SIGNED WORDS AND PERMUTATIONS, I: A FUNDAMENTAL TRANSFORMATION

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This paper is dedicated to the memory of Percy Alexander MacMahon

Abstract. The statistics major index and inversion number, usually defined on ordinary words, have their counterparts in signed words, namely the so-called flag-major index and flag-inversion number. We give the construction of a new transformation on those signed words that maps the former statistic onto the latter one. It is proved that the transformation also preserves two other set-statistics: the inverse ligne of route and the lower records.

1. Introduction

The second fundamental transformation, as it was called later on (see [16, chap. 10], or [15, ex. 5.1.1.19]), was described in these proceedings [8]. Let $w = x_1x_2\ldots x_m$ be a (finite) word, whose letters $x_1$, $x_2$, $\ldots$, $x_m$ are integers. The integer-valued statistics Inversion Number “inv” and Major Index “maj” attached to the word $w$ are defined by

\begin{align}
\text{inv } w &:= \sum_{1 \leq i \leq m-1} \sum_{i<j} \chi(x_i > x_j), \\
\text{maj } w &:= \sum_{1 \leq i \leq m-1} i \chi(x_i > x_{i+1}),
\end{align}

making use of the $\chi$-notation that maps each statement $A$ to the value $\chi(A) = 1$ or 0, depending on whether $A$ is true or not.

If $m = (m_1, m_2, \ldots, m_r)$ is a sequence of $r$ nonnegative integers, the rearrangement class of the nondecreasing word $1^{m_1}2^{m_2}\ldots r^{m_r}$, that is, the class of all the words than can be derived from $1^{m_1}2^{m_2}\ldots r^{m_r}$ by permutation of the letters, is denoted by $R_m$. The second fundamental transformation, denoted by $\Phi$, maps each word $w$ on another word $\Phi(w)$ and has the following properties:

(a) $\text{maj } w = \text{inv } \Phi(w)$;
(b) $\Phi(w)$ is a rearrangement of $w$, and the restriction of $\Phi$ to each rearrangement class $R_m$ is a bijection of $R_m$ onto itself.
Further properties were proved later on by Foata, Schützenberger [10] and Björner, Wachs [9]. In particular, when the transformation is restricted to act on rearrangement classes \( R_m \) such that \( m_1 = \cdots = m_r = 1 \), that is, on symmetric groups \( S_r \).

The purpose of this paper is to construct an analogous transformation not simply on words, but on signed words, so that new equidistribution properties on classical statistics, such as the (Coxeter) length function (see [8] p. 9, [14] p. 12), defined on the group \( B_m \) of the signed permutations can be derived. By a signed word we mean a word \( w = x_1 x_2 \ldots x_m \) whose letters are positive or negative integers. If \( m = (m_1, m_2, \ldots, m_r) \) is a sequence of nonnegative integers such that \( m_1 + m_2 + \cdots + m_r = m \), let \( B_m \) be the set of all rearrangements \( w = x_1 x_2 \ldots x_m \) of the sequence \( 1^{m_1} 2^{m_2} \ldots r^{m_r} \), with the convention that some letters \( i \) may be replaced by their opposite values \(-i\). For typographical reasons we shall use the notation \( \overline{i} := -i \) in the sequel. The class \( B_m \) contains \( 2^m (m_1, m_2, \ldots, m_r) \) signed words. When \( m_1 = m_2 = \cdots = m_r = 1 \), \( m = r \), the class \( B_m \) is simply the group \( B_m \) of the signed permutations of order \( m \).

Next, the statistics “inv” and “maj” must be adapted to signed words and correspond to classical statistics when applied to signed permutations. Let

\[
(\omega; q)_n := \begin{cases} 1, & \text{if } n = 0, \\ (1 - \omega)(1 - \omega q) \ldots (1 - \omega q^{n-1}), & \text{if } n \geq 1, \end{cases}
\]

denote the usual \( q \)-ascending factorial in a ring element \( \omega \) and let

\[
\left[ \frac{m_1 + \cdots + m_r}{m_1, \ldots, m_r} \right]_q := \frac{(q; q)_{m_1 + \cdots + m_r}}{(q; q)_{m_1} \cdots (q; q)_{m_r}}
\]

