

## ON MATROIDS ON EDGE SETS OF GRAPHS WITH CONNECTED SUBGRAPHS AS CIRCUITS

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**ABSTRACT.** It is proved that if  $\mathcal{F}$  is a finite family of connected, finite graphs, then a graph  $G$  exists such that the subgraphs of  $G$  isomorphic to a member of the family cannot be regarded as the circuits of a matroid on the edge set of  $G$ .

1. In a recent paper [1] we have proved that there are only two matroids on the edge set of any graph  $G$  (let us call them edge set matroids), whose circuits are connected subgraphs which form homeomorphic equivalent classes. These matroids are the polygon-matroid, whose circuits are the cycles, and the matroid of bi-circular subgraphs, where a bi-circular graph is a graph formed by two cycles which either have a path in common, or a vertex in common, or are disjoint but linked by a path; these graphs are homeomorphic to those pictured in Figure 1.

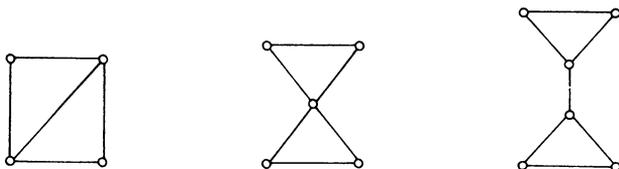


FIGURE 1

The hypothesis concerning homeomorphism is essential to the arguments in [1]. If we drop this hypothesis, the problem of finding all edge set matroids seems to be a very difficult one. As an unknown referee pointed out to me, a matroid of this kind is the matroid whose circuits are: (i) all cycles of even length; (ii) all graphs consisting of two cycles

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of odd length, having only one vertex in common; (iii) all graphs consisting of two cycles of odd length, joined by a path. In any graph  $G$ , the subgraphs of these kinds are the circuits of a matroid on the edge set of  $G$  but a cycle of odd length, although homeomorphic to a cycle of even length, is not a circuit of the matroid.

In this note we prove a theorem concerning edge set matroids. Our terminology is now slightly different from that used in [1]: we reserve the word "circuit" for the matroid-circuits and use "cycle" for simple closed paths in a graph. Moreover a matroid is defined as follows (see Whitney [2]):

Let  $E$  be a set of elements and  $\mathcal{K}$  a family of subsets of  $E$  (circuits).  $\mathcal{K}$  defines a matroid on  $E$  if and only if the following axioms hold:

AXIOM 1. No circuit is properly contained in another circuit.

AXIOM 2. If  $K$  and  $K'$  are distinct circuits,  $a \in K \cap K'$  and  $b \in K' - K$ , then a circuit  $K''$  exists such that  $b \in K'' \subset K \cup K' - \{a\}$ .

2. Let  $\mathcal{F}$  be a family of abstract connected graphs such that in any graph  $G$  the subgraphs isomorphic to members of  $\mathcal{F}$  are the circuits of a matroid. Call the members of  $\mathcal{F}$  circuits. Then

LEMMA 1. *No circuit has a pendant edge.*

PROOF. Let  $K$  be a circuit with a pendant edge, say  $x$ . Take another circuit  $K'$ , equal to  $K$ , and let  $K \cup K'$  be such that  $x$  is the only edge common to  $K$  and  $K'$  and the pendant vertex of  $x$  in each circuit coincides with the vertex of higher degree in the other circuit. Clearly,  $x$  is a bridge in  $K \cup K'$ . By Axiom 2,  $K \cup K' - \{x\}$  contains a circuit. But since all circuits must be connected, the existence of such a circuit contradicts Axiom 1.

Thus the lemma is proved.

THEOREM 1. *Let  $\mathcal{F}$  be a finite family of connected, finite graphs. Then a graph  $G$  exists such that the subgraphs of  $G$  isomorphic to a member of  $\mathcal{F}$  (or, for brevity's sake belonging to  $\mathcal{F}$ ) cannot be regarded as the circuits of a matroid on the edge set of  $G$ .*

PROOF. Let  $\mathcal{F}$  be a finite family of finite, connected graphs. The members of this family may eventually be regarded as the circuits of a matroid on the edge set of some graphs. However a graph  $G$  always exists with a subgraph which, according to the definition of a matroid, must also be a circuit but which does not belong to the family. This is a consequence from the fact that, for the members of a family  $\mathcal{F}$  of connected, finite graphs to be circuits of an edge-set matroid defined on any graph  $G$ , there must always exist a member of  $\mathcal{F}$  with a pair of edges of minimal distance arbitrarily large.

