

## THE JUMP INVERSION THEOREM FOR $Q_{2n+1}$ -DEGREES

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ABSTRACT. Assuming projective determinacy we extend Friedberg's Jump Inversion theorem to  $Q_{2n+1}$ -degrees, after noticing that it fails for  $\Delta^1_{2n+1}$ -degrees.

**0. Preliminaries.** We list some results from the theory of countable analytical sets and  $Q$ -theory. For a more extensive account, including proofs, see [2 and 5]. Some familiarity with forcing in the analytical hierarchy is assumed; consult [3 and 4].

DEFINITION 0.1 (PD).  $C_{2n+1}$  is the largest countable  $\Pi^1_{2n+1}$  set of reals.

DEFINITION 0.2 (PD).  $C_{2n+2}$  is the largest countable  $\Sigma^1_{2n+2}$  set of reals.

We mention some of their properties:  $C_{2n+2}$  is the set of reals that are recursive in some element of  $C_{2n+1}$ . The set  $C_m$  is made up of  $\Delta^1_m$ -degrees (a  $\Delta^1_m$ -degree is a set of reals that is an equivalence class for the equivalence relation  $\alpha \equiv_{\Delta^1_m} \beta \Leftrightarrow \alpha \in \Delta^1_m(\beta)$  and  $\beta \in \Delta^1_m(\alpha)$ ). The  $\Delta^1_m$ -degrees in the set  $C_m$  are well-ordered by  $\alpha \leq_{\Delta^1_m} \beta \Leftrightarrow \alpha \in \Delta^1_m(\beta)$ .

DEFINITION 0.3. Given  $S \subset \omega^\omega$  let  $H_{2n+1}(S) = \{\alpha: \forall \beta \in S (\alpha \in \Delta^1_{2n+1}(\beta))\}$ ; we call it the *hull* of  $S$ . If  $S$  is a nonempty  $\Sigma^1_{2n+1}$  set then  $H_{2n+1}(S)$  is called a  $\Sigma^1_{2n+1}$ -hull. We let  $Q_{2n+1} =$  the union of all  $\Sigma^1_{2n+1}$ -hulls.

We have, assuming PD: The set  $Q_{2n+1}$  is  $\Pi^1_{2n+1}$ . Every  $\Sigma^1_{2n+1}$ -hull is  $\Pi^1_{2n+1}$ -bounded (this means that if  $R(\alpha, x)$  is  $\Pi^1_{2n+1}$  this so is  $\exists \alpha \in H_{2n+1}(S) R(\alpha, x)$ ). The set  $Q_{2n+1}$  is the largest  $\Sigma^1_{2n+1}$ -hull, and the largest  $\Pi^1_{2n+1}$ -bounded set. Relativizing to an arbitrary real  $\beta$  we may define the set  $Q_{2n+1}(\beta)$ . We define also  $\alpha \leq_{Q_{2n+1}} \beta \Leftrightarrow \alpha \in Q_{2n+1}(\beta)$ , and  $\alpha \equiv_{Q_{2n+1}} \beta \Leftrightarrow \alpha \in Q_{2n+1}(\beta)$  and  $\beta \in Q_{2n+1}(\alpha)$ . This is an equivalence relation, and the equivalence classes are called  $Q_{2n+1}$ -degrees. The set  $C_{2n+1}$  consists of such degrees. The set  $Q_{2n+1}$  is the largest initial segment of  $C_{2n+1}$  closed under  $\leq_{\Delta^1_{2n+1}}$ ; it consists of the  $\Delta^1_{2n+1}$ -degrees in  $C_{2n+1}$  up to and not including the degree of the first nontrivial (i.e. non- $\Delta^1_{2n+1}$ )  $\Pi^1_{2n+1}$  singleton  $y_0^{2n+1}$ . Relativizing to  $\alpha$  we have  $y_\alpha^{2n+1}$ . If  $\alpha \leq_{Q_{2n+1}} \beta$  then  $y_\alpha^{2n+1} \leq_{\Delta^1_{2n+1}} y_\beta^{2n+1}$ , and  $y_\alpha^{2n+1}$  plays the role of the *jump* for  $Q_{2n+1}$ -degrees. The set  $Q_{2n+1}$  is closed under the  $\Delta^1_{2n+1}$ -jump.

To obtain an *ordinal assignment* for the  $Q_{2n+1}$ -degrees we proceed as follows.

