

STATISTICS FOR SPECIAL q, t -KOSTKA POLYNOMIALS

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ABSTRACT. Kirillov and Reshetikhin introduced rigged configurations as a new way to calculate the entries $K_{\lambda\mu}(t)$ of the Kostka matrix. Macdonald defined the two-parameter Kostka matrix whose entries $K_{\lambda\mu}(q, t)$ generalize $K_{\lambda\mu}(t)$. We use rigged configurations and a formula of Stembridge to provide a combinatorial interpretation of $K_{\lambda\mu}(q, t)$ in the case where μ is a partition with no more than two columns. In particular, we show that in this case, $K_{\lambda\mu}(q, t)$ has nonnegative coefficients.

1. INTRODUCTION

In [Mac2], Macdonald defined a basis $P_\lambda(q, t)$ of the ring of symmetric functions. Hall–Littlewood symmetric functions, Jack polynomials, Schur functions, and zonal polynomials are all either limiting or special cases of the $P_\lambda(q, t)$'s. He also defined a transition matrix, whose entries are denoted $K_{\lambda\mu}(q, t)$, between a renormalized version of the $P_\lambda(q, t)$'s and another basis S_λ of the ring of symmetric functions. $(K_{\lambda\mu}(q, t))_{\lambda, \mu \vdash n}$ generalizes the Kostka matrix $(K_{\lambda\mu})_{\lambda, \mu \vdash n}$.

Macdonald conjectured that the entries in the two-parameter Kostka matrix are polynomials in q and t with nonnegative integer coefficients. All that is known a priori is that the entries are rational functions of q and t . Garsia and Haiman have constructed, for each partition μ of n , a finite-dimensional bigraded S_n -module whose irreducible multiplicities they conjecture to be rescaled versions of the entries $K_{\lambda\mu}(q, t)$. In [GH], they give several constructions of S_n -modules conjectured to have this property, together with an announcement of the special cases for which they can prove their conjecture. The main special cases correspond to the entries $K_{\lambda\mu}(q, t)$ in which μ is either a hook, or has at most two rows or two columns. Their results do not provide any explicit combinatorial interpretation of the entries.

In the paper [Ste], Stembridge gave a direct proof of the hook case of Macdonald's conjecture, and gave a formula for the polynomial in the two-column

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case, which proves the entries are polynomials. His formula is

$$(1) \quad K_{\lambda_{2r}1^{n-2r}}(q, t) = \sum_{s=0}^r q^{r-s} (t^{n-r}q; t^{-1})_s \begin{bmatrix} r \\ s \end{bmatrix}_t K_{\lambda_{2s}1^{n-2s}}(t)$$

where $K_{\lambda_{2s}1^{n-2s}}(t)$ is the Kostka (charge) polynomial [Mac1] and $|\lambda| = n$. We use (1) to show that there are statistics c_r and cut_r defined on the set \mathcal{M}_0^0 of Kirillov and Reshetikhin rigged configurations which correspond to standard Young tableaux of shape λ such that

Theorem 1.1.

$$(2) \quad K_{\lambda_{2r}1^{n-2r}}(q, t) = \sum_{(\alpha(0), L) \in \mathcal{M}_0^0} q^{cut_r(\alpha(0), L)} t^{c_r(\alpha(0), L)}.$$

Theorem 1.1 proves that $K_{\lambda_{2r}1^{n-2r}}(q, t)$ has nonnegative coefficients.

This paper is divided into five sections. In Section 2, we explain the necessary Kirillov and Reshetikhin material and introduce notation. In Section 3 we rewrite $K_{\lambda_{2r}1^{n-2r}}(q, t)$ as a sum of “difference” polynomials. In Section 4 we show the difference polynomials are nonnegative by showing they are generating functions for sets \mathcal{M}_m^d of rigged configurations. In Section 5 we finish proving Theorem 1.1.

2. KIRILLOV AND RESHETIKHIN’S RIGGED CONFIGURATIONS

Kirillov and Reshetikhin [KR1, KR2] introduced rigged configurations as a new way to calculate $K_{\lambda\mu}(t)$ for any pair of partitions λ and μ of n . Fix n , a positive integer, and λ , a partition of n . Let $\alpha = (\alpha^1, \alpha^2, \dots, \alpha^x)$ be a sequence of partitions such that $|\alpha^i| = \lambda_{i+1} + \lambda_{i+2} + \dots$. For any such sequence, if μ is any partition of n , let $\alpha(\mu)$ be the sequence of partitions $(\mu' = \alpha^0, \alpha^1, \dots, \alpha^x)$, and if m is a nonnegative integer, let $\alpha(m) = ((n - m, m), \alpha^1, \alpha^2, \dots, \alpha^x)$. $\alpha(\mu)$ is called a configuration.

