A First Order Method for Differential Equations of Neutral Type

By R. N. Castleton and L. J. Grimm*

Abstract. A first order method is presented for solution of the initial-value problem for a differential equation of neutral type with implicit delay in the critical case where the time-lag is zero and the method of stepwise integration does not apply. A convergence theorem is proved, and numerical examples are given.

1. Introduction. In this note, we present a first order method for the numerical solution of the initial-value problem (IVP) for a neutral-type functional-differential equation without previous history:

(1)
$$x'(t) = f(t, x(t), x(g(t, x(t))), x'(g(t, x(t)))),$$

(2)
$$x(a) = x_0, \quad x'(a) = z_0,$$

where z_0 is a real root of the algebraic equation

(3)
$$z = f(a, x_0, x_0, z).$$

Here, x(t) is a scalar function to be determined on some finite interval [a, b]. We shall make the following assumptions regarding f and g:

(H1) f and g are continuous and satisfy uniform Lipschitz conditions of the form

$$|f(t, x_1, y_1, z_1) - f(t, x_2, y_2, z_2)| \le L\{|x_1 - x_2| + |y_1 - y_2|\} + L_z |z_1 - z_2|,$$

$$|g(t, x_1) - g(t, x_2)| \le L_g |x_1 - x_2|$$

in their respective domains E and E', where

$$E = \{(t, x, y, z): a \le t \le b, |x - x_0| \le c, |y - x_0| \le c, |z| \le M\}$$

and E' is the projection of E in the (t, x) space; c, M, L, L_v, L_z are constants, with $L_z < 1$, M is such that $\sup_{(t, x, y, z) \in E} |f(t, x, y, z)| < M$, and M(b - a) < c. (H2) $a \le g(t, x) \le t$ for $(t, x) \in E'$.

Our hypotheses, together with additional smoothness and growth conditions on f and g, ensure the local existence of a solution of the IVP (1)–(2). Furthermore, x(t) is the only solution having a bounded derivative on [a, b]; see [2], [4]. Our result extends a method developed by Feldstein [3] for the equation of retarded type

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$$x'(t) = f(t, x(t), x(g(t)))$$

to the neutral-type equation with implicit delay (1). Other methods for implicitdelay equations are given in [1].

2. The Algorithm \mathfrak{A} . Let y(t) = x(g(t, x(t))); z(t) = x'(g(t, x(t))). Let N be a positive integer, and let h = (b - a)/N. For each nonnegative integer $n \le N$, let $t_n = a + nh$. Let [s] denote the integer part of s. Define the algorithm \mathfrak{A} as follows:

(4)
$$f_n = f(t_n, x_n, y_n, z_n), \quad g_n = g(t_n, x_n),$$

(5)
$$q(n) = [(g_n - a)/h], \quad r(n) = (g_n - a)/h - q(n),$$

(6)
$$y_0 = x_0, y_n = x_{q(n)} + hr(n)f_{q(n)},$$

$$(7) z_n = f_{q(n)},$$

(8)
$$x_{n+1} = x_n + h f_n$$

Note that condition (H2) implies $q(n) \le n$, thus, the algorithm is well defined. For n=0, $g_0=a$, q(0)=0, and r(0)=0. Thus, $y_0=x_0$ and $z_0=f(a,x_0,x_0,z_0)$. Let u_0 , an approximation of the root z_0 , be chosen independently of h. It is of interest to note that such an approximation does not destroy the order h convergence of the algorithm. It is of further interest that (6) may be simplified to $y_n=x_{q(n)}$. The error bound established in the convergence theorem for this "simplified" algorithm is larger but still of order h, as noted following the proof of convergence of the algorithm $\mathfrak A$. The second numerical example of Section 4 demonstrates both the algorithm $\mathfrak A$ and the simplified algorithm.

