INFINITE GAMES WITH IMPERFECT INFORMATION(1)

BY

MICHAEL ORKIN

ABSTRACT. We consider an infinite, two person zero sum game played as follows: On the nth move, players A, B select privately from fixed finite sets, A_n , B_n , the result of their selections being made known before the next selection is made. After an infinite number of selections, a point in the associated sequence space, Ω , is produced upon which B pays A an amount determined by a payoff function defined on Ω . In this paper we extend a result of Blackwell and show that if the payoff function is the indicator function of a set in the Boolean algebra generated by the C_{δ} 's (with respect to a natural topology on Ω) then the game in question has a value.

1. Introduction. Infinite games with imperfect information have been studied by several writers, notably Blackwell [1], [2], and Shapley [5]. Before proceeding with the main result of this paper, we will discuss the structure and admissible strategies of these games.

Let $\{A_n\}$, $\{B_n\}$ be sequences of nonempty finite sets. Let $Z_n = A_n \times B_n$, and let Ω be the space $\prod_{n=1}^{\infty} Z_n$ of infinite sequences $\omega = (z_1, z_2, \cdots)$ where $z_n \in Z_n$. Let Ω be topologized as follows (for a related discussion, see [3]):

Suppose X is the set of all positions, i.e. finite sequences, $x=(z_1,z_2,\cdots,z_n),\ z_i\in Z_i,\ n=0,1,2,\cdots$. Then if $\omega\in\Omega,\ x\in X$, we define x to be a neighborhood of ω if ω passes through x. If the positions are thus considered as sets, they form a base for a Hausdorff, disconnected topology for Ω in which Ω is compact.

In this topology any open set is defined by a subset of X (a countable collection of positions). Any set defined by a finite collection of positions is both open and closed. It is shown by Wolfe [7] that if G is a G_{δ} then there exists a collection of positions T such that $G = \{\omega \in \Omega \mid \omega \text{ passes through infinitely many members of } T\}$, which we will henceforth denote by G = T i.o.

Now, suppose f is a bounded Baire function on Ω . Then we define a zero sum two person game G_f , played as follows:

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First, player A selects $a_1 \in A_1$ while player B simultaneously selects $b_1 \in B_1$. The result, $z_1 = (a_1, b_1) \in Z_1$, is announced to both players, upon which A selects $a_2 \in A_2$ while B selects $b_2 \in B_2$, etc. The result of this infinite sequence of moves is a point $\omega = (z_1, z_2, \dots) \in \Omega$ and B pays A the amount $f(\omega)$.

A strategy α (β) for A (B) gives for each position x (of length n, say) a probability distribution on A_{n+1} (B_{n+1}) with the stipulation that if the current position is x, A (B) will make his next choice according to α (β). A pair of strategies (α , β) defines a probability distribution $P_{\alpha\beta}$ on Ω and, hence, an expected payoff to A in G_f when A uses α and B uses β :

$$E(f, \alpha, \beta) = \int f(\omega) dP_{\alpha\beta}(\omega).$$

The lower and upper values of G_f are, respectively,

$$L(G_f) = \sup_{\alpha} \inf_{\beta} E(f, \alpha, \beta), \qquad U(G_f) = \inf_{\beta} \sup_{\alpha} E(f, \alpha, \beta).$$

It is always true that $L(G_f) \leq U(G_f)$; if $L(G_f) = U(G_f)$, this common value is called the value of G_f and will be denoted by $Val(G_f)$.

2. Main Result. We will show that if $f = I_G$, where $G \in B(G_{\delta})$ (the Boolean algebra generated by the G_{δ} 's), then G_f has a value. In [1], Blackwell proved this result if G is a G_{δ} . Before proving this result, we give two examples of games of this type and mention a related open question.