be the \( q \)-multinomial coefficient. Back to MacMahon [17], [18], [19], it was known that the above \( q \)-multinomial coefficient, which is the true \( q \)-analog of the cardinality of \( R_m \), was the generating function for the class \( R_m \) by either one of the statistics “inv” or “maj.” Consequently, the generating function for \( B_m \) by the new statistics that are to be introduced on \( B_m \) must be a plausible \( q \)-analog of the cardinality of \( B_m \). The most natural \( q \)-analog we can think of is certainly

\[
(-q; q)_m \left[ \frac{m_1 + \cdots + m_r}{m_1, \ldots, m_r} \right]_q = \sum_{w \in B_m} q^{\text{finv } w},
\]

which tends to \( 2^m \left( \frac{m_1 + \cdots + m_r}{m_1, \ldots, m_r} \right) \) when \( q \) tends to 1. As a substitute for “inv” we are led to introduce the following statistic “finv,” called the flag-inversion number, which will be shown to meet our expectation, that is,

\[
\sum_{w \in B_m} q^{\text{finv } w} = \sum_{w \in B_m} q^{\text{inv } w}.
\]

This identity is easily proved by induction on \( r \). Let \( w = x_1 x_2 \ldots x_m \) be a signed word from the class \( B_m \). To define \( \text{finv } w \) we use “inv” defined in (1.1), together with

\[
\overline{\text{inv } w} := \sum_{1 \leq i \leq m-1} \sum_{i < j} \chi(x_i > x_j),
\]

and define

\[
\text{finv } w := \text{inv } w + \overline{\text{inv } w} + \sum_{1 \leq j \leq m} \chi(x_j < 0).
\]

The salient feature of this definition of “finv” is the fact that it does not involve the values of the letters, but only the comparisons between letters, so that it can be applied to each arbitrary signed word. Moreover, the definition of “finv” is similar
to that of “fmaj” given below in (1.7). Finally, its restriction to the group $B_m$ of the signed permutations is the traditional length function:

$$\text{finv}|_{B_m} = \ell.$$  

This is easily shown, for instance, by using the formula derived by Brenti [7] for the length function $\ell$ over $B_n$:

$$\ell w = \text{inv} w + \sum_{1 \leq j \leq m} |x_j| \chi(x_j < 0).$$

Next, the statistic “maj” is to be replaced by “fmaj”, the flag-major index, introduced by Adin and Roichman [1] for signed permutations (see also [2]) showed that “fmaj” was equidistributed with the length function $\ell$ over $B_n$. Their definition of “fmaj” can be used verbatim. For a signed word $w$, we have Iligne $\Psi(w) = \text{fmaj} w$.

For example, with $w = x_1x_2 \ldots x_m$ those definitions read:

$$\text{fmaj } w := 2 \text{maj } w + \sum_{1 \leq j \leq m} \chi(x_j < 0),$$

$$\text{fdes } w := 2 \text{des } w + \chi(x_1 < 0),$$

where “des” is the usual number of descents $\text{des } w := \sum_i \chi(x_i > x_{i+1})$. We postpone the construction of our transformation $\Psi$ on signed words to the next section. The main purpose of this paper is to prove the following theorem.

**Theorem 1.1.** The transformation $\Psi$ constructed in section 2 has the following properties:

(a) $\text{fmaj } w = \text{finv } \Psi(w)$ for every signed word $w$;

(b) the restriction of $\Psi$ to each rearrangement class $B_m$ of signed words is a bijection of $B_m$ onto itself, so that “fmaj” and “finv” are equidistributed over each class $B_m$.

**Definition 1.1.** Let $w = x_1x_2 \ldots x_m \in B_m$ be a signed word. We say that a nonnegative integer $i$ belongs to the inverse ligne of route, Iligne $w$, of $w$, if one of the following two conditions holds:

1. $i = 0$, $m_1 \geq 1$ and the rightmost letter $x_k$ satisfying $|x_k| = 1$ is equal to $\overline{1}$;
2. $i \geq 1$, $m_i = m_{i+1} = 1$ and the rightmost letter that belongs to $\{i, \overline{i}, i + 1, \overline{i} + 1\}$ is equal to $i$ or $i + 1$.