Define, for $k \geq 1$,

$$(3) \quad P_l^k(\alpha(\mu)) = \sum_{i=1}^l (\alpha_i^{k-1} - 2\alpha_i^k + \alpha_i^{k+1}).$$

A rigged configuration is a pair $(\alpha(\mu), L)$, where L labels the columns of the partitions in α . In particular,

- (1) $0 \leq L_i^k \leq P_l^k(\alpha(\mu))$ for $k \geq 1$ and $1 \leq i \leq \alpha_i^k$ where the length of column i of α^k is l ,
- (2) If column i of α^k has the same length as column $i + 1$, then $L_i^k \leq L_{i+1}^k$.

Not all sequences $\alpha(\mu)$ will have labels. Kirillov and Reshetikhin call a sequence $\alpha(\mu)$ a μ -admissible λ configuration when there is at least one labelling function L , that is, when $P_i^k(\alpha(\mu)) \geq 0$. They have defined a bijection between μ -admissible λ rigged configurations and column strict tableaux of shape λ and content μ .

Further, Kirillov and Reshetikhin define the charge of a rigged configuration $c((\alpha(\mu), L))$. Let

$$c(\alpha(\mu)) = n(\mu) - \sum_{i \geq 1} \mu'_i \alpha_i^1 + \sum_{k, i \geq 1} \alpha_i^k (\alpha_i^k - \alpha_i^{k+1}),$$

where $n(\mu) = \sum (i - 1)\mu_i$. Then

$$c((\alpha(\mu), L)) = c(\alpha(\mu)) + \sum_{i, k \geq 1} L_i^k.$$

The Kirillov and Reshetikhin theorem is now

$$K_{\lambda\mu}(t) = \sum_{(\alpha(\mu), L)} t^{c((\alpha(\mu), L))}$$

where the sum is over all μ -admissible λ rigged configurations.

In this paper, $\mu = (2^m 1^{n-2m})$, so that $\alpha(\mu) = \alpha(m)$. We need several properties which are peculiar to this case.

(1) Let

$$(4) \quad c'(\alpha) = \sum_{k, i \geq 1} \alpha_i^k (\alpha_i^k - \alpha_i^{k+1}).$$

Then

$$(5) \quad c((\alpha(m))) = \binom{n-m}{2} + \binom{m}{2} - (n-m)\alpha_1^1 - m\alpha_2^1 + c'(\alpha).$$

(2) By the definition of $P_i^k(\alpha(m))$

$$(6) \quad P_i^k(\alpha(m+1)) = \begin{cases} P_i^k(\alpha(m)) - 1 & \text{if } i = k = 1, \\ P_i^k(\alpha(m)) & \text{otherwise.} \end{cases}$$

Finally, let $\mathcal{M}_m^0 = \{(\alpha(m), L) | \alpha(m) \text{ is a } (n-m, m)\text{-admissible } \lambda \text{ configuration and } L \text{ is a label}\}$. Call \mathcal{M}_0^0 the set of standard rigged configurations.

3. DIFFERENCE POLYNOMIALS

In this section, we rewrite Stembridge's formula (1), changing it to (8). Equation (8) is crucial because we will show the polynomials $M_{r-k}^k(t)$ are generating functions for sets of rigged configurations.

Lemma 3.1. *The coefficient of q^k in $K_{\lambda 2^r 1^{n-2r}}(q, t)$ is*

$$(7) \quad \sum_{s=0}^r (-1)^{k-(r-s)} t^{(k-r+s)(n-r-s+1) + \binom{k-(r-s)}{2}} \begin{bmatrix} s \\ r-k \end{bmatrix}_t \begin{bmatrix} r \\ s \end{bmatrix}_t K_{\lambda 2^s 1^{n-2s}}(t).$$

Proof. Lemma 3.1 is a consequence of the q -binomial theorem [And, 3.3.6]. \square

Definition 3.1. *Define the polynomials $M_m^d(t)$ recursively by $M_m^0(t) = K_{\lambda 2^m 1^{n-2m}}(t)$ and $M_m^{d+1}(t) = M_m^d(t) - t^{n-2m-(d+1)} M_{m+1}^d(t)$.*