Example 1. On each move, players A and B choose, simultaneously, a 0 or 1. The winning set S of the form $G_{\delta} \cup F_{\sigma}$, is defined as follows:

 $S = G \cup F$ where $G = \{\omega | \omega_n = (0, 0) \text{ for infinitely many } n \text{ and } \omega_n = (1, 1) \text{ for infinitely many } n\},$

 $F = \{\omega | \omega_n = (0, 0) \text{ for at most finitely many } n \text{ and } \omega_n = (1, 1) \text{ for at most finitely many } n\}.$

The value of this game is 1, which can be achieved by A with a nonrandom strategy; he starts by saying 1 on each move. If B says 0 on every move, F is hit. If B ever says 1, A then starts saying 0's. If B then says 1's forever, F is still hit. If B ever says 0 again, A switches back to 1's, etc. If there are an infinite number of changes G is hit, otherwise F is hit.

Example 2. The winning set is a G_{\S} . On each move, the players choose simultaneously a 0 or 1. If player A ever says 1, the game is over on that move; if B also said 1, A wins; if B said 0, B wins. If A never says 1, the game continues and A wins if there are infinitely many moves with outcome (0, 0). (In other words, A tries to predict B's choice. See [2] for a related game.)

The value of this game is 1, but there are no optimal strategies for A. Here is a strategy for A, due to David Blackwell, which, for fixed N, guarantees A at least 1-1/N:

Define $N_j = 2^j N$, $j = 1, 2, \dots$, so that

$$\sum_{i=1}^{\infty} \frac{1}{N_i} = \frac{1}{N}.$$

Player A divides the trials into successive blocks of lengths N_1, N_2, \cdots . If he has not yet stopped the game, i.e. played 1 when block j is reached, he selects X_j at random from $\{1, 2, \cdots, N_j\}$. He then plays 1 at the X_j th trial of block j if B's previous $X_j - 1$ plays in the block are all 1's; otherwise, he plays 0 throughout the block. Then, clearly,

$$P(A \text{ loses on } j \text{th block } | j \text{th block is reached}) \leq 1/N_{j}$$

Thus, $P(A \text{ loses on } j\text{th block}) \leq 1/N_j$, and $P(A \text{ loses by failing to match}) \leq \sum (1/N_j) = 1/N$. However, by the nature of this strategy, if the game goes on forever, A will win, since there would then be (0, 0)'s in each block. \square

The following question remains unsolved in general: Do games with payoffs which are simple functions based on sets in $B(G_{\delta})$ have a value, i.e. games with payoff of the form $f = c_1 I_{B_1} + \dots + c_n I_{B_n}$, where $B_i \in B(G_{\delta})$, c_i are constants. In fact, we do not even know whether or not the much simpler games with payoffs of the form $I_{Q_1} - I_{Q_2}$ have a value, where Q_1 , Q_2 are open and disjoint. In another paper we will discuss some special cases of these kinds of games and show that they have a value.

We are now ready to prove the main result.

Lemma 1. Consider the class of sets of the form

$$G_1 \cup F_1 \cup (G_2 \cap F_2) \cup \cdots \cup (G_n \cap F_n)$$

where $G_i \in G_{\delta}$, $F_i \in F_{\sigma}$. This class of sets is precisely $B(G_{\delta})$.

Proof. By the fact that a finite union of G_{δ} 's is a G_{δ} , a finite intersection of F_{σ} 's is an F_{σ} , and by the standard results for generating Boolean algebras (e.g. see [4, Proposition I. 2.2, p. 7]) it is easily shown that every set in $B(G_{\delta})$ is of the form $\bigcup_{i=1}^{n} (G_{i} \cap F_{i})$, $G_{i} \in G_{\delta}$, $F_{i} \in F_{\sigma}$. Thus, every set in $B(G_{\delta})$ is of the form

$$\left(\bigcup_{i=1}^{n} (G_{i} \cap F_{i})\right)^{c} = \bigcap_{i=1}^{n} (G_{i}^{c} \cup F_{i}^{c})$$

$$= \left(\bigcap_{i=1}^{n} F_{i}^{c}\right) \cup \left(\bigcap_{i=1}^{n} G_{i}^{c}\right) \cup \left(\bigcup_{i=1}^{n} \left(F_{i}^{c} \cap \left(\bigcap_{j \neq i} G_{j}^{c}\right)\right)\right)$$

which, again using the fact about finite unions (intersections) of G_{δ} 's $(F_{\sigma}$'s) is easily seen to be of the required form.

