

SELF-DUAL AXIOMS FOR MANY-DIMENSIONAL PROJECTIVE GEOMETRY

BY

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ABSTRACT. Proposed and compared are four equivalent sets R, S, T, D of self-dual axioms for projective geometries, using points, hyperplanes and incidence as primitive elements and relation. The set R is inductive on the number of dimensions. The sets S, T, D all include the axiom "on every n points there is a plane", the dual of this axiom, one axiom on the existence of a certain configuration, and one or several axioms on the impossibility of certain configurations. These configurations consist of $(n + 1)$ points and $(n + 1)$ planes for sets S, T , but of $(n + 2)$ points and $(n + 2)$ planes for set D . Partial results are obtained by a preliminary study of self-dual axioms for simplicial spaces (spaces which may have fewer than 3 points per line).

1. Definitions and results. Projective spaces, with finite dimension n , are defined classically by nondual axioms, using undefined points and certain sets called lines. We propose several equivalent sets of self-dual axioms, using undefined points, planes (= hyperplanes) and incidence. With further details and definitions to be given later, these axioms for *Projective* spaces P [or in brackets for *simplicial* spaces N] are

The classic [1, p. 24] *Set LP* [or *LN*].

Axiom I. *Any two distinct points are in exactly one line.*

Axiom IIP [or **IIN**]. *There exist at least three [or two] points in each line.*

Axiom III. *If five distinct points P, Q, R, S, T satisfy the collineations $\overline{PQR}, \overline{PST}$, then there exists a point U satisfying $\overline{QSU}, \overline{RTU}$.*

Axiom IV (and V). *There exists at least (and at most) $n + 1$ independent points.*

A *Set RP* [or *RN*] inductive on dimensions.

Axiom a. *For $n = 1$ the relation "on" is one-to-one.*

Axiom bP[or **bN**]. *For $n = 1$ there exist at least three [or two] points.*

Axiom c. *For $n \geq 2$ there exist at least one point and one plane not on each other.*

Axiom d. *For $n \geq 2$ any reduction $[R, r]$ is an $(n - 1)$ -dimensional space.*

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A Set TP [or TN] of axioms by configurations.

Axiom 1 (and 2). *On every n points (planes) there exists at least one plane (point).*

Axiom 3TP [or **3TN**]. *There exists a semidoublex [or a semisimplex].*

Axiom 4 (and 5). *Distinct planes (points) are distinguishable.*

Axiom 6T. *If T is a point on the n planes, and if t is a plane on the n points, of a semisimplex, then T is not on t .*

If we replace Axiom 3TP [or 3TN] by the stronger

Axiom 3D [or **3SN**]. *There exists a doublex [or a simplex],*

then we shall replace Axioms 4, 5, 6T by weaker axioms given later.

We consider also *Weak spaces* W defined equivalently by later configuration axioms or by the following inductive *Set RW*.

Axiom aW. *For $n = 0$ all incidences are "not on".*

Axioms bW and cW. *For $n = 0$ and for $n \geq 1$ there exist a point and a plane not on each other.*

Axiom dW. *For $n \geq 1$ any reduction $[R, r]$ is an $(n - 1)$ -dimensional space.*

We repeat that we start with undefined objects called *points*, undefined objects called *planes*, and an undefined relation *on* between single points and single planes. The symbols \circ and \sim mean "on" and "not on". Thus $A \circ a$, $A \sim b$ means that the point A is on the plane a but not on the plane b . *Independent* is defined in §5. By a *reduction* $[R, r]$, where always we assume $R \sim r$, we mean the set of all points on the plane r and of all planes on the point R . Two planes (or points) are said to be *distinguishable* if there exists a point (or plane) on exactly one of the two planes (or points). *Duality* means to interchange points and planes.

Axioms and configurations are identified by letters: P = projective, N = simplicial, W = weak, U = undistinguished, L = lines, R = reduction, S = square, T = triangle, D = doublex. For any positive integer α , a set of α points A_i and of α planes a_i , with $i = 1, 2, \dots, \alpha$, is called an α -square if the corresponding (square) incidence table is completely specified, and is called an α -triangle if the incidences $A_i ? a_j$ are specified for at least all $i \geq j$. We define various configurations (giving notation = name = properties)

TNa = *singlenot α -triangle* = α -triangle with $A_i \circ a_j$ for $i > j$; $A_i \sim a_i$.

$SN\alpha$ = *singlenot α -square* = α -square with $A_i \circ a_j$ for $i \neq j$; $A_i \sim a_i$.

$TP\alpha$ = *doublenot α -triangle* = α -triangle with $A_i \circ a_j$ for $i > j$; $A_i \sim a_i$, $A_i \sim a(i + 1)$.

$SP\alpha$ = *doublenot α -square* = α -square with $A_i \sim a_i$, $A_i \sim a(i + 1)$; $A_i \circ a_j$ otherwise.

$D(n + 2)$ = *doublex* = $(n + 2)$ -square with $A_i \sim a_i$, $A_i \sim a(i + 1)$, $A(n + 2) \sim a_1$; $A_i \circ a_j$ otherwise.

In particular

Semisimplex = TN_n , *Simplex* = $SN(n+1)$, *Semidoublex* = $TP(n+1)$.

It is allowed to take $\alpha = 1$ in a doublenot triangle or square, the hypothesis $A_i \sim a(i+1)$ being then vacuously true. Indeed for $\alpha = 1$ we have $TN_1 = SN_1 = TP_1 = SP_1$. We use capitals for points, lower cases for planes, script capitals for squares. For instance the α -square \mathcal{B} has automatically the points B_i and the planes b_i with $1 \leq i \leq \alpha$.

The axioms by configurations, including axioms given earlier, and to be considered in various later sets, are

Axioms 3TN, 3SN, 3TP, 3SP or 3D. *There exists a configuration TN_n , $SN(n+1)$, $TP(n+1)$, $SP(n+1)$ or $D(n+2)$ respectively.*

Axiom 4N. *There is only one plane on the first n points of any simplex.*

Axioms 6T, 6W or 6P. *If a point T is on the n planes, and if a plane t is on the n points, of a configuration TN_n , SN_n or SP_n respectively, then $T \sim t$.*

Axiom 6D. *Doublenot $(n+2)$ -squares are impossible.*

Axioms 7N, 7P or 7D. *No plane is on the $(n+1)$ points of a configuration $SN(n+1)$, $SP(n+1)$ or on the $(n+2)$ points of a doublex $D(n+2)$.*

Axioms 5N, 8N, 8P or 8D. *Duals of Axioms 4N, 7N, 7P, 7D.*

We consider the following sets of axioms:

LN, LP, RW, RN, RP defined earlier,

$TN = \{1, 2, 3TN, 4, 5, 6T\}$,

$TU = \{1, 2, 3TP, 6T\}$, $TP = \{TU, 4, 5\}$,

$SW = \{1, 2, 3SN, 6W\}$,

$SN = \{1, 2, 3SN, 4N, 5N\}$,

$SU = \{1, 2, 3SP, 7P, 8P\}$, $SP = \{SU, 4, 5\}$,

$DU = \{1, 2, 3D, 6D, 7D, 8D\}$, $DP = \{DU, 4, 5\}$.

