## ON PRINCIPAL SOLUTIONS OF LINEAR DIFFERENTIAL EQUATIONS

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1. **Introduction.** In this paper we investigate the asymptotic behavior of solutions of certain nth order nonhomogeneous linear ordinary differential equations,  $\Omega(y) = \phi$ , near a singular point at  $\infty$ . The class of nth order linear differential operators,  $\Omega$ , treated here consists roughly of those whose coefficients are complex functions, defined and analytic in unbounded sectorial regions, and have asymptotic expansions as  $x \rightarrow \infty$  in terms of real (but not necessarily integral) powers of x and/or functions (called trivial) which are of smaller rate of growth  $(\prec)$  than all powers of x as  $x \to \infty$ . (We are using here the concept of asymptotic equivalence ( $\sim$ ) as  $x \to \infty$  and the order relations "\(\sigma\)" introduced in [3, \(\xi\)13]. However, it should be noted (see [3, §128(g)]) that the class of operators treated here includes as a special case those operators where no requirement is imposed except that each coefficient be analytic and have an asymptotic expansion (in the customary sense) of the form  $\sum C_i x^{-\lambda_i}$  with  $\lambda_i$  real and  $\lambda_i \to +\infty$  as  $j \to \infty$ . (A summary of the necessary definitions from [3] appears in §2 below.)

In [5], Strodt showed that if  $\phi$  is a nontrivial analytic function which also possesses, as  $x\to\infty$ , an asymptotic expansion in terms of real powers of x and/or trivial functions, then the equation  $\Omega(y)=\phi$  has at least one solution  $y_0$  which is  $\sim$  to a logarithmic monomial (i.e., a function of the form  $Kx^{\alpha_0}(\log x)^{\alpha_1}(\log \log x)^{\alpha_2} \cdot \cdot \cdot (\log_q x)^{\alpha_q}$  for complex  $K\neq 0$  and real  $\alpha_j$ ) and such that if  $f \prec y_0$ , then  $\Omega(f) \prec \phi$ . (A solution with these two properties is called a *principal solution* in [3, §69] and is clearly of minimal rate of growth at  $\infty$ .)

In this paper we consider the case where  $\phi$  is any function  $\sim$  to a logarithmic monomial, and we show in §3 that the equation  $\Omega(y) = \phi$  always has a principal solution. As a corollary (§5) we apply a result proved in [1, §12] to obtain a representation theorem for those solutions of  $\Omega(y) = \phi$  which are  $\prec$  some power of x. Our method of proving §3 consists of first obtaining a sufficiently close approximate solution by successive integrations of a factored equation, and then using the approximate solution to transform the equation into one in which an exact solution can be obtained using [4, §99].

2. Concepts from [3]. (a) [3, §94]. Let  $-\pi \le a \le b \le \pi$ . For each non-

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negative real-valued function g on (0, (b-a)/2), let V(g) be the union (over  $\delta \in (0, (b-a)/2)$ ) of all sectors  $a+\delta < \arg(x-h(\delta)) < b-\delta$  where  $h(\delta) = g(\delta) \exp(i(a+b)/2)$ . The set of all V(g) (for all choices of g) is denoted F(a, b) and is a filter base which converges to  $\infty$ . A statement is said to hold except in finitely many directions (briefly, e.f.d.) in F(a, b) if there are finitely many points  $r_1 < \cdots < r_q$  in (a, b) such that the statement holds in each of  $F(a, r_1)$ ,  $F(r_1, r_2)$ ,  $\cdots$ ,  $F(r_q, b)$  separately.