Lemma 2. Consider the class of sets $\mathcal{C} = \bigcup_{n=1}^{\infty} \mathcal{C}_n$ generated in the following way: $\mathcal{C}_1 = G_{\delta}$; if n > 1, $\mathcal{C}_n = sets$ of the form $G_{\delta} \cup A_{n-1}$, where $A_n = complements$ of sets in \mathcal{C}_n (e.g. $\mathcal{C}_2 = G_{\delta} \cup F_{\sigma}$). We claim $\mathcal{C} = B(G_{\delta})$.

Proof. By using De Morgan's law and the fact that $G_{\delta} \cup G_{\delta} = G_{\delta}$, $F_{\sigma} \cap F_{\sigma} = F_{\sigma}$, it is easily seen that \mathcal{C}_4 = sets of the form $G_1 \cup F_1 \cup (F_2 \cap G_2)$, \mathcal{C}_{2n} = sets of the form $G_1 \cup F_1 \cup (\bigcup_{i=2}^n (F_i \cap G_i))$ which by Lemma 1 gives the sets in $B(G_{\delta})$. \square

The next lemma is the first step in an induction which will yield the main result.

Lemma 3. Let ϕ be upper semicontinuous, $0 \le \phi \le 1$. Suppose $G \in G_{\delta}$. Then the game with payoff $\overline{\phi} = \min(\phi, I_G)$ has a value.

Proof. The first part of this proof and A's method of play is the same as in [1]. Suppose G=T i.o., where T is a collection of positions. For any position x, let G_x^* be the game, starting from x with payoff $U_t(G_{\overline{\phi}})$ if T is hit for the first time after x at t, with payoff 0 if T is never hit after x, where $U_t(G_{\overline{\phi}})$ is the upper value of the original game starting from t. This payoff is lower semicontinuous, so, by [6], G_x^* has a value and player B has an optimal strategy. We claim $\operatorname{Val}(G_x^*) \geq U_x(G_{\overline{\phi}})$; for fixed $\epsilon > 0$, we present a strategy for B starting from x such that no matter what A does, $E_x(\overline{\phi}) \leq \operatorname{Val}(G_x^*) + \epsilon$:

Let B, starting from x, play optimally in G_x^* until T is hit for the first time after x, say at t. Then B plays, starting from t, to keep $E_t(\overline{\phi}) \leq U_t(G_{\overline{\phi}}) + \epsilon$, so

$$E_{x}(\overline{\phi}) = \sum_{t \in T} p(t) E_{t}(\overline{\phi}) \leq \sum_{t \in T} p(t) U_{t}(G_{\overline{\phi}}) + \epsilon \leq \operatorname{Val}(G_{x}^{*}) + \epsilon.$$

Now, for $\epsilon>0$, we describe a strategy for A such that no matter what B does, $E(\overline{\phi})\geq U(G_{\overline{\phi}})-\epsilon$, and the lemma will be proved. First, A plays $\epsilon/4$ optimally in G_e^* (e denotes the empty position). If T is hit after e, say at t_1 , A then plays $\epsilon/8$ optimally in $G_{t_1}^*$, etc. (If T is hit for the nth time at t_n , A then plays $\epsilon/2^{n+1}$ optimally in $G_{t_n}^*$.) Let the resulting sequence of moves be denoted by $z=(z_1,z_2,\ldots)$.

