REVIEWS AND DESCRIPTIONS OF TABLES AND BOOKS

The numbers in brackets are assigned according to the American Mathematical Society classification scheme. The 1980 Mathematics Subject Classification (1985 Revision) can be found in the December index volumes of Mathematical Reviews.

24[41–02, 65Dxx].—GÜNTHER NÜRNBERGER, *Approximation by Spline Functions*, Springer, Berlin, 1989, xi + 243 pp., 25 cm. Price $39.50.

In the past 30 years, there has been an explosive development in the theory of spline functions (piecewise polynomials and generalizations thereof) and their applications. This great interest is due in no small part to the fact that splines are ideal tools for use in approximating functions, and for designing methods for solving problems numerically. While a number of monographs on splines have appeared in the past 10 years, much of the theory has not yet appeared in book form. This book helps fill this gap by providing a comprehensive development of *best approximation* by splines, along with some related topics.

The book is divided into three parts. The first part (approx. 80 pp.) is devoted primarily to various types of Chebyshev systems and their properties. The topics discussed include divided differences, interpolation, existence, uniqueness, strong uniqueness, characterization, and computation of best approximations in each of the L_1, L_2, and L_∞ norms. Much of this material has been treated in other books, but the unified treatment given here (most of the results are given full and detailed proofs) provides a useful basis for the study of the rest of the book.

The second part of the book (approx. 100 pp.) deals with various types of Weak Chebyshev systems, and in particular with polynomial splines as their most important example. First, the basic properties of splines (including B-splines, zero properties, etc.) are developed. Then several special interpolation methods are discussed. The most interesting sections deal with best approximation by splines with fixed knots. In addition to the usual questions of existence, uniqueness, strong uniqueness, and characterization, the problems of continuity of the metric projection, the existence of a continuous selection for it, one-sided L_1 approximation and its connection with Gauss quadrature formulae, and the optimal approximation of linear functionals are all treated in detail. A Remez algorithm for finding best spline approximants in the uniform norm is also presented. Many of the results in this part of the book are quite new.

The third part of the book is in the form of an extended Appendix (approx. 30 pp.), divided into three sections. The first section sketches the theory of best

approximation by splines with free knots. The second section outlines ways of defining bivariate splines, including tensor-products, blending, simplex splines, and piecewise polynomial spaces on triangulations and other partitions. The final section deals with the solution of ODE's by collocation. This part of the book contains no proofs, but does cite a considerable number of references. A bibliography of approximately 400 references is included.

This book should be of interest to researchers in the theory of splines, as well as to users of splines in approximation and numerical analysis. It can be read by anyone with a good background in elementary analysis. The material is well organized, and the text reads very smoothly.

L.L.S.

25[94–01, 65T05, 68Qxx, 94A11].—RICHARD TOLIMIERI, MYOUNG AN & CHAO LU, *Algorithms for Discrete Fourier Transform and Convolution*, Springer, New York, 1989, xv + 350 pp., 24 cm. Price $59.00.

Discrete Fourier transforms and finite convolutions form a mainstay of digital signal processing algorithms. Ever since the discovery of the Cooley-Tukey fast Fourier transform, there has been a flurry of activity in designing efficient algorithms for finite harmonic analysis, and these algorithms have found applications far beyond the realm of digital signal processing. There is now a growing list of monographs devoted to these algorithms; we mention the widely used references Blahut [1] and Nussbaumer [3].

The book under review presents a unified approach to many fast Fourier transform and convolution algorithms, using matrix factorizations and tensor products of matrices as the common themes. The individual steps in such an algorithm are viewed as matrix operations, and the full algorithm amounts then to a matrix factorization, typically involving diagonal matrices, permutation matrices, and tensor products of relatively simple matrices. Conversely, matrix factorizations of an appropriate type can be translated into algorithms. This provides a systematic and mathematically appealing framework for the design of discrete Fourier transform and convolution algorithms. This approach also makes it easy to switch from parallelized to vectorized algorithms and vice versa, as one can move from one to the other essentially by transposition of the matrix factorization and by application of a commutation theorem for tensor products of matrices. This allows an adaptation of the algorithms to the available computer architecture.