- (b) [3, §13]. If f is analytic in some V(g), then  $f \to 0$  in F(a, b) means that for any  $\epsilon > 0$ , there is a  $g_1$  such that  $|f(x)| < \epsilon$  for all  $x \in V(g_1)$ .  $f \prec 1$  means that in addition to  $f \to 0$ , all functions  $\theta_j^k f \to 0$  where  $\theta_j f = x \log x \cdot \cdot \cdot \log_{j-1} x f'$  and where  $\theta_j^k$  is the kth iteration of the operator  $\theta_j$ . Then  $f_1 \prec f_2$ ,  $f_1 \sim f_2$ ,  $f_1 \approx f_2$  mean respectively  $f_1/f_2 \prec 1$ ,  $f_1 f_2 \prec f_2$  and  $f_1 \sim cf_2$  for some constant  $c \neq 0$ . If  $f \sim c$ , we write  $f(\infty) = c$ , while if  $f \prec 1$ , we write  $f(\infty) = 0$ . If  $M = x^{\alpha_0}(\log x)^{\alpha_1} \cdot \cdot \cdot \cdot (\log_r x)^{\alpha_r}$  for some r and M is not constant, then by [3, §28]  $f \prec M$  implies  $f' \prec M'$ . If  $f \approx M$ , then  $\delta_k(f)$  will denote  $\alpha_k$ . If  $j \geq i$ , then  $s_{ji}(\alpha)$  will denote the elementary symmetric function of degree i in  $\alpha$ ,  $\alpha 1$ ,  $\cdots$ ,  $\alpha j + 1$ .
- (c) [3, §49]. A logarithmic domain of rank zero (briefly an  $LD_0$ ) over F(a, b) is a complex vector space E of functions (each analytic in some V(g)) which contains the constants and such that any finite linear combination of elements of E, with coefficients which for some  $q \ge 0$  are functions of the form  $cx^{\alpha_0}(\log x)^{\alpha_1} \cdot \cdot \cdot (\log_q x)^{\alpha_q}$  (for real  $\alpha_j$ ), is either  $\sim$  to a function of this latter form or is trivial.
- 3. The main theorem. Consider the equation  $\Omega(y) = \phi$ , where  $\Omega(y)$  is an nth order linear differential polynomial with coefficients in an  $LD_0$  over F(a, b), and where  $\phi$  is a function which in F(a, b) is  $\sim$  to a logarithmic monomial. If  $\theta$  is the operator  $\theta y = xy'$ ,  $\Omega(y)$  may be written  $\Omega(y) = \sum_{i=0}^n B_i(x)\theta^{iy}$ , where the functions  $B_i$  belong to an  $LD_0$ . We assume  $B_n$  is nontrivial. By dividing the equation  $\Omega(y) = \phi$  through by the highest power of x which is  $\sim$  to a coefficient  $B_i$ , we may assume that for some  $m \ge 0$ ,  $B_m \approx 1$  while for each j,  $B_j < 1$  or  $B_j \approx 1$ . Let  $F(\alpha) = \sum_{i=0}^n B_i(\infty)\alpha^i$ . Let Q be the logarithmic monomial such that  $\phi \sim Q$  and let  $\delta_i(Q) = \sigma_i$  for each j. Define a logarithmic monomial M as follows. If  $F(\sigma_0) \ne 0$ , let  $M = (F(\sigma_0))^{-1}Q$ . If  $\sigma_0$  is a root of F of multiplicity r, then let

$$M = (F^{(r)}(\sigma_0)/r!)^{-1}(s_{rr}(\sigma_1 + r))^{-1}(\log x)^r Q$$
 if  $\sigma_1 \in \{-1, -2, \cdots, -r\}$ , while if  $\sigma_1 \in \{-1, -2, \cdots, -r\}$ , let

$$M = c(\log x)^r(\log_2 x \cdot \cdot \cdot \log_k x)Q$$

where  $k = \min\{j: j \ge 2, \sigma_j \ne -1\}$ , and

$$c = (s_{r,r-1}(\sigma_1 + r))^{-1}(\sigma_k + 1)^{-1}(F^{(r)}(\sigma_0)/r!)^{-1}.$$

Then: (1) The equation  $\Omega(y) = \phi$  possesses at least one solution  $y_0 \sim M$  e.f.d. in F(a, b). (2) If  $y_0$  is a solution of  $\Omega(y) = \phi$  such that  $y_0 \sim M$  in some  $F(a_1, b_1)$ , then for any function f which is  $\forall y_0$  in  $F(a_1, b_1)$ , we have  $\Omega(f) \prec \phi$  in  $F(a_1, b_1)$ . In particular, among all solutions of  $\Omega(y) = \phi$  in  $F(a_1, b_1)$ ,  $y_0$  is of minimal rate of growth at  $\infty$ .