For example, with $w = \overline{7}4\overline{3}2\overline{5}6\overline{7}$ we have Iligne $w = \{0, 2, 6\}$.

**Remark.** The expression “line of route” was used by Foulkes [11], [12]. We have added the letter “g” making “ligne of route,” thus bringing a slight touch of French. Note that 0 may or may not belong to the inverse ligne of route. The ligne of route of a signed word $w = x_1x_2 \ldots x_m$ is defined to be the set, denoted by Ligne $w$, of all the $i$’s such that either $1 \leq i \leq m - 1$ and $x_i > x_{i+1}$, or $i = 0$ and $x_1 < 0$. In particular, maj $w = \sum_{0 \leq i \leq m-1} i \chi(i \in \text{Ligne } w)$. Finally, if $w$ is a signed permutation, then Iligne $w = \text{Ligne } w^{-1}$. For ordinary permutations, some authors speak of descent set and descent set of the inverse, instead of ligne of route and inverse ligne of route, respectively.

**Theorem 1.2.** The transformation $\Psi$ constructed in section 2 preserves the inverse ligne of route:

(c) Iligne $\Psi(w) = \text{Iligne } w$ for every signed word $w$. 
Definition 1.2. Let \( w = x_1 x_2 \ldots x_m \) be a signed word of length \( m \). A letter \( x_i \) is said to be a lower record of \( w \), if either \( i = m \), or \( 1 \leq i \leq m - 1 \) and \( |x_i| < |x_j| \) for all \( j \) such that \( i + 1 \leq j \leq m \). When reading the lower records of \( w \) from left to right, we get a signed subword \( x_i x_{i+1} \ldots x_{i+k} \), called the lower record subword, denoted by Lower \( w \), which has the property that \( \min x_i = |x_i| < |x_{i+1}| < \cdots < |x_{i+k}| = |x_m| \). The notion of lower record is classical in the statistical literature. In combinatorics the expression “strict right-to-left minimum” is also used.

With our previous example \( w = \overline{441325367} \) we get Lower \( w = \overline{25367} \). Our third goal is to prove the following result.

Theorem 1.3. The transformation \( \Psi \) constructed in section 2 preserves all the lower records:

(d) Lower \( \Psi(w) = \text{Lower } w \) for every signed word \( w \).

For each signed permutation \( w = x_1 x_2 \ldots x_m \) let

\[
\text{ifmaj } w := 2 \sum_{1 \leq j \leq m} i \chi(j \in \text{Iligne } w) + \sum_{1 \leq j \leq m} \chi(x_j < 0),
\]

\[
\text{ifdes } w := 2 \sum_{1 \leq j \leq m} \chi(j \in \text{Iligne } w) + \chi(x_i = -1 \text{ for some } i).
\]

It is immediate to verify that

\( \text{finv } w = \text{finv } w^{-1} \), \( \text{ifmaj } w = \text{ifmaj } w^{-1} \), \( \text{ifdes } w = \text{ifdes } w^{-1} \),

where \( w^{-1} \) denotes the inverse of the signed permutation \( w \) (written as a linear word \( w^{-1} = w^{-1}(1) \ldots w^{-1}(m) \)).

Let \( w := w^{-1} \). Then the chain

\[
\left( \begin{array}{c} w \\ \text{finv} \\ \text{ifdes} \\ \text{ifmaj} \\ \text{finv} \end{array} \right) \xrightarrow{1} \left( \begin{array}{c} w_1 \\ \text{finv} \\ \text{ifdes} \\ \text{ifmaj} \\ \text{finv} \end{array} \right) \xrightarrow{\Psi} \left( \begin{array}{c} w_2 \\ \text{finv} \\ \text{ifdes} \\ \text{ifmaj} \\ \text{finv} \end{array} \right) \xrightarrow{1} \left( \begin{array}{c} w_3 \\ \text{finv} \\ \text{ifdes} \\ \text{ifmaj} \\ \text{finv} \end{array} \right)
\]

shows that the four generating polynomials \( \sum t^{\text{ifdes } w} q^{\text{ifmaj } w} \), \( \sum t^{\text{ifdes } w} q^{\text{ifmaj } w} \), \( \sum t^{\text{ifdes } w} q^{\text{ifmaj } w} \), and \( \sum t^{\text{ifdes } w} q^{\text{ifmaj } w} \) \((w \in B_m)\) are identical. Their analytic expression will be derived in a forthcoming paper [9].

