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## SPHERE-PACKING IN THE HAMMING METRIC ${ }^{1}$

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Let $V_{n}(2)$ be the $n$-dimensional vector space over GF(2), with vectors represented as $n$-tuples of 0 's and 1 's. The Hamming metric $d(x, y)$ is defined to be the number of coordinates in which $x$ and $y$ disagree. If $A=\left\{a_{1}, a_{2}, \cdots, a_{M}\right\}$ is a set of $M$ vectors, we define $d(A)=\min _{i \neq j} d\left(a_{i}, a_{j}\right)$, and $\bar{d}(A)=$ mean $_{i \neq j} d\left(a_{i}, a_{j}\right)$. Finally define

$$
D(n, M)=\max _{|A|=M} d(A)
$$

We present in this paper a method of obtaining an upper bound on $D(n, M)$ which is always at least as good as the well-known bounds, and which is frequently better. At the same time, the method gives a satisfactory explanation of the relationship between the various known upper bounds on $D(n, M)$ (Hamming [1], Plotkin [1], and Elias [2]). The weakness of the method seems to be that for the most part it deals only with the average distance between vectors, and further progress probably awaits a technique which is able to deal more directly with the minimum distance.

We need three theorems. Throughout $A=\left\{a_{1}, a_{2}, \cdots, a_{M}\right\}$ is a set of $M$ vectors from $V_{n}(2)$.

Theorem 1. Let $S_{r}(x)$ be the sphere of radius $r$ centered at $x$. Then the mean value of $\left|S_{r}(x) \cap A\right|$ as $x$ varies over $V_{n}(2)$ is

$$
M_{r}=\frac{M}{2^{n}} \sum_{k \leq r}\binom{n}{k} .
$$

[^0]Proof. Each $a_{i}$ appears in exactly

$$
\sum_{k \leq r}\binom{n}{k}
$$

spheres of radius $r$, so that

$$
\sum_{x}\left|S_{r}(x) \cap A\right|=M \sum_{k \leq r}\binom{n}{k}
$$

Theorem 2 (Plotkin). Suppose $A \subseteq S_{r}(x)$, and let the mean distance of vectors in $A$ to $x$ be $\overline{\bar{r}}$. Then

$$
\bar{d}(A) \leqq \min (2 r, 2(M /(M-1)) \bar{r}(1-\bar{r} / n)) .
$$

Proof. The value $2 r$ is obvious. We assume $x=0$, and arrange the $M$ vectors in an $M \times n$ array ( $a_{i j}$ ) with column sums $s_{k}$. In column $k$, a pair of entries ( $a_{i k}, a_{j k}$ ) contribute 1 to $d\left(a_{i}, a_{j}\right)$ if and only if $a_{i k} \neq a_{j k}$. Hence

$$
\binom{M}{2} \bar{d}=\sum d\left(a_{i}, a_{j}\right)=\sum s_{k}\left(M-s_{k}\right)=M \sum s_{k}-\sum s_{k}^{2} .
$$

But $\sum s_{k}=M \bar{r}$ and by Schwarz's inequality, $\sum s_{k}^{2} \geqq 1 / n\left(\sum s_{k}\right)^{2}$ $=M^{2} \bar{r}^{2} / n$, so that

$$
\binom{M}{2} \bar{d} \leqq M^{2} \bar{r}-M^{2} \bar{r}^{2} / n
$$

and the theorem follows.
Theorem 3. $D(n, M) \leqq D\left(n-t,\left\{M / 2^{t}\right\}\right), t=0,1,2, \cdots$.
Proof. There must be a set of at least $\left\{M / 2^{t}\right\}(\{M\}$ is the smallest integer $\geqq M$ ) vectors from $A$ which agree on the first $t$ coordinates.

Using Theorems 1, 2, and 3, it is possible to obtain a two-parameter ( $r$ and $t$ ) family of upper bounds on $D(n, M)$, as follows. For each $r$, Theorem 1 guarantees that we can find a sphere of radius $r$ which contains at least $\left\{M_{r}\right\}$ vectors from $A$; Theorem 2 (with $\bar{r}$ replaced by $\min (r, n / 2)$ ) then gives an upper bound on the average distance of this subset which is also an upper bound on $d(A)$. And Theorem 3 allows us to repeat this procedure for the parameters $\left(n-t,\left\{M / 2^{t}\right\}\right)$, $t=1,2, \cdots$.

The explanation of the relationship between this procedure and the other known bounds is easily stated: If we locate the smallest $r$ for which Theorems 1 and 2 give any upper bound at all ( $M_{r}>1$ ) and apply the $2 r$ part of Theorem 2, the result is numerically the same as

Hamming's bound. If we apply Theorems 1 and 2 with the largest allowable $r(r=n)$ to the sequence of pairs $\left(n-t,\left\{M / 2^{t}\right\}\right)$ as per Theorem 3, the result is Plotkin's bound. (We conjecture that only Plotkin's bound is improved by an application of Theorem 3.) Finally, if instead of spheres of radius $r$ we use shells of radius $r \leqq n / 2$, we obtain a somewhat weaker bound. This bound is the same as the version of Elias' bound given in [2].

This procedure improves known bounds on $D(n, M)$ for even modest values of the parameters. For example $D\left(22,2^{14}\right) \leqq 6$ is given by the Hamming, Plotkin, and Elias bounds, while the procedure of this paper gives $D\left(22,2^{14}\right) \leqq 5$. Another interesting example is $D\left(53,2^{23}\right) \leqq 18$ (Hamming, Plotkin), $\leqq 17$ (Elias), and $D\left(53,2^{23}\right) \leqq 16$ by our methods. It is known [2] that Elias' bound is asymptotically better than both the Hamming and the Plotkin bounds. The bound of this paper is not asymptotically better than Elias'.

## References

1. W. W. Peterson, Error-correcting codes, Wiley, New York, 1961.
2. A. D. Wyner, On coding and information theory, J. SIAM (to appear).

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