PROOF OF PART (2). We consider  $\Omega(y)-Q$  and apply the algorithm introduced in [3, §66] which produces the set of those logarithmic monomials N (called principal monomials in [3, §67]) such that  $\Omega(N) \sim Q$  and  $\Omega(f) \prec Q$  whenever  $f \prec N$ . For  $\Omega(y) - Q$  we find by applying the algorithm that M is the unique principal monomial. Hence if  $f \prec M$ , then  $\Omega(f) \prec Q$ . Since  $y_0 \sim M$  and  $\phi \sim Q$ , part (2) clearly follows.

The proof of part (1) will be based on a sequence of lemmas and will be concluded in §6.

- 4. Lemma. Let  $\gamma$  be a complex number and let  $\psi$  be a function which in F(a, b) is  $\sim$  to a logarithmic monomial R. Let  $\delta_i(R) = \lambda_i$ . Define a logarithmic monomial N as follows:
  - (a) If  $\lambda_0 \neq \gamma$ , let  $N = (\lambda_0 \gamma)^{-1} R$ .
- (b) If  $\lambda_0 = \gamma$ , let  $N = (\lambda_q + 1)^{-1} (\log x \cdot \cdot \cdot \log_q x) R$  where  $q = \min \{j: j \ge 1, \lambda_j \ne -1\}$ .

Then in F(a, b), the equation  $xy' - \gamma y = \psi$  has at least one solution  $y^* \sim N$ .

PROOF. Under the change of variable  $y = x^{\lambda_0}z$  and multiplication by  $x^{-\lambda_0}$ , the equation  $xy' - \gamma y = \psi$  becomes

$$(1) xz' + (\lambda_0 - \gamma)z = \psi_0$$

where  $\psi_0 = x^{-\lambda_0} \psi$ .

Let  $N_0 = x^{-\lambda_0}N$ . The proof is divided into three cases.

Case A.  $\operatorname{Re}(\lambda_0 - \gamma) \neq 0$ .

In this case, under  $z = N_0 + N_0 w$  and division by  $(\lambda_0 - \gamma) N_0$ , equation (1) becomes,

(2) 
$$x(\lambda_0 - \gamma)^{-1}w' + f(x)w = g(x),$$

where  $f \sim 1$  (since  $xN_0' \prec N_0$  by a simple calculation) and where  $g \prec 1$  since  $(\lambda_0 - \gamma)N \sim \psi$ . Thus (2) is normal in the sense of [3, §83] with divergence monomial  $(\lambda_0 - \gamma)x^{-1}$ . Since  $d = \text{Re}(\lambda_0 - \gamma) \neq 0$ , it follows

from [3, §111] (when d>0) and [3, §117] (when d<0) that (2) possesses a solution  $w_0 \prec 1$  in F(a, b). Then clearly  $y^* = x^{\lambda_0}(N_0 + N_0 w_0)$  is  $\sim N$  and satisfies the equation  $xy' - \gamma y = \psi$ .

Case B.  $\lambda_0 = \gamma$ .

Thus (1) is of the form  $z' = x^{-1}\psi_0$ . With N as defined in (b) above, it is proved in [2, p. 272] that for some constant A,  $z_0 = A + \int_{x_0}^x x^{-1}\psi_0$  is  $\sim N_0$  in F(a, b). Hence if  $y^* = x^{\lambda_0}z_0$ , then  $y^*$  satisfies the conclusion. Case C.  $\text{Re}(\lambda_0 - \gamma) = 0$  and  $\lambda_0 \neq \gamma$ .

In this case, (1) may be written  $xz'-\sigma iz=\psi_0$  where  $\sigma=i(\lambda_0-\gamma)$  is a nonzero real number. Under  $z=-(\sigma i)^{-1}\psi_0+w$ , this becomes  $xw'-\sigma iw=\psi_1$  where  $\psi_1=(\sigma i)^{-1}x\psi_0'$ . Since  $\psi_0 < (\log x)^{\lambda_1+1/2}$ ,  $\psi_1 < (\log x)^{\lambda_1-1/2}$  by §2(b), and hence  $\psi_1 < \psi_0$  since  $(\log x)^{\lambda_1-\epsilon} < \psi_0$  for all  $\epsilon>0$ . Under  $w=-(\sigma i)^{-1}\psi_1+u$ , we obtain  $xu'-\sigma iu=\psi_2$  where by §2(b),  $\psi_2 < (\log x)^{\lambda_1-3/2}$  (thus  $\psi_2 < \psi_0$  since  $\psi_0 > (\log x)^{\lambda_1-\epsilon}$  for all  $\epsilon>0$ ). Clearly this process can be repeated so as to make the constant term  $< (\log x)^{\alpha}$  for  $\alpha$  as small as desired. Hence there is a function  $f \sim -(\sigma i)^{-1}\psi_0$  in F(a,b) such that under z=f+v, equation (1) becomes

$$xv' - \sigma iv = \phi_1$$

where  $\phi_1$  is chosen so that,

$$\phi_1 \prec (\log x)^{-1-t}$$

where  $t = 1 + \max\{0, -2\lambda_1\}$ .