After the necessary background on ring and field theory and on tensor products of matrices, a detailed discussion of the Cooley-Tukey FFT algorithm from the viewpoint described above is given. Several variants of the Cooley-Tukey algorithm, including those of Gentleman-Sande, Pease, and Korn-Lambiotte, are also presented. These algorithms make use of the additive structure of residue class rings of integers. The Good-Thomas algorithm is described as an exam-

ple of an algorithm based on the multiplicative structure of such rings. Two chapters are devoted to linear and cyclic convolution. Linear convolution is identified with polynomial multiplication, and cyclic convolution is identified with polynomial multiplication modulo a binomial $x^N - 1$. The basic tools for the design of convolution algorithms are the convolution theorem and the Chinese remainder theorem. The Cook-Toom algorithm, the Winograd small convolution algorithm, and the Agarwal-Cooley algorithm are the principal convolution algorithms that are discussed. Convolution algorithms can be applied to the design of multiplicative Fourier transform algorithms. This was first observed by Rader who showed that for a prime p, a p-point Fourier transform can be computed by a $(p - 1)$-point cyclic convolution. This principle can be extended to N-point Fourier transforms for various values of N, such as prime powers N or squarefree N. The duality between periodic and decimated data established by the Fourier transform is studied and applied to the computation of N-point Fourier transforms, where N is an odd prime power. Fourier transforms of multiplicative characters $\bmod N$ and orthogonal bases diagonalizing Fourier transform matrices are also discussed.

The underlying idea of the book, namely to present the algorithms for finite harmonic analysis from the systematic viewpoint of matrix factorizations, is certainly attractive. However, the execution of this idea leaves a lot to be desired. The writing is repetitious and tedious, and the approach is numbingly slow. For instance, an algorithm might first be presented for $N = 15$, then for N a product of two distinct primes, then for N a product of three distinct primes, and finally for the general case of squarefree N, always with the same arguments and usually with the same wording. The book could easily be streamlined to half its size without any loss of information. A further irritant is the sloppy attitude that pervades the book. On p. 79 the authors speak of the "following diagram" and on pp. 271 and 274 of the " following table", but the reader will look in vain for any of these. On p. 279, Problems 6 and 7 refer to tables that do not exist. The definition of a subfield on p. 23 is worth repeating: "F is a subfield of K in the sense that F is a subset of K containing 1, and it is closed under the addition and multiplication in K". According to this "definition", the natural numbers would form a subfield of the reals! The distinction between "relatively prime" and "pairwise relatively prime" is often not made. In the statement of the unique factorization theorem for polynomials on p. 22 one has to assume that the irreducible factors are monic in order to guarantee uniqueness up to the ordering of factors. On p. 16 the authors speak of the "Euler quotient function" instead of the Euler totient function. There are many misprints and even some systematic spelling errors. No serious attempt is made to assign proper credit for the results that are presented. For instance, with respect to the construction of orthogonal bases diagonalizing Fourier transform matrices only papers from the 1980's are quoted, although it is well known that this was first achieved by McClellan and Parks [2] and that the basic ideas go back to Schur [4]. Overall, the careless and repetitious style takes away a lot of

the pleasure one might have had in reading this book.

H. N.

1. R. E. Blahut, *Fast algorithms for digital signal processing*, Addison-Wesley, Reading, MA, 1985.
2. J. H. McClellan and T. W. Parks, *Eigenvalue and eigenvector decomposition of the discrete Fourier transform*, IEEE Trans. Audio Electroacoust. **20** (1972), 66–74.
3. H. J. Nussbaumer, *Fast Fourier transform and convolution algorithms*, Springer, Berlin, 1981.
4. I. Schur, *Über die Gaußschen Summen*, Göttinger Nachr. **1921**, 147–153.

