PSEUDORANDOM VECTOR GENERATION
BY THE COMPOUND INVERSIVE METHOD

FRANK EMMERICH

Abstract. Pseudorandom vectors are of importance for parallelized simulation methods. In this paper a detailed analysis of the compound inversive method for the generation of \( k \)-dimensional uniform pseudorandom vectors, a vector analog of the compound inversive method for pseudorandom number generation, is carried out. In particular, periodicity properties and statistical independence properties of the generated sequences are studied based on the discrete discrepancy of \( s \)-tuples of successive terms in the sequence. The results show that the generated sequences have attractive statistical independence properties for pseudorandom vectors of dimensions \( k \leq 4 \).

1. Introduction

The generation of uniform pseudorandom numbers in the interval \([0, 1)\) and of uniform pseudorandom vectors in the interval \([0, 1)^k\) is a basic and crucial task in any stochastic simulation. A review of several methods is given in Niederreiter’s monograph [12]. Since the simple nature of the classical linear congruential method for the generation of pseudorandom numbers and of the matrix method for the generation of pseudorandom vectors implies undesirable regularities (cf. [1]), several nonlinear methods for the generation of pseudorandom numbers and of pseudorandom vectors have been introduced [1, 2, 4, 5, 10, 11, 13]. A particularly attractive nonlinear approach is the inversive congruential method for the generation of pseudorandom numbers (cf. [1, 2, 5]) and the inversive method for the generation of pseudorandom vectors, which was introduced in [9] and is analyzed in detail in [14]. In [3] a compound version of the inversive congruential method for the generation of pseudorandom numbers was introduced in order to achieve a very long period length and an algorithm which allows a simple parallelized implementation. The analog of this approach for the generation of pseudorandom vectors, the compound inversive method, will be introduced and studied in this paper.

First, the generation of pseudorandom vectors by the (ordinary) inversive method is described, which is due to Niederreiter [14]. Let \( \mathbb{Z}_n = \{0, 1, \ldots, n-1\} \) for integers \( n \geq 1 \). Further, let \( k \geq 1 \) be an integer, the given dimension of the vectors to be generated. Choose a (large) prime \( p \), and put \( q = p^k \). Denote by \( F_q \) and \( F_q^* \) the finite field with \( q \) elements and its multiplicative group of nonzero elements, respectively. For \( \gamma \in F_q^* \) define \( \overline{\gamma} \in F_q^* \) by \( \overline{\gamma} = \gamma^{-1} \), i.e., \( \overline{\gamma} \) is the multiplicative inverse of \( \gamma \).
inverse of $\gamma$ in $F_q^*$, and put $\overline{\gamma} = 0$. Now, parameters $\alpha, \beta \in F_q$ with $\alpha \neq 0$ are selected and a sequence $\gamma_0, \gamma_1, \ldots$ of elements of $F_q$ is generated by choosing an initial value $\gamma_0$ and using the recursion

$$\gamma_{n+1} = \alpha \overline{\gamma}_n + \beta$$

for $n \geq 0$. Note that $F_q$ can be viewed as a $k$-dimensional vector space over $F_p$ (cf. [7, Chapter 1.4]). Let $B$ be an ordered basis of $F_q$ over $F_p$, and denote by $c_n \in F_p^k$ for $n \geq 0$ the coordinate vector of $\gamma_n \in F_q$ relative to $B$. Since $F_p = \mathbb{Z}/p\mathbb{Z}$ can be identified with the set $\mathbb{Z}_p = \{0, 1, \ldots, p-1\}$ of integers, each vector $c_n$ can be viewed as an element of $\mathbb{Z}_p^k$. Then

$$u_n = \frac{1}{p} c_n \in [0, 1)^k$$

for $n \geq 0$ defines an inverse sequence $(u_n)_{n \geq 0}$ of pseudorandom vectors.

Now, the compound inversive method for the generation of pseudorandom vectors is introduced. Let $m = \prod_{i=1}^r p_i$, for an integer $r \geq 1$ and arbitrary distinct primes $p_1, \ldots, p_r$. Let $m_i = m/p_i$ and $q_i = p_i^k$ for $1 \leq i \leq r$. Then for $1 \leq i \leq r$ sequences $(\gamma_n^{(i)})_{n \geq 0}$ of elements of $F_{q_i}$ are generated by choosing an initial value $\gamma_0^{(i)}$ and using the recursion

$$\gamma_{n+1}^{(i)} = \alpha_i \overline{\gamma}_n^{(i)} + \beta_i$$

for $n \geq 0$, where $\alpha_i \in F_{q_i}^*$ and $\beta_i \in F_{q_i}$ are the parameters of the $i$th generator.