DEFINITION 0.4.

$$\begin{aligned} \lambda_{2n+1} &= \sup \{ \xi: \xi \text{ is the length of a } \Sigma^1_{2n+1} \text{ wellfounded relation on } \omega^\omega \} \\ &= \sup \{ \xi: \xi \text{ is the length of a } \Delta^1_{2n+1} \text{ prewellordering of } \omega^\omega \}. \end{aligned}$$

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Relativizing to  $\alpha$  we obtain  $\lambda_{2n+1}(\alpha)$ . Finally

$$k_{2n+1}(\alpha) = \sup\{\lambda_{2n+1}(\langle \alpha, \beta \rangle) : \lambda_{2n+1}(\langle \alpha, \beta \rangle) < \lambda_{2n+1}(y_\alpha^{2n+1})\}.$$

Of course,  $\lambda_{2n+1}$  is the ordinal assignment for the  $\Delta^1_{2n+1}$ -degrees, e.g. the Spector Criterion holds:  $\mathbf{d} \leq_{\Delta_{2n+1}} \mathbf{e} \Rightarrow [\mathbf{d}' \leq_{\Delta_{2n+1}} \mathbf{e} \Leftrightarrow \lambda_{2n+1}(d) < \lambda_{2n+1}(e)]$ . Now we have  $\lambda_{2n+1}(\alpha) < k_{2n+1}(\alpha) < \lambda_{2n+1}(y_\alpha^{2n+1})$ ,  $k_{2n+1}(\alpha)$  is invariant under  $\equiv_{Q_{2n+1}}$ ,  $\alpha \leq_{Q_{2n+1}} \beta \Rightarrow k_{2n+1}(\alpha) \leq k_{2n+1}(\beta)$ , and the Spector Criterion is true for  $Q_{2n+1}$ -degrees:  $\mathbf{d} \leq_{Q_{2n+1}} \mathbf{e} \Rightarrow [\mathbf{d}' \leq_{Q_{2n+1}} \mathbf{e} \Leftrightarrow k_{2n+1}(d) \leq k_{2n+1}(e)]$ . Naturally  $\mathbf{d}'$  is the degree of  $y_d^{2n+1}$ .

The relation  $k_{2n+1}(\alpha) \leq k_{2n+1}(\beta)$  is  $\Sigma^1_{2n+1}$ .

**1. Background and definitions.** One of the early results in the theory of Turing degrees was the following:

**FRIEDBERG JUMP INVERSION THEOREM [1].** *If  $\mathbf{b} \geq \mathbf{0}'$  then there exists an  $\mathbf{a}$  such that  $\mathbf{a}' = \mathbf{a} \vee \mathbf{0}' = \mathbf{b}$ .*

Of course  $\mathbf{0}$  denotes the degree of the recursive sets, and  $'$  denotes the Turing jump operation.

Next, the question was considered in the context of hyperdegrees. Let  $\mathbf{0}$  denote the hyperdegree of the hyperarithmetical sets and  $'$  the hyperjump. Does the above theorem hold? The answer is yes [6]:

**JUMP INVERSION THEOREM FOR  $\Delta^1_1$ -DEGREES.** *If  $\mathbf{b} \geq \mathbf{0}'$  then there exists an  $\mathbf{a}$  such that  $\mathbf{a}' = \mathbf{a} \vee \mathbf{0}' = \mathbf{b}$ .*

A natural question now is: does the inversion theorem hold for  $\Delta^1_{2n+1}$ -degrees? (We are assuming PD, needless to say). By a well-known argument Determinacy implies that there exists *some* cone on which inversion holds (a *cone*, by definition, is  $\{\mathbf{a} : \mathbf{a} \geq \mathbf{b}\}$ , and  $\mathbf{b}$  is called the base of the cone). But what is the base of the cone? Is it again  $\mathbf{0}'$ ? (I.e. the  $\Delta^1_{2n+1}$ -jump of the degree of  $\Delta^1_{2n+1}$  sets.) Surprisingly, the answer is no.

**THEOREM (KECHRIS, UNPUBLISHED) (PD).** *If  $n \geq 1$ , then no real in  $C_{2n+2}$  can be a base for a cone of inversion of the  $\Delta^1_{2n+1}$ -jump. ("Cone of inversion" of course means that every member of the cone is the  $\Delta^1_{2n+1}$ -jump of some  $\Delta^1_{2n+1}$ -degree.)*

**PROOF.** For notational simplicity we let  $2n + 1 = 3$ . If a member of  $C_4$  were a base then it would be recursive in a member of  $C_3$ , so without loss of generality assume a base  $b$  is in  $C_3$ . Consider the set  $C = \{\alpha : \exists \beta \in Q_3(\alpha) (\beta \in C_3 \text{ and } \alpha \leq_{\Delta_3} \beta)\}$ . It is a subset of  $C_4$ , and it is  $\Pi^1_3$ , because the quantification is bounded. So it is countable, and hence a subset of  $C_3$ . Since  $b \in C_3$  everything  $\geq b$  in  $C_3$  is the  $\Delta^1_3$ -jump of a member of  $C$ , thus a member of  $C_3$ . However the  $\Delta^1_3$ -degrees in  $C_3$  are wellordered with successor steps taken by the  $\Delta^1_3$ -jump, so that a limit stage of this wellordering gives immediately a contradiction. ( $C_3$  is closed under  $\equiv_{Q_3}$ , hence  $\alpha' \in C_3 \Rightarrow \alpha \in C_3$ , hence no limit level of  $C_3$  is a  $\Delta^1_3$ -jump.)