Lemma 3.2. *The coefficient of q^k in $K_{\lambda 2^r 1^{n-2r}}(q, t)$ is $\begin{bmatrix} r \\ k \end{bmatrix}_t M_{r-k}^k(t)$, so that*

$$(8) \quad K_{\lambda 2^r 1^{n-2r}}(q, t) = \sum_{k=0}^r \begin{bmatrix} r \\ k \end{bmatrix}_t M_{r-k}^k(t) q^k.$$

Proof. We use induction on $r + k$ and (7) to show that

$$M_{r-(k+1)}^{k+1}(t) = \begin{bmatrix} r \\ k+1 \end{bmatrix}_t^{-1} (\text{coeff. of } q^{k+1} \text{ in } K_{\lambda 2^r 1^{n-2r}}(q, t)).$$

The lemma is true if $r = 0$ or $k = 0$. Assume

$$M_{b-a}^a = \begin{bmatrix} b \\ a \end{bmatrix}_t^{-1} (\text{coeff. of } q^a \text{ in } K_{\lambda 2^b 1^{n-2b}}(q, t))$$

if a and b are nonnegative integers such that $a + b \leq r + k$.

$$\begin{aligned} M_{r-(k+1)}^{k+1}(t) &= M_{(r-1)-k}^{k+1}(t) \\ &= M_{(r-1)-k}^k(t) - t^{n-2(r-1-k)-(k+1)} M_{(r-1)-k+1}^k(t) \\ &= \begin{bmatrix} r-1 \\ k \end{bmatrix}_t^{-1} (\text{coeff. of } q^k \text{ in } K_{\lambda 2^{(r-1)} 1^{n-2(r-1)}}(q, t)) \\ &\quad - t^{n-2(r-1-k)-(k+1)} \begin{bmatrix} r \\ k \end{bmatrix}_t^{-1} (\text{coeff. of } q^k \text{ in } K_{\lambda 2^r 1^{n-2r}}(q, t)) \\ &= \begin{bmatrix} r-1 \\ k \end{bmatrix}_t^{-1} \sum_{s=0}^{r-1} (-1)^{k-(r-1-s)} t^{(k-(r-1)+s)(n-(r-1)-s+1)+(k-\binom{r-1}{2}-s)} \\ &\quad \times \begin{bmatrix} s \\ r-1-k \end{bmatrix}_t \begin{bmatrix} r-1 \\ s \end{bmatrix}_t K_{\lambda 2^s 1^{n-2s}}(t) \\ &\quad - t^{n-2(r-1-k)-(k+1)} \begin{bmatrix} r \\ k \end{bmatrix}_t^{-1} \sum_{s=0}^r (-1)^{k-(r-s)} t^{(k-r+s)(n-r-s+1)+(k-\binom{r-s}{2})} \\ &\quad \times \begin{bmatrix} s \\ r-k \end{bmatrix}_t \begin{bmatrix} r \\ s \end{bmatrix}_t K_{\lambda 2^s 1^{n-2s}}(t). \end{aligned}$$

We will finish proving the lemma assuming $0 \leq s \leq r - 1$. The case $s = r$ is a degenerate special case of what follows. Now we need to show the coefficient of $K_{\lambda 2^s 1^{n-2s}}(t)$, $0 \leq s \leq r - 1$, in the last expression in this string of equalities is equal to

$$(-1)^{k+1-r+s} t^{(k+1-r+s)(n-r-s+1)+(k+1-\binom{r-s}{2})} \begin{bmatrix} r \\ k+1 \end{bmatrix}_t^{-1} \begin{bmatrix} s \\ r-1-k \end{bmatrix}_t \begin{bmatrix} r \\ s \end{bmatrix}_t,$$

which is the coefficient of $K_{\lambda 2^s 1^{n-2s}}(t)$ in $\begin{bmatrix} r \\ k+1 \end{bmatrix}_t^{-1} (\text{coeff. of } q^{k+1} \text{ in } K_{\lambda 2^r 1^{n-2r}}(q, t))$.