To prove it let  $K$  be a circuit,  $\alpha=(a_1, a_2)$ ,  $\beta=(b_1, b_2)$  two edges of  $K$ . Consider the four distances  $d(a_i, b_j)$  for  $i, j=1, 2$ . Let  $r$  be the minimal distance between  $\alpha$  and  $\beta$ , and suppose we choose a pair  $\alpha, \beta$  in  $K$  for which this distance is maximal among all edge pairs. Moreover, without loss of generality, we may suppose  $d(a_1, b_1)=r$ . There are 6 distinct cases which are summarized in Table I (columns 1 to 5).

Cases	$d(a_1, b_1)$	$d(a_1, b_2)$	$d(a_2, b_1)$	$d(a_2, b_2)$	$d(a_1, b_1)$	$d(a_1, b_2)$	$d(a_2, b_1)$	$d(a_2, b_2)$
I	$r (>0)$	$r$	$r$	$r$	$r$	$r+1$	$r$	$r+1$
II	$r (>0)$	$r+1$	$r$	$r$	$r$	$r+1$	$r$	$r+1$
III	$r (>0)$	$r+1$	$r$	$r+1$	$r+1$	$r+1$	$r+1$	$r+1$
IV	$r (>0)$	$r+1$	$r+1$	$r$	$r$	$r+1$	$r+1$	$r+1$
V	$r (\geq 0)$	$r+1$	$r+1$	$r+1$	$r$	$r+1$	$r+1$	$r+2$
VI	$r (\geq 0)$	$r+1$	$r+1$	$r+2$	$r+1$	$r+1$	$r+2$	$r+2$

TABLE I

Take another circuit  $K'$  and let  $K \cup K'$  be such that  $\beta$  is the only edge common to  $K$  and  $K'$ . For simplicity let us say the edges of  $K$  are black and those of  $K'$  are blue. Now, by Axiom 2, a circuit  $K''$  exists such that  $\alpha \in K'' \subset K \cup K' - \{\beta\}$ . By Axiom 1,  $K''$  contains both black and blue edges. Since  $K''$  must be connected and as a consequence of Lemma 1, either  $K''$  contains at least one blue path  $P(b_1, b_2)$  with length  $s \geq 2$ , or at least one of the vertices  $b_1$  and  $b_2$  is a cut-point of  $K''$  and there exists at least one cycle in the blue block of  $K''$  relative to this cut-point.

If blue paths exist, then take one with minimum length  $s \geq 2$ . We distinguish two possibilities:

(a)  $s \geq 3$ . Let  $\beta'$  be an edge of  $P(b_1, b_2)$  incident to neither  $b_1$  nor  $b_2$ . The minimal distance between  $\alpha$  and  $\beta'$ , which both belong to  $K''$ , is  $\geq r+1$ , that is to say, we obtain a new circuit  $K''$  from a given circuit  $K$  with a pair of edges  $\alpha$  and  $\beta'$  whose minimal distance is greater than the minimal distance between the edges  $\alpha$  and  $\beta$  of  $K$ .

(b)  $s=2$ . Let  $(b_1, x), (x, b_2)$  be the edges in  $P(b_1, b_2)$ . We have to examine the 6 cases of Table I. In cases I, II, IV and V, we set  $\beta'=(b_1, x)$  and  $x$  plays now the role of  $b_2$ . In cases III and VI, we set  $\beta'=(x, b_2)$  and  $x$  plays the role of  $b_1$ . The new distances between the endpoints of  $\alpha$  and  $\beta'$  are given in the columns 6 to 9 of Table I. With this operation we obtain, in cases III and VI, a pair of edges in  $K''$ , namely  $\alpha$  and  $\beta'$ , whose minimal distance is larger than the distance between  $\alpha$  and  $\beta$ . In the remaining cases, to obtain a circuit with a pair of edges satisfying this condition, one or two iterations of this operation may be required, each time with  $K''$  and  $\beta'$  in the roles of  $K$  and  $\beta$ , respectively. In fact, cases I and II yield case III, case V yields case VI and case IV yields in a first iteration case V

which in turn yields case VI. Now from cases III and VI, a new iteration allows us to achieve our aim.

If no blue path exists, then take the above mentioned blue cycle. Suppose the cycle belongs to the blue block of  $b_1$ . (The same argument holds a fortiori with  $b_2$  instead of  $b_1$ .) Let  $\beta'$  be an edge of the cycle non-incident to  $b_1$ . Obviously, the minimal distance between  $\alpha$  and  $\beta'$  is  $\geq r+1$ .

Hence it is always possible to obtain from a pair of edges  $\alpha, \beta$  in a circuit  $K$ , whose distance is  $r$ , a new pair  $\alpha, \beta'$  in a circuit  $K''$ , whose distance is  $\geq r+1$ . By repeating the argument, the theorem is proved.

Theorem 1 may also be stated more briefly as follows.

**THEOREM 1'.** *No edge-set matroid (on an arbitrary graph) may exist with a finite number of connected, finite graphs as circuits.*

#### REFERENCES

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