So the inversion theorem is a property of hyperdegrees that fails to generalize to  $\Delta^1_{2n+1}$ -degrees,  $n \geq 1$ . Usually in such cases the validity of the property is restored if instead of  $\Delta^1_{2n+1}$ -degrees we work with  $Q_{2n+1}$ -degrees. Indeed, it is the case that the jump inversion theorem holds for  $Q_{2n+1}$ -degrees, i.e. the base is again  $\mathbf{0}'$ . Moreover we can establish that the  $Q_{2n+1}$ -jump is never one-to-one.

**JUMP INVERSION THEOREM FOR  $Q_{2n+1}$ -DEGREES (PD).** *If  $\mathbf{c}$  is a  $Q_{2n+1}$ -degree  $\geq \mathbf{0}'$  then there exist  $Q_{2n+1}$ -degrees  $\mathbf{a}, \mathbf{b}$  such that  $\mathbf{a} \vee \mathbf{b} = \mathbf{a}' = \mathbf{b}' = \mathbf{c}$ .*

The rest of the paper is devoted to the proof of this theorem.

**2. The proof.** For notational simplicity we work with  $2n + 1 = 3$ . First we establish a lemma.

**LEMMA 2.1.** *If  $\mathbf{0}' \not\leq \mathbf{b}$  (i.e.  $k_3^0 = k_3^b$ ) then  $\mathbf{b}' = \mathbf{b} \vee \mathbf{0}'$ .*

**PROOF.** By the Spector Criterion  $\mathbf{0}' \not\leq \mathbf{b}$  iff  $k_3^0 = k_3^b$ . Now  $k_3^0 < k_3^{b \vee \mathbf{0}'}$ , so again by the Spector Criterion  $\mathbf{b}' \leq \mathbf{b} \vee \mathbf{0}'$ . The opposite inequality is obvious.

**PROOF OF THE THEOREM.** The set  $\{\alpha: k_3^\alpha = k_3^0 \text{ and } \alpha \notin Q_3\}$  is  $\Sigma_3^1$  and comeager. In fact there is a sequence  $D_0, D_1, \dots$  of dense open sets,  $\{D_i\} \in \Delta_3^1(y_0)$ , such that  $\bigcap D_i \subset \{\alpha: k_3^\alpha = k_3^0 \text{ and } \alpha \notin Q_3\}$ . This is implicit in [3]; briefly, comeagerness is characterized by the Banach-Mazur game. Use the Game Formula to unfold it and make it  $\Pi_2^1$ ; then the set of winning strategies is also  $\Pi_2^1$ , so there is a winning strategy recursive in  $y_0$ , by the Martin-Solovay basis theorem [5]. This gives the dense open sets.

We describe an inductive construction of reals  $a$  and  $b$ . Set  $a_{-1} = b_{-1} = \emptyset$ .

*Inductive step.* Suppose  $a_n, b_n$  have been constructed (they are finite sequences of integers). Consider the dense, open set  $D_{n+1}$  and extend  $a_n$  by a finite segment  $s$ , least in some fixed enumeration, so that the basic neighborhood defined by  $\hat{a}_n s$  is contained in  $D_{n+1}$ . Extend  $\hat{b}_n s$  by a finite segment  $t$ , least again, so that the basic neighborhood defined by  $\hat{b}_n s t$  is contained in  $D_{n+1}$ . Set now  $a_{n+1} = \hat{a}_n s \hat{t}\{c(n)\}$ ,  $b_{n+1} = \hat{b}_n s \hat{t}\{c(n) + 1\}$ .

This completes the inductive step.

Let now  $a = \bigcup a_n, b = \bigcup b_n$ . Since  $a, b \in \bigcap D_i$  we have by Lemma 2.1 that  $\mathbf{a}' = \mathbf{a} \vee \mathbf{0}', \mathbf{b}' = \mathbf{b} \vee \mathbf{0}'$ . Now  $\mathbf{a} \vee \mathbf{0}' \geq \mathbf{c}$ , because using  $y_0$  we may trace the construction of  $a$  and find all  $c(n)$ 's. Likewise  $\mathbf{b} \vee \mathbf{0}' \geq \mathbf{c}$ . However  $\mathbf{a} \vee \mathbf{0}' \leq \mathbf{c}$ , too, because  $\mathbf{0}' \leq \mathbf{c}$  and the construction of  $a$  only needs  $y_0$  and  $c$ . The same holds for  $b$ , and therefore we have  $\mathbf{a}' = \mathbf{b}' = \mathbf{a} \vee \mathbf{0}' = \mathbf{b} \vee \mathbf{0}' = \mathbf{c}$ . Finally note that  $\mathbf{a} \vee \mathbf{b} \geq \mathbf{c}$ , because if both  $a$  and  $b$  are available then considering the points where they differ  $c$  may be obtained. So we have  $\mathbf{a}' = \mathbf{b}' = \mathbf{a} \vee \mathbf{b} = \mathbf{c}$ , and  $a, b$  cannot have the same degree.

**REMARK.** The real  $a, b$  may also be chosen to be of minimal degree by using perfect trees in  $Q_3$  instead of finite sequences.

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