If $g_n = t_n$ for any $n, 1 \le n \le N$, then q(n) = n, r(n) = 0, and (7) becomes $z_n = f(t_n, x_n, y_n, z_n)$ which has exactly one root z in the interval [-M, M] under the conditions (H1)-(H2) together with the smoothness and growth conditions mentioned in Section 1. We must in general include a procedure for finding this root, and this in turn will affect the error estimate. As before, such an estimate does not destroy the order h convergence of the algorithm. For simplicity, we do not take this into account, since our aim is to show the convergence of the algorithm \mathfrak{A} .

Thus, we shall assume in the convergence proof that (7) will not reduce to $z_n = f(t_n, x_n, y_n, z_n), n \ge 1$.

3. Convergence.

THEOREM. Let f and g satisfy (H1)-(H2) and suppose, in addition, that there exists a unique solution x(t) of (1)-(2) with $\sup_{\{a,b\}} |x''(t)| \le B$. Then, for each $t_n \in [a,b]$, $0 < n \le N$,

$$|x_n - x(t_n)| \leq h \left\{ L_z |z_0 - u_0| e^{s(b-a)} + \frac{B}{2s} \left(\frac{1 + L_z}{1 - L_z} \right) (e^{s(b-a)} - 1) \right\} + O(h^2)$$

where

$$s = L(1 + c_0) + L_1c_1$$

$$c_0 = 1 + ML_v,$$

 $c_1 = (L(2 + ML_v) + BL_v)/(1 - L_v),$

 u_0 is the approximation to z_0 mentioned above, and x_n is given by algorithm \mathfrak{A} .

Proof. Let $e_n = |x_n - x(t_n)|$; $e_n^* = |y_n - y(t_n)|$; $e_n^{**} = |z_n - z(t_n)|$. From (8) and Taylor's formula, we obtain

(9)
$$e_{n+1} \leq e_n + h(L(e_n + e_n^*) + L_2 e_n^{**}) + h^2 B/2.$$

Equation (5) implies that $g_n = t_{q(n)} + hr(n)$, and hence, in a similar manner, we have (after replacing n by (n + 1))

(10)
$$e_{n+1}^* \leq ML_0 e_{n+1} + e_{q(n+1)} + hr(n+1) \{ L(e_{q(n+1)} + e_{q(n+1)}^* + L_2 e_{q(n+1)}^{**} \} + h^2 r^2 (n+1) B/2,$$

$$(11) e_{n+1}^{**} \leq BL_{q}e_{n+1} + L(e_{q(n+1)} + e_{q(n+1)}^{*}) + L_{\ell}e_{q(n+1)}^{**} + hr(n+1)B.$$

We then have two cases to consider:

Case 1. q(n + 1) = n + 1 and r(n + 1) = 0. Under these conditions, (9) is unchanged:

(9a)
$$e_{n+1} \leq e_n(1 + hL) + e_n^*hL + e_n^{**}hL_z + h^2B/2.$$

(10) becomes

(10a)
$$e_{n+1}^* \leq e_{n+1}(1 + ML_q) = e_{n+1}c_0.$$

And (11) becomes

$$e_{n+1}^{**} \leq (L + BL_n)e_{n+1} + Le_{n+1}^* + L_ne_{n+1}^{**}$$

or

(11a)
$$e_{n+1}^{**} \leq \left(\frac{L + BL_g + L(1 + ML_g)}{1 - L_z}\right) e_{n+1} = e_{n+1}c_1.$$

Define the partial ordering for vectors: $v_1 = (v_1^1, \dots, v_1^k) \le v_2 = (v_2^1, \dots, v_2^k)$ if $v_1^i \le v_2^i$, $i = 1, \dots, k$. Then, in vector form, (9a), (10a), and (11a) become

$$\begin{bmatrix} e_{n+1} \\ e_{n+1}^* \\ e_{n+1}^{**} \end{bmatrix} \leq \begin{bmatrix} 1 + hL & hL & hL_z \\ (1 + hL)c_0 & hLc_0 & hL_zc_0 \\ (1 + hL)c_1 & hLc_1 & hL_zc_1 \end{bmatrix} \begin{bmatrix} e_n \\ e_n^* \\ e_n^{**} \end{bmatrix} + hB \begin{bmatrix} h/2 \\ hc_0/2 \\ hc_1/2 \end{bmatrix}$$

which is of the form $d_{n+1} \leq A_1 d_n + b_1$.