Again, let $c_n^{(i)}$ be the coordinate vector of $\gamma_n^{(i)}$ relative to an ordered basis $B_i$ of $F_{q_i}$ over $F_{p_i}$, and let

$$u_n^{(i)} = \frac{1}{p_i} c_n^{(i)} \in [0, 1)^k$$

for $n \geq 0$. Then a compound inversive sequence $(u_n)_{n \geq 0}$ of pseudorandom vectors of $[0, 1)^k$ is defined by

$$u_n \equiv u_n^{(1)} + \cdots + u_n^{(r)} \pmod{1}.$$  

For $r = 1$ one has the (ordinary) inversive method, which is due to Niederreiter [14]. If $k = 1$, the compound inversive congruential method of Eichenauer-Herrmann [3] for the generation of pseudorandom numbers is obtained.

In §2 a criterion for the sequence $(u_n)_{n \geq 0}$ having the maximum possible period length $m^k$ will be established. A very important property that should be asked of pseudorandom numbers and of pseudorandom vectors for stochastic simulations is the statistical independence of successive terms in the generated sequence. A reliable theoretical approach for assessing statistical independence properties is based on the notion of discrepancy of $s$-tuples of successive terms in the sequence. Unfortunately, for $k \geq 3$ the discrepancy of an inversive sequence of pseudorandom vectors is dominated by the discretization error. So the discrete discrepancy and the discrete star discrepancy are introduced. These discrete versions of the corresponding ordinary discrepancies are natural quantities, since the considered pseudorandom vectors are rational points with fixed denominator $m$ (cf. [14]). The definition of the discrete discrepancy and of the discrete star discrepancy, as well as general discrepancy estimates, are stated in §3. In §§5 and 6 upper and lower bounds for the discrete (star) discrepancy are established. These results rely heavily on bounds for certain rational exponential sums which are stated in §4.
2. Period length

In order to obtain a criterion for the maximum possible period length \( m^k \) of a compound inversive sequence of pseudorandom vectors, a corresponding result for the underlying inversive sequences is recalled from [14]. For \( 1 \leq i \leq r \) a key role is played by the polynomial

\[
G_i(x) = x^2 - \beta_i x - \alpha_i \in F_q[x]
\]

associated with the underlying recursion in the inversive method. This polynomial has the factorization

\[
G_i(x) = (x - \sigma_i)(x - \tau_i)
\]

with nonzero roots \( \sigma_i, \tau_i \in F_q \), and the quotient \( \sigma_i \tau_i^{-1} \) of these roots is decisive in the criterion below, which is due to Niederreiter [14, Theorem 1].

**Lemma 1.** Let \( 1 \leq i \leq r \). The sequence \( u_n^{(i)}, u_1^{(i)}, \ldots \) of pseudorandom vectors generated by the inversive method has period length \( q_i = p_i^k \) if and only if the order of \( \sigma_i \tau_i^{-1} \) in the multiplicative group \( F_q^* \) is equal to \( q_i + 1 \).

In order to establish a criterion for the period length of a sequence of pseudorandom vectors generated by the compound inversive method, a sequence \( (c_n)_{n \geq 0} \) of vectors in \( \mathbb{Z}_m^k \) is defined by

\[
c_n = m_1 c_1^{(i)} + \cdots + m_r c_r^{(i)} \pmod{m},
\]

i.e., \( u_n = c_n/m \) for \( n \geq 0 \). Then the sequences \( (c_n)_{n \geq 0} \) and \( (u_n)_{n \geq 0} \) have the same period length. Obviously, the sequence \( (c_n)_{n \geq 0} \) has maximum possible period length \( m^k \) if and only if the sequences \( (c_n^{(i)})_{n \geq 0} \) have maximum possible period length \( q_i \) for \( 1 \leq i \leq r \). Since for \( 1 \leq i \leq r \) the recursion \( t_{n+1}^{(i)} = \alpha_i t_n^{(i)} + \beta_i \) is bijective in \( F_q \), the sequence \( (u_n^{(i)})_{n \geq 0} \) and also the sequence \( (u_n^{(i)})_{n \geq 0} \) are always purely periodic. Theorem 1 summarizes these properties.