The technique we now employ to prove the existence of a solution  $v_0 \prec f$  of (3) is similar to the technique used by Strodt in the proof of [6, Section 107].

Let  $E_1 \in F(a, b)$  be such that on  $E_1$ , we have  $|x| \ge 2$  and

$$|\phi_1(x)| \leq (\log |x|)^{-1-t/2}.$$

For x and  $x_1$  in  $E_1$ , let  $B(x, x_1) = \exp \int_x^{x_1} (-\sigma i/u) du$ , where the contour is any rectifiable path from x to  $x_1$  in  $E_1$ . Then clearly, if we put  $L(x, \rho) = B(x, \rho x)$  for  $1 \le \rho < \infty$ , we have

(6) 
$$|L(x, \rho)| \equiv 1 \text{ and } \partial L(x, \rho)/\partial x \equiv 0.$$

Hence,

$$|L(x, \rho)| \rho^{-1} |\phi_1(\rho x)| \le (\rho \log 2\rho)^{-1} (\log 2\rho)^{-t/2}$$

for  $x \in E_1$  and  $1 \le \rho < \infty$ ; and since the right side is

$$(-2/t)d((\log 2\rho)^{-t/2})/d\rho$$
,

we have by the M-test [7, p. 22] that the integral

(7) 
$$v_0(x) = -\int_1^{\infty} L(x,\rho)\rho^{-1}\phi_1(\rho x)d\rho$$

is uniformly convergent on  $E_1$  and thus represents an analytic function there (e.g., [7, p. 100]) whose derivative may be calculated by differentiating under the integral sign. In view of (5), clearly

(8) 
$$|v_0(x)| \leq (2/t)(\log |x|)^{-t/2}$$

on  $E_1$  and hence  $v_0 \rightarrow 0$  in F(a, b). Differentiating (7), we see easily that  $v_0' - \sigma i x^{-1} v_0 = x^{-1} \phi_1$  in  $E_1$ , so  $v_0$  is a solution of (3). Successively differentiating (7) and using (6), we see that for all j

(8a) 
$$\theta^{j}v_{0}(x) = -\int_{1}^{\infty} L(x,\rho)\rho^{-1}(\theta^{j}\phi_{1})(x\rho)d\rho$$

(where, for example,  $(\theta\phi_1)(x\rho) = x\rho\phi_1'(x\rho)$  etc.) in  $E_1$ . Since  $\phi_1 \prec (\log x)^{-1-t/2}$ , it follows (see §2(b)) that  $\theta^j\phi_1 \prec (\log x)^{-(j+1)-t/2}$  in F(a, b), and so, for each j, there is an  $S_j \in F(a, b)$  and a constant  $c_j$  such that

$$\mid \theta^{j} \phi_{1}(x) \mid \leq c_{j} (\log \mid x \mid)^{-(j+1)-t/2} \text{ on } S_{j}.$$

Hence by (8a), there is a  $C_i$  such that

(9) 
$$|\theta^{j}v_{0}(x)| \leq C'_{j}(\log |x|)^{-j-t/2} \text{ in } S_{j}.$$

Thus  $\theta^{j}v_0 \rightarrow 0$  in F(a, b) for each j. Now let p > 1. Then, by the definition of the operator  $\theta_p$ ,  $\theta_p v_0 = G\theta v_0$  where  $G = \log x \cdot \cdot \cdot \log_{p-1} x$ . It is routine to verify by induction on j that for  $j = 1, 2, \cdots$ 