Case 2.
$$q(n + 1) \le n \text{ and } 0 \le r(n + 1) < 1.$$

Let

$$\delta_n = \max_{1 \leq i \leq n} e_i, \qquad \delta_n^* = \max_{1 \leq i \leq n} e_i^*, \qquad \delta_n^{**} = \max_{1 \leq i \leq n} e_i^{**}.$$

Then, (9) becomes

(9b)
$$\delta_{n+1} \leq \delta_n (1 + hL) + \delta_n^* hL + \delta_n^{**} hL_1 + h^2 B/2.$$

And (10) becomes

$$\delta_{n+1}^* \leq M L_{\sigma} \delta_{n+1} + \delta_n (1 + hL) + h L \delta_n^* + h L_{z} \delta_n^{**} + h^2 B/2.$$

Using (9b), we have

$$\delta_{n+1}^* \leq (\delta_n(1+hL) + \delta_n^*hL + \delta_n^{**}hL_z + h^2B/2)(1+ML_q)$$

or

(10b)
$$\delta_{n+1}^* \leq \delta_n (1 + hL)c_0 + \delta_n^* hLc_0 + \delta_n^{**} hL_z c_0 + h^2 c_0 B/2.$$

Finally, (11) becomes

$$\delta_{n+1}^{**} \leq \delta_{n+1}BL_q + \delta_nL + \delta_n^*L + \delta_n^{**}L_z + hB.$$

Further, enlarging δ_n to δ_{n+1} and δ_n^* to δ_{n+1}^* on the right, and using $1 - L_z > 0$, we find

$$\delta_{n+1}^{**} \leq \delta_{n+1} \left(\frac{L + BL_0}{1 - L_1} \right) + \delta_{n+1}^{*} \frac{L}{1 - L_2} + \frac{hB}{1 - L_2}$$

Using (9b) and (10b), we have

$$\delta_{n+1}^{**} \leq \left(\frac{L + BL_o + Lc_0}{1 - L_z}\right) \left(\delta_n(1 + hL) + \delta_n^* hL + \delta_n^{**} hL_z + \frac{h^2 B}{2}\right) + \frac{hB}{1 - L_z}$$

or

(11b)
$$\delta_{n+1}^{**} \leq \delta_n (1 + hL)c_1 + \delta_n^* hLc_1 + \delta_n^{**} hL_2 c_1 + \frac{hB}{1 - I} + \frac{h^2 c_1 B}{2}.$$

Then, as a vector system, (9b), (10b), and (11b) become

$$\begin{bmatrix}
\delta_{n+1} \\
\delta_{n+1}^{*} \\
\delta_{n+1}^{**}
\end{bmatrix} \leq \begin{bmatrix}
1 + hL & hL & hL_{z} \\
(1 + hL)c_{0} & hLc_{0} & hL_{z}c_{0} \\
(1 + hL)c_{1} & hLc_{1} & hL,c_{1}
\end{bmatrix} \begin{bmatrix}
\delta_{n} \\
\delta_{n}^{*} \\
\delta_{n}^{**}
\end{bmatrix} + hB \begin{bmatrix}
h/2 \\
hc_{0}/2 \\
hc_{1}/2 + 1/(1 - L_{z})
\end{bmatrix}$$

which is of the form $d_{n+1} \le A_2 d_n + b_2$. Comparing this with the result obtained in Case 1, we find that A_1 and A_2 are identical and that $b_1 \le b_2$. Thus, any bound obtained here in Case 2 for d_{n+1} will also bound d_{n+1} in Case 1.