26[60-02, 49D37, 60G15, 60G17, 60G35, 60G60, 62C10].—Jonas Mockus, *Bayesian Approach to Global Optimization—Theory and Applications*, Mathematics and Its Applications (Soviet Series), Kluwer Academic Publishers, Dordrecht, 1989, xiv + 254 pp., 24 $\frac{1}{2}$ cm. Price $59.00 /Dfl 190.00.

In the Bayesian approach to global optimization an objective function f is a priori modelled as a realization of a stochastic process (also called random function), which can be viewed as a probability distribution on a class of functions. The objective function is evaluated in certain points, and the posterior stochastic process is computed, conditional on the observed function values. The posterior information is used to determine the location of the next point where f will be evaluated.

The stochastic processes are taken to be Gaussian. A Gaussian process is characterized by a mean and covariance function. The covariance function specifies how the correlation of the function values $f(x)$ and $f(y)$ depends on the originals x and y. Gaussian stochastic processes have the very attractive property that the posterior process, conditional on a number of observed function values, is Gaussian as well. However, the determination of the posterior mean and covariance function involves the time-consuming operation of the inversion of a matrix, whose size is equal to the number of observations.

On page 12 the author argues that the Bayesian approach and a stochastic model of f were first applied to global optimization in his reference of 1963. However, H. J. Kushner already in 1962 published a paper on this subject [1].

The author first defines an optimal n-step optimization strategy which minimizes the expected deviation from the global optimum. The strategy can be computed by solving n Bellman equations. However, it is well known that serious problems arise from a practical point of view, even for moderate values of n. Therefore, a one-step approximation is introduced, where the next observation is considered as the last one. Even this approximation, however, is hard to implement, since at each step the matrix referred to above has to be inverted, and since the next observation point is the optimum value of a function, say ϕ, which also will have different local optima. Several simplifications are carried out to achieve a method which is able to process a reasonable number of function evaluations. Unfortunately, these simplifications all ruin

the consistency of the model. It also has to be remarked that no exact result is known about the distribution of the global maximum of a Gaussian random function. This means that, unless the number of function evaluations is fixed in advance, the question of how to evaluate the error of an approximation to the global optimum cannot be answered satisfactorily.

An outline of the book is as follows. Chapter 1 contains a discussion of the main advantages of the Bayesian approach. Chapter 2 presents a general definition of Bayesian methods of global optimization. Chapter 3 contains an axiomatic justification of the Bayesian approach. In Chapter 4 the Gaussian class of prior random functions is derived from the conditions of homogeneity, independence of partial differences, and continuity of sample functions. Chapter 5 provides the expressions for the one-step approximation of the dynamic programming equations. This chapter also discusses the replacement of the Kolmogorov consistency conditions by the weaker condition of the risk function continuity. Chapter 6 discusses methods to reduce the dimensionality of global optimization problems. In Chapter 7 the Bayesian approach is applied to find local optima of objective functions with noise. In Chapter 8 a number of real-life applications is described. Chapter 9 provides a description of the portable FORTRAN package which is contained in the book.

Although the Bayesian approach to global optimization, in my opinion, did not yet yield efficient algorithms which are fully theoretically justified, this new book shows that the approach is very appealing, and that the approximations work well. Also, the book contains all the relevant theorems, proofs, and computer programs. Hence, although the book is not very clearly written, and contains very many typos (7 in the preface), it can serve well for investigators who want to pursue the approach. In addition, the programs of the methods, which are based on approximations, can be used by practitioners to solve real-life problems.

C. G. E. BOENDER

Econometric Institute
Erasmus University Rotterdam
3000 DR Rotterdam, The Netherlands

1. H. J. Kushner, *A versatile stochastic model of a function of unknown and time varying form*, J. Math. Anal. Appl. **5** (1962), 150–167.

