

# DUALITY FOR BOOLEAN EQUALITIES

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**Introduction.** Boolean equalities occur naturally in Logic and especially in the theory of polyadic algebras. For instance, they have been used in [2] under the name of *Champ Logique*. The purpose of this paper is to develop a duality theory for Boolean equalities and by doing so establish a representation theory for them.

**0. Notation.** The basic notation for Boolean algebras is the standard one (for instance, see [1]). Throughout this paper, the letter  $O$  will denote the two-element Boolean algebra; whenever a topology is assumed to be defined on  $O$ , then this topology is the discrete one. If  $B$  is a Boolean algebra, then a *valuation* of  $B$  is a homomorphism from  $B$  onto  $O$ . We make precise the terminology we shall be using in connection with the Stone duality theory for Boolean algebras. A *Boolean space* is a compact Hausdorff space whose topology is generated by the clopen sets. If  $X$  is a Boolean space, then the *dual algebra* of  $X$  is the Boolean algebra of clopen sets of  $X$ . If  $B$  is a Boolean algebra, then the *dual space* of  $B$  is the Boolean space  $X$  of all valuations of  $B$  (the topology on  $X$  is the topology generated by all sets of the form  $\{v: vp = 1\}$  where  $p$  belongs to  $B$ ).

**1. Definition and generalities.** A *Boolean equality* on a set  $\Delta$  is a pair  $(B, E)$  where  $B$  is a Boolean algebra and  $E$  is a mapping from  $\Delta \times \Delta$  into  $B$  so that, for all  $s, t, u$  in  $\Delta$ ,

$$(1.1) \quad E(s, s) = 1,$$

$$(1.2) \quad E(s, t) = E(t, s),$$

$$(1.3) \quad E(s, u) \wedge E(u, t) \leq E(s, t),$$

and

$$(1.4) \quad B \text{ is generated by the range of } E, \text{ i.e., by } E(\Delta \times \Delta).$$

In the sequel, it will be convenient to view an *equivalence relation* on  $\Delta$  as a mapping  $\phi$  from  $\Delta \times \Delta$  into  $O$  so that  $\phi$  satisfies (1.1), (1.2) and (1.3). Clearly, if  $\phi$  is an equivalence relation on  $\Delta$ , then  $(O, \phi)$  is a Boolean equality on  $\Delta$ . Suppose  $(B, E)$  and  $(\tilde{B}, \tilde{E})$  are Boolean equalities on a set  $\Delta$ . A homomorphism  $f$  from  $B$  into  $\tilde{B}$  shall be called an *equality homomorphism* if  $fE = \tilde{E}$ , i.e., if  $fE(s, t) = \tilde{E}(s, t)$  for all  $s$  and  $t$  in  $\Delta$ . It is easy to see that an equality homomorphism from  $B$  into  $\tilde{B}$  is necessarily onto. The Boolean equalities  $(B, E)$  and  $(\tilde{B}, \tilde{E})$

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are said to be *isomorphic* if there exists an equality isomorphism from  $B$  onto  $\bar{B}$ . If  $M$  is an ideal in  $B$ ,  $\bar{B} = B/M$ ,  $f$  is the natural projection from  $B$  onto  $\bar{B}$  and  $\bar{E}(s, t) = fE(s, t)$  for all  $s, t$  in  $\Delta$ , then  $(\bar{B}, \bar{E})$  is a Boolean equality on  $\Delta$ ;  $(\bar{B}, \bar{E})$  shall be called a *quotient* of  $(B, E)$ .

**2. Duality.** Let  $\Delta$  be a nonempty set. In this section, we shall establish a one-to-one correspondence between Boolean equalities on  $\Delta$  and certain sets of equivalence relations on  $\Delta$ . As we shall see, this correspondence obeys all the laws of a genuine duality theory. We begin by introducing a topology on the set  $\Phi$  of all equivalence relations on  $\Delta$ . Note first that the set of all functions from  $\Delta \times \Delta$  into  $O$  with the product topology is a Boolean space. The topology induced on  $\Phi$  shall also be referred to as the *product topology*. This topology is generated by the sets of the form  $\{\phi: \phi \in \Phi, \phi(s, t) = 1\}$  together with their complements. We have the following basic fact.

