

PROOF OF A CONJECTURE OF FRIEDMAN

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ABSTRACT. We prove that every uncountable hyperarithmetical set has a member of each hyperdegree $\geq \mathbf{0}$, the hyperdegree of Kleene's $\mathbf{0}$.

We improve the main result of Friedman [1] by proving his conjecture that every uncountable hyperarithmetical set has a member of each hyperdegree $\geq \mathbf{0}$, the hyperdegree of Kleene's $\mathbf{0}$. This result has been obtained independently by Friedman by a different method. Friedman's proof uses ideas employed by L. Harrington to obtain a partial result.

In [2] it is shown that there is a function $f: \omega \rightarrow \omega$ of hyperdegree $\mathbf{0}$ and a Gödel number e such that $(\forall k)(f(k) \leq g(k))$ implies that f is hyperarithmetical in g with Gödel number e .

By [1] it suffices to prove that each recursive tree of finite sequences of natural numbers with uncountably many branches has a branch of each hyperdegree $\geq \mathbf{0}$. Let T be such a tree and let f be as above. Let $x: \omega \rightarrow 2$ have hyperdegree $\geq \mathbf{0}$. Let A be the Cantor-Bendixson perfect subtree of T . $A \in \Sigma_1^1$. A point in A is *good* if it has at least two immediate successors in A .

We shall define a branch h through A . Let $\sigma_0 < \sigma_1 < \dots$ be the good points on our branch. $h(\ln(\sigma_n))$ will always be given either the smallest or the next to smallest possible value. $h(\ln(\sigma_{2n}))$ and $h(\ln(\sigma_{2n+1}))$ will have the minimal possible value unless there are exactly k numbers $m < n$ such that

$$h(\ln(\sigma_{2m})) \text{ is not minimal, and } f(k) \leq \ln(\sigma_{2n}).$$

When this happens, $h(\ln(\sigma_{2n}))$ will not be minimal and $h(\ln(\sigma_{2n+1}))$ will be minimal just in case $x(k) = 0$.

For $g: \omega \rightarrow \omega$ let $g \in C$ if and only if, for each k , there are at least $2k + 1$ good points σ_n of length $\leq g(k)$ such that $h(\ln(\sigma_n))$ is not minimal. C is Σ_1^1 in h and $g \in C \Rightarrow (\forall k)(g(k) \geq f(k))$. f is hyperarithmetical in h since

$$f^* = f \Leftrightarrow (\exists g \in C) (f^* \text{ is hyperarithmetical in } g \text{ with Gödel number } e).$$

Since both f and A are hyperarithmetical in each of x and h , x and h are hyperarithmetical in one another.

REFERENCES

1. H. Friedman, *Borel sets and hyperdegrees*, J. Symbolic Logic **38** (1973), 405–409.
2. C. Jockusch and R. Soare, *Encodability of Kleene's $\mathbf{0}$* , J. Symbolic Logic **38** (1973), 437–440.

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