**Theorem 1.** The sequence \( u_0, u_1, \ldots \) of pseudorandom vectors generated by the compound inversive method is always purely periodic and has period length \( m^k \) if and only if the underlying inversive sequences \( (u_n^{(i)})_{n \geq 0} \) of pseudorandom vectors have period length \( q_i \) for \( 1 \leq i \leq r \).

In this paper it will always be assumed that the compound inversive sequence \( (u_n)_{n \geq 0} \) has the maximum possible period length \( m^k \). This standing hypothesis will not be mentioned explicitly in the results below.

3. General estimates for discrete discrepancies

As mentioned above, a reliable theoretical approach for assessing the statistical independence properties of a sequence of pseudorandom vectors is based on the discrete (star) discrepancy of \( s \)-tuples of successive terms in the sequence. Let \( x_n = M^{-1} y_n \in [0, 1]^d \) for \( 0 \leq n < N \), where \( d \geq 1 \) and \( M \geq 2 \) are integers and \( y_0, y_1, \ldots, y_{N-1} \in \mathbb{Z}_M^d \) are \( N \) arbitrary lattice points. For any subinterval \( J \) of \( [0, 1]^d \), denote by \( L(J) \) the number of points among \( x_0, x_1, \ldots, x_{N-1} \) falling into \( J \) and let \( \text{Vol}(J) \) be the \( d \)-dimensional volume of \( J \). Then the discrete discrepancy \( E_{N,M} \) of the points \( x_0, x_1, \ldots, x_{N-1} \) is defined by

\[
E_{N,M} = E_{N,M}(x_0, x_1, \ldots, x_{N-1}) = \sup_{J \in \mathcal{J}} \left| \frac{L(J)}{N} - \text{Vol}(J) \right|,
\]
where $\mathcal{J}_M$ is the family of all subintervals $J$ of $[0,1)^d$ of the form

$$J = \prod_{j=1}^d \left( \frac{a_j}{M}, \frac{c_j}{M} \right)$$

with integers $0 \leq a_j < c_j \leq M$ for $1 \leq j \leq d$. The discrete star discrepancy $E_{N,M}^*$ of the points $x_0, x_1, \ldots, x_{N-1}$ is defined by

$$E_{N,M}^* = E_{N,M}^*(x_0, x_1, \ldots, x_{N-1}) = \sup_{J \in \mathcal{J}_M} \left| \frac{L(J)}{N} - \text{Vol}(J) \right|,$$

where $\mathcal{J}_M$ is the family of all subintervals $J$ of $[0,1)^d$ of the form

$$J = \prod_{j=1}^d \left( 0, \frac{c_j}{M} \right)$$

with integers $0 < c_j \leq M$ for $1 \leq j \leq d$. It is obvious that $E_{N,M}^* \leq E_{N,M}$.

Before general upper and lower estimates for the discrete (star) discrepancy can be stated, some further notation is necessary. Let $C_d(M)$ be the set of points $h = (h_1, \ldots, h_d)$ with integer coordinates satisfying $-M/2 < h_j \leq M/2$ for $1 \leq j \leq d$, and let $C_d^+(M) = C_d(M) \setminus \{0\}$. Further, let

$$r(h, M) = \begin{cases} M \sin(\pi|h|/M) & \text{for } h \in C_d^+(M), \\ 1 & \text{for } h = 0, \end{cases}$$

and

$$r(h, M) = \prod_{j=1}^d r(h_j, M)$$

for $h = (h_1, \ldots, h_d) \in C_d(M)$. For real $t$ the abbreviation $e(t) = e^{2\pi it}$ is used, and $\langle x, y \rangle$ stands for the standard inner product of $x, y \in \mathbb{R}^d$.

In the following, three general results for estimating discrete (star) discrepancies are stated. Lemmas 2 and 3 follow from [14, Lemmas 1 and 3], and Lemma 4 is due to Niederreiter [8, Lemma 2.3].