(2.1) LEMMA. *The set  $\Phi$  of all equivalence relations on  $\Delta$  together with the product topology is a Boolean space.*

PROOF. It suffices to show that  $\Phi$  is closed in the Boolean space of all functions from  $\Delta \times \Delta$  into  $O$ . Let  $\Phi_1, \Phi_2$  and  $\Phi_3$  be the sets of functions from  $\Delta \times \Delta$  into  $O$  satisfying (1.1), (1.2) and (1.3) respectively. Clearly  $\Phi = \Phi_1 \cap \Phi_2 \cap \Phi_3$ . It suffices to show that  $\Phi_1, \Phi_2$  and  $\Phi_3$  are closed. For instance,  $\Phi_3$  is the intersection of all sets of the form  $\{\phi: \phi(s, u) = 0\} \cup \{\phi: \phi(u, t) = 0\} \cup \{\phi: \phi(s, t) = 1\}$  as  $s, u$  and  $t$  run over  $\Delta$ ; since these sets are closed, it follows that  $\Phi_3$  is closed. The proof of the fact that  $\Phi_1$  and  $\Phi_2$  are closed is similar.

We are now in a position to make more precise the comment made at the beginning of this section. We shall establish a one-to-one correspondence between nonempty compact subspaces of  $\Phi$  and Boolean equalities on  $\Delta$ . First, we define the dual Boolean equality of a nonempty compact subspace  $\Psi$  of  $\Phi$ . The *dual Boolean equality* of  $\Psi$  is the pair  $(B, E)$  where  $B$  is the dual Boolean algebra of  $\Psi$  and  $E$  is defined by

$$(2.2) \quad E(s, t) = \{\phi: \phi \in \Psi, \phi(s, t) = 1\} \quad \text{for all } s$$

and  $t$  in  $\Delta$ . It is easy to check that  $(B, E)$  is indeed a Boolean equality on  $\Delta$ . We now define the dual equivalence relation space of a Boolean equality  $(B, E)$ . The *dual equivalence relation space* of  $(B, E)$  is a subspace  $\Psi$  of  $\Phi$  so that  $\phi \in \Psi$  if and only if there exists a valuation  $v$  of  $B$  for which

$$(2.3) \quad \phi(s, t) = vE(s, t)$$

for all  $s$  and  $t$  in  $\Delta$ . Note that if  $v$  is a valuation of  $B$  and  $\phi$  is defined by (2.3), then  $\phi \in \Phi$  and therefore  $\phi \in \Psi$ . We turn now to the proof of the fact that  $\Psi$  is a Boolean space.

(2.4) THEOREM. *If  $\Psi$  is the dual equivalence relation space of a Boolean equality  $(B, E)$ , then  $\Psi$  is homeomorphic with the dual space of  $B$ ; in particular, this implies that  $\Psi$  is a Boolean space.*

PROOF. Let  $X$  be the dual space of  $B$  and define a one-to-one mapping  $\chi$  from  $X$  onto  $\Psi$  by

$$(2.5) \quad \chi(v)(s, t) = vE(s, t)$$

for all valuations  $v$  in  $X$  and all  $s, t$  in  $\Delta$ . Now observe that the topology on  $X$  is generated by the sets of the form  $\{v: vE(s, t) = 1\}$  together with their complements as  $s$  and  $t$  run over  $\Delta$ . The proof of the fact that  $\chi$  is a homeomorphism is then completed after observing that the following identity holds:

$$\chi\{v: vE(s, t) = 1\} = \{\phi: \phi \in \Psi, \phi(s, t) = 1\}$$

for all  $s, t$  in  $\Delta$ . Since  $X$  is a Boolean space, it follows that  $\Psi$  is a Boolean space.

We now show that the correspondence we have just established obeys the laws of duality.