27[33–04, 65D20].—UNITED LABORATORIES, INC., *Mathematical Function Library for Microsoft-FORTRAN*, Wiley, New York, 1989, xvii + 341 pp., $25 \frac{1}{2}$ cm., loose leaflets in 3-hole-punched binder, including three $5\frac{1}{4}''$ diskettes. Price $295.00.

In 1964 the National Bureau of Standards (now the National Institute of Standards and Technology) issued a massive handbook of formulas, graphs and numerical tables of the elementary mathematical functions and the so-called

higher transcendental functions or special functions of mathematical physics [6]. This N.B.S. Handbook immediately filled, and continues to fill, a tremendous need in scientific work; according to the Science Citation Index (published by the Institute for Scientific Information, Inc., Philadelphia, PA), the current rate at which it is cited in the mathematical and scientific literature is of the order of 1,300 entries per year.

The need for the numerical tables of the elementary functions in the N.B.S. Handbook has now largely disappeared: library routines for generating exponentials, logarithms, and trigonometric and hyperbolic functions are available in all major scientific software libraries as well as being required by the FORTRAN Standard [1]. Also, these routines are incorporated in hand-held calculators designed for scientific calculations. The loose-leaf manual and accompanying diskettes under review, which we shall refer to as the "UL Library", may be regarded as an attempt to replace the numerical tables of the higher transcendental functions supplied in the N.B.S. Handbook by a comprehensive software package. The functions treated include Bessel and related functions, hypergeometric and confluent hypergeometric functions, elliptic functions and integrals, exponential integral and related functions, error function and related functions, Gamma and incomplete Gamma functions, orthogonal polynomials, probability functions and random number generators. This list is not quite isomorphic with the list of functions tabulated in the N.B.S. Handbook; among the omissions are parabolic cylinder functions, Mathieu functions, spheroidal wave functions and the Riemann Zeta function. The UL Library is designed to be used on Personal Computers of IBM type and equipped with a Microsoft FORTRAN 77 Compiler. For efficiency, a numeric co-processor is recommended.[1] The operating precision is IEEE double precision (53 bits in the floating-point mantissa), but the accuracy of the computed function values is generally less, sometimes well below single precision.

The project is an extremely ambitious enterprise, especially as it appears that all of the programs for the library routines have been constructed *ab initio*. For such a project to be completed successfully its authors need to have a thorough knowledge of, and experience in, several areas of classical and numerical analysis, including analytic properties of the higher transcendental functions, asymptotic analysis, approximation theory, and error and stability analyses. How well have the present authors (who are nameless) succeeded?

A comprehensive answer to the question just posed would necessitate a tremendous amount of numerical testing. We concentrated our testing on just a few functions with which we have had previous software experience, namely Airy, Bessel, hypergeometric, confluent hypergeometric and Legendre functions. Usually the UL Library performed in accordance with the specifications for each routine. However, we studied the documentations in detail, looking for major

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[1] The library diskettes provide FORTRAN programs, which could be compiled and used on a Personal Computer without a co-processor, as well as coprocessor assembly code for each subroutine. For the purposes of this review, attention is restricted to the co-processor assembly code.

ways in which the algorithms used might go astray. Unfortunately, we were successful.

The first, and very serious, type of failure arise when the UL Library delivers completely wrong answers without any error message. This occurs with the routine BESJR in §8.8, which is designed to generate the Bessel function $J_\nu(x)$ for real values of the argument x and order ν. For example, with $x = 9\pi$ and $\nu = 2\frac{1}{2}$, $4\frac{1}{2}$, $6\frac{1}{2}$, $8\frac{1}{2}$, $10\frac{1}{2}$, BESJR yields the following values of $J_\nu(x)$ to ten decimal places:

$$-0.00670\ 05817,\ 0.03764\ 71379,\ -0.05418\ 00781,$$
$$0.31114\ 25305,\ -0.06567\ 89923.$$

These should be compared with the correct values[2]

$$0.01592\ 10880,\quad -0.05237\ 32560,\quad 0.10316\ 94874,$$
$$-0.14727\ 18330,\quad 0.14363\ 19977.$$