The coefficient of $K_{\lambda^{2s}1^{n-2s}}(t)$ in the last expression in the string of equalities is equal to

$$\begin{aligned} & (-1)^{k+1-r+s} t^{(k+1-r+s)(n-r-s+1) + \binom{k+1-(r-s)}{2}} \\ & \times \left(\left[\begin{matrix} r-1 \\ k \end{matrix} \right]_t^{-1} t^{k+1-r+s} \left[\begin{matrix} s \\ r-1-k \end{matrix} \right]_t \left[\begin{matrix} r-1 \\ s \end{matrix} \right]_t \right. \\ & \left. + t^{(n-2(r-1-k)-(k+1))-(n-r-s+1)-(k-r+s)} \left[\begin{matrix} r \\ k \end{matrix} \right]_t^{-1} \left[\begin{matrix} s \\ r-k \end{matrix} \right]_t \left[\begin{matrix} r \\ s \end{matrix} \right]_t \right). \end{aligned}$$

Using the q -factorial definition of the q -binomial coefficient, this turns into

$$\begin{aligned} & (-1)^{k+1-r+s} t^{(k+1-r+s)(n-r-s+1) + \binom{k+1-(r-s)}{2}} \left[\begin{matrix} r \\ k+1 \end{matrix} \right]_t^{-1} \left[\begin{matrix} s \\ r-1-k \end{matrix} \right]_t \left[\begin{matrix} r \\ s \end{matrix} \right]_t \\ & \times \left(\frac{r_t}{(k+1)_t} t^{k+1-r+s} \times \frac{(r-s)_t}{r_t} + \frac{(r-k)_t}{(k+1)_t} \times \frac{(s-r+k+1)_t}{(r-k)_t} \right), \end{aligned}$$

where $n_t = 1 - t^n$. The quantity in the brackets boils down to one, so we are done. \square

4. DIFFERENCE POLYNOMIALS ARE NONNEGATIVE

In this section, we show that the polynomials $M_m^k(t)$, defined in the last section, are the generating functions for sets of rigged configurations, thus showing that $K_{\lambda^{2r}1^{n-2r}}(q, t)$ has nonnegative coefficients.

Definition 4.1. Let $\mathcal{M}_m^k = \{(\alpha(m), L) \in \mathcal{M}_m^0 \mid L_{\alpha_2^1+1}^1 = L_{\alpha_2^1+2}^1 = \dots = L_{\alpha_2^1+k}^1 = 0\}$. Note that if $\alpha_1^1 - \alpha_2^1 < k$, then $(\alpha(m), L) \notin \mathcal{M}_m^k$ for any L .

Lemma 4.1. The generating function for the set \mathcal{M}_m^k is $M_m^k(t)$; that is,

$$\sum_{(\alpha(m), L) \in \mathcal{M}_m^k} t^{c((\alpha(m), L))} = M_m^k(t).$$

Proof. The proof is by induction. The lemma is true if $k = 0$, by the original Kirillov and Reshetikhin result. The definition of $M_m^{k+1}(t)$ is

$$M_m^k(t) - t^{n-2m-(k+1)} M_{m+1}^k(t).$$

In order to prove the lemma we need an injection $\phi_k : \mathcal{M}_{m+1}^k \rightarrow \mathcal{M}_m^k$ such that

- (1) $c(\phi_k((\alpha(m+1), L))) = c((\alpha(m+1), L)) + n - 2m - (k+1)$ and
- (2) $(\alpha(m), L) \in \mathcal{M}_m^k$ is not in the image of ϕ_k if and only if $(\alpha(m), L) \in \mathcal{M}_m^{k+1}$.

Let $\phi_k(\alpha(m+1), L) = (\alpha(m), \hat{L})$, where

$$\hat{L}_j^i = \begin{cases} L_j^i + 1 & \text{if } i = 1 \text{ and } j \geq \alpha_2^1 + k + 1, \\ L_j^i & \text{otherwise.} \end{cases}$$

Please note $(\alpha(m), \hat{L}) \in \mathcal{M}_m^k$ by (6) and also that (2) above is satisfied.

To see that (1) above is satisfied,

$$c((\alpha(m+1), L)) = \binom{n-(m+1)}{2} + \binom{m+1}{2} - (n-(m+1))\alpha_1^1 - (m+1)\alpha_2^1 + c'(\alpha) + \sum_{i,j \geq 1} L_j^i$$

and

$$\begin{aligned} c(\phi_k(\alpha(m+1), L)) &= c(\alpha(m), \hat{L}) \\ &= \binom{n-m}{2} + \binom{m}{2} - (n-m)\alpha_1^1 - m\alpha_2^1 + c'(\alpha) + \sum_{i,j \geq 1} \hat{L}_j^i \\ &= \binom{n-m}{2} + \binom{m}{2} - (n-m)\alpha_1^1 - m\alpha_2^1 + c'(\alpha) + \sum_{i,j \geq 1} L_j^i + (\alpha_1^1 - (\alpha_2^1 + k)) \\ &= \binom{n-(m+1)}{2} + n-m + \binom{m+1}{2} - (m+1) - (n-(m+1))\alpha_1^1 \\ &\quad - (m+1)\alpha_2^1 + \sum_{i,j \geq 1} L_j^i - k \\ &= c(\alpha(m+1), L) + n - 2m - k - 1. \quad \square \end{aligned}$$