(2.6) THEOREM. *If  $(B, E)$  is a Boolean equality on  $\Delta$  and  $\Psi$  is the dual equivalence relation space of  $(B, E)$ , then the dual Boolean equality of  $\Psi$  is isomorphic with  $(B, E)$ ; if  $\Psi$  is a Boolean subspace of  $\Phi$  and  $(B, E)$  is the dual Boolean equality of  $\Psi$ , then the dual equivalence relation space of  $(B, E)$  is  $\Psi$ .*

PROOF. To prove the first part of the theorem, let  $(\tilde{B}, \tilde{E})$  be the dual Boolean equality of  $\Psi$  and let  $X$  be the dual space of  $B$ . Let  $\chi$  be the natural mapping from  $X$  onto  $\Psi$  defined by (2.5). If  $C$  is the dual algebra of  $X$ ,  $f$  the natural isomorphism from  $B$  onto  $C$ ,  $\chi_0$  the isomorphism from  $C$  onto  $\tilde{B}$  induced by  $\chi$ , then  $\chi_0 f$  is an isomorphism from  $B$  onto  $\tilde{B}$  so that  $(\chi_0 f)E(s, t) = \tilde{E}(s, t)$  for all  $s$  and  $t$  in  $\Delta$ . It follows that  $(B, E)$  and  $(\tilde{B}, \tilde{E})$  are isomorphic. To prove the second part of the theorem, assume  $\Psi$  is a Boolean subspace of  $\Phi$  and  $(B, E)$  is the dual Boolean equality of  $\Psi$ . Let  $X$  be the dual space of  $B$  and let  $g$  be the natural homeomorphism from  $\Psi$  onto  $X$ , i.e., if  $\phi \in \Psi$ , then  $g(\phi)$  is the unique valuation of  $B$  so that, for all  $p$  in  $B$ ,  $g(\phi)p = 1$  if and only if  $\phi$  belongs to the clopen set  $p$ . We have, for every  $\phi$  in  $\Psi$  and every  $s, t$  in  $\Delta$ ,  $g(\phi)E(s, t) = \phi(s, t)$  which shows, by (2.3), that  $\Psi$  is the dual equivalence relation space of  $(B, E)$ .

(2.7) COROLLARY. *Two Boolean equalities on a set  $\Delta$  are isomorphic if and only if they have the same dual equivalence relation space.*

Our next and final task in this section is to find necessary and sufficient conditions (in term of dual spaces) for a Boolean equality to be a quotient of another one.

(2.8) THEOREM. *Suppose  $(B, E)$  and  $(\tilde{B}, \tilde{E})$  are Boolean equalities on  $\Delta$  and suppose  $\Psi$  and  $\tilde{\Psi}$  are their respective dual equivalence relation spaces. Then  $(\tilde{B}, \tilde{E})$  is a quotient of  $(B, E)$  if and only if  $\tilde{\Psi}$  is a subset of  $\Psi$ .*

PROOF. Suppose there exists an equality homomorphism  $f$  from  $B$  onto  $\tilde{B}$ . Let  $\phi$  in  $\tilde{\Psi}$  and let  $\bar{v}$  be the unique valuation of  $\tilde{B}$  so that  $\phi(s, t) = \bar{v}\tilde{E}(s, t)$  for all  $s, t$  in  $\Delta$ . Let  $v = \bar{v}f$ . Then  $v$  is a valuation of  $B$  and  $vE(s, t) = \phi(s, t)$  which shows that  $\phi \in \Psi$ . Hence  $\tilde{\Psi}$  is a subset of  $\Psi$ . Conversely, assume  $\tilde{\Psi}$  is a subset of  $\Psi$ . In view of the preceding theorem we may and do assume that  $(\tilde{B}, \tilde{E})$  and  $(B, E)$  are the dual Boolean equalities of  $\tilde{\Psi}$  and  $\Psi$  respectively. Let  $g$  be the inclusion mapping of  $\tilde{\Psi}$  into  $\Psi$ ; since  $g$  is continuous, it induces a Boolean homomorphism  $\bar{g}$  from  $B$  onto  $\tilde{B}$ . But