Instead of Set X of axioms, often we write Axioms $\{X\}$. The Axioms 6P, 7N, 8N do not appear in any previous set, but will be used in proofs.

We shall prove

Theorem A1. *Sets RW and SW are equivalent (and define weak spaces). Proof in §3.*

Theorem B1. *Sets RN, SN, TN, LN are equivalent (and define simplicial spaces). Proof in §§3, 4, 5.*

Theorem C1. *Sets TU, SU, DU are equivalent. Proof in §§6, 8.*

Theorem D1. *Sets LP, RP, TP, SP, DP are equivalent (and define projective spaces). Proof in §§5, 6, 8.*

In the last section we study some examples, and discuss independencies within sets of axioms.

Theorem D is our main achievement. Its proof requires all results in this article, including the theory of W and N spaces. The longest proofs are $\text{Set } TP \Rightarrow \text{Set } RP$ and $\text{Set } DU \Rightarrow \text{Axiom } 6T$. These two proofs have a large common part, in particular §§2 and 7. The Axioms 4, 5 have no interest in discussions involving only configuration axioms (because an identification of undistinguishable elements is trivial), but they are fundamental when studying reduction axioms (because distinguishability in the whole space does not imply distinguishability in the reductions).

We have also

Theorem E1. *Any $(n + 2)$ points are the points of a doublex if and only if no $(n + 1)$ of them are coplanar.*

Proof. The proof is easy. Axiom 1 implies the "if" and Lemma B7 gives the "only if". \square

Theorem F1. *For two-dimensional spaces Axioms 1, 2, 3TP, 6P imply Axiom 6T.*

Proof. The proof of this innocuous theorem is difficult and involves most results developed in this article. See Theorem A9. The proof given there shows that Axioms (1, 2, 3TP, 6P) \Rightarrow Axiom 3SP for $n = 2$ only, and that Axioms (1, 2, 3SP, 6P) \Rightarrow Axiom 6T for all n .

The Set SN appeared already in [2] and the set (1, 2, 3D, 6T), for $n = 3$ only, appeared in [5].

2. Moving a simplex. We assume Axioms $\{SW\}$ or $\{SN\}$ and prove

Theorem A2. *For every point R and plane r with $R \sim r$ there exists a simplex \mathcal{B} such that $B(n + 1) = R$, $b(n + 1) = r$.*

This theorem states the existence of a simplex \mathcal{B} containing a given set \mathcal{S} of points and planes. We find \mathcal{B} by "moving" a simplex from the starting position \mathcal{A} , which is the simplex of Axiom 3SN, into the desired terminal position \mathcal{B} , by successive replacements of one point or one plane of \mathcal{A} by one point or one plane of \mathcal{S} .

The notation $\{\dots \uparrow \dots\}$ means the finite sequence starting with the term on the left and ending with the term on the right. Thus $\{B6 \uparrow B8\} = \{A2 \uparrow A4\}$ means $B6 = A2$, $B7 = A3$, $B8 = A4$.

Lemma B2. *A permutation of the points of a simplex and the same permutation of the planes yields a new simplex.*

The proof is obvious.

Lemma C2. *No plane is on all points of a simplex (reworded: Axiom 6W or 4N \Rightarrow Axiom 7N).*

Proof. The existence of a plane t on all points of a simplex \mathcal{U} contradicts Axiom 6W with $T = A(n+1)$, $t = t$, $SNn = \{A1 \uparrow An; a1 \uparrow an\}$, and contradicts Axioms 4N because $t \neq a(n+1)$. \square

Given a simplex \mathcal{U} , a point P , and two integers α, β with $1 \leq \alpha < \beta \leq n+1$, the notation $a(\alpha, \beta; P)$ means any plane on the point P and on the $(n-1)$ points A_i with $i \neq \alpha, i \neq \beta$. We use also the dual notation $A(\alpha, \beta; p)$.

Lemma D2. *Under the above hypotheses, at least one plane $b = a(\alpha, \beta; P)$ exists, and any such plane satisfies either $A\alpha \sim b$ or $A\beta \sim b$.*

Proof. Trivial consequence of Axiom 1 and Lemma C. \square

Lemma E2. *Given a simplex \mathcal{U} , a plane r and two integers β, γ with $1 \leq \beta < \gamma \leq (n+1)$, there exists a simplex \mathcal{C} such that either $C\beta \circ r$ or $C\gamma \circ r$ and such that $C_i = A_i, c_i = a_i$ for all $i \neq \beta, i \neq \gamma$.*

Proof. We consider the point $B = A(\beta, \gamma; r)$ and the plane $b = a(\beta, \gamma; B)$. Using Lemmas C and C dual, there are four cases:

Case 1. $B \sim a\beta, B \circ a\gamma$.

Case 2. $B \circ a\beta, B \sim a\gamma$.

Case 3. $B \sim a\beta, B \sim a\gamma, A\gamma \sim b$.

Case 4. $B \sim a\beta, B \sim a\gamma, A\beta \sim b, A\gamma \circ b$.

In Case 1 we obtain \mathcal{C} from \mathcal{U} by replacing $A\beta$ by B . In Case 3 we obtain \mathcal{C} from \mathcal{U} by replacing $A\beta$ and $a\gamma$ by B and b . Cases 2, 4 differ from Cases 1, 3 only by an interchange of β and γ . \square

Lemma F2. *Given a simplex \mathcal{U} and a plane r there exists a simplex \mathcal{D} such that $d(n+1) = r$. Moreover $D_i = A_i, d_i = a_i$ for all i such that $A_i \circ r, 1 \leq i \leq n$.*

Lemma G2. *And if $A(n+1) \sim r$ then $D(n+1) = A(n+1)$.*

(Incidentally, Axioms 1, 2, 7N, 8N imply Lemma F but not G.)

Proof of Lemma F. We consider three cases:

Case 5. $A_i \sim r$ only for $i = n+1$.