2. The construction of the transformation

For each signed word \( w = x_1 x_2 \ldots x_m \) the first or leftmost (resp. last or rightmost) letter \( x_1 \) (resp. \( x_m \)) is denoted by \( F(w) \) (resp. \( L(w) \)). Next, define \( s_1 w := x_1 x_2 \ldots x_m \). The transformation \( s_1 \) changes the sign of the first letter. Together with \( s_1 \) the main ingredients of our transformation are the bijections \( \gamma_x \) and \( \delta_y \) defined for each integer \( x, y \), as follows.

If \( L(w) \leq x \) (resp. \( L(w) > x \)), then \( w \) admits the unique factorization

\[
(v_1 y_1, v_2 y_2, \ldots, v_p y_p),
\]

called its \( x \)-right-to-left factorisation having the following properties:

(i) each \( y_i \) \((1 \leq i \leq p)\) is a letter verifying \( y_i \leq x \) (resp. \( y_i > x \));
(ii) each \( v_i \) \((1 \leq i \leq p)\) is a factor which is either empty or has all its letters greater than (resp. smaller than or equal to) \( x \).
Then, $\gamma_x$ is defined to be the bijection that maps $w = v_1 y_1 v_2 y_2 \ldots v_p y_p$ onto the signed word

\begin{equation}
\gamma_x(w) := y_1 v_1 y_2 v_2 \ldots y_p v_p.
\end{equation}

In a dual manner, if $F(w) \geq x$ (resp. $F(w) < x$) the signed word $w$ admits the unique factorization

\[(z_1 w_1, z_2 w_2, \ldots, z_q w_q)\]
called its $x$-left-to-right factorisation having the following properties:

(i) each $z_i$ ($1 \leq i \leq q$) is a letter verifying $z_i \geq x$ (resp. $z_i < x$);
(ii) each $w_i$ ($1 \leq i \leq q$) is a factor which is either empty or has all its letters less than (resp. greater than or equal to) $x$.

Then, $\delta_x$ is defined to be the bijection that sends $w = z_1 w_1 z_2 w_2 \ldots z_q w_q$ onto the signed word

\begin{equation}
\delta_x(w) := w_1 z_1 w_2 z_2 \ldots w_q z_q.
\end{equation}

Next, if $(v_1 y_1, v_2 y_2, \ldots, v_p y_p)$ is the $x$-right-to-left factorization of $w$, we define

\begin{equation}
\beta_x(w) := \begin{cases} 
\delta_{\tau} \gamma_x(w), & \text{if either } \tau \leq y_1 \leq x, \text{ or } x < y_1 < \tau, \\
\delta_{\tau} s_1 \gamma_x(w), & \text{otherwise}.
\end{cases}
\end{equation}

The fundamental transformation $\Psi$ on signed words that is the main object of this paper is defined as follows: if $w$ is a one-letter signed word, let $\Psi(w) := w$; if it has more than one letter, write the word as $wx$, where $x$ is the last letter. By induction determine $\Psi(w)$, then apply $\beta_x$ to $\Psi(w)$ and define $\Psi(wx)$ to be the juxtaposition product:

\begin{equation}
\Psi(wx) := \beta_x(\Psi(w)) x.
\end{equation}

The proof of Theorem 1.1 is given in section 3. It is useful to note the following relation:

\begin{equation}
y_1 \leq x \iff L(w) \leq x
\end{equation}

and the identity

\begin{equation}
\Psi(wx) = \Psi(w) x, \quad \text{whenever } x < -\max\{|x_i|\} \text{ or } x \geq \max\{|x_i|\}.
\end{equation}