**Lemma 2.** Let $M \geq 2$ be an integer and $y_0, y_1, \ldots, y_{N-1} \in \mathbb{Z}_M^d$. Then the discrete discrepancy $E_{N,M}$ of the points $x_n = M^{-1}y_n \in [0,1)^d$, $0 \leq n < N$, satisfies

$$E_{N,M} \leq \sum_{h \in C_d^+(M)} \frac{1}{r(h, M)} \left| \frac{1}{N} \sum_{n=0}^{N-1} e(h \cdot x_n) \right|.$$  

**Lemma 3.** Let $M \geq 2$ be an integer and $y_0, y_1, \ldots, y_{N-1} \in \mathbb{Z}_M^d$. Then the discrete star discrepancy $E_{N,M}^*$ of the points $x_n = M^{-1}y_n \in [0,1)^d$, $0 \leq n < N$, satisfies

$$\left| \sum_{n=0}^{N-1} e(h \cdot x_n) \right| \leq \frac{2}{\pi} \left( \prod_{j=1}^d (2\pi|h_j| + 1) - 1 \right) NE_{N,M}^*$$

for any $h = (h_1, \ldots, h_d) \in \mathbb{Z}^d$ for which not all coordinates are divisible by $M$.

**Lemma 4.** Let $M \geq 2$ be an integer. Then

$$\sum_{h \in C_d^+(M)} \frac{1}{r(h, M)} < \frac{2}{\pi} \log M + \frac{2}{5}.$$
4. Exponential sums

Lemmas 2 and 3 show that a crucial role for estimating the discrete (star) discrepancy of \( s \)-tuples of successive terms in a compound inversive sequence \((u_n)_{n \geq 0}\) is played by certain exponential sums which are defined below. Subsequently, the \( ks \)-dimensional points

\[
\mathbf{v}_n = (u_n, \ldots, u_{n+s-1}) \in [0,1)^ks
\]

for \( n \geq 0 \) and

\[
\mathbf{v}^{(i)}_n = (u^{(i)}_n, \ldots, u^{(i)}_{n+s-1}) \in [0,1)^ks
\]

for \( n \geq 0 \) and \( 1 \leq i \leq r \) are considered, and the abbreviations

\[
S(h) = \sum_{n=0}^{m^k-1} e(h \cdot \mathbf{v}_n)
\]

and

\[
S_i(h) = \sum_{n=0}^{p_i^k-1} e(h \cdot \mathbf{v}^{(i)}_n)
\]

for \( 1 \leq i \leq r \) and \( h \in \mathbb{Z}^{ks} \) will be used. Some properties of these exponential sums are collected in the following two results. Lemma 6 follows from [14, Proof of Theorem 2].

**Lemma 5.** Let \( h \in \mathbb{Z}^{ks} \). Then

\[
S(h) = \prod_{i=1}^{r} S_i(h).
\]

**Proof.** First, it follows from \( \mathbf{v}_n \equiv \sum_{i=1}^{r} \mathbf{v}^{(i)}_n \mod 1 \) for \( n \geq 0 \) that

\[
S(h) = \sum_{n=0}^{m^k-1} e \left( \sum_{i=1}^{r} h \cdot \mathbf{v}^{(i)}_n \right) = \sum_{n=0}^{m^k-1} \prod_{i=1}^{r} e(h \cdot \mathbf{v}^{(i)}_n).
\]

Now, the Chinese Remainder Theorem implies that

\[
S(h) = \sum_{(n_1, \ldots, n_r) \in \mathbb{Z}^{ks}_{p_i k}} \prod_{i=1}^{r} e(h \cdot \mathbf{v}^{(i)}_{n_i}).
\]

Since the sequence \((\mathbf{v}^{(i)}_n)_{n \geq 0}\) has period length \( q_i = p_i^k \) for \( 1 \leq i \leq r \), one finally obtains

\[
S(h) = \sum_{(n_1, \ldots, n_r) \in \mathbb{Z}^{ks}_{p_i k}} \prod_{i=1}^{r} e(h \cdot \mathbf{v}^{(i)}_{n_i}) = \prod_{i=1}^{r} \sum_{n_i \in \mathbb{Z}_{p_i k}} e(h \cdot \mathbf{v}^{(i)}_{n_i}) = \prod_{i=1}^{r} S_i(h). \quad \blacksquare
\]