We define a sequence of random variables: $X_0 = U(G_{\overline{\phi}})$; for $k \ge 1$, $X_k = U_{t_k}(G_{\overline{\phi}})$ if T was hit for the kth time at t_k , $X_k = 0$ if T was hit less than k times. Thus, we have

(1)
$$E(X_k | X_{k-1}, \dots, X_0) \ge X_{k-1} - \epsilon/2^{k+1}.$$

This is obvious if $X_{k-1}=0$. If not, T was hit for the (k-1)st time at t_{k-1} , say, after which A played $\epsilon/2^{k+1}$ optimally in $G_{t_k}^*$ to get at least $\operatorname{Val}(G_{t_k}^*) - \epsilon/2^{k+1} \ge X_{k-1} - \epsilon/2^{k+1}$. Since the payoff in $G_{t_k}^*$ is X_k , (1) follows. Taking expectations on both sides, we get

(2)
$$E(X_k) \ge E(X_{k-1}) - \epsilon/2^{k+1} \Longrightarrow E(X_k) \ge U(G_{\overline{\phi}}) - \epsilon/2.$$

Now, by the definition of upper semicontinuity and the nature of the topology on Ω , for every point $z=(z_1, z_2, \dots)$ and every $\epsilon > 0$, there exists k such that any point $\omega=(\omega_1, \omega_2, \dots)$ with $\omega_i=z_i$ for $i \leq k$ has the property

$$\phi(\omega) \le \phi(z) + \epsilon/2 \implies \text{if } z \in G, \ \overline{\phi}(\omega) \le \overline{\phi}(z) + \epsilon/2$$

$$\implies (\text{still if } z \in G) \ U_{(z_1, \dots, z_k)} (G_{\overline{\phi}}) \le \overline{\phi}(z) + \epsilon/2$$

$$\implies (\text{for any } z) \ \lim \sup_{n} X_n \le \overline{\phi}(z) + \epsilon/2.$$

The last implication is obvious if $z \in G$. If $z \notin G$, T is only hit say N times, so for $n \ge N$, $X_n = 0$. Using Fatou's lemma on the last inequality, we get

$$E(\vec{\phi}) \ge \lim \sup_{n} E(X_n) - \epsilon/2 \ge U(G_{\vec{\phi}}) - \epsilon. \quad \Box$$

Theorem 1. Suppose $H \in \mathcal{C}_n$, i.e. $H = G \cup S$, $G \in G_{\delta}$, $S^c \in \mathcal{C}_{n-1}$. Suppose, also, that ϕ is upper semicontinuous with the property $0 \le \phi \le 1$, $\phi = 1$ on S. Then the game with payoff $\overline{\phi} = \min(\phi, I_H)$ has a value.

Proof. Lemma 3 shows that the theorem is true for sets in \mathcal{C}_1 . Suppose the theorem is true for sets in \mathcal{C}_{n-1} . Let $H \in \mathcal{C}_n$, $H = G \cup S$, where $G \in G_\delta$, $S^c \in \mathcal{C}_{n-1}$ (assume, without loss of generality, that $G \neq \emptyset$). Suppose G = T i.o. for some collection of positions T. For any position x, let H_x^* be the game starting at x which continues until the first time T is hit after x, say at t, with A getting $U_t(G_{\overline{\phi}})$ when this happens. Otherwise, the game continues and A gets I_S . We claim H_x^* has a value (the payoff in H_x^* may be neither upper nor lower semicontinuous).

