We fix a finite alphabet $A$ and consider languages and grammars with this alphabet.

**Definition 1.** A formal grammar $G$ with prohibition consists of rules that are divided into two parts: positive $PG$ and negative $NG$. These rules generate in a conventional manner, i.e., by derivation or recursive inference, two languages $L(PG)$ and $L(NG)$.

**Definition 2.** The language $L(G)$ of the grammar $G$ with prohibition is equal to $L(PG) \setminus L(NG)$.

**Definition 3.** A language $L$ is recursive/recursive if both $L(PG)$ and $L(NG)$ are recursive.

**Theorem 1.** A language $L$ is recursive/recursive if and only if it is recursive.

**Remark.** This property is not always true because in general grammars with prohibition are more powerful than conventional formal grammars from a given class. For instance, enumerable/enumerable language is not always enumerable. There are recursive languages with arbitrarily high derivational complexity $K$ where $K$ can be time $T$, space $S$, dispersion $D$ or index $I$ [A.V. Gladkii, Formal Grammars and Languages, Nauka, 1973].

**Theorem 2.** For any recursive language $L$, there is grammar $G$ with prohibition such that $L = L(G)$ and $T(G, n) = O(n), S(G, n) = n, D(G, n) = 1$, and $I(G, n) = 1$. (Received January 18, 2005)