**Lemma 6.** Let \( h \in \mathbb{Z}^{ks} \) and \( 1 \leq i \leq r \). Then \( |S_i(h)| = p_i^k \) for \( h \equiv 0 \mod p_i \) and

\[
|S_i(h)| \leq (s-1)(2p_i^{k/2}+1)
\]

for \( h \not\equiv 0 \mod p_i \).
5. Upper bounds for the discrete discrepancy

**Theorem 2.** The discrete discrepancy \( E_{m^k,m}^{(ks)} = E_{m^k,m}(v_0, v_1, \ldots, v_{m^k-1}) \) satisfies

\[
E_{m^k,m}^{(ks)} \leq m^{-k/2} \left( \frac{2}{\pi} \log m + \frac{7}{5} \right)^{ks} \left( \prod_{i=1}^{r} \left( (s-1)(2 + p_i^{-k/2}) + p_i^{-k/2} - m^{(k-4)/2} \right) \right)
\]

for any sequence of \( k \)-dimensional pseudorandom vectors generated by the compound inversive method and all \( s \geq 2 \).

**Proof.** First, Lemma 2 is applied with \( M = m, N = m^k, d = ks \), and \( x_n = v_a \) for \( 0 \leq n < m^k \). This yields

\[
E_{m^k,m}^{(ks)} \leq \frac{1}{m^k} \sum_{h \in C_{ks}^k(m)} \frac{1}{r(h,m)} |S(h)|
\]

\[
= \frac{1}{m^k} \sum_{h \in C_{ks}^k(m)} \frac{1}{r(h,m)} \prod_{i=1}^{r} |S_i(h)|
\]

\[
= \frac{1}{m^k} \sum_{I \subseteq \{1, \ldots, r\}} \left| \prod_{|I| < r} \right| \sum_{h \in C_{ks}^k(m)} \prod_{\substack{h \equiv 0 (mod\ p_i), i \notin I \\ h \not\equiv 0 (mod\ p_i), i \notin I}} \frac{1}{r(h,m)} \prod_{i=1}^{r} |S_i(h)|,
\]

where in the second step Lemma 5 has been used. Now, denote by \( A(h) \) the number of nonzero coordinates of \( h \in \mathbb{Z}^{ks} \). Observe that always \( A(h) \geq 1 \) for \( h \in C_{ks}^k(m) \) and that \( S(h) = 0 \) for \( A(h) = 1 \). Hence, it follows that

\[
E_{m^k,m}^{(ks)} \leq \frac{1}{m^k} \sum_{I \subseteq \{1, \ldots, r\}} \left| \prod_{|I| < r} \right| \sum_{h \in C_{ks}^k(m)} \prod_{\substack{h \equiv 0 (mod\ p_i), i \notin I \\ h \not\equiv 0 (mod\ p_i), i \notin I \\ A(h) \geq 2}} \frac{1}{r(h,m)} \prod_{i=1}^{r} |S_i(h)|.
\]

Now, Lemma 6 can be applied in order to obtain

\[
E_{m^k,m}^{(ks)} \leq \frac{1}{m^k} \sum_{I \subseteq \{1, \ldots, r\}} \left| \prod_{|I| < r} \right| \sum_{h \in C_{ks}^k(m)} \prod_{\substack{h \equiv 0 (mod\ p_i), i \notin I \\ h \not\equiv 0 (mod\ p_i), i \notin I \\ A(h) \geq 2}} \frac{1}{r(h,m)} \prod_{i=1}^{r} |S_i(h)|.
\]

\[
\leq \frac{1}{m^k} \sum_{I \subseteq \{1, \ldots, r\}} \left| \prod_{|I| < r} \right| \sum_{h \in C_{ks}^k(m)} \prod_{\substack{h \equiv 0 (mod\ p_i), i \notin I \\ h \not\equiv 0 (mod\ p_i), i \notin I \\ A(h) \geq 2}} \frac{1}{r(h,m)}.
\]
where \( m_I = \prod_{i \in I} p_i \) for subsets \( I \) of \( \{1, \ldots, r\} \). Straightforward calculations show that

\[
\sum_{\substack{h \in C^*_m(m) \\ \mathbf{h} \equiv 0 \pmod{m_I} \atop A(h) \geq 2}} \frac{1}{r(h, m)} = \left( \sum_{\substack{h \in C^*_m(m) \\ \mathbf{h} \equiv 0 \pmod{m_I}}} \frac{1}{r(h, m)} + 1 \right)^{k_s} - k_s \sum_{\substack{h \in C^*_m(m) \\ \mathbf{h} \equiv 0 \pmod{m_I}}} \frac{1}{r(h, m)} - 1
\]