Not only are the numerical values totally incorrect, all the signs are wrong, too. Similar gross errors occur for $\nu = 22\frac{1}{2}(2)50\frac{1}{2}$, for example. On the other hand, BESJR generates correct values (within the prescribed error tolerance) for the intermediate values $\nu = 12\frac{1}{2}(2)20\frac{1}{2}$, and also for $\nu = \frac{1}{2}$ and $1\frac{1}{2}(2)49\frac{1}{2}$, x again being 9π. The reason for the errors appears to be that J. C. Miller's backward recurrence algorithm has been used, with the trial values normalized on the value of a single Bessel function, in fact $J_{1/2}(9\pi)$. Since $J_{1/2}(9\pi)$ is zero, this procedure is bound to lead to meaningless answers. Yet this cannot be the entire explanation, otherwise all values in the range $\nu = 21\frac{1}{2}(1)50\frac{1}{2}$ would be incorrect and not merely alternate ones. Thus the documentation must be in error, too. And there may be further inaccuracies here. For example, according to the documentation the value for $\nu = 2\frac{1}{2}$ is computed from the power series expansion of $J_\nu(x)$: either the Miller algorithm was used instead, or, perhaps less likely, the power series was summed incorrectly.

Similar errors occur for numerous other values of ν, both integer and non-integer, that we tested. For example, with $x = 8.6537279129\ldots$ (the third positive zero of $J_0(x)$), BESJR generates accurate values for $\nu = 0, 1, 2, 3$ and $7(1)50$, but grossly inaccurate values for $\nu = 4, 5, 6$. Furthermore, erroneous values of $J_\nu(x)$ are generated when the value of x is merely moderately close to one of the critical values, the magnitudes of the errors being inversely proportional to the distances of x from the critical value.

A companion routine to BESJR is BESYR (§8.10), which is designed to generate the Bessel function $Y_\nu(x)$ for real values of x and ν. Since one of the algorithms used in BESYR draws upon values of $J_\nu(x)$, we expected—and found—some difficulties. For example, BESYR computed $Y_\nu(9\pi)$ with wrong signs and incorrect numerical values for $\nu = -50\frac{1}{2}(2)-22\frac{1}{2}$ and $-10\frac{1}{2}(2)-2\frac{1}{2}$,

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[2]Obtained by use of D. E. Amos' package [2]. See also comments made below on validation.

while the results were correct at intermediate values of ν. (This similarity of behavior with BESJR reflects the identity $Y_\nu(x) = (-1)^n J_{-\nu}(x)$ when $\nu = n - \frac{1}{2}$, n being an integer.)

A second type of failure is the generation of results which, while not completely inaccurate, contain errors greatly in excess of the accuracy claimed in the documentation. One example is provided by the routine BESYR mentioned in the preceding paragraph. The documentation claims at least 12-digit accuracy when the order ν is an integer. In itself this claim is careless because it takes no account of the inevitable loss of relative precision in the neighborhoods of zeros. But the situation is actually much worse. With $\nu = 119$ and 120, for example, we found that in the range $x = 848(0.1)852$ the accuracy often fell to 8 or 9 digits, even well away from the zeros of $Y_{119}(x)$ and $Y_{120}(x)$. Another example is provided by the routines AIRYA and AIRYAD (§§8.1 and 8.3) for the Airy function $\mathrm{Ai}(x)$ and its first derivative. In both routines at least 13-digit accuracy is claimed when $-10^{10} \leq x \leq 100$. But we found many values of x in the (zero-free) range 5.2 to 5.8 for which they yielded only 8 or 9 correct digits.

The third type of failure arises when the UL Library generates an inaccurate value, but the user is warned by an error message such as "unable to compute the ... function with acceptable accuracy" or "numeric overflow in the ... function". However, if these failures occur too frequently, then there will be huge gaps in the effective ranges claimed for the variables. Such is the case with the routines CHGFU (§9.2) for generating the confluent hypergeometric function $U(a, b, x)$, and HPRGMT (§20.1) for generating the hypergeometric function $F(a, b; c; x)$.