**Example.** Let $w = 3 \overline{2} 1 \overline{3} 3$. The factorizations used in the definitions of $\gamma_x$ and $\delta_x$ are indicated by vertical bars. First, $\Psi(3) = 3$. Then

\begin{align*}
|3| & \xrightarrow{\tau} 3 |s_1| \xrightarrow{\tau} 3, \text{ so that } \Psi(3 \overline{2}) = 3 \overline{2} ; \\
|3 \overline{2}| & \xrightarrow{\tau} 3 \overline{2} |s_1|, \text{ so that } \Psi(3 \overline{2} 1) = 3 \overline{2} 1 ; \\
|3 \overline{2} 1 \overline{3} 3| & \xrightarrow{\tau} 3 \overline{2} 1 \overline{3} 3 |s_1|, \text{ so that } \Psi(3 \overline{2} 1 \overline{3} 3) = 3 \overline{2} 1 \overline{3} 3 ;
\end{align*}

and $\Psi(3 \overline{2} 1 \overline{3} 3) = 3 \overline{2} 1 \overline{3} 3$, because of (2.6);

\begin{align*}
|3 \overline{2} 1 \overline{3} 3| & \xrightarrow{\tau} 3 \overline{2} 1 \overline{3} 3 |s_1| \xrightarrow{\tau} 3 \overline{2} 1 \overline{3} 3.
\end{align*}

Thus, with $w = 3 \overline{2} 1 \overline{3} 3$ we get $\Psi(w) = 3 \overline{2} 1 \overline{3} 3$. We verify that $\text{fmaj } w = \text{finv } \Psi(w) = 19$, $\text{liligne } w = \text{liligne } \Psi(w) = \{1\}$, $\text{Lower } w = \text{Lower } \Psi(w) = 13$. 

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3. Proof of Theorem 1.1

Before proving the theorem we state a few properties involving the above statistics and transformations. Let \(|w|\) be the number of letters of the signed word \(w\) and let \(|w|_x\) be the number of its letters greater than \(x\) with analogous expressions involving the subscripts \(\geq x\), \(< x\) and \(\leq x\). We have:

\[
(3.1) \quad \text{fmaj } wx = \text{fmaj } w + \chi(x < 0) + 2 |w|\chi(L(w) > x),
\]

\[
(3.2) \quad \text{finv } wx = \text{finv } w + |w|_x + |w|_{< x} + \chi(x < 0),
\]

\[
(3.3) \quad \text{finv } \gamma_x(w) = \text{finv } w + |w|_{\leq x} - |w|\chi(L(w) \leq x),
\]

\[
(3.4) \quad \text{finv } \delta_x(w) = \text{finv } w + |w|_x - |w|\chi(F(w) \geq x).
\]

Next, let \(y_1\) denote the first letter of the signed word \(w''\). Then

\[
(3.5) \quad \text{finv } s_1 w'' = \text{finv } w'' + \chi(y_1 > 0) - \chi(y_1 < 0),
\]

\[
(3.6) \quad |s_1 w''|_x = |w''|_x + \chi(x > 0)(\chi(y_1 < x) - \chi(x < y_1)) + \chi(x < 0)(\chi(y_1 < x) - \chi(x \leq y_1)).
\]

Theorem 1.1 is now proved by induction on the word length. Assume that \(\text{fmaj } w = \text{finv } \Psi(w)\) for a given \(w\). Our purpose is to show that

\[
(3.7) \quad \text{fmaj } wx = \text{finv } \Psi(wx)
\]

holds for all letters \(x\). Let \(w' = \Psi(w)\), so that by (2.4) the words \(w\) and \(w'\) have the same rightmost letter. Denote the \(x\)-right-to-left factorization of \(w'\) by \((v_1y_1, \ldots, v_py_p)\). By (2.3) the signed word \(v := \beta_x(w')\) is defined by the chain

\[
(3.8) \quad w' = v_1y_1 \ldots v_py_p \xrightarrow{\gamma_x} w'' = y_1v_1 \ldots y_pv_p = z_1w_1 \ldots zqw_q
\]

\[
\xrightarrow{s_1} v = w_1z_1 \ldots w_qz_q
\]

if either \(x \leq y_1 \leq x\), or \(x < y_1 < x\), and by the chain

\[
(3.9) \quad w' = v_1y_1 \ldots v_py_p \xrightarrow{\gamma_x} w'' = y_1v_1 \ldots y_pv_p
\]

\[
\xrightarrow{s_1} v = y_1v_1 \ldots y_pv_p = z_1w_1 \ldots zqw_q
\]

\[
\xrightarrow{s_p} v = w_1z_1 \ldots w_qz_q,
\]

otherwise. Note that \((z_1w_1, \ldots, zqw_q)\) designates the \(\overline{x}\)-left-to-right factorization of \(w''\) in chain (3.8) and of \(w''\) in chain (3.9).