Case 6. $A_i \sim r$ only for $i = \beta$ where $1 \leq \beta \leq n$.

Case 7. $A_i \sim r$ for at least two values of i ; say $i = \beta, i = \gamma$.

In Case 5 we obtain \mathcal{D} from \mathcal{U} by replacing $a(n+1)$ by r . In Case 6 we obtain \mathcal{D} from \mathcal{U} by replacing $A\beta, A(n+1), a\beta, a(n+1)$ by $A(n+1), A\beta, a(n+1), r$ respectively. In Case 7 we apply Lemma E repeatedly, until we get out of this case. \square

Proof of Lemma G. Let β, γ be the smallest values of i such that $A_i \sim r$. We consider three cases:

Case 8. $\beta = n+1, \gamma$ does not exist.

Case 9. $\beta < \gamma = n+1$.

Case 10. $\beta < \gamma \leq n$.

Then Case 8 is the same as Case 5. In Case 9 let $B = A(\beta, n + 1; \tau)$. Then $B \sim a\beta$ by Axiom 6W with $T = B$, $t = a\beta$, $SNn = \{A1 \uparrow A(\beta - 1), A(\beta + 1) \uparrow A(n + 1); a1 \uparrow a(\beta - 1), a(\beta + 1) \uparrow an, r\}$, or by Axiom 5N with $A(n + 1)$, B distinguished by τ . Thus, similarly to Cases 1 and 3, the desired simplex is

$$\mathbb{D} = \{A1 \uparrow A(\beta - 1), B, A(\beta + 1) \uparrow A(n + 1); a1 \uparrow an, r\}.$$

In Case 10 we apply Lemma F repeatedly until we get into Case 8 or 9. \square

Lemma H2. *For any point R there exists a simplex \mathbb{E} such that $E(n + 1) = R$.*

Proof. We use the simplex \mathbb{A} of Axiom 3SN and apply the dual of Lemma F. \square

Proof of Theorem A2. We apply Lemmas F, G to the simplex of Lemma H. \square

3. Reductions for W and SN spaces. We prove Theorems A1 and the equivalence $RN \Leftrightarrow SN$ in Theorem B1. For a W-space which is not an N-space, see Example A9.

We use the *notations* z (= zero), y , n or m affixed to axioms to indicate that the axioms hold for zero dimension, for one dimension, for the original (n -dimensional) space or for all reductions (they are $m = n - 1$ dimensional) of the original space respectively.

Lemma A3. *Axiom 1n \Leftrightarrow Axiom 1m.*

Proof. The implication Axiom 1n \Rightarrow Axiom 1m is trivial. For the converse we make an induction on α , with $1 \leq \alpha \leq n$, in the statement "on every α points there exists a plane". Given α points $\{P1 \uparrow P\alpha\}$, by the inductive hypothesis there exists a plane p on $\{P1 \uparrow P(\alpha - 1)\}$. Either $P\alpha \circ p$ and p is as desired, or $P\alpha \sim p$, but then the reduction $[P\alpha, p]$ contains a plane q on $\{P1 \uparrow P(\alpha - 1)\}$, and this plane q is as desired. \square

Lemma B3. *Axioms (3SNm, cW or c) \Rightarrow Axiom 3SNn. Conversely Axioms {SWn} or {SNn} \Rightarrow Axiom 3SNm.*

Proof. The first implication is trivial. The second implication results from Theorem A2. \square

Lemma C3. *Axiom 6Wn \Leftrightarrow Axiom 6Wm.*

Proof. The singlenot n -square \mathbb{A} in Axiom 6Wn and the singlenot m -square \mathbb{B} in Axiom 6Wm are related by $\{A1 \uparrow An, a1 \uparrow an\} = \{B1 \uparrow B(n - 1), R; b1 \uparrow b(n - 1), r\}$. \square

Lemma D3. *Axiom 4Nn \Leftrightarrow Axiom 4Nm.*

Proof. The n -simplex \mathbb{A} in Axiom 4Nn is related to the m -simplex \mathbb{B} in Axiom 4Nm by $\{A1 \uparrow A(n + 1); a1 \uparrow a(n + 1)\} = \{B1 \uparrow B(n - 1), R, Bn; b1 \uparrow b(n - 1), r, bn\}$. \square

Lemma E3. *Axioms (aW, bW) \Leftrightarrow Axioms {SWz}.*

Proof. Axiom $aW =$ Axiom $6Wz$. Axiom $bW =$ Axiom $3SNz$. Axiom $bW \Rightarrow$ Axioms $(1z, 2z)$. \square

Lemma F3. *Axioms $(a, bN) \Leftrightarrow$ Axioms $\{SNy\}$.*

Proof. Axiom $a \Rightarrow$ Axioms $(1, 2, 4N, 5N)$. Axioms $(a, bN) \Rightarrow$ Axiom $3SN$. (Lemma H2, Axiom $4N) \Rightarrow$ (at most one plane on each point). (This, its dual, Axioms $1, 2) \Rightarrow$ Axiom a . Axiom $3SN \Rightarrow$ Axiom bN .

Theorem G3. *Axioms $\{SW\} \Leftrightarrow$ Axioms $\{RW\}$.*

Proof. Use $(3SN \Rightarrow c)$, Lemmas A, A dual, B, C, E. \square

Theorem H3. *Axioms $\{SN\} \Leftrightarrow$ Axioms $\{RN\}$.*

Proof. Use $(3SN \Rightarrow c)$, Lemmas A, A dual, B, D, F. \square

4. **The SN, TN equivalence.** To prove this equivalence we use lemmas, including a Theorem D4 which requires the moving of a simplex.

Lemma A4. *Axioms $\{TN\} \Rightarrow$ Axiom $3SN$.*

Proof. Let \mathcal{U} be the semisimplex of Axiom $3TN$ and let $A(n+1)$ be a point on the n planes of \mathcal{U} . For each integer α with $1 \leq \alpha \leq n+1$, let $b\alpha$ be a plane on all points A_i except possibly $A\alpha$. Then $A\alpha \sim b\alpha$ (Axiom $6T$). The simplex $\mathcal{B} = \{A1 \uparrow A(n+1); b1 \uparrow b(n+1)\}$ satisfies Axiom $3SN$. \square

Lemma B4. *Axioms $\{TN\} \Rightarrow$ Axiom $4N$.*

Proof. Axiom $6T \Rightarrow$ (no distinguishable planes are on the first n points of a simplex). Hence $(4, 6T) \Rightarrow 4N$. \square