CHGFU employs two algorithms. The first is evaluation of the asymptotic expansion of $U(a, b, x)$ for large x, which is quite sound. The second is based on a formula that expresses $U(a, b, x)$ as a difference of two M-type confluent hypergeometric functions, which is unsound because of the potential for massive numerical cancellation. In consequence, although the documentation claims that the effective ranges of the variables are given by

$$-50 \leq a \leq 50, \quad -50 \leq b \leq 50, \quad -100 \leq x \leq 100,$$

with the exclusion of integer values of b and nonpositive integer values of a, extensive regions are inadmissible. These include, for example,

$$\begin{cases} a = 0.5, \ b = 0.5, \ 7.4 \leq x \leq 19.9; \ a = 2.7, \ b = 5.4, \ 10.1 \leq x \leq 437.4; \\ a = 20, \ b = 2.5, \ 0.34 \leq x \leq 20,000. \end{cases}$$

For HPRGMT the documentation states that a and b can have any real values between -10^{10} and 10^{10}, c can have any real value between -10^{20} and 10^{20} (other than a negative integer) and x can have any real value between -10^{10} and 1. But, again, these claims are misleading. For example, when x is

in the interval $[0, 1)$ HPRGMT sums the hypergeometric series

$$F(a, b; c; x) = \sum_{n=0}^{\infty} \frac{a(a+1)\cdots(a+n-1)b(b+1)\cdots(b+n-1)}{c(c+1)\cdots(c+n-1)} \frac{x^n}{n!}.$$

Since the radius of convergence is 1 the algorithm must fail for values of x sufficiently close to 1. Sample intervals of failure were found to be:

$$a = b = 100, \quad c = 1, \quad 0.95 \leq x < 1;$$
$$a = b = 10000, \quad c = 1, \quad 0.0013 \leq x < 1;$$
$$a = b = 10^8, \quad c = 1, \quad 10^{-10} \leq x < 1.$$

Even when failure does not occur, execution can be extremely slow. Thus for $a = b = c = 1$ and $x = 0.99999$, 20 minutes elapsed on an IBM PS2 computer before the answer was produced.

The weaknesses in the algorithms used for $J_\nu(x)$, $Y_\nu(x)$ and $U(a, b, x)$ were avoidable, if only because robust software for generating these functions is already available; see, for example, [3], [8]. We also observe that instead of normalizing the trial values of $J_\nu(x)$ obtained in the Miller algorithm via a single value of $J_\nu(x)$, the identity

$$(\tfrac{1}{2}x)^\nu = \sum_{k=0}^{\infty} \frac{(\nu + 2k)\Gamma(\nu + k)}{k!} J_{\nu+2k}(x), \quad \nu \neq 0, -1, -2, \ldots,$$

[6, equation (9.1.87)], could have been used instead. This is the appropriate generalization of the identity

$$1 = J_0(x) + 2J_2(x) + 2J_4(x) + \cdots$$

that the authors use in a routine BESJ (§8.7) for generating functions of integer order. Again, a stable way of generating $U(a, b, x)$ is backward integration of the confluent hypergeometric equation, with initial values derived from the asymptotic expansions of $U(a, b, x)$ and $\partial U(a, b, x)/\partial x$ for large x.

In addition to occasional poor choices of algorithm, we noticed instances of poor choices of the actual functions being generated. Thus, functions that exhibit exponential growth or decay when the argument x is large would have been better replaced by their logarithms. This would greatly increase the threshold at which overflow or underflow occurs. Such functions include the Gamma function, $\Gamma(x)$, the exponential integral, $\text{Ei}(x)$, the complementary error function, $\text{erfc}\, x$, the Airy functions, $\text{Ai}(x)$ and $\text{Bi}(x)$, the incomplete Gamma function $\Gamma(a, x)$, the modified Bessel functions $I_\nu(x)$ and $K_\nu(x)$, and the confluent hypergeometric function $M(a, b, x)$. In the case of $\Gamma(x)$, a separate routine *is* given for $\ln|\Gamma(x)|$ in §13.3 but since this is constructed simply by taking logarithms of the values obtained by the library routine for $\Gamma(x)$ there is no increase in the overflow threshold. It is the logarithm that should have been generated first!