(i) Suppose that one of the conditions \(x \leq y_1 \leq x\), \(x < y_1 < x\) holds, so that (3.8) applies. We have

\[
\text{finv } \Psi(wx) = \text{finv } wx = \text{finv } v + |v|_x + |v|_{< x} + \chi(x < 0) \quad \text{[by (3.2)]},
\]

\[
\text{finv } v = \text{finv } s_1(w'') = \text{finv } w'' + |w''|_{\geq x} - |w''|\chi(F(w'') \geq x) \quad \text{[by (3.4)]},
\]

\[
\text{finv } w'' = \text{finv } \gamma_x(w') = \text{finv } w' + |w'|_{\leq x} - |w'|\chi(L(w') \leq x) \quad \text{[by (3.3)]},
\]

\[
\text{finv } w' = \text{fmaj } w \quad \text{[by induction]},
\]

\[
\text{fmaj } w = \text{fmaj } wx - \chi(x < 0) - 2 |w|\chi(L(w) > x) \quad \text{[by (3.1)]}.
\]

By induction,

\[
(3.10) \quad L(w) = L(w') \quad \text{and} \quad \chi(L(w') > x) = 1 - \chi(L(w') \leq x).
\]
Also $F(w''') = y_1$. As $w'$, $w''$, $v$ are true rearrangements of each other, we have

$$|v|_{>x} + |w'|_{<x} = |w|, |v|_{<\bar{x}} + |w''|_{\geq \bar{x}} = |w|.$$ Hence,

$$\text{finv } \Psi(wx) = \text{fmaj } wx + |w|[\chi(L(w') \leq x) - \chi(y_1 \geq \bar{x})].$$

By (2.5), if $\bar{x} \leq y_1 \leq x$ holds, then $L(w') \leq x$ and the expression between brackets is null. If $x < y_1 < \bar{x}$ holds, then $L(w') > x$ and the same expression is also null. Thus (3.7) holds.

(ii) Suppose that none of the conditions $\bar{x} \leq y_1 \leq x$, $x < y_1 < \bar{x}$ holds, so that (3.9) applies. We have

$$\text{finv } \Psi(wx) = \text{finv } vx = \text{finv } v + |v|_{>x} + |v|_{<\bar{x}} + \chi(x < 0) \quad \text{[by (3.2)]},$$

$$\text{finv } v = \text{finv } \delta_x(w''') = \text{finv } w'' + |w''|_{\geq \bar{x}} - |w'''| \chi(F(w''') \geq \bar{x}) \quad \text{[by (3.4)]},$$

$$\text{finv } w''' = \text{finv } s_1 y'' = \text{finv } w''' + \chi(y_1 > 0) - \chi(y_1 < 0) \quad \text{[by (3.5)]},$$

$$\text{finv } w'' = \text{finv } \gamma_x(w') = \text{finv } w' + |w'|_{\leq x} - |w'| \chi(L(w') \leq x) \quad \text{[by (3.3)]},$$

$$\text{finv } w = \text{fmaj } w \quad \text{[by induction]},$$

$$\text{fmaj } w = \text{fmaj } wx - \chi(x < 0) - 2|w|\chi(L(w) > x) \quad \text{[by (3.1)]}.$$