The following formula is a consequence of Lemma 3.1 and Lemma 3.2:

$$M_m^k(t) = t^{\binom{n-m}{2} + \binom{m}{2}} \sum_{j=0}^k (-1)^j t^{-\binom{j}{2} - \binom{n-(m+j)}{2} - \binom{m+j}{2}} \begin{bmatrix} k \\ j \end{bmatrix}_{t^{-1}} K_{\lambda_{2^{m+j} 1^{n-2(m+j)}}}(t).$$

Lemma 4.1 therefore has the following corollary.

Corollary 4.1. *The polynomial*

$$t^{\binom{n-m}{2} + \binom{m}{2}} \sum_{j=0}^k (-1)^j t^{-\binom{j}{2} - \binom{n-(m+j)}{2} - \binom{m+j}{2}} \begin{bmatrix} k \\ j \end{bmatrix}_{t^{-1}} K_{\lambda_{2^{m+j} 1^{n-2(m+j)}}}(t)$$

has nonnegative coefficients.

Let $\tilde{K}_{\lambda\mu}(t)$ be the cocharge polynomial; that is, $\tilde{K}_{\lambda\mu}(t) = t^{n(\mu)} K_{\lambda\mu}(t^{-1})$, where $n(\mu) = \sum (i-1)\mu_i$. Then we rewrite the sum in Corollary 4.1 in terms of cocharge polynomials and Lemma 4.1 has a second corollary.

Corollary 4.2. *The polynomial*

$$\sum_{j=0}^k (-1)^j t^{\binom{j}{2}} \begin{bmatrix} k \\ j \end{bmatrix}_t \tilde{K}_{\lambda_{2^{m+j} 1^{n-2(m+j)}}}(t)$$

has nonnegative coefficients.

5. PROOF OF THEOREM 1.1

In this section we finish proving the main theorem of this paper. We define the statistics c_r and cut_r on \mathcal{M}_0^0 , the standard rigged configurations. Then we construct a surjection from \mathcal{M}_0^0 onto $\bigcup_{k=0}^r \mathcal{M}_{r-k}^k$ which respects the statistics.

Definition 5.1. Let $(\alpha(0), L) \in \mathcal{M}_0^0$ and let $a_i = L_{\alpha_2^1+i}^1$, that is, a_i is the label of the i th column of length 1 in $(\alpha(0), L)$. Let $a_{\alpha_1^1-\alpha_2^1+1} = \infty$ and let $a_0 = 0$. We define $cut_r(\alpha(0), L)$ to be the least j , $0 \leq j \leq r$, such that $j + 1 + a_{j+1} > r$. If there is no such j , we let $cut_r(\alpha(0), L) = r$. Note that if $\alpha_1^1 - \alpha_2^1 \leq r$, then $cut_r(\alpha(0), L) \leq \alpha_1^1 - \alpha_2^1$.

Definition 5.2. Let $(\alpha(0), L) \in \mathcal{M}_0^0$. Let $k = cut_r(\alpha(0), L)$. Then we define $c_r((\alpha(0), L)) = c((\alpha(0), L)) - (n - r)(r - k)$.

Definition 5.3. Let $k = cut_r(\alpha(0), L)$. Define $\Psi((\alpha(0), L)) = (\alpha(r - k), \tilde{L})$, where

$$\tilde{L}_j^i = \begin{cases} 0, & i = 1 \text{ and } \alpha_2^1 + 1 \leq j \leq \alpha_2^1 + k, \\ L_j^i - (r - k), & i = 1 \text{ and } \alpha_2^1 + k + 1 \leq j \leq \alpha_1^1, \\ L_j^i, & \text{otherwise.} \end{cases}$$

Since $P_1^1(\alpha(m + 1)) = P_1^1(\alpha(m)) - 1$, if $L_j^1 \leq P_1^1(\alpha(0))$, then $\tilde{L}_j^1 \leq P_1^1(\alpha(r - k))$, so that $\Psi(\alpha(0), L) \in \mathcal{M}_{r-k}^0$. Since $\tilde{L}_j^1 = 0$ for $\alpha_2^1 + 1 \leq j \leq \alpha_2^1 + k$, $\Psi((\alpha(0), L)) \in \mathcal{M}_{r-k}^k$. Also note that the image of Ψ is $\bigcup_{j=0}^r \mathcal{M}_{r-j}^j$.