(10) 
$$\theta_p^j v_0 = \sum_{\alpha=1}^j G_{\alpha j} \theta^{\alpha} v_0,$$

where

$$(11) G_{\alpha j} = \sum_{i} m(i_1, \cdots, i_j, \alpha) G^{i_1}(\theta G)^{i_2} \cdots (\theta^{j-1} G)^{i_j}$$

in which the *m*'s are constants,  $i_1 + \cdots + i_j = j$  and  $i_2 + 2i_3 + \cdots + (j-1)i_j = j - \alpha$  for each term in (11). Now for all  $\epsilon > 0$ ,  $G < (\log x)^{1+\epsilon}$  so (see §2(b))  $\theta^j G < (\log x)^{1-j+\epsilon}$  for each *j*. Hence, by (11), for each  $\alpha$  and *j* we have  $G_{\alpha j} < (\log x)^{\alpha+\epsilon j}$  for all  $\epsilon > 0$ . Now *t* is a fixed positive number and so for each given  $\alpha$  and *j*, we have, by taking  $\epsilon = t/5j$ , that  $G_{\alpha j} < (\log x)^{\alpha+t/4}$ . Hence there exist  $S_{\alpha j} \in F(a, b)$  and constants  $d(\alpha, j)$  such that on  $S_{\alpha j}$ ,

$$|G_{\alpha i}(x)| \leq d(\alpha, j)(\log |x|)^{\alpha+t/4}$$
.

Thus by (9) and (10), for each p and j,

$$|\theta_{p}^{i}v_{0}(x)| \leq m_{pj}(\log |x|)^{-t/4}$$

in some element of F(a, b) for some constant  $m_{pj}$ . Thus  $\theta_p^j v_0 \rightarrow 0$  for all p and j since t > 0 and so

$$(12) v_0 \prec 1 \text{ in } F(a,b).$$

Since  $v_0$  solves (3), we have

(13) 
$$v_0 = (\sigma i)^{-1} (x v_0' - \phi_1).$$

Since  $v_0 \prec 1$ ,  $xv_0' \prec (\log x)^{-1}$ . Thus since  $\phi_1 \prec (\log x)^{-1-t/2}$ , we have by (13) that  $v_0 \prec (\log x)^{-1}$ . Hence  $xv_0' \prec (\log x)^{-2}$ , and so if -1-t/2 < -2, we have  $v_0 \prec (\log x)^{-2}$ . Continuing this way, if m is the greatest integer <1+t/2, then  $v_0 \prec (\log x)^{-m}$ , and so since  $m+1 \ge 1+t/2$ , we have by (13) that  $v_0 \prec (\log x)^{-1-t/2}$  in F(a, b). Thus by (4),  $v_0 \prec (\log x)^{\lambda_1-1}$  and so  $v_0 \prec f$  in F(a, b). Hence if  $z_0 = f + v_0$ , then  $z_0 \sim f$  and  $z_0$  solves (1). Hence  $y^* = x^{\lambda_0} z_0$  is a solution of  $xy' - \gamma y = \psi$  and  $y^* \sim N$  in F(a, b) concluding the proof.

5. Lemma. Assume the hypothesis and notation of §3. Let  $\Phi(y) = \sum_{i=0}^{n} B_{i}(\infty)\theta^{i}y$ . Then there exists a function  $y^{*}$  such that  $\Phi(y^{*}) = \phi$  and  $y^{*} \sim M$  in F(a, b) (where M is as in §3).

PROOF. Let  $F(\alpha) = \sum B_i(\infty)\alpha^i$  be of degree p. If p=0, take  $y^* = (B_0(\infty))^{-1}\phi$ . Hence we may assume p>0. It is easy to verify that if  $F(\alpha) = b_p(\alpha - \alpha_1) \cdot \cdot \cdot \cdot (\alpha - \alpha_p)$  (where  $b_p = B_p(\infty)$ ), then  $b_p^{-1}\Phi = (\theta - \alpha_1) \circ \cdot \cdot \cdot \circ (\theta - \alpha_p)$  where the order of the factors is immaterial. Let  $\phi^* = b_p^{-1}\phi$  and let  $Q^* = b_p^{-1}Q$ . We solve  $\Phi(y) = \phi$  by successive integrations on  $b_p^{-1}\Phi = \phi^*$  using §4, and we adopt the following notation. We let  $y_1$  be any solution of  $xy' - \alpha_1 y = \phi^*$  given by §4. Since  $y_1$  is  $\sim$  to a logarithmic monomial, we let  $y_2$  be any solution of  $xy' - \alpha_2 y = y_1$  given by §4. In general,  $y_{j+1}$  is any solution of  $xy' - \alpha_{j+1}y = y_j$  ( $1 \le j \le p-1$ ) given by §4. Then clearly  $y^* = y_p$  solves  $\Phi(y) = \phi$ . We will show  $y^* \sim M$ .