Moreover, $F(w''') = \bar{y}_1$, so that $\chi(F(w''') \geq \bar{x}) = \chi(\bar{y}_1 \geq \bar{x}) = \chi(y_1 \leq x) = \chi(L(w') \leq x) = \chi(L(w) \leq x)$. As $v$ and $w'''$ are rearrangements of each other, we have $|v|_{<\bar{x}} + |w'''|_{\geq \bar{x}} = |w|$. Using (3.6) since $|v|_{>x} = |w'''|_{>x} = |s_1 w'''|_{>x}$ we have

$$\text{finv } \Psi(wx) = |w'''|_{>x} + \chi(x > 0)(\chi(y_1 < \bar{x}) - \chi(y_1 < y_1))$$

$$+ \chi(x < 0)(\chi(y_1 \leq x) - \chi(\bar{x} \leq y_1)) + \chi(x < 0)$$

$$+ |w| - |w| \chi(y_1 \leq x) + \chi(y_1 > 0) - \chi(y_1 < 0)$$

$$+ |w'|_{\leq x} - |w'| \chi(y_1 \leq x)$$

$$+ \text{fmaj } wx - \chi(x < 0) - 2|w| + 2|w| \chi(y_1 \leq x)$$

$$= \text{fmaj } wx + \chi(x > 0)(\chi(y_1 < \bar{x}) - \chi(x < y_1))$$

$$+ \chi(x < 0)(\chi(y_1 \leq x) - \chi(\bar{x} \leq y_1))$$

$$+ \chi(y_1 > 0) - \chi(y_1 < 0)$$

$$= \text{fmaj } wx,$$

for, if none of the conditions $\bar{x} \leq y_1 \leq x$, $x < y_1 < \bar{x}$ holds, then one of the following four holds: (a) $y_1 > x > 0$; (b) $\bar{y}_1 > x > 0$; (c) $y_1 \leq x < 0$; (d) $\bar{y}_1 \leq x < 0$; and in each case the sum of the factors in the above sum involving $\chi$ is zero.

The construction of $\Psi$ is perfectly reversible. First, note that $s_1$ is an involution and the maps $\gamma_x$, $\delta_{\bar{x}}$ send each class $B_m$ onto itself, so that their inverses are perfectly defined. They can also be described by means of left-to-right and right-to-left factorizations. Let us give the construction of the inverse $\Psi^{-1}$ of $\Psi$. Of course, $\Psi^{-1}(v) := v$ if $v$ is a one-letter word. If $vx$ is a signed word whose last letter is $x$, determine $v' := \delta_x^{-1}(v)$ and let $z_1$ be its first letter. If one of the conditions $\bar{x} \leq z_1 \leq x$ or $x < z_1 < \bar{x}$ holds, the chain (3.8) is to be used in reverse order, so that $\Psi^{-1}(vx) := (\Psi^{-1}(v) \delta_x^{-1}(v))x$. If none of those two conditions holds, then $\Psi^{-1}(vx) := (\Psi^{-1}(v) s_1 \delta_{\bar{x}}^{-1}(v))x$. □
4. Proofs of Theorems 1.2 and 1.3

Before proving Theorem 1.2 we note the following two properties.

**Property 4.1.** Let \( I_x = \{ i \in \mathbb{Z} : i < -|x| \} \) (resp. \( J_x = \{ i \in \mathbb{Z} : -|x| < i < |x| \} \), resp. \( K_x = \{ i \in \mathbb{Z} : |x| < i \} \) and let \( w \) be a signed word. Then, the bijections \( \gamma_x \) and \( \delta_x \) do not modify the mutual order of the letters of \( w \) that belong to \( I_x \) (resp. \( J_x \), resp. \( K_x \)).

**Proof.** Let \( y \in I_x \) and \( z \in I_x \) (resp. \( y \in J_x \) and \( z \in J_x \), resp. \( y \in K_x \) and \( z \in K_x \)) be two letters of \( w \) with \( y \) to the left of \( z \). In the notations of (2.1) (resp. of (2.2)) both \( y, z \) are either among the \( y_i \)'s (resp. the \( z_i \)'s) or letters of the \( v_i \)'s (resp. the \( w_i \)'s). Accordingly, \( y \) remains to the left of \( z \) when \( \gamma_x \) (resp. \( \delta_x \)) is applied to \( w \).

**Property 4.2.** Let \( w = x_1 x_2 \ldots x_m \in B_m \) be a signed word and let \( i \) be a positive integer such that \( m_i = m_{i+1} = 1 \). Furthermore, let \( x \) be an integer such that \( x \not \in \{ i, i+1, i+1 \} \). Then the following conditions are equivalent: (a) \( i \in \text{Iligne } w \); (b) \( i \in \text{Iligne } s_{i} w \); (c) \( i \in \text{Iligne } \gamma_x w \); (d) \( i \in \text{Iligne } \delta_x w \).