Troublesome singularities can sometimes be avoided by introduction of appropriate factors. Thus each of the functions $M(a, c, x)$ and $F(a, b; c; x)$ has poles at $c = 0, -1, -2, \cdots$, but both $M(a, c, x)/\Gamma(c)$ and $F(a, b; c; x)/\Gamma(c)$ are entire functions of c; this was pointed out many years ago in [7]. On the other hand, there is a nonsensical statement in §9.2 that $U(a, b, x)$ is not well defined when a and b are negative integers and $|b| \geq |a|$. In fact, as noted in [7, p. 258], $U(a, b, x)$ is entire in a and b. It is the poor algorithm used to compute $U(a, b, x)$ that may fail when a or b is an integer or close to an integer.

We could continue in this vein and we could also provide a substantial list of typographical errors in the manual and operational defects in the software.[3] Instead, let us turn now to the topic of validation. Several articles have been written on the difficulty of checking software for the generation of mathematical functions; see, for example, [3]. How were the UL Library routines validated by the authors? The only clue supplied in the manual appears to be the statement on p. 36 that the least accurate results in the output of an algorithm always occur at the interface with another algorithm. Presumably this means that there has been substantial cross-checking at these interfaces. This is indeed a powerful type of check, but one that does not guard against all types of algorithmic and programming errors, as we have already observed with the routines BESJR and BESYR.

In fairness to the authors there are two subsections (§§6.2, 6.3) in which they encourage users to check output by using identities satisfied by the higher transcendental functions. For example, the error and Bessel functions are both special cases of confluent hypergeometric functions. However, these kinds of identities are rather specialized, and there is always the danger that they may have already been used in constructing the library routine. For example, there are identities that relate the Airy functions $\mathrm{Ai}(x)$, $\mathrm{Bi}(x)$ and their derivatives to Bessel functions and modified Bessel functions of orders $\pm 1/3$, $\pm 2/3$. But if these identities are used as cross-checks, then the inaccuracies we noted above for the routine, AIRYA and AIRYAD may not show up because the same kind of inaccuracy is present in a (parent) routine BESKR (§8.14) for the modified Bessel function $K_\nu(x)$. In contrast, a powerful form of cross-check that is *not* mentioned in the manual, but which is widely applicable, is to employ identities of Wronskian or Casoratian type, for example,

$$\mathrm{Ai}(x)\mathrm{Bi}'(x) - \mathrm{Ai}'(x)\mathrm{Bi}(x) = 1/\pi,$$
$$J_{\nu+1}(x)Y_\nu(x) - J_\nu(x)Y_{\nu+1}(x) = 2/(\pi x).$$

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[3]For example, there is a statement on p. 24 that the computer will halt when a library routine is called and a co-processor is not present. We found that is not always true; furthermore, instead of identifying the absence of a co-processor as the problem, the error messages misleadingly report errors in the library routine.

These were the checks that we used ourselves to detect errors in the UL Library routines as well as to ascertain whether or not the UL Library output is correct in cases of discrepancy with output from other libraries.

Let us summarize our findings. The UL Library provides PC users with a new set of routines for generating an extensive collection of higher transcendental functions, indeed most of the functions tabulated in the N.B.S. Handbook. The library is relatively easy to install and its price is reasonable. It will provide a useful tool for mathematical physicists and other scientists, especially those engaged in calculations of exploratory type. However, in our comparatively small sample we encountered failures of various kinds, including some extremely serious ones. Because of this, we conclude that the authors may have lacked some experience and expertise, or perhaps just the necessary time, for the mammoth task of constructing a robust library of the kind intended. In consequence, users will need to exercise great care with any output from the library, applying independent checks wherever possible. Furthermore, users must also be prepared for disappointments: the viable ranges of a routine may turn out to be a good deal less than is claimed in the documentation, especially in the case of functions that have not been treated by earlier software workers.