Case I.  $F(\sigma_0) \neq 0$ . By  $\S4(a)$ ,  $y_1 \sim (\sigma_0 - \alpha_1)^{-1}Q^*$ . Similarly  $y_2 \sim (\sigma_0 - \alpha_2)^{-1}(\sigma_0 - \alpha_1)^{-1}Q^*$ . Continuing by  $\S4(a)$ ,  $y_p \sim (F(\sigma_0))^{-1}Q$  so  $y_p \sim M$ .

Case II.  $\sigma_0$  is a root of F of multiplicity r and  $\sigma_1 \notin \{-1, \dots, -r\}$ . Let  $\alpha_1 = \dots = \alpha_r = \sigma_0$ . By  $\{4(b), y_1 \sim (\sigma_1 + 1)^{-1}(\log x)Q^*$ . Similarly by  $\{4(b), \dots, y_n \in A(b)\}$ 

$$y_j \sim (\sigma_1 + j)^{-1} \cdot \cdot \cdot (\sigma_1 + 1)^{-1} (\log x)^j Q^*$$
 for  $2 \le j \le r$ .

Then by §4(a),

$$y_{r+1} \sim (\sigma_0 - \alpha_{r+1})^{-1} (\sigma_1 + r)^{-1} \cdot \cdot \cdot (\sigma_1 + 1)^{-1} (\log x)^r Q^*$$

and by continuing to use §4(a), clearly

$$y_p \sim K(\log x)^r Q$$
 where  $K = (F^{(r)}(\sigma_0)/r!)^{-1}(s_{rr}(\sigma_1 + r))^{-1}$ 

so  $y_p \sim M$ .

Case III.  $\sigma_0$  is a root of F of multiplicity r and  $\sigma_1 = -1$ . Thus,  $\min\{j: j \ge 1, \sigma_j \ne -1\} = k$  (as in §3). Thus, by §4(b), (assuming  $\alpha_1 = \cdots = \alpha_r = \sigma_0$ ),

$$y_1 \sim (\sigma_k + 1)^{-1} (\log x \cdot \cdot \cdot \log_k x) Q^*.$$

Since  $\delta_1(y_1) = 0$ , we find by continuing up to r using §4(b) that  $y_r \sim ((r-1)!)^{-1} (\log x)^{r-1} y_1$ . We now continue using §4(a) and find that  $y_p \sim b_p(F^{(r)}(\sigma_0)/r!)^{-1} y_r$ . Since  $(r-1)! = s_{r,r-1}(\sigma_1+r)$ , clearly  $y_p \sim M$ .

Case IV.  $\sigma_0$  is a root of F of multiplicity r and  $\sigma_1 \in \{-2, \dots, -r\}$ . Let  $s = -\sigma_1$ . Since  $\sigma_1 \neq -1$ , by §4(b), (assuming  $\alpha_1 = \dots = \alpha_r = \sigma_0$ ),  $y_1 \sim (\sigma_1 + 1)^{-1} (\log x) Q^*$ . Continuing up to s - 1, we find

$$y_{s-1} \sim [(\sigma_1 + 1) \cdot \cdot \cdot (\sigma_1 + s - 1)]^{-1} (\log x)^{s-1} Q^*.$$

Since  $\delta_1(y_{s-1}) = -1$ , we have by §4(b),

$$y_s \sim (\sigma_k + 1)^{-1} \log x \cdot \cdot \cdot \log_k x y_{s-1}$$
.

Since  $\delta_1(y_s) = 0$ , we have, using §4(b), that  $y_r \sim ((r-s)!)^{-1} (\log x)^{r-s} y_s$ . Now, using §4(a), we find  $y_p \sim b_p(F^{(r)}(\sigma_0)/r!)^{-1} y_r$ . Since  $(r-s)! = (\sigma_1 + r) \cdot \cdot \cdot \cdot (\sigma_1 + s + 1)$ , it follows that  $y_p \sim M$ .