\[
= \left( \frac{1}{m_I} \sum_{g \in C^*_m(m/m_I)} \frac{1}{r(g, m/m_I) + 1} \right)^{k_s} - k_s \sum_{g \in C^*_m(m/m_I)} \frac{1}{r(g, m/m_I)} - 1
\]

\[
= \sum_{n=2}^{k_s} \frac{1}{m_I} \left( \sum_{g \in C^*_m(m/m_I)} \frac{1}{r(g, m/m_I)} \right)^n
\]

\[
\leq \frac{1}{m_I} \sum_{n=2}^{k_s} \left( \sum_{g \in C^*_m(m/m_I)} \frac{1}{r(g, m/m_I)} \right)^n
\]

\[
< \frac{1}{m_I} \left( 1 + \sum_{g \in C^*_m(m/m_I)} \frac{1}{r(g, m/m_I)} \right)^{k_s}.
\]

Hence, Lemma 4 implies that

\[
\sum_{\substack{h \in C^*_m(m) \\ \mathbf{h} \equiv 0 \pmod{m_I} \atop A(h) \geq 2}} \frac{1}{r(h, m)} < \frac{1}{m_I} \left( \frac{2}{\pi} \log(m/m_I) + \frac{7}{5} \right)^{k_s} \leq \frac{1}{m_I} \left( \frac{2}{\pi} \log m + \frac{7}{5} \right)^{k_s}.
\]

Altogether, one obtains

\[
E^{(k_s)}_{m^k, m} < \frac{1}{m^k} \left( \frac{2}{\pi} \log m + \frac{7}{5} \right)^{k_s} \sum_{I \subseteq \{1, \ldots, r\}} \prod_{i \in I} (s-1)(2p_i^{k/2} + 1)
\]

\[
= \frac{1}{m^k} \left( \frac{2}{\pi} \log m + \frac{7}{5} \right)^{k_s} \left( \sum_{I \subseteq \{1, \ldots, r\}} \prod_{i \in I} (s-1)(2p_i^{k/2} + 1) - m^{k-2} \right)
\]

\[
= \frac{1}{m^k} \left( \frac{2}{\pi} \log m + \frac{7}{5} \right)^{k_s} \left( \prod_{i=1}^r ((s-1)(2p_i^{k/2} + 1) + p_i^{k-2}) - m^{k-2} \right)
\]

\[
= m^{-k/2} \left( \frac{2}{\pi} \log m + \frac{7}{5} \right)^{k_s} \left( \prod_{i=1}^r ((s-1)(2 + p_i^{-k/2}) + p_i^{(k-4)/2}) - m^{(k-4)/2} \right),
\]

which is the desired result. \( \square \)

For a fixed number \( r \) of prime factors of \( m \), Theorem 2 shows that \( E^{(k_s)}_{m^k, m} = O(m^{-k/2}(\log m)^{k_s}) \) for \( k \leq 4 \) or \( r = 1 \). For \( k \geq 5 \) and \( r \geq 2 \) the order of magnitude increases, since \( \prod_{i=1}^r ((s-1)(2 + p_i^{-k/2}) + p_i^{(k-4)/2}) - m^{(k-4)/2} \) has an order
of magnitude $m_1^{(k-4)/2}$ in this case, where $m_1 = m/p_1$ and $p_1 = \min_{1 \leq i \leq r} p_i$ is assumed. Then it follows that $E^{(ks)}_{m^+ m^-} = O(m_1^{(k-4)/2} m^{-k/2} (\log m)^k)$. Lower bounds for $E^{(ks)}_{m^+ m^-}$, which will be established in the next section, show that the result of Theorem 2 is best possible up to the logarithmic factor. Further, it should be observed that this bound is independent of the specific choice of the parameters $\alpha_i$ and $\beta_i$ in the underlying recursions provided the parameters are chosen in such a way that the quotient $\sigma_i \tau_i^{-1}$ has order $q_i + 1$ in the multiplicative group $F_q^*$. It should also be observed that the bound is independent of the ordered basis $B_i$ of $F_q$ over $F_{p_i}$, which is used for the representation of the sequence $(c_n^{(i)})_{n \geq 0}$.