$$g^{-1}\{\phi: \phi \in \Psi, \phi(s, t) = 1\} = \{\phi: \phi \in \tilde{\Psi}, \phi(s, t) = 1\}$$

for all  $s$  and  $t$  in  $\Delta$  which is the same as saying that  $\bar{g}E(s, t) = \tilde{E}(s, t)$  for all  $s, t$  in  $\Delta$ . Hence  $\bar{g}$  is an equality homomorphism from  $B$  onto  $\tilde{B}$ , and that completes the proof of the theorem.

(2.9) COROLLARY. *There exists a free Boolean equality on  $\Delta$  i.e., there exists a unique Boolean equality  $(B, E)$  on  $\Delta$  having the property that any other Boolean equality on  $\Delta$  is a quotient of  $(B, E)$ .*

PROOF. The dual Boolean equality of the space  $\Phi$  of all equivalence relations on  $\Delta$  is the required Boolean equality.

**3. Boolean equalities and transformations.** In this section, as always,  $\Delta$  shall be a fixed nonempty set and  $\Phi$  shall be the Boolean space of all equivalence relations on  $\Delta$ . A function from  $\Delta$  into  $\Delta$  shall be called a *transformation* (on  $\Delta$ ). The purpose of this section is to study the effect of transformations on Boolean equalities. We begin by extending transformations  $\eta$  on  $\Delta$  to transformations  $\bar{\eta}$  on  $\Phi$  by

$$(3.1) \quad \bar{\eta}(\phi)(s, t) = \phi(\eta s, \eta t)$$

for all  $s, t$  in  $\Delta$ . The mapping  $\eta \rightarrow \bar{\eta}$  sends the identity onto the identity and  $\eta\xi \rightarrow \bar{\xi}\bar{\eta}$  for all transformations  $\eta$  and  $\xi$  on  $\Delta$ . Moreover,

$$(3.2) \quad \bar{\eta}^{-1}\{\phi: \phi \in \Phi, \phi(s, t) = 1\} = \{\phi: \phi \in \Phi, \phi(\eta s, \eta t) = 1\}$$

for all  $s, t$  in  $\Delta$ . This implies that  $\bar{\eta}$  is continuous on  $\Phi$  for all transformations  $\eta$  on  $\Delta$ . Suppose  $\Psi$  is a subset of  $\Phi$  and  $G$  is a set of transformations on  $\Delta$ . We shall say that  $\Psi$  is *invariant* under  $G$  or that  $G$  *leaves  $\Psi$  invariant* if  $\bar{\eta}(\Psi) \subseteq \Psi$  for all  $\eta$  in  $G$ . This last concept allows us to formulate the main theorem of this section.

(3.3) **THEOREM.** *Suppose  $(B, E)$  is a Boolean equality on  $\Delta$  and suppose  $\eta$  is a transformation on  $\Delta$ . Then there exists a (necessarily unique) endomorphism  $S(\eta)$  of  $B$  so that*

$$(3.4) \quad S(\eta)E(s, t) = E(\eta s, \eta t)$$

*for all  $s, t$  in  $\Delta$ , if and only if the dual equivalence relation space of  $(B, E)$  is invariant under  $\eta$ .*

