

On Stieltjes' Continued Fraction for the Gamma Function

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Abstract. The first forty-one coefficients of a continued fraction for $\ln \Gamma(z) + z - (z - \frac{1}{2}) \ln z - \ln \sqrt{2\pi}$, are given. The computation, based on Wall's algorithm for converting a function's power series representation to a continued fraction representation, was run on the algebraic manipulation system MACSYMA.**

1. Introduction. Recall Stirling's formula for the gamma function:

$$\ln \Gamma(z) = -z + (z - \frac{1}{2}) \ln z + \ln \sqrt{2\pi} + J(z)$$

where, for $\text{real}(z) > 0$,

$$J(z) = \frac{1}{\pi} \int_0^{\infty} \ln \frac{1}{1 - e^{-2\pi u}} \cdot \frac{z}{z^2 + u^2} du.$$

Furthermore, asymptotically

$$(1) \quad J(z) = \sum_{p=0}^{\infty} (-1)^p \frac{c_p}{z^{2p+1}},$$

where

$$c_p = \frac{B_{2p+2}}{(2p+1)(2p+2)}, \quad p = 0, 1, 2, \dots,$$

and $B_2 = 1/6$, $B_4 = 1/30$, $B_6 = 1/42$, \dots , are the Bernoulli numbers. Henrici [2] refers to $J(z)$ as the *Binet function*, and gives the details for the derivation of the above formulae.

Wall [6, pp. 192-202] gives an algorithm for constructing a continued fraction development of power series such as (1), which we summarize below:

Using the symbolic operation on polynomials of *formal integration* with respect to a variable u , and an infinite sequence of numbers c_0, c_1, c_2, \dots , in which the i th

Received June 12, 1979.

AMS (MOS) subject classifications (1970). Primary 65A05, 68A15.

Key words and phrases. Gamma function, continued fraction, symbolic computation, approximation.

*Work supported in part by the U.S. Department of Energy under contract DE-AT03-76SF00034 and PADE-AS03-79ER10358. Preparation of this paper supported in part by the National Science Foundation under grant MCS #78-07291.

**Developed by the Matlab group at the Massachusetts Institute of Technology's Laboratory for Computer Science, currently supported, in part, by the U.S. Department of Energy under Contract Number E(11-1)-3070 and by the National Aeronautics and Space Administration under Grant NSG 1323.

power of u is replaced by c_i :

$$\begin{aligned} \int (k_0 + k_1u + \cdots + k_nu^n) d\phi_c(u) \\ \equiv k_0c_0 + k_1c_1 + \cdots + k_nc_n, \end{aligned}$$

one computes a_i , $i = 0, 1, 2, \dots$, of

$$\sum_{p=0}^{\infty} \frac{c_p}{z^{p+1}} = \frac{a_0}{z + \frac{a_1}{z + \frac{a_2}{z + \cdots}}}$$

by defining the auxiliary polynomials

$$(2) \quad q_{-1}(u) = 0, \quad q_0(u) = c_0,$$

initializing

$$(3) \quad a_0 = c_0$$

and using the recurrence for $p = 1, 2, 3, \dots$:

$$(4) \quad q_p(u) = uq_{p-1}(u) - a_{p-1}q_{p-2}(u),$$

$$(5) \quad a_p = \frac{\int u^p q_p(u) d\phi_c(u)}{\prod_{i=0}^{p-1} a_i},$$

where

$$(6) \quad e_p = \begin{cases} \text{if } p \text{ is even} & c_{p/2}, \\ \text{if } p \text{ is odd} & 0, \end{cases} \quad p = 1, 2, \dots$$

Stieltjes [5, pp. 520–521] gives the first five a_i for $J(z)$, noting that “Le calcul des $[a_i]$ est très pénible . . . la loi de ces nombres paraît étrangement compliqué.” However, advances of the last decade in the power of algebraic manipulation languages and systems have made it easy to use the recurrence (2)–(6) as the basis for a computer program. The MACSYMA system [3], [4] was used to compute the first forty-one a_i coefficients.

2. The Coefficients. The first seven coefficients computed via MACSYMA agree with those given by Stieltjes, and by Wall [6, p. 365]:

$$a_0 = \frac{1}{12}, \quad a_1 = \frac{1}{30}, \quad a_2 = \frac{53}{210}, \quad a_3 = \frac{195}{371}, \quad a_4 = \frac{22999}{22737},$$

$$a_5 = \frac{29944523}{19733142}, \quad a_6 = \frac{109535241009}{48264275462}.$$

where

$$b_0 = .4, \quad b_i = 12(a_{2i+1} + a_{2i}), \quad c_i = 144(a_{2i}a_{2i-1}), \quad i = 1, 2, 3, \dots$$

3. Details of the Computation. The computation was carried out on the MACSYMA Consortium Decsystem 10 (KL model) computer, located at the Massachusetts Institute of Technology's Laboratory for Computer Science in Cambridge, Massachusetts. The programming was done in the MACSYMA language [4], whose mathematical features allowed a few lines of code to completely specify the computation. The computation of the coefficients a_0 through a_{40} , as well as their 40-digit approximations, took approximately twenty minutes of central processing unit time (including the time spent on list "garbage collection"). The limiting constraint to continuing the computation is that the space requirements of the simple data representations used in the program exhaust available memory (approximately 1 million bytes) after the computation of a_{40} .

A listing of the MACSYMA program and output is available upon request from the author.

Acknowledgment. I gratefully acknowledge the illuminating commentary of W. Kahan, who brought this problem to my attention.

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