Lemma C4. *Axioms 1 and $4N$ imply: Given a simplex \mathcal{D} and a point R with $R \sim d1$, there exists a simplex \mathcal{B} such that $B1 = R, B_i = D_i$ for $2 \leq i \leq n+1, b1 = d1$.*

Proof. For each integer α with $2 \leq \alpha \leq n+1$ we define $b\alpha = d(1, \alpha; R)$. Then $D\alpha \sim b\alpha$ because otherwise $d1$ and $b\alpha$ would be two distinct (because distinguished by R) planes on the last n points of the simplex \mathcal{D} , contrary to Axiom $4N$. The desired simplex is $\mathcal{B} = \{R, D2 \uparrow D(n+1); d1, b2 \uparrow b(n+1)\}$. \square

Theorem D4. *Axioms $\{SN\}$ imply: For every singlenot α -triangle \mathcal{Q} (the plane $a\alpha$ may be missing) there exists a simplex \mathcal{B} such that $B_i = A_i$ for $1 \leq i \leq \alpha$.*

Proof. We make an induction on α . For $\alpha = 1$ the theorem is true, by Lemma H2. For $\alpha > 1$ the inductive hypothesis and a circular permutation gives a simplex \mathcal{C} such that $C_i = A_i$ for $2 \leq i \leq \alpha$. Lemmas B2, F2 applied to \mathcal{C} and $r = a1$ give a simplex \mathcal{D} such that $d1 = a1$ and $D_i = C_i$ for $2 \leq i \leq \alpha$. Lemma C with $R = A1$ gives the desired simplex \mathcal{B} . \square

Lemma E4. *Axioms* $\{SN\} \Rightarrow$ *Axiom* 6*T*.

Proof. Let $T, t, TNn = \{A1 \uparrow An; a1 \uparrow an\}$ be as in *Axiom* 6*T*. By *Theorem* D there exists a simplex \mathcal{B} such that $\{B1 \uparrow B(n+1)\} = \{A1 \uparrow An, T\}$. Then $t = b(n+1)$ (*Axiom* 4*N*), hence $T \sim t$. \square

Lemma F4. *Axioms* $\{SN\} \Rightarrow$ *Axiom* 4.

Proof. Let r, s be two distinct planes. There exists a simplex \mathcal{E} with $e(n+1) = r$ (*Lemma* H2 dual). At least one of the points $\{E1 \uparrow En\}$ is not on s (*Axiom* 4*N*), hence distinguishes r, s . \square

Theorem G4. *Axioms* $\{SN\} \Leftrightarrow$ *Axioms* $\{TN\}$.

Proof. Use *Lemmas* A, B, B dual, E, F, F dual. \square

5. **The nondual theory.** We have discussed the *Axioms* $\{SN\}$ in an earlier publication [2], and gave there a self-dual treatment of flats, but we did not go far enough to compare *Axioms* $\{SN\}$ and $\{LN\}$. We do it now. We abandon our emphasis on self-duality. The present results are not used in later sections.

We assume *Axioms* $\{SN\}$. A *set-plane* p is defined as the set of all points on the plane p . A *line* is defined as the intersection of all set-planes containing two distinct points. If these points are P, Q , the line is denoted by \overline{PQ} .

Lemma A5. *In any simplex* \mathcal{U} , *the line* $\overline{A1A2}$ *is the intersection of the planes* $\{a3 \uparrow a(n+1)\}$.

Proof. The two points $P = A1, Q = A2$, and therefore all points of \overline{PQ} , are on $\{a3 \uparrow a(n+1)\}$. Conversely, if R is a point on these planes, and if r is a plane on the two points P, Q , then $R \circ r$, because otherwise we contradict *Axiom* 6*T* with $T = Q, t = a3, TNn = \{A4 \uparrow A(n+1), R, P; a4 \uparrow a(n+1), r, a1\}$. Thus R , being on all such planes r , is in \overline{PQ} . \square

Lemma B5. *If* R, S *are distinct points in a line* \overline{PQ} , *then* $\overline{PQ} = \overline{RS}$.

Proof. There exists a simplex \mathcal{U} with $A1 = P, A2 = Q$ (*Theorem* D4). We may assume $P \neq S$. Then $S \sim a2$ (*Axiom* 5*N*). Let $s = a(1, 2; S)$. Then $s \neq a2$, hence $P \sim s$ (*Axiom* 4*N*). Thus $\mathcal{B} = \{P, S, A3 \uparrow A(n+1); s, a2 \uparrow a(n+1)\}$ is a simplex. Comparing the planes of \mathcal{U} and \mathcal{B} , we get $\overline{PQ} = \overline{PS}$ (*Lemma* A). Similarly $\overline{PS} = \overline{RS}$. \square

Lemma C5. *Lines and reductions to one dimension are identical.*

Proof. Each line \overline{PQ} is the set of points of the one-dimensional space obtained by the succession of the $(n-1)$ reductions $[A3, a3] \uparrow [A(n+1), a(n+1)]$, with \mathcal{U} defined as above. Conversely each reduction to one dimension has two points P, Q (*Axiom* b*N*) and is identical to \overline{PQ} . \square

Lemma D5. *Axioms* $\{SN\} \Rightarrow$ *Axioms* (I, IIN, III).

Proof. The definition of a line and Lemma B imply Axioms I and IIN. Let P, Q, R, S, T be defined as in Axiom III. Excluding the trivial case of five collinear points, there exists a singlenot 3-triangle with points Q, S, P , hence (Theorem D4) there exists a simplex \mathcal{Q} with $A1 = Q, A2 = S, A3 = P$. Then $\mathcal{B} = \{R, T, P, A4 \uparrow A(n+1); a1, a2, b3, a4 \uparrow a(n+1)\}$, for an appropriate plane $b3$, also is a simplex. There exists a point U on the n planes $\{a3, b3, a4 \uparrow a(n+1)\}$. This point U is in \overline{QS} and in \overline{RT} , as shown by \mathcal{Q} and \mathcal{B} . This proves Axiom III. \square

In any plane satisfying Axioms (I, IIN, III) Dembowski defines: A set s of points is a *subspace* if for any two distinct points P, Q in s every point of the line \overline{PQ} is in s . A point A is *independent* of a set \mathcal{S} if there exists a subspace containing \mathcal{S} but not A . A set \mathcal{S} is *independent* if every point in \mathcal{S} is independent of the remaining points in \mathcal{S} .

Lemma E5. *Axioms* $\{SN\} \Rightarrow$ *Axiom* IV.

