did. With respect to the second problem, he said that he "would give a method", but none appears in [11]. Fortunately, he does provide us with his method in a later paper [12]. Indeed, he even states in [12] that he used this idea in his work on factoring \( F_6 \). (Actually, he does not mention \( F_6 \) explicitly, but refers to [1], which is an announcement of his factorization of \( F_6 \).)

The idea is very simple and rather neat. Suppose, for a given pair of integers

\(^1\)Landry singled 6 out as a special modulus in his treatment in [11], but it is not really necessary to do this.
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$P, Q \geq 0$, we want to find positive integers $x', x$ such that $x' \leq x$ and

$$x' + x = P,$$
$$x'x = Q.$$ 

Suppose further that

$$x' = d_010^k + d_110^{k-1} + \cdots + d_k \quad (0 \leq d_i \leq 9, \ d_0 \neq 0),$$

and put $a = d_010^k$. If

$$x_1 = x - a \ (\geq 0),$$
$$x_1' = x' - a \ (\geq 0),$$
$$P_1 = P - 2a \ (\geq 0),$$
$$Q_1 = Q - (P - a)a \ (\geq 0),$$

then

$$x_1 + x_1' = P_1,$$
$$x_1x_1' = Q_1.$$ 

Notice that $a$ is the largest possible multiple of $10^k$ such that the four inequalities above can hold. Also, since $x' = Q/(P - x')$ and $0 < x'/P \leq 1/2$, we get

$$Q/P < x' \leq 2Q/P;$$

thus not many trials are needed to find $a$. Indeed, as Landry notes, we have

$$\frac{P}{P - x'} = 1 + \frac{x'}{x},$$

$$\frac{P_1}{P_1 - x_1'} = \frac{P_1}{x_1} = \frac{x_1 + x_1'}{x_1} = 1 + \frac{x_1'}{x_1} = 1 + \frac{x' - a}{x - a};$$

thus, $P_1/(P_1 - x_1')$ is closer to 1 than $P/(P - x')$. It follows that the values for successive values of $a_n$ are closer to $Q_n/P_n$ as $n$ increases. Thus, the value of $x'$ can be computed very simply, one digit at a time. If there are no possible values for $x, x'$, this will become clear when no value for an $a_n$ can be determined, i.e., no $Q_n = 0$.

Consider the following example:

$$x + x' = 3842 = P,$$
$$x'x = 1330945 = Q.$$ 

Since $[Q/P] = 346$, try $a = 300$.

$$P = 3842 \quad 1330945 = Q$$
$$2a = 600 \quad 1062600 = a(P - a)$$
$$P_1 = P - 2a = 3242 \quad 268345 = Q_1$$
$$2a_1 = 160 \quad 259360 = a_1(P_1 - a_1)$$
$$P_2 = P_1 - 2a_1 = 3082 \quad 15385 = Q_2$$
$$2a_2 = 10 \quad 15385 = a_2(P_2 - a_2)$$
$$P_3 = P_2 - 2a_2 = 3072 \quad 0 = Q_3.$$ 

Hence, $x' = 385$ and $x = 3457$. 

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In the large, then, Landry's factoring idea is the following: (1) Test for all possible values of \( K \) whether or not (3.5) has a solution. If it does, we have a factor of \( N \). If not, (2) test (3.4) for all possible values of \( n_1 \leq B \). We either find a factor of \( N \), or \( N \) is a prime.

4. A POSSIBLE SETUP FOR FACTORING \( F_6 \)

Since Landry seems never to have published anything more on the factorization problem than what we have mentioned above, the remainder of this paper must be somewhat speculative. We can only describe a technique for factoring \( F_6 \) which is consistent with all of Landry's work that has been published. Whatever method he did use would most likely not differ significantly from what we shall now describe.

We can (in view of Lucas [17] and Landry [9]) use \( P = 256 \) and \( A = 1 \); hence, \( R = R' = 0 \), \( Q = Q' = 2^{48} \), and we get

\[
\begin{align*}
(4.1) & \quad n_1 + n_2 = 256h, \\
(4.2) & \quad n_1 n_2 = 2^{48} - h, \\
(4.3) & \quad h = \frac{2^{48} + n_1^2}{256n_1 + 1} \in \mathbb{Z}.
\end{align*}
\]

Also, in order for (4.1) and (4.2) to hold, we require that

\[
(4.4) \quad 2^{14}t^2 + h - 2^{48} = t^2 \quad (t \in \mathbb{Z}).
\]

If we use

\[
G = 2^4 \cdot 3 \cdot 5 \cdot 7 \cdot 11 \cdot 13 \cdot 17 \cdot 19 = 77597520,
\]
we find from (4.3) that if \( n_1 > 14170 \), then \( h' = 0 \) for \( h = Gh' + K \) \((K > 0)\). If \( h' = 0 \), we get \( h = K \) and

\[
(4.5) \quad 2^{14}K^2 + K - 2^{48} = t^2 \quad (t \in \mathbb{Z})
\]

from (4.4). We also get

\[
\begin{align*}
(4.6) & \quad n_1 + n_2 = 256K, \\
(4.7) & \quad n_1 n_2 = 2^{48} - K
\end{align*}
\]

from (4.1) and (4.2).