Lemma 5.1. Let $(\alpha(r - k), \tilde{L}) \in \mathcal{M}_{r-k}^k$. Then

$$\sum_{(\alpha(0), L) \in \Psi^{-1}((\alpha(r-k), \tilde{L}))} t^{c_r((\alpha(0), L))} = \begin{bmatrix} r \\ k \end{bmatrix}_t t^{c((\alpha(r-k), \tilde{L}))}.$$

Proof. First we show that

$$c(\Psi((\alpha(0), L))) = c(\alpha(0), L) - \left(\sum_{i=\alpha_2^1+1}^{\alpha_2^1+k} L_i^1 \right) - (n - r)(r - k)$$

so that

$$\begin{aligned} (9) \quad c_r(\alpha(0), L) &= \sum_{i=\alpha_2^1+1}^{\alpha_2^1+k} L_i^1 + c(\Psi(\alpha(0), L)) \\ &= \sum_{i=\alpha_2^1+1}^{\alpha_2^1+k} L_i^1 + c(\alpha(r - k), \tilde{L}), \end{aligned}$$

$$\begin{aligned}
 &c(\Psi((\alpha(0), L))) \\
 &= \sum_{i=k+1+\alpha_2^1}^{\alpha_1^1} L_i^1 - (\alpha_1^1 - \alpha_2^1 - k)(r - k) + \sum_{i>1, j \geq 1} L_j^i + c(\alpha(r - k)) \\
 &= \sum_{i \geq 1, j \geq 1} L_j^i - \sum_{j=\alpha_2^1+1}^{\alpha_2^1+k} L_j^1 - (\alpha_1^1 - \alpha_2^1 - k)(r - k) \\
 &\quad + \left[\binom{n - (r - k)}{2} + \binom{r - k}{2} - (n - (r - k))\alpha_1^1 - (r - k)\alpha_2^1 + c'(\alpha) \right] \text{ by (5)} \\
 &= \sum_{j, i \geq 1} L_j^i - \sum_{j=\alpha_2^1+1}^{\alpha_2^1+k} L_j^1 + \binom{n - (r - k)}{2} + \binom{r - k}{2} + n\alpha_1^1 + c'(\alpha) + k(r - k) \\
 &= c(\alpha(0), L) - \binom{n}{2} - \sum_{j=\alpha_2^1+1}^{\alpha_2^1+k} L_j^1 + \binom{n - (r - k)}{2} + \binom{r - k}{2} + k(r - k) \\
 &= c(\alpha(0), L) - \sum_{j=\alpha_2^1+1}^{\alpha_2^1+k} L_j^1 - (n - r)(r - k).
 \end{aligned}$$

Also note that $\psi^{-1}((\alpha(r - k), \tilde{L})) = \{(\alpha(0), L)\}$ such that

$$L_j^i \begin{cases} \leq r - k & \text{for } \alpha_2^1 + 1 \leq j \leq \alpha_2^1 + k \text{ and } i = 1, \\ = \tilde{L}_j^1 + (r - k) & \text{for } j \geq \alpha_2^1 + k + 1 \text{ and } i = 1, \\ = \tilde{L}_j^i & \text{otherwise.} \end{cases}$$

Thus

$$\begin{aligned}
 &\sum_{(\alpha(0), L) \in \Psi^{-1}(\alpha(r - k), \tilde{L})} t^{c_r(\alpha(0), L)} \\
 &= \sum_{0 \leq L_{\alpha_2^1+1}^1 \leq \dots \leq L_{\alpha_2^1+k}^1 \leq r - k} t^{\sum_{j=1}^k L_{\alpha_2^1+j}^1 + c(\alpha(r - k), \tilde{L})} \text{ by (9)} \\
 &= \begin{bmatrix} r \\ k \end{bmatrix}_t t^{c(\alpha(r - k), \tilde{L})}.
 \end{aligned}$$

This last lemma, the fact that the image of Ψ is $\bigcup_{j=0}^r \mathcal{M}_{r-j}^j$, and Lemma 3.2 and Lemma 4.1, finish the proof of Theorem 1.1. \square

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