Proof. Let \mathcal{Q} be the simplex of Axiom $3SN$. For each integer α with $1 \leq \alpha \leq n+1$ the point $A\alpha$ is independent of the other points of \mathcal{Q} because $a\alpha$ is a subspace containing these other points but not $A\alpha$. \square

Lemma F5. *For any simplex* \mathcal{Q} *and point* R , *the point* R *depends on* $\{A1 \uparrow A(n+1)\}$.

Proof. We move a simplex \mathcal{B} from the initial position \mathcal{Q} into a final position \mathcal{E} such that $E(n+1) = R$ (Lemma H2). Each new point introduced during this motion has the form $C = B(\beta, \gamma; r)$, hence is on $\overline{B\beta B\gamma}$ (Lemma A), thus dependent on $\{B1 \uparrow B(n+1)\}$. In particular R is dependent on $\{B1 \uparrow B(n+1)\}$ and each $B\alpha$ is dependent on $\{A1 \uparrow A(n+1)\}$. \square

Lemma G5. *Axioms* $\{SN\} \Rightarrow$ *Axiom* V.

Proof. We make an induction on dimensions, by proving that Axioms $(\{SN\}, Vm) \Rightarrow$ *Axiom* Vn . Given $(n+2)$ points $\{A1 \uparrow A(n+2)\}$, if $\{A1 \uparrow A(n+1)\}$ are independent then they are not coplanar (Axiom Vm), hence they form a simplex, and $A(n+2)$ depends on them (Lemma F). This establishes Axiom Vn . \square

Dembowski defines a hyperplane as a maximal proper subspace.

Lemma H5. *The points and hyperplanes of any space* LN *satisfy Axioms* $\{SN\}$.

We do not prove this well-known result.

Theorem J5. *Axioms* $\{LN\} \Leftrightarrow$ *Axioms* $\{SN\}$ *and* *Axioms* $\{LP\} \Leftrightarrow$ *Axioms* $\{RP\}$.

Proof. The first \Leftrightarrow holds by Lemmas D, E, G, H. The second \Leftrightarrow holds by Theorem H3 and because Axiom IIP \Leftrightarrow Axiom bP (Lemma C). \square

Incidentally, the following theorem is true for simplicial spaces, but not for weak spaces.

Theorem K5. *Any two reductions $[R, r]$ and $[S, r]$ for the same plane r consist of the same point and the same set-planes.*

We may denote this reduction by $[r]$. See Example A9.

6. The $RP \Rightarrow TP \Rightarrow DP$ implications. We prove

Theorem A6. *The sets of axioms satisfy $\{RP\} \Rightarrow (\{SN\}, 3SP) \Rightarrow \{TP\} \Rightarrow \{TU\} \Rightarrow \{SU\} \Rightarrow \{DU\}$.*

We use several lemmas.

Lemma B6. *Axioms $\{RP\} \Rightarrow$ Axiom $3SP$.*

Proof. We may use Axioms $\{SN\}$ (by $RP \Rightarrow RN$ and Theorem H3). Let \mathcal{Q} be the simplex of Axiom $3SN$. The succession of the $(n - 1)$ reductions $[A1, a1] \uparrow [A(n - 1), a(n - 1)]$ is a one-dimensional space containing the two points $An, A(n + 1)$ and (Axiom bP) a third point P . Then $P \neq A(n + 1)$ and Axiom $5N$ gives $P \sim an$, while $P \neq An$ gives $P \sim a(n + 1)$. Thus we have found a point $Bn = P$ satisfying $Bn \circ aj$ for and only for $j \neq n, j \neq n + 1$. Similarly, for each integer α with $1 \leq \alpha < n$, after a permutation (Lemma B2), we can find a point $B\alpha$ satisfying $B\alpha \circ aj$ for and only for $j \neq \alpha, j \neq \alpha + 1$. Then $\mathcal{B} = \{B1 \uparrow Bn, A(n + 1); a1 \uparrow a(n + 1)\}$ satisfies Axiom $3SP$. \square

Lemma C6. *Axioms $\{TU\} \Rightarrow$ Axiom $3SP$.*

Proof. Let \mathcal{Q} be the semidoublex Axiom $3TP$. We construct a doublenot square \mathcal{C} with the same points as \mathcal{Q} . For each integer α with $3 \leq \alpha \leq n + 1$, let $B\alpha$ be a point on the n planes aj with $j \neq \alpha - 1$, and let $c\alpha = a(\alpha - 1, \alpha; B\alpha)$ be a plane on the one point $B\alpha$ and on the $(n - 1)$ points Ai with $i \neq \alpha - 1, i \neq \alpha$. The point $B\alpha$ and the plane $c\alpha$ exist. Moreover $A(\alpha - 1) \sim c\alpha$ and $A\alpha \sim c\alpha$ because otherwise we would contradict Axiom $6T$ by taking $T = B\alpha, t = c\alpha, TNn = \{A1 \uparrow A(\alpha - 2), \text{either } A(\alpha - 1) \text{ or } A\alpha, A(\alpha + 1) \uparrow A(n + 1); a1 \uparrow a(\alpha - 2), a\alpha \uparrow a(n + 1)\}$. Then $C = \{A1 \uparrow A(n + 1); a1, a2, c3 \uparrow c(n + 1)\}$ satisfies Axiom $3SP$. \square

Lemma D6. *Axioms $\{SU\} \Rightarrow$ Axiom $3D$.*

Proof. Let A be the doublenot square of Axiom $3SP$. Let b be a plane on $\{A1 \uparrow An\}$ and B be a point on $\{a2 \uparrow a(n + 1)\}$. Then $A(n + 1) \sim b$ and $B \sim a1$, by Axioms $(7P, 8P)$ respectively. Also $B \sim b$ by Axiom $7P$ with $SP(n + 1) = \{An \downarrow A1, B; a(n + 1) \downarrow a1\}$. The desired doublex is $\{A1 \uparrow A(n + 1), B; a1 \uparrow a(n + 1), b\}$. \square

Proof of Theorem A6. It is obvious that $\{RP\} \Rightarrow \{RN\}, 3SP \Rightarrow 3TP, \{TP\} \Rightarrow \{TU\}, 6T \Rightarrow (7P, 8P) \Rightarrow (6D, 7D, 8D)$. The other needed results are Theorems H3, G4 and Lemmas B6, C6, D6. \square

7. Moving a doublex. We assume Axioms $\{DU\}$ and prove:

Theorem A7. Every doublet α -square \mathcal{Q} (or \mathcal{Q} plus a point $A(\alpha + 1)$ on $\{a1 \uparrow a\alpha\}$) can be completed into a doublet.