**PROOF.** We may as well assume that  $(B, E)$  is the dual Boolean equality of some Boolean subspace  $\Psi$  of  $\Phi$ . Assume first that  $\Psi$  is invariant under  $\eta$  i.e.,  $\bar{\eta}(\Psi) \subseteq \Psi$ . Since  $\bar{\eta}$  is continuous, it follows from the Stone duality theory that  $\bar{\eta}$  induces an endomorphism  $S(\eta)$  of  $B$ . By definition,  $S(\eta)p = \bar{\eta}^{-1}p$  for every  $p$  in  $B$  (remember that the elements of  $B$  are the clopen subsets of  $\Psi$ ). Therefore, for all  $s$  and  $t$  in  $\Delta$ ,  $S(\eta)E(s, t) = E(\eta s, \eta t)$  which completes one part of the theorem. Conversely, assume there exists a (necessarily unique) endomorphism  $S(\eta)$  of  $B$  so that  $S(\eta)E(s, t) = E(\eta s, \eta t)$  for all  $s$  and  $t$  in  $\Delta$ . Let  $\chi$  be the mapping defined by (2.5). If  $\phi \in \Phi$  and  $v$  is a valuation of  $B$  so that  $\chi(v) = \phi$ , then  $\chi(vS(\eta)) = \bar{\eta}\phi$  which shows that  $\bar{\eta}\phi \in \Psi$  and therefore  $\Psi$  is invariant under  $\eta$ .

(3.5) **COROLLARY.** *Suppose  $(B, E)$  is a Boolean equality on  $\Delta$ ,  $\Psi$  the dual equivalence relation space of  $(B, E)$  and  $G$  the semi-group of transformations on  $\Delta$  which leave  $\Psi$  invariant. Then there exists a unique mapping  $S$  from  $G$  into endomorphisms  $S(\eta)$  of  $B$  so that  $S(\eta)E(s, t) = E(\eta s, \eta t)$  for all  $\eta$  in  $G$  and all  $s, t$  in  $\Delta$ . Moreover, the mapping  $S$  has the following two properties: (3.6) If  $\delta$  is the identity on  $\Delta$ , then  $S(\delta)$  is the identity on  $B$ ; (3.7) If  $\eta$  and  $\xi$  are in  $G$ , then  $S(\eta\xi) = S(\eta)S(\xi)$ .*

**4. Boolean equalities and operations.** If  $\Delta$  is a nonempty set and  $n$  is a positive integer, then a mapping from  $\Delta^n$  into  $\Delta$  shall be called an  $n$ -operation or simply an operation on  $\Delta$ . The purpose of this section is to establish the connection between operations on  $\Delta$  and Boolean equalities on  $\Delta$ .

Suppose  $(B, E)$  is a Boolean equality on  $\Delta$  and  $T$  is an  $n$ -operation. Then  $(B, E)$  is said to be *compatible* with  $T$  if

$$E(s_1, t_1) \wedge \cdots \wedge E(s_n, t_n) \leq E(T(s_1, \cdots, s_n), T(t_1, \cdots, t_n))$$

for all  $s_1, \cdots, s_n, t_1, \cdots, t_n$  in  $\Delta$ ;  $E$  itself is said to be compatible with  $T$  if  $(B, E)$  is compatible with  $T$ .

It is easy to see that if  $T$  is an operation on  $\Delta$ , then the space of equivalence relations compatible with  $T$  is compact. This, combined with the following theorem, implies that there exists a free Boolean equality on  $\Delta$  which is compatible with  $T$  (i.e., free with respect to the property of being compatible with  $T$ ).

(4.1) THEOREM. *Suppose  $(B, E)$  is a Boolean equality on a set  $\Delta$  and suppose  $\Psi$  is the dual equivalence relation space of  $(B, E)$ . If  $T$  is an operation on  $\Delta$ , then  $(B, E)$  is compatible with  $T$  if and only if every equivalence relation  $\phi$  in  $\Psi$  is compatible with  $T$ .*

PROOF. Using (2.2) and (2.3), the proof reduces to a straightforward verification.

We conclude by making a few remarks concerning Boolean equalities on familiar algebraic structures. Suppose, for instance,  $(B, E)$  is a Boolean equality on a ring  $\Delta$  and assume moreover that  $(B, E)$  is compatible with the ring operations. By (4.1), if  $\phi$  is in the dual of  $(B, E)$ , then  $\phi$  corresponds to a uniquely determined ideal in  $\Delta$ . In other words, to every Boolean equality on  $\Delta$  compatible with the ring operations there corresponds a uniquely determined set of ideals in  $\Delta$ . Such sets of ideals are easily characterized.

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