**Proof.** (a)⇔(b) holds by Definition 1.1, because \( s_{i} \) has no action on the rightmost letter belonging to \( \{ i, i+1, i+1 \} \). For the other equivalences we can say the following. If the two letters of \( w \) that belong to \( \{ i, i+1, i+1 \} \) are in \( I_x \) (resp. \( J_x \), resp. \( K_x \)), Property 4.1 applies. Otherwise, if \( i \in \text{Iligne } w \), then \( w \) is either of the form \( \ldots i+1 \ldots i \ldots \) or \( \ldots i \ldots i+1 \ldots \) and the order of those two letters is immaterial.

Theorem 1.2 holds for each one-letter signed word. Let \( w = x_1 x_2 \ldots x_m \in B_m \) be a signed word, \( x \) a letter and \( i \) a positive integer. Assume that \( \text{Iligne } w = \text{Iligne } \Psi(w) \).

If \( x = i \), then \( i \in \text{Iligne } w x \) if and only if \( w \) contains no letter equal to \( \pm i \) and exactly one letter equal to \( \pm (i+1) \). As \( \beta_x \Psi(w) \) is a rearrangement of \( w \) with possibly sign changes for some letters, the last statement is equivalent to saying that \( \beta_x \Psi(w) \) has no letter equal to \( \pm i \) and exactly one letter equal to \( \pm (i+1) \). This is also equivalent to saying that \( i \in \text{Iligne } \beta_x \Psi(w)x = \text{Iligne } \Psi(wx) \). In the same manner, we can show that:

- if \( x = i+1 \), then \( i \in \text{Iligne } w x \) if and only if \( i \in \text{Iligne } \Psi(wx) \);
- if \( x = i \), then \( i \not \in \text{Iligne } w x \) and \( i \not \in \text{Iligne } \Psi(wx) \);
- if \( x = i+1 \), then \( i \not \in \text{Iligne } w x \) and \( i \not \in \text{Iligne } \Psi(wx) \).

Now, let \( i \) be such that none of the integers \( i, i+1, i, i+1 \) is equal to \( x \). There is nothing to prove if \( m_i = m_{i+1} = 1 \) does not hold, as \( i \) does not belong to any of the sets \( \text{Iligne } w \), \( \text{Iligne } \beta_x \Psi(w) \). Otherwise, the result follows from Property 4.2 because \( \Psi \) is a composition product of \( \beta_x, s_{i} \) and \( \gamma_x \). Finally, the equivalence \( [0 \in \text{Iligne } \beta_x(w)x] \iff [0 \in \text{Iligne } w x] \) follows from Property 4.1 when \( |x| > 1 \) and the result is evident when \( |x| = 1 \).

The proof of Theorem 1.3 also follows from Property 4.1. By definition the lower records of \( w x \), other than \( x \), belong to \( J_x \). As the bijections \( \gamma_x \) and \( \delta_x \) do not modify the mutual order of the letters of \( w \) that belong to \( J_x \), we have \( \text{Lower } w x = \text{Lower } \beta_x(w)x \). When the chain (3.8) is used. When (3.9) is applied, so that \( y_i \not \in J_x \), we also have \( z_i = \overline{y}_i \not \in J_x \). Thus, neither \( y_i \), nor \( z_i \) can be lower records for each word ending with \( x \). Again, \( \text{Lower } w x = \text{Lower } \beta_x(w)x \).  

\[\square\]
5. Concluding remarks

Since the works by MacMahon, much attention has been given to the study of statistics on the symmetric group or on classes of word rearrangements, in particular by the M.I.T. school (23, 20, 27, 13, 7). It was then natural to extend those studies to other classical Weyl groups, as was done by Reiner (20, 21, 22, 23, 24) for the signed permutation group. Today the work has been pursued by the Israeli and Roman schools [1, 2, 3, 4]. The contribution of Adin, Roichman [1] has been essential with their definition of the flag major index for signed permutations. In our forthcoming paper [9] we will derive new analytical expressions, in particular for several multivariable statistics involving “fmaj,” “finv” and the number of lower records.

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