Integers are taken in the additive algebra modulus $(n + 2)$. The inequality $\alpha < \beta$ means $1 \leq \alpha < \beta \leq n + 2$, and the arrows \uparrow or \downarrow mean in the increasing or decreasing cyclical order, starting with the term on the left and ending with the term on the right. Thus for $n = 4$ we have

$$\begin{aligned} \{A2 \uparrow A4\} &= \{A2, A3, A4\}, & \{A2 \downarrow A4\} &= \{A2, A1, A6, A5, A4\}, \\ \{A4 \uparrow A2\} &= \{A4, A5, A6, A1, A2\}, & \{A4 \downarrow A2\} &= \{A4, A3, A2\}. \end{aligned}$$

When $\alpha = \beta + 1$ the context will make it clear whether $\{A\alpha \uparrow A\beta\}$ represents the empty set or the set of all Ai .

Lemma B7. Any doublet \mathcal{Q} and plane r satisfy $Ai \sim r$ for at least two values of i .

Proof. Otherwise there would exist an integer α such that $Ai \circ r$ for $i \neq \alpha$. Then $A\alpha \circ r$ is contradicted by Axiom 7D, and $A\alpha \sim r$ is contradicted by Axiom 6D with $SP(n + 2) = \{A\alpha \uparrow A(\alpha - 1); r, a(\alpha + 1) \uparrow a(\alpha - 1)\}$. \square

Given a doublet \mathcal{Q} and two distinct integers α, β , the notation $a(\alpha, \beta)$ means any plane on the n points Ai with $i \neq \alpha, i \neq \beta$. Given a doublet \mathcal{Q} , a point B , and three distinct integers α, β, γ , the notation $a(\alpha, \beta, \gamma; B)$ means any plane on the point B and on the $(n - 1)$ planes Ai with $i \neq \alpha, i \neq \beta, i \neq \gamma$.

Given a doublet \mathcal{Q} and two distinct integers α, γ , a permutation of planes gives (Lemma C, below) a new doublet $\mathcal{B} = \mathcal{Q}(B\alpha, B\gamma)$ defined by

$$\begin{aligned} B\alpha &= A(\alpha, \gamma), & B\gamma &= A(\alpha + 1, \gamma + 1), \\ \{B(\gamma + 1) \uparrow B(\alpha - 1)\} &= \{A(\gamma + 1) \uparrow A(\alpha + 1)\}, & \{B(\alpha + 1) \uparrow B(\gamma - 1)\} &= \{A(\gamma - 1) \downarrow A(\alpha + 1)\}, \\ \{b(\gamma + 1) \uparrow b\alpha\} &= \{a(\gamma + 1) \uparrow a\alpha\}, & \{b(\alpha + 1) \uparrow b\gamma\} &= \{a\gamma \downarrow a(\alpha + 1)\}. \end{aligned}$$

If for the point $B\alpha$ or $B\gamma$ we can use a previously defined point C then we write $\mathcal{U}(B\alpha = C, B\gamma)$ or $\mathcal{U}(B\alpha, B\gamma = C)$.

Given a doublet \mathcal{Q} , a point B and two integers β, γ such that $\gamma + 1 \neq \beta \neq \gamma, B \sim a\gamma, B \sim a(\gamma + 1)$, but $B \circ aj$ for $j \neq \beta, j \neq \gamma, j \neq \gamma + 1$, we obtain (Lemma C, below) a new doublet $\mathcal{C} = \mathcal{Q}(C\gamma = B, c\beta)$ defined by

$$\begin{aligned} C\gamma &= B, & c\beta &= a(\beta - 1, \beta, \gamma; B), \\ C_i &= A_i \text{ for } i \neq \gamma, & c_j &= a_j \text{ for } j \neq \beta. \end{aligned}$$

The notations $A(\alpha, \beta), A(\alpha, \beta, \gamma; b), \mathcal{U}(b\alpha, b\gamma), \mathcal{U}(C\beta, c\gamma = b)$ are defined by duality.

Lemma C7. The two $(n + 2)$ -squares $\mathcal{B} = \mathcal{U}(B\alpha, B\gamma), \mathcal{C} = \mathcal{U}(C\gamma = B, c\beta)$ defined above exist and are doublets.

Proof. The points $B\alpha$, $B\gamma$ exist (Axiom 2) and satisfy the needed \sim (Lemma B dual). The plane $c\beta$ exists (Axiom 1), and satisfies $A(\beta - 1) \sim c\beta$, $A\beta \sim c\beta$, as we show by eliminating the other alternatives: $A(\beta - 1) \circ c\beta$, $A\beta \circ c\beta$ is contradicted by Lemma B. $A(\beta - 1) \circ c\beta$, $A\beta \sim c\beta$ is contradicted by Axiom 6D with $SP(n + 2) = \{C\beta \uparrow C(\beta - 1); c\beta \uparrow c(\beta - 1)\}$. $A(\beta - 1) \sim c\beta$, $A\beta \circ c\beta$ is contradicted by Axiom 6D with $SP(n + 2) = \{C(\beta - 1) \downarrow C\beta; c\beta \downarrow c(\beta + 1)\}$. \square

Given a doublex \mathfrak{A} and a plane r , we denote by $\beta(\mathfrak{A}), \gamma(\mathfrak{A})$, with $\beta < \gamma$, the largest two integers i such that $Ai \sim r$.

Lemma D7. *Given a doublex \mathfrak{A} and a plane r , let $\beta = \beta(\mathfrak{A})$ and assume $\beta \geq 2$. Then there exists a doublex \mathfrak{B} satisfying $\beta(\mathfrak{B}) < \beta(\mathfrak{A}), \gamma(\mathfrak{B}) \leq \gamma(\mathfrak{A}), Bi = Ai$ for $1 \leq i < \beta - 1, bj = aj$ for $1 \leq j < \beta$.*

Lemma E7. *Moreover either $\gamma(\mathfrak{B}) < \gamma(\mathfrak{A})$ or $B(\beta - 1) = A(\beta - 1)$.*

Proofs. Let $\gamma = \gamma(\mathfrak{A}), B = A(\beta, \gamma, \gamma + 1; r)$. We consider three cases.

Case 1. $B \sim a\gamma, B \sim a(\gamma + 1)$,

Case 2. $B \sim a\beta, B \sim a\gamma, B \circ a(\gamma + 1)$,

Case 3. $B \sim a\beta, B \circ a\gamma, B \sim a(\gamma + 1)$.

The doublex \mathfrak{B} satisfying Lemma D is

Case 1. $\mathfrak{B} = \mathfrak{A}(B\gamma = B, b\beta)$,

Case 2. $\mathfrak{B} = \mathfrak{A}(B\beta = B, B\gamma)$,

Case 3. $\mathfrak{B} = \mathfrak{A}(B(\beta - 1), B\gamma = B)$.