6. LOWER BOUNDS FOR THE DISCRETE DISCREPANCY

Theorem 1 shows that a sequence of pseudorandom vectors generated by the compound inversive method has the maximum possible period length $m^k$ if and only if all underlying (ordinary) inversive sequences of pseudorandom vectors have maximum possible period length $q_i$, which depends on the quotient of the roots $\sigma_i$, $\tau_i \in F_{p_i}$ of the polynomial

$$G_i(x) = x^2 - \beta_i x - \alpha_i = (x - \sigma_i)(x - \tau_i)$$

according to Lemma 1. Since $\sigma_i + \tau_i = \beta_i$ and $\sigma_i \tau_i = -\alpha_i$, an easy calculation shows that

$$\left( x - \frac{\sigma_i}{\tau_i} \right) \left( x - \frac{\tau_i}{\sigma_i} \right) = x^2 + \left( \frac{\beta_i^2}{\alpha_i} + 2 \right) x + 1.$$ 

Therefore, the desired property that $\sigma_i \tau_i^{-1}$ has order $q_i + 1$ in $F_q^*$ just depends on the value of $\beta_i^2 \alpha_i^{-1} \in F_q$. Now, put $\beta_i^j \alpha_i = \zeta_i$ and let $P_{q_i}$ be the set of $\zeta$’s for which the desired property is satisfied, i.e., for which $x^2 + (\zeta_i^{-1} + 2)x + 1$ has roots of order $q_i + 1$ in $F_q^*$. Further, for a fixed nontrivial additive character $\eta_i$ of $F_{q_i}$ and any $\beta_i, \zeta_i \in F_{q_i}$, put

$$K_i(\beta_i, \zeta_i) = \sum_{\gamma \in F_{q_i}} \eta_i(\beta_i \gamma + \beta_i \zeta_i \gamma).$$

The following result, which is due to Niederreiter [14, Lemma 6], will be used in the proof of a lower bound for the discrete star discrepancy in Theorem 3 below. To avoid trivial cases, it will be assumed that $q_i = p_i^k \geq 4$ for $1 \leq i \leq r$.

**Lemma 7.** Let $1 \leq i \leq r$ and $\eta_i$ be a nontrivial additive character of $F_{q_i}$. Then for any $\zeta_i \in F_{q_i}$ and $0 < t \leq \sqrt{\frac{q_i - 3}{q_i - 1}}$, there exist more than $H_{q_i}(t)(q_i - 1)$ values of $\beta_i \in F_{q_i}$ such that

$$|K_i(\beta_i, \zeta_i)| \geq t q_i^{1/2},$$

where

$$H_{q_i}(t) = \frac{1 - t^2 - 2(q_i - 1)^{-1}}{4 - t^2 + 4q_i^{1/2} + q_i^{-1}}.$$ 

**Theorem 3.** (i) Let $k \leq 4$. For $1 \leq i \leq r$ let $\zeta_i \in F_{q_i}$ and $0 < t_i \leq \sqrt{\frac{q_i - 3}{q_i - 1}}$, and let $H_{q_i}(t_i)$ be defined as in Lemma 7. Then for fixed underlying ordered bases of $F_q$ over $F_{p_i}$, there exist more than $\prod_{i=1}^r H_{q_i}(t_i)(q_i - 1)$ values of $(\beta_1, \ldots, \beta_r) \in F_{q_1} \times \cdots \times F_{q_r}$ such that the discrete star discrepancy $E_{m^+ m^-}^{(ks)} = E_{m^+ m^-}^{(ks)}(v_0, v_1, \ldots, v_{m^k-1})$
for any compound inversive sequence of \( k \)-dimensional pseudorandom vectors with \( \alpha_i = \beta_i^2 \zeta_i \) satisfies

\[
E^{*\pm(k)}_{\mu^k, m} \geq \frac{1}{8(\pi + 1)} \left( \prod_{i=1}^r t_i \right) m^{-k/2}
\]

for all \( s \geq 2 \).

(ii) Let \( k \geq 5 \) and suppose that \( p_1 = \min_{1 < i < r} p_i \). Let \( \zeta_1 \in F_{q_1} \), and let \( h_{q_1} \) be defined as in Lemma 7. Then for a fixed underlying ordered basis of \( F_{q_1} \) over \( F_{p_1} \), there exist more than \( H_{q_1}(t)/(q_1 - 1) \) values of \( \alpha \in F_{q_1} \) such that the discrete star discrepancy \( E^{*\pm(k)}_{\mu^k, m} = E^{*\pm(k)}_{\mu^k, m}(v_0, v_1, \ldots, v_{m^k-1}) \) for any compound inversive sequence of \( k \)-dimensional pseudorandom vectors with \( \alpha_1 = \beta_1^2 \zeta_1 \) satisfies

\[
E^{*\pm(k)}_{\mu^k, m} \geq \frac{t}{8(\pi + 1)m^{(k-4)/2}} m^{-k/2}
\]

for all \( s \geq 2 \).