To complete the proof, we shall use the following lemmas [3] which may be verified by induction:

LEMMA 1. Suppose A is a $k \times k$ real matrix and b is a real k-vector. Let $\{d_n\}$ $(n = 0, 1, \dots)$ satisfy $d_{n+1} \leq Ad_n + b$. Then

$$d_{n+1} \leq A^{n+1}d_0 + \left(\sum_{i=0}^n A^i\right)b.$$

LEMMA 2. Let $p = (p_1, \dots, p_k)$, $q = (q_1, \dots, q_k)$. Suppose the $k \times k$ matrix A has the form $A = p^T q$. Then

$$A^{n} = \left(\sum_{i=1}^{k} p_{i}q_{i}\right)^{n-1}A.$$

By Lemma 1,

$$d_{n+1} \leq A_2^{n+1} d_0 + \left(\sum_{i=0}^n A_2^i \right) b_2,$$

where

$$d_0 = \begin{bmatrix} e_0 \\ e_0^* \\ e_0^{**} \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ |z_0 - u_0| \end{bmatrix}.$$

Then, because

$$A_2 = \begin{bmatrix} 1 \\ c_0 \\ c_1 \end{bmatrix} (1 + hL, hL, hL_t),$$

we can make use of Lemma 2 to obtain

$$A_2^i = (1 + hL + hLc_0 + hL_2c_1)^{i-1}A_2 = (1 + hs)^{i-1}A_2.$$

Two results follow from this: $A_2^{n+1} = (1 + hs)^n A_2 \le e^{s(b-a)} A_2$, and

$$\sum_{i=1}^{n} A_{2}^{i} = A_{2} \sum_{i=1}^{n} (1 + hs)^{i-1} = \frac{((1 + hs)^{n} - 1)}{hs} A_{2} \le \frac{1}{hs} (\exp(s(b - a)) - 1) A_{2}.$$

Finally,

$$d_{n+1} \leq A_2^{n+1} d_0 + \left(\sum_{i=0}^n A_2^i\right) b_2$$

$$\leq h \begin{cases} |z_0 - u_0| \ L_z e^{s(b-a)} \begin{bmatrix} 1 \\ c_0 \\ c_1 \end{bmatrix}$$

$$+ \frac{B}{2s} \left(hs + \frac{1+L_{z}}{1-L_{z}} \right) \left(e^{s(h-a)} - 1 \right) \begin{bmatrix} 1 \\ c_{0} \\ c_{1} \end{bmatrix} + B \begin{bmatrix} \frac{h}{2} \\ \frac{hc_{0}}{2} \\ \frac{hc_{1}}{2} + \frac{1}{1-L_{z}} \end{bmatrix}$$

which gives

$$e_{n+1} \leq \delta_{n+1} \leq h \left\{ |z_0 - u_0| \ L_z e^{s(b-a)} + \frac{B}{2s} \left(hs + \frac{1+L_z}{1-L_z} \right) (e^{s(b-a)} - 1) + \frac{hB}{2} \right\}$$

and the theorem follows.

For the simplified algorithm, where (6) is replaced by $y_n = x_{q(n)}$ the following bound is possible:

$$d_{n+1} \leq h \begin{cases} |z_0 - u_0| \ L_z e^{s(b-a)} \begin{bmatrix} 1 \\ c_0 \\ c_{1.} \end{bmatrix} \end{cases}$$

(13)
$$+ \left(\frac{B}{2s}\left(hs + \frac{1+L_z}{1-L_z}\right) + \frac{1}{s}\left(\frac{ML}{1-L_z}\right)\right)(e^{s(b-a)} - 1)\begin{bmatrix}1\\c_0\\c_1\end{bmatrix} \\ + B\begin{bmatrix}\frac{h}{2}\\\frac{hc_0}{2}\\\frac{hc_1}{2} + \frac{1}{1-L}\end{bmatrix} + \begin{bmatrix}0\\M\\\frac{ML}{1-L_z}\end{bmatrix},$$

and hence

$$e_{n+1} \le h \left\{ |z_0 - u_0| \ L_z e^{s(b-a)} + \left(\frac{B}{2s} \left(hs + \frac{1+L_z}{1-L_z} \right) + \frac{1}{s} \left(\frac{ML}{1-L_z} \right) \right) (e^{s(b-a)} - 1) + \frac{hB}{2} \right\}.$$