Moreover we get $\gamma(\mathfrak{B}) < \gamma(\mathfrak{A})$ in Cases 1 and 3, and we get $Bi = Ai$ for $1 \leq i < \beta$ in Cases 1 and 2. \square

Lemma F7. *Given a doublex \mathfrak{A} , a plane r and an integer δ , there exists a doublex \mathfrak{C} such that $\beta(\mathfrak{C}) \leq \delta, Ci = Ai$ for $1 \leq i < \delta, cj = aj$ for $1 \leq j \leq \delta$.*

Lemma G7. *Moreover either $\gamma(\mathfrak{C}) < n + 2$ or $C\delta = A\delta$.*

Proofs. If $\beta(\mathfrak{A}) > \delta$ we apply Lemma D repeatedly until we get a doublex \mathfrak{C} with $\beta(\mathfrak{C}) \leq \delta$. This proves Lemma F. Moreover Lemma E implies Lemma G. \square

Lemma H7. *Given a doublex \mathfrak{A} , a plane r and an integer α such that $Ai \circ r$ for $1 \leq i < \alpha$, there exists a doublex \mathfrak{D} such that $d(\alpha + 1) = r, Di = Ai$ for $1 \leq i < \alpha, dj = aj$ for $2 \leq j \leq \alpha$.*

Lemma J7. *Moreover if $A\alpha \sim r$ then $D\alpha = A\alpha$.*

Lemma K7. *Or if $A(\alpha + 1) \sim r$ then $d1 = a1$.*

In the following three proofs, let $\beta = \beta(\mathfrak{C}), \gamma = \gamma(\mathfrak{C})$.

Proof of Lemma H. Let \mathfrak{C} be the doublex of Lemma F for $\delta = \alpha$. Then $\beta = \alpha, r = c(\alpha, \gamma)$. The desired doublex is $\mathfrak{D} = \mathfrak{C}(d(\alpha + 1) = r, d(\gamma + 1))$. \square

Proof of Lemma J. Let \mathcal{C} be the doublex of Lemma F for $\delta = \alpha + 1$. If $\beta = \alpha$ the desired doublex \mathcal{D} is as defined in the proof of Lemma H. If $\beta = \alpha + 1$ then the desired doublex is $\mathcal{D} = \mathcal{C}(D\gamma, d(\alpha + 1) = r)$. \square

Proof of Lemma K. If $A\alpha \circ r$ we look at the proof of Lemma H. We have $C\alpha \sim r$, hence $C\alpha \neq A\alpha$, and Lemma G gives $\gamma < n + 2$. Then $\gamma + 1 \neq 1$ and $d1 = c1 = a1$. If $A\alpha \sim r$ we look at the proof of Lemma J. Either $C(\alpha + 1) \circ r$, then $C(\alpha + 1) \neq A(\alpha + 1)$, hence again $\gamma < n + 2$, or $C(\alpha + 1) \sim r$, then $d_j = c_j$ for $j \neq \gamma$. \square

Lemma L7. *Given a doublex \mathcal{E} , a point R and an integer α such that $R \circ e_j$ for $1 \leq j < \alpha$, there exists a doublex \mathcal{F} such that $F\alpha = R$, $F_i = E_i$ and $f_i = e_i$ for $1 \leq i < \alpha$.*

Lemma M7. *Moreover if $R \sim e\alpha$ then $f\alpha = e\alpha$.*

Proofs. These two lemmas are obtained from Lemmas H, J by the duality $A_i \rightarrow e_i, a(i + 1) \rightarrow E_i, r \rightarrow R$. \square

Lemma N7. *For any point R and plane r with $R \sim r$ there exists a doublex \mathcal{D} with $D1 = R, d1 = r$.*

Proof. To the doublex \mathcal{E} of Axiom 3D we apply Lemma L with $\alpha = 1$ to obtain a doublex \mathcal{F} with $F1 = R$, and to the doublex \mathcal{F} we apply Lemmas H, J with $\alpha = 1$ to obtain the desired doublex \mathcal{D} . \square

Proof of Theorem A7. We refer to this theorem by $A\alpha$ without the parentheses and by $A\alpha +$ with the parentheses. We make an induction on α . Theorem A1 is true, being the same as Lemma N. Now we assume Theorem $A(\beta - 1)$. Let \mathcal{A} be a doublerot β -square, and let i mean all i in $1 \leq i < \beta$. The inductive hypothesis gives a doublex \mathcal{B} with $B_i = A_i, b_i = a_i$. Lemma L for $\alpha = \beta$ gives a doublex \mathcal{C} with $C_i = B_i, C\beta = A\beta, c_i = b_i$. Lemmas H, J, K for $\alpha = \beta - 1$ give a doublex \mathcal{D} with $D_i = C_i, d_i = c_i, d\beta = a\beta$ (but $D\beta \neq A\beta$). Lemmas L, M for $\alpha = \beta$ give a doublex \mathcal{E} with $E_i = D_i, E\beta = A\beta, e_i = d_i, e\beta = d\beta$. The existence of \mathcal{E} proves Theorem $A\beta$. This completes the induction. The existence of \mathcal{C} proves Theorem $A\alpha +$. \square

8. The $DP \Rightarrow RP$ implications. We prove

Theorem A8. *Axioms $\{DP\} \Rightarrow$ Axioms $\{RP\}$.*

Theorem B8. *Axioms $\{DU\} \Rightarrow$ Axiom 6T.*

We use most previous results and additional lemmas.

Lemma C8. *Axioms $\{DU\} \Rightarrow$ Axiom 6P.*

Proof. The doublerot n -square $\mathcal{A} = SPn$ plus the point $A(n + 1) = T$ of Axiom 6P can be completed into a doublex \mathcal{B} (Theorem A7). Hence $T \sim t$ (Lemma B7). \square

We recall that n means for the original space, m means for all reductions $[R, r]$, and y means for one dimension. We assume $n \geq 2$.