6. Conclusion of main theorem (§3). For each i,  $B_i = b_i + w_i$  where  $b_i = B_i(\infty)$  and  $\delta_0(w_i) < 0$ . Letting  $\Phi(y) = \sum_{i=0}^n b_i \theta^i y$  and  $\Gamma(y) = \sum_{i=0}^n w_i \theta^i y$ , we have by §5 that there exists a function  $y^* \sim M$  in F(a, b) such that  $\Phi(y^*) = \phi$ . Under  $y = y^* + z$ ,  $\Omega(y) = \phi$  becomes  $\Omega(z) = -\Gamma(y^*)$ . Now if  $\delta_0(y^*) = \lambda$ , then it is easily verified that  $\delta_0(\theta^j y^*) \leq \lambda$  for each j. Letting  $\epsilon > 0$  be such that  $\delta_0(w_i) < -\epsilon$  for each i, we have

(1) 
$$\delta_0(\Gamma(y^*)) < \lambda - \epsilon.$$

We now utilize a technique employed by Strodt in [5] which we outline here for the reader's convenience. Let  $H = \{\alpha \colon F(\alpha) = 0\}$ . Then if q is a real number not in H and we let  $k_q = (F(q))^{-1}$ , it is easily seen that the principal monomial of  $\Omega(y) - x^q$  is  $k_q x^q$ . Hence if we let  $\Lambda_q(\omega) = x^{-q}\Omega(k_q x^q \omega)$ , then by the properties of a principal monomial we have  $\Lambda_q(1) \sim 1$  and  $\Lambda_q(E) < 1$  if E < 1. (Thus  $\Lambda_q$  is unimajoral in the terminology of [4, Section 13]). Further, it is easily seen that  $\Lambda_q$  has coefficients in an  $LD_0(F(a, b))$  and that  $\partial \Lambda_q/\partial \omega^{(n)}$  is a nontrivial function. Thus by [4, Section 27],  $\Lambda_q$  possesses at least one principal fac-

torization sequence, that is, a sequence  $(V_1, \dots, V_n)$  of logarithmic monomials such that  $\Lambda_q$  may be written

$$\Lambda_q = \dot{V}_n \cdot \cdot \cdot \dot{V}_1 + \sum_{j=0}^n E_j \dot{V}_j \cdot \cdot \cdot \dot{V}_1$$

where  $\dot{V}_j$  is the operator  $\dot{V}_j(y) = y - y'/V_j$  and where each  $E_j < 1$ . Now by definition of  $\Lambda_q$ , it is easily verified that

$$\Lambda_q(\omega) = k_q \sum_{j=0}^n B_j (q + \theta)^j \omega,$$

and so it follows from [4, Section 44] that all principal factorization sequences for  $\Lambda_{\sigma}(\omega)$  can be obtained as follows. If we let

$$C_1^* \Lambda_q(y) = k_q \sum_{j=0}^n B_j(q + xy)^j$$

and if  $N_1, \dots, N_n$  are the logarithmic monomials such that the zeros  $y_1, \dots, y_n$  of  $C_1^*\Lambda_q(y)$  satisfy  $y_j/N_j \rightarrow 1$  for each j, then  $(V_1, \dots, V_n)$  is a principal factorization sequence for  $\Lambda_q$  if and only if  $(V_1, \dots, V_n)$  is a permutation of  $(N_1, \dots, N_n)$  and for each j,  $V_j$  is either  $\prec$  or  $\approx$  to  $V_{j+1}$ . Since  $\{B_0, \dots, B_n\}$  is contained in an  $LD_0(F(a, b))$ , it easily follows that if  $(V_1, \dots, V_n)$  is a principal factorization sequence for  $\Lambda_q$ , then for each j,  $V_j$  has the form

$$(2) V_j = c_i x^{-1+t_i}$$

for some constant  $c_j$  and some  $t_j \ge 0$ .  $V_j$  is called nonexceptional if either  $t_j > 0$  or  $c_j$  is not purely imaginary, and  $(V_1, \dots, V_n)$  is called nonexceptional if each  $V_j$  is nonexceptional. From the definition of  $C_1^*\Lambda_q(y)$ , it follows that for q and s not in H, we have

