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Mercurial Signatures.

A canonical digital signature scheme consists of three algorithms: key generation, signing, and verifying. It needs to be (1) correct: verification accepts (PK, M, σ) if σ is the output of the signing algorithm on input (SK, M) , for SK corresponding to PK , and (2) unforgeable: (informally) a signature that verifies under PK can only be produced by PK 's owner.

In a *mercurial* signature, public keys and messages are partitioned into equivalence classes using relations \equiv_k and \equiv_m , and there are additional algorithms:

* **ConvertSig**: On input (PK, M, σ) where σ is a valid signature on M under public key PK , output (PK', M, σ') , where $PK' \equiv_k PK$, and σ' is a valid signature on M under public key PK' .

* **ChangeRep**: On input (PK, M, σ) where σ is a valid signature on M under public key PK , output (PK, M', σ') , where $M' \equiv_m M$, and σ' is a valid signature on M' under public key PK .

Further, for an appropriate choice of message space and public key space, the new message M' cannot be linked to the original M , and the new public key PK' cannot be linked to PK .

In this talk we will go over constructions and applications of mercurial signatures and open problems related to them. (Received September 15, 2020)