Lemma D8. *Axiom 6Pn* \Rightarrow *Axioms (6Dm, 7Dm, 8Dm).*

Proof. The $SP(m + 2) = \{A1 \uparrow A(n + 1); a1 \uparrow a(n + 1)\}$ in Axiom 6Dm cannot exist, by Axiom 6Pn with $T = A(n + 1)$, $t = r$, $SPn = \{A1 \uparrow An; a1 \uparrow an\}$. A plane p in $[R, r]$ on the m -doublex \mathcal{D} is Axiom 7Dm cannot exist, by Axiom 6Pn with $T = R$, $t = p$, $SPn = \{D1 \uparrow Dn, d1 \uparrow dn\}$. Duality gives Axiom 8Dm. \square

Lemma E8. *Axioms {DU n }* \Rightarrow *Axiom 3Dm.*

Proof. Given R, r with $R \sim r$ there exists a doublex \mathcal{U} with $A1 = R$, $a1 = r$ (Lemma N7). Let $B = A(2, n + 2)$, $b = a(2, n + 2)$. The desired m -doublex is $\{A2 \uparrow A(n + 1), B; b, a3 \uparrow a(n + 2)\}$, noticing that $B \sim b$ by Axiom 6Dm. \square

Lemma F8. *Axioms {DP n }* \Rightarrow *Axiom 5m.*

Proof. Let P, Q be two distinct points in a reduction $[R, r]$, where distinct means distinguishable by a plane s in the original space (Axiom 5n). The case $R \circ s$ being trivial, we assume $R \sim s$. Then the doublenot rectangle $\{R, P, Q; r, s\}$ can be completed into a doublex \mathcal{B} (Theorem A7). The plane $b4$ is in $[R, r]$ and distinguishes P, Q . \square

Lemma G8. *Axioms {DP y }* \Rightarrow *Axioms (a, bP).*

Proof. Lemma C for $n = 1$ becomes $\{DUy\} \Rightarrow 6Ty$. We have also $3D \Rightarrow 3TN$. Hence $\{DPy\} \Rightarrow \{TNy\} \Rightarrow (a, bN)$. Moreover $3Dy \Rightarrow bP$. \square

Proof of Theorems A8 and B8. We get Theorem A8 from Lemmas C, D, E, F, F dual, G. We get Theorem B8 from Theorems A8, H3, E4 and by defining distinct to mean distinguishable in the original space. \square

9. Counterexamples and weaker axioms. First we give one example of a non-simplicial weak space and two examples of nonprojective simplicial spaces.

Example A9. We consider the space of five points and five planes with the following incidence table

	a	b	c	d	e
A	•	•	~	~	•
B	•	~	•	~	~
C	~	•	•	•	~
D	~	~	•	~	•
E	•	~	~	•	•

This space satisfies Axioms $\{SW\}$ for $n = 2$, but not Axiom 6T (take $T = A$, $t = a$,

$TN2 = \{B, E; e, b\}$). The distinguishable points A, E are not distinguishable in the reduction $[B, e]$. The reduction $[B, e]$ satisfies neither Axiom $3TP$ nor Axiom 5, while the reduction $[C, e]$ is projective. Thus Theorem K5 fails. \square

Example B9. The space consisting of a simplex only is simplicial but not projective. \square

Example C9. Let \mathcal{P} and \mathcal{Q} be two projective spaces. We declare that each point of \mathcal{P} is on all planes of \mathcal{Q} , and that each point of \mathcal{Q} is on all planes of \mathcal{P} . We form thus a new space \mathcal{N} , which we call the *join* of \mathcal{P} and \mathcal{Q} . This space \mathcal{N} is a nonprojective simplicial space, with dimension $n = \alpha + \beta + 1$, where α and β are the dimensions of \mathcal{P} and \mathcal{Q} .

Proof. The join \mathcal{N} satisfies Axioms (1, 2, $3SN$, 4, 5, $6T$), but the lines joining any one point of \mathcal{P} to any one point of \mathcal{Q} have only two points. Thus Axiom bP is not satisfied. \square

We say that an Axiom A can be weakened within a set $\{A, B, C, \dots\}$ of axioms if there exists an Axiom A' , which does not imply A , such that $\{A', B, C, \dots\} \Rightarrow A$.

Theorem A9. *Axiom $6T$ can be weakened within Axioms $\{TP\}$ by an additional incidence specification.*

The added specification is $s1 \sim s2$ for the semisimplex \mathcal{S} in Axiom $6T$.

Proof. Using this weaker Axiom $6T'$, the proof of Lemma C6 remains valid with $c3 = a(2, 3; B3)$ if $A1 \sim a3$ and with $c3 = a3$ if $A1 \circ a3$. Thus we proved Axiom $3SP$, hence Axioms $\{SU\}, \{DU\}$ (Theorem A6) and Axioms $6T$ (Theorem B8). \square

Theorem B9. *Axiom $3TP$ cannot be weakened within Axioms $\{TP\}$ by deletion of an incidence specification.*

Lemma C9. *Given two integers α, β with $1 \leq \alpha < \beta \leq n + 2$, Axioms $\{DU\}$ imply the existence of an $(n + 2)$ -square \mathcal{B} which is doublet except that $B\beta \sim b\alpha$.*

Proof of Lemma C. Let \mathcal{A} be a doublet, let $P = A(1, \alpha + 1)$, $p = a(1, \beta)$. Then $P \sim p$ by Axiom $6T$ (proved in Theorem B8) with $T = P$, $t = p$, $TPn = \{A2 \uparrow A\alpha, A(\beta - 1) \downarrow A(\alpha + 1), A(\beta + 1) \uparrow A(n + 2); a2 \uparrow a\alpha, a\beta \downarrow a(\alpha + 2), a(\beta + 1) \uparrow a(n + 2)\}$. The desired $(n + 2)$ -square $\mathcal{B} = \{A(\alpha - 1) \downarrow A1, P, A(\alpha + 1) \uparrow A(n + 2); a\alpha \downarrow a2, p, a(\alpha + 1) \uparrow a(n + 2)\}$. \square

Proof of Theorem B. When \mathcal{P} and \mathcal{Q} are projective spaces or zero-dimensional weak spaces, then their join \mathcal{N} in Example C satisfies the Axiom $3SP'$ obtained from Axiom $3SP$ by deletion of any one of the three specifications $A(\alpha + 1) \sim a(\alpha + 1)$, $A(\alpha + 1) \sim a(\alpha + 2)$, $A(\alpha + 2) \sim a(\alpha + 2)$. The Axiom $3SP''$ obtained from Axiom $3SP$ by deletion of any one $A\beta \circ a\alpha$, for $1 \leq \alpha < \beta \leq n + 1$, is satisfied by any $(n - 1)$ -dimensional projective space \mathcal{P} (Lemma C), and hence by the join \mathcal{N} of \mathcal{P} with a zero-dimensional weak space \mathcal{Q} . Thus neither Axioms ($\{TN\}, 3SP'$) nor Axioms ($\{TN\}, 3SP''$) imply Axiom bP . \square

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