$$C_1^* \Lambda_q(y) = k_q(k_s)^{-1} C_1^* \Lambda_s((q-s)x^{-1}+y).$$

We now fix s and we fix a principal factorization sequence for  $\Lambda_s$  (which may be exceptional). By the above relation,  $y^{\sharp}$  is a zero of  $C_1^*\Lambda_q(y)$  if and only if  $(q-s)x^{-1}+y^{\sharp}$  is a zero of  $C_1^*\Lambda_s(y)$ , and so it easily follows from the previous discussion that except for finitely many real q,  $\Lambda_q(\omega)$  possesses a nonexceptional principal factorization sequence. Thus we have outlined the proof given in [5] that there is a finite set G of real numbers such that for any real number  $q \notin G$ ,  $\Lambda_q(\omega)$  is unimajoral and possesses a nonexceptional principal factorization sequence.

In our case here, we choose a real number  $q \notin G$  such that  $\lambda - \epsilon < q < \lambda$ , and let  $(V_1, \dots, V_n)$  be a nonexceptional principal fac-

torization sequence for this  $\Lambda_q$ . If  $V_j$  has the form in (2), then its indicial function (as defined in [4, Section 61]) is the function defined on (a, b) given by  $f_j(\alpha) = \cos(t_j\alpha + \arg c_j)$ . (Thus if  $t_j = 0$ ,  $f_j$  is the constant function  $\cos(\arg c_j)$ .) Since each  $V_j$  is nonexceptional, each  $f_j$  has only finitely many zeros in (a, b). Hence the union of the sets of roots of the  $f_j$  is a finite set  $\gamma_1 < \cdots < \gamma_d$  in (a, b). Thus if  $I = (a_1, b_1)$  is any of the intervals  $(a, \gamma_1)$ ,  $(\gamma_1, \gamma_2)$ ,  $\cdots$ ,  $(\gamma_d, b)$ , then no indicial function vanishes anywhere in I, so by definition [4, §98],  $(V_1, \cdots, V_n)$  is unblocked in  $(a_1, a_2, b_1)$  for any  $a_2$ . Furthermore, by (1) and choice of q,  $x^{-q}\Gamma(y^*) < 1$ , and so by definition [4, §88],  $(V_1, \cdots, V_n)$  is a strong factorization sequence for  $\Lambda_q(\omega) + x^{-q}\Gamma(y^*)$  in F(I). Hence by [4, §99, Theorem II], the equation  $\Lambda_q(\omega) + x^{-q}\Gamma(y^*) = 0$  possesses a solution  $\omega_0 < 1$  in F(I). Letting  $z_0 = k_q x^q \omega_0$ , we have  $\Omega(z_0) = -\Gamma(y^*)$ ; and since  $q < \lambda$ , we have  $z_0 < y^*$ . Letting  $y_0 = y^* + z_0$ , we then have  $\Omega(y_0) = \phi$  and  $y_0 \sim M$  in F(I), which concludes the proof.

7. Corollary. Under the hypothesis and notation of §3, let I be a subinterval of (a, b) such that in F(I) there is a solution  $y_0 \sim M$  of  $\Omega(y) = \phi$  (as just proved) and such that a complete logarithmic set of solutions  $\{g_1, \dots, g_p\}$  of  $\Omega(y) = 0$  exists in F(I). (It was shown in  $[1, \S11]$  that if  $p = \max\{j: B_j(\infty) \neq 0\}$ , then e.f.d. in F(a, b) there exist solutions  $g_1, \dots, g_p$  of  $\Omega(y) = 0$  such that  $g_j \sim x^{\alpha_j} (\log x)^{\beta_j}$  for some complex  $\alpha_j$  and integer  $\beta_j$  and such that  $(\alpha_k, \beta_k) \neq (\alpha_j, \beta_j)$  if  $k \neq j$ .)

Then if  $y^{\sharp}$  is any solution of  $\Omega(y) = \phi$  which in F(I) is  $\langle x^{\delta}$  for some constant  $\delta$ , then e.f.d. in F(I) there exist constants  $c_1, \dots, c_p$  and a trivial function T(x) such that  $y^{\sharp} = y_0 + \sum_{j=1}^p c_i g_i + T$ .

PROOF.  $y^{\sharp} - y_0$  is a solution of  $\Omega(y) = 0$  and is  $\prec x^{\alpha}$  for some  $\alpha$  so the result follows immediately from [1, §12].

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