We next investigate the possible values of \( h \). By using (4.4), we see that the only possible values of \( h \) are given by

\[
\begin{align*}
h & \equiv 0^*, 4^*, 1, 9 \pmod{16}, \\
h & \equiv 1 \pmod{3}, \\
h & \equiv 0, 1 \pmod{5}, \\
h & \equiv 1, 4, 6 \pmod{7}, \\
h & \equiv 1, 3, 4, 9, 10 \pmod{11}, \\
h & \equiv 0, 1, 2, 3, 5^*, 8, 11^* \pmod{13}, \\
h & \equiv 0, 1, 2, 3^*, 10^*, 11, 12, 13, 15 \pmod{17}, \\
h & \equiv 5, 7, 8, 9, 11, 13, 14, 15, 17 \pmod{19}.
\end{align*}
\]
Residues with an asterisk (*) lead to similar factors with respect to the given modulus or in the case of 16 a divisor of that modulus. In his letter to Lucas, Landry said that he left certain forms aside; in view of [11], these would likely be those types of forms. This is simply because for these forms the corresponding modulus (or a divisor of it) can be absorbed into the $P$ value; and, as this increases the $P$ value, the resulting calculations can be accomplished more quickly. Thus, if we eliminate these forms from consideration, we find that the total number of possible values of $h \mod G$ is

$$2 \times 1 \times 2 \times 3 \times 5 \times 5 \times 7 \times 9 = 18900.$$ 

We must now deal with the problem of the values of $n_1$ which are conjugate to the $h$-values. We first point out that if $h \equiv 1 \pmod{16}$, then $n_1 \equiv 1, 7 \pmod{8}$; if $h \equiv 9 \pmod{16}$, then $n_1 \equiv 3, 5 \pmod{8}$. The other conjugate values modulo $p$, where $p | G$, are given in the tables below.

<table>
<thead>
<tr>
<th>$p$</th>
<th>$h$</th>
<th>$n_1$</th>
</tr>
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<tbody>
<tr>
<td>3</td>
<td>1</td>
<td>0, 1</td>
</tr>
<tr>
<td>7</td>
<td>1</td>
<td>4, 0</td>
</tr>
<tr>
<td>11</td>
<td>1</td>
<td>12, 1</td>
</tr>
<tr>
<td>13</td>
<td>1</td>
<td>5, 8</td>
</tr>
<tr>
<td>17</td>
<td>1</td>
<td>4, 13</td>
</tr>
<tr>
<td>19</td>
<td>1</td>
<td>6, 1, 12, 13, 16, 18, 9, 15, 0, 4, 8, 14, 5, 7, 10, 11, 3, 17</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>$p$</th>
<th>$h$</th>
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<tbody>
<tr>
<td>5</td>
<td>0</td>
<td>2</td>
</tr>
<tr>
<td>3</td>
<td>1</td>
<td>6, 4, 8, 9, 11, 13, 14, 15, 17</td>
</tr>
</tbody>
</table>

Notice that if $n_1 \leq 14170$, we cannot have $n_1 \equiv 0 \pmod{2}$, $n_1 \equiv 2 \pmod{3}$, $n_1 \equiv 4 \pmod{5}$, $n_1 \equiv 5 \pmod{7}$, $n_1 \equiv 7 \pmod{11}$, $n_1 \equiv 3, 4, 10 \pmod{13}$, $n_1 \equiv 5, 10, 16 \pmod{17}$, or $n_1 \equiv 2 \pmod{19}$. Up to 14170 we would expect, then, to find about

$$14170 \times \frac{1}{2} \times \frac{2}{3} \times \frac{4}{5} \times \frac{6}{7} \times \frac{10}{11} \times \frac{14}{13} \times \frac{18}{19} \approx 1767$$

values of $n_1$ which could be admissible. In fact there are exactly 1773 such values.

We now turn to the algorithm for factoring $F_6$.

5. The algorithm

One of the difficulties one encounters on using Landry’s ideas to factor $F_6$ is keeping track of the large number (18900 here) of possible values of $K$. Landry may have just done this in a very ordinary (and tedious) fashion, but it is possible that he might have done this as we describe here. We first note that if $G = G_1G_2$, where $\gcd(G_1, G_2) = 1$ and

$$K \equiv K_1 \pmod{G_1}, \quad K \equiv K_2 \pmod{G_2},$$
then
\[ K \equiv W_1 + W_2 \pmod{G}, \]
where
\[ W_1 \equiv G_2 Z_2 K_1, \quad W_2 \equiv G_1 Z_1 K_2 \pmod{G} \]
and
\[ G_1 Z_1 \equiv 1 \pmod{G_2}, \quad G_2 Z_2 \equiv 1 \pmod{G_1}. \]

This is just a simple case of what today is called the Chinese Remainder Theorem. In its more general form it can be found in Gauss [4, art. 36]. Since Landry had read [4] (see [6]), it is entirely possible that he was aware of this and that he utilized it in his calculations. We will show how below.