TABLE I. $x_n(h)$ denotes the value of x_n for step size h.

t_n	$x(t_n)$	$x_n(2^{-4})$	$x_n(2^{-6})$	$x_n(2^{-8})$	$x_n(2^{-10})$
0	0	0	0	0	0
. 0625	.0039	0	. 0029	. 0034	. 0039
. 1250	. 0158	. 0078	. 0138	. 0153	. 0157
. 1875	. 0360	. 0238	. 0329	. 0352	. 0358
. 2500	. 0653	. 0484	. 0610	. 0642	. 0650
. 3125	. 1048	. 0825	. 0990	. 1032	. 1044
. 3750	. 1562	. 1275	. 1485	. 1541	. 1556
. 4375	. 2224	. 1853	. 2119	. 2196	. 2217
. 5000	. 3078	. 2593	. 2942	. 3043	. 3069
. 5625	. 4206	. 3547	. 4026	. 4159	. 4194
. 6250	. 5771	. 4856	. 5518	. 5705	. 5754
. 6875	. 8185	. 6707	. 7778	. 8080	. 8159
. 7500	1.3244	. 9860	1.2205	1.2968	1.3174

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t_n	$x(t_n)$	$x_n^{(1)}(2^{-2})$	$x_n^{(2)}(2^{-2})$	$x_n^{(1)}(2^{-4})$	$x_n^{(2)}(2^{-4})$			
. 25	. 2474	. 2500	. 2500	. 2483	. 2478			
. 50	. 4794	. 4930	. 4892	. 4838	. 4759			
. 75	. 6816	. 7180	. 6866	. 6942	. 6739			
1.00	. 8414	. 9228	. 8569	. 8697	. 8273			
t_n	$x(t_n)$	$x_n^{(1)}(2^{-8})$	$x_n^{(2)}(2^{-8})$	$x_n^{(1)}(2^{-12})$	$x_n^{(2)}(2^{-12})$			
. 25	. 2474	. 2475	. 2471	. 2474	. 2474			
. 50	. 4794	. 4797	. 4787	. 4794	. 4794			
. 75	. 6816	. 6825	. 6802	. 6817	. 6815			
1.00	. 8414	. 8435	. 8390	. 8416	. 8413			

TABLE II. $x_n^{(1)}(h)$ denotes the value of x_n for step size h by algorithm \mathfrak{A} ; $x_n^{(2)}(h)$ denotes the value of x_n for step size h by the simplified algorithm.

4. Examples. (a) We solve the IVP

$$x'(t) = \frac{-4tx^2(t)}{4 + \log^2 \cos t} + \tan 2t + \frac{1}{2} \tan^{-1} z$$

 $(z_0 = 0, x_0 = 0, z = x'(g(t, x(t))) \equiv x'(tx^2(t)/(1 + x^2(t))))$ on the interval [0, .75]. The existence and uniqueness of the solution is guaranteed by the results of [2] mentioned earlier. The only solution is $x(t) = -\frac{1}{2} \log \cos 2t$.

The results of the computation by algorithm A are given in Table I.

(b) Consider the IVP

$$x'(t) = \cos t(1 + y) + xz - \sin(t(1 + \sin^2 t)),$$

with $y = x(tx^2(t))$, $z = x'(tx^2(t))$, $z_0 = 1$, $x_0 = 0$, on the interval [0, 1]. As in example (a), existence and uniqueness of the solution are guaranteed by the results of [2]. Here, the solution is $x(t) = \sin t$.

The results of the computation by the algorithm $\mathfrak A$ and by the simplified algorithm are given in Table II.

Computation Division Lawrence Livermore Laboratory Livermore, California 94551

Department of Mathematics University of Missouri, Rolla Rolla, Missouri 65401

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