Observe that if C is closed, $J \in G_{\delta} \Rightarrow C \cap J \in G_{\delta}$; $J \in F_{\sigma} \Rightarrow C \cap J \in F_{\sigma}$; therefore $J \in \mathcal{C}_n \Rightarrow C \cap J \in \mathcal{C}_n$. Let \mathcal{O}_x be the open set defined by the collection of positions passing through x which later hit T. Let $C_x = \mathcal{O}_x^c$. Then $S^c \cap C_x \in \mathcal{C}_{n-1}$ since S^c is. Define $\phi^* = 1 - f$, where f is the payoff in H_x^* . Also, define the upper semicontinuous function $g: g = \phi^*$ on \mathcal{O}_x , $g \equiv 1$ elsewhere. Thus g satisfies the conditions of the theorem and $\phi^* = \min(g, I_{C_x \cap S^c})$, so by the induction hypothesis the game starting at x with payoff ϕ^* has a value. Therefore, H_x^* has a value since its payoff is $f = 1 - \phi^*$ (the method of proof in Lemma 3 allows negation of the payoff since the same proof can be used by reversing the role of the players; it clearly allows the addition of a constant to the payoff).

By reasoning similar to that in Lemma 3, it can be shown that $\operatorname{Val}(H_x^*) \geq U_x(G_{\overline{\phi}})$. Now, for fixed $\epsilon > 0$, we will exhibit a strategy for player A which guarantees that $E(\overline{\phi}) \geq U(G_{\overline{\phi}}) - \epsilon$.

As in Lemma 3, A starts out playing $\epsilon/4$ optimally in H_e^* , etc. (If T is hit for the nth time at t_n , A then plays $\epsilon/2^{n+1}$ optimally in $H_{t_n}^*$.) Let the resulting play be $z = (z_1, z_2, \cdots)$.

Define the random variables: $X_0 = U(G_{\overline{\phi}})$; for $k \ge 1$, $X_k = U_{t_k}(G_{\overline{\phi}})$ if T is hit for the kth time at t_k , $X_k = I_S$ if T is hit less than k times. By reasoning similar to that in Lemma 3, we get

(3)
$$E(X_k) \ge U(G_{\overline{\phi}}) - \epsilon/2.$$

Again, by the definition of upper semicontinuity, there exists $k_{(\epsilon,z)}$ such that any point $\omega=(\omega_1,\,\omega_2,\,\cdots)$ agreeing with z up to $z_{k_{(\epsilon,z)}}$ has the property

$$\phi(\omega) \le \phi(z) + \epsilon/2 \implies \text{if } z \in H, \ \overline{\phi}(\omega) \le \overline{\phi}(z) + \epsilon/2$$

$$\implies \text{if } z \in H, \ U_{(z_1, \dots, z_k)}(G_{\overline{\phi}}) \le \overline{\phi}(z) + \epsilon/2$$

$$\implies \text{for any } z, \ \lim \sup_{n} X_n \le \overline{\phi}(z) + \epsilon/2.$$

Again, the last step is obvious if $z \in G$. If not, $S \cap G^c$ is hit and $\overline{\phi} = 1$ or $S^c \cap G^c$ is hit and $\limsup_n X_n = 0$. Thus, by Fatou, $E(\overline{\phi}) \ge \limsup_n E(X_n) \ge U(G_{\overline{\phi}}) - \epsilon$. \square

Corollary 1. If $H \in B(G_{\delta})$, the game with payoff I_H has a value.

Proof. Let $\phi = 1$ and use the theorem and Lemma 2.

Corollary 2. G, has a value if f satisfies the following conditions:

- (a) There is a collection T of nonoverlapping positions (nonoverlapping means $x \in T \implies x$ is not an initial segment of any other member of T) such that if $x \in T$ then f is constant on all sequences passing through x.
 - (b) $0 \le f \le 1$ on T.
 - (c) There exists $H \in B(G_{\delta})$ such that $f = I_H$ if T is never hit.

Proof. The function $\phi = 1$ off T, $\phi = f$ otherwise, is upper semicontinuous, $\phi = 1$ on H, and $f = \min(\phi, I_H)$ so the theorem applies. \square

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DEPARTMENT OF STATISTICS, UNIVERSITY OF CALIFORNIA, BERKELEY, CALIFORNIA 94720

Current address: Department of Mathematics, Case Western Reserve University, Cleveland, Ohio 44106