We also know that Landry was aware of the idea of using exclusion moduli to eliminate possible candidates as solutions for certain equations. Thus he might very well have used this technique to eliminate values of \( K \) which fail to satisfy (4.5). Suppose that \( \mathcal{E} = \{23, 29, \ldots\} \) (no prime in \( \mathcal{E} \) can divide \( G \)) is the set of prime exclusion moduli which Landry used. Note that |\( \mathcal{E} \)| need not be very large and could have even been 0.

A simple idea for keeping track of the \( K \)-values (and for eliminating many of them) can now be implemented by making use of three relatively short lists: \( \mathcal{L}_1, \mathcal{L}_2, \mathcal{L}_3 \). We point out that there are 135 possible values of \( K \pmod{G_1} \) when \( G_1 = 1463 = 7 \cdot 11 \cdot 19 \). For each of these compute a list made of the corresponding \( W_1 \)-values given above, together with a sequence of the values of each \( W_1 \) modulo the exclusion moduli in \( \mathcal{E} \). There are 140 possible values of \( K \pmod{G_2} \) when \( G_2 = 53040 = 16 \cdot 3 \cdot 5 \cdot 13 \cdot 17 \). For each of these compute a list of \( \mathcal{L}_2 \) made up of the corresponding \( W_2 \)-values, together with a sequence of the values of each \( W_2 \) modulo the elements in \( \mathcal{E} \). Finally, compute a list \( \mathcal{L}_3 \) of all values of \( Y \pmod{E} \) such that
\[
\left( \frac{2^{14} Y^2 + Y - 2^{48}}{E} \right) = -1
\]
for \( E \in \mathcal{E} \).

We now have the following

Algorithm.

1. To each \( W_1 \in \mathcal{L}_1 \) add every element \( W_2 \in \mathcal{L}_2 \) (and reduce \( \pmod{G} \)). Check if the resulting value of \( K \) can be excluded from satisfying (4.5) by using the easily computed value of \( K \pmod{E} \) for \( E \in \mathcal{E} \) and \( \mathcal{L}_3 \). If this does not exclude the \( K \)-value, use the method of [12] to solve the simultaneous equations (4.6), (4.7).

2. Test if any of the 1773 possible values for \( n_1 \leq 14170 \) is such that (4.3) holds.

Even if |\( \mathcal{E} \)| = 0, this algorithm could probably have been executed by hand in the several months that Lucas said that Landry required. However, it is useful to see how effective the use of exclusion moduli is in this problem. The total number of \( K \)-values here is 18900. In the table below we give the number
of $K$-values remaining in this problem as various exclusion moduli are used.

<table>
<thead>
<tr>
<th>$\mathcal{E}$</th>
<th>new modulus</th>
<th>number of $K$-values remaining</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>—</td>
<td>18900</td>
</tr>
<tr>
<td>1</td>
<td>23</td>
<td>9041</td>
</tr>
<tr>
<td>2</td>
<td>29</td>
<td>4656</td>
</tr>
<tr>
<td>3</td>
<td>31</td>
<td>2259</td>
</tr>
<tr>
<td>4</td>
<td>37</td>
<td>1126</td>
</tr>
<tr>
<td>5</td>
<td>41</td>
<td>542</td>
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<td>8</td>
<td>53</td>
<td>67</td>
</tr>
<tr>
<td>9</td>
<td>59</td>
<td>37</td>
</tr>
<tr>
<td>10</td>
<td>61</td>
<td>19</td>
</tr>
<tr>
<td>11</td>
<td>67</td>
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<td>12</td>
<td>71</td>
<td>6</td>
</tr>
<tr>
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<td>2</td>
</tr>
<tr>
<td>14</td>
<td>79</td>
<td>1</td>
</tr>
</tbody>
</table>

In view of this rapid rate of decrease (a decrease rate with which Landry would have to have been familiar after doing the work reported in [6] and [7]), it is difficult to accept that he would not have made use of at least a few exclusion moduli in order to lessen his work load.

Note that this algorithm would find the factor 274177 toward the end of all the work, as most of the work is done in Step (1) of the algorithm and $n_1 = 1071$ is the 137th value in the list (in ascending order) of the 1773 possible values for $n_1 \leq 14170$ in Step (2).

In summary, then, Landry seems to have been most unlucky in his attempts to factor $F_6$. He first tried to use Fermat's difference of squares method; but, as the factors of $F_6$ are not at all close in value, he failed. This failure seems to have left him for awhile with the impression that $F_6$ is prime. When Lucas later stated that it is composite, Landry decided to try again. As simple trial division would (possibly) be very time-consuming, he elected (contrary to Lucas's advice, it appears) to use a more sophisticated technique, which grew out of his early primality testing ideas. This method would, after a large but manageable amount of labor, ultimately yield the factorization. When, after a great deal of work, Landry finally discovered the factorization, one of the factors turned out to be so small that it could have been found considerably more quickly by the simple trial division process. It must, however, be emphasized that his ingenious method would still have worked with about the same amount of effort even if the small factor had been much larger. Unfortunately for Landry, it was not.

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