For heavy systematic computations many users will find the more robust IMSL and NAG libraries [4], [5] to be preferable. The variety of functions covered in these libraries is not as large, but the viable ranges of the variables are considerably more extensive and the precision is often higher. Moreover, both IMSL and NAG provide many desirable features, such as linear algebra packages, in addition to routines for generating higher transcendental functions.

In preparing this review I am heavily indebted to Dr. Daniel Lozier of the National Institute of Standards and Technology who helped devise and carry out the numerical testing, and also to Sang Chin, a student at Duke University, who assisted.

F.W.J.O.

1. American National Standard Programming Language FORTRAN, ANSI X3., 9–1978, American National Standards Institute, 1430 Broadway, New York, NY 10018.
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3. W. J. Cody, *Software for special functions*, Rend. Sem. Mat. Univ. Politec. Torino, Special issue: *Special functions: theory and computation* (1985), 91–116.
4. IMSL *Library Reference Manual*, IMSL, 7500 Bellaire Boulevard, Houston, TX 77036-5085, Edition 9, June 1982.
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28[17–01, 17–04, 22–01].—JOHAN G. F. BELINFANTE & BERNARD KOLMAN, *A Survey of Lie Groups and Lie Algebras with Applications and Computational Methods*, Classics in Applied Mathematics, Vol. 2, SIAM, Philadelphia, PA, 1989, x + 164 pp., 23 cm. Price $24.50 paperback.

The book under review is the second edition of a text which appeared first as book in 1972 and which is an expanded version of three review papers which appeared in the mid-sixties. It has now appeared as volume 2 of a series with the promising title "Classics in Applied Mathematics". The book consists of three chapters, entitled Lie groups and Lie algebras, Representation theory, Constructive methods, and an extensive bibliography of pre-1972 references.

In the first two chapters the authors explain the notions of Lie groups and Lie algebras and their basic properties. All proofs are omitted. Instead, extensive references are given. As an example of the narrative style we quote the definition of a Lie algebra: "*Abstractly, a Lie algebra L is a vector space equipped with a product $[x, y]$ satisfying certain axioms* (references to Freudenthal-de Vries, Dynkin, Jacobson and van der Waerden). *We shall continue to use the bracket notation for products when we deal with any Lie algebra. One of the axioms for a Lie algebra is that the product $[x, y]$ be bilinear, that is, linear in x and y separately. We also assume that the Lie product is anticommutative, $[x, y] = -[y, x]$. Finally we assume that the Jacobi identity, $[x, [y, z]] + [y, [z, x]] + [z, [x, y]] = 0$, holds for all vectors x, y, z in the Lie algebra.*" As examples, the applications of Lie groups and algebras in classical and quantum mechanics are described in a few pages. The description of classical mechanics looks incomprehensible for readers without a background in physics. Such a background is not necessary to understand that the basic definitions of pure states and transition probabilities in the section on quantum symmetries are wrong.

The third chapter focusses on the calculation of weight multiplicities of irreducible highest weight representations and the decomposition of tensor products. This chapter is without doubt the most interesting. Physicists and applied mathematicians will find here the basic ideas behind the tables and the computer algebra packages which they use for their calculations.

Many books have been written on Lie groups and algebras. Some are called "Classics". The omission of the beautiful proofs of, e.g., the Weyl character formula, and the sloppy treatment of applications, suggest that this book hardly deserves this title. The third chapter, however, is still well worth reading for everyone who occasionally uses Lie algebra representations.

HUIB-JAN IMBENS

Max-Planck-Institut für Mathematik
Gottfried-Claren-Strasse 26
D-5300 Bonn
West Germany