Proof. First, Lemma 3 is applied with \( M = m \), \( N = m^k \), \( d = ks \), \( x_n = v_n \) for \( 0 < n < m^k \), and

\[
h = (h_0, \ldots, h_{s-1}) \in \mathbb{Z}^{ks},
\]

where \( h_i = (h_{i1}, \ldots, h_{is}) \in \mathbb{Z}^k \) for \( 0 \leq l \leq s - 1 \) with \( h_{0i} = h_{i1} = m_I = \prod_{i \in I} p_i \) for an arbitrary subset \( I \) of \( \{1, \ldots, r\} \) with \( |I| < r \) and all other \( h_{ij} = 0 \). This yields

\[
|S(h)| \leq \frac{2}{\pi} \left( \prod_{i=0}^{s-1} \prod_{j=1}^k (2\pi|h_{ij}| + 1) - 1 \right) m^k E^{\pm(k)}_{\mu^k, m} = S(\pi + m_I^{-1}) m_I^2 m^k E^{\pm(k)}_{\mu^k, m}.
\]

Thus, it follows from Lemma 5 that

\[
E^{\pm(k)}_{\mu^k, m} \geq \frac{1}{8(\pi + 1)m_I^2 m^k} \prod_{i=1}^r |S_i(h)| \geq \frac{1}{8(\pi + 1)m_I^2 m^k} \prod_{i=1}^r |S_i(h)|.
\]

Since \( p_i \) divides \( h_{0i} = h_{i1} = m_I \) for \( i \in I \) and all other \( h_{ij} = 0 \), it follows that \( |S_i(h)| = q_i \) for \( i \in I \) according to Lemma 6. This yields

\[
E^{\pm(k)}_{\mu^k, m} \geq \frac{1}{8(\pi + 1)m_I^2 m^k} \prod_{i \in \{1, \ldots, r\} \setminus I} |S_i(h)|.
\]

Now, let \( i \in \{1, \ldots, r\} \) with \( i \notin I \) be fixed. Let \( \zeta_i \in F_{q_i} \) and let \( \alpha_i = \beta_i^2 \zeta_i \). Subsequently, a lower bound for \( |S_i(h)| \) with \( h_{0i} = h_{i1} = m_I \) and all other \( h_{ij} = 0 \) will be established. Let \( (\delta^{(i)}_1, \ldots, \delta^{(i)}_k) \) be the dual basis \( B_i^\ast \) of the fixed underlying ordered basis \( B_i \) of \( F_{q_i} \) over \( F_{p_i} \). Further, let \( \alpha_i, \beta_i \in F_{q_i} \), be the parameters of the \( i \)th recursion. Define a nontrivial additive character \( \chi_i \) of \( F_{q_i} \) by \( \chi_i(\gamma) = e(\frac{1}{m_I} \text{Tr}(\gamma)) \) for \( \gamma \in F_{q_i} \), where \( \text{Tr} \) denotes the trace function from \( F_{q_i} \) in \( F_{p_i} \). Then

\[
|S_i(h)| = \frac{q_i - 1}{\prod_{j=1}^{s-1} p_I^{(i)}(m_I^{(i)}h)j}.
\]
where $\mu^{(i)}_l = \sum_{j=1}^k h_{ij} \delta^{(i)}_j \in F_q$, for $0 \leq l \leq s - 1$ (cf. [14, Proof of Theorem 2]). Since $h_{0,1} = h_{1,1} = m_1$ and all other $h_{i,j} = 0$, it follows that $\mu^{(i)}_0 = \mu^{(i)}_1 = m_1 \delta^{(i)}_1$ and $\mu^{(i)}_l = 0$ for $2 \leq l \leq s - 1$. Hence, one obtains

\[
S_i(h) = \sum_{n=0}^{q_i-1} \chi_i(m_1 \delta^{(i)}_1) + m_1 \delta^{(i)}_1 (\gamma^{(i)}_{n+1})
\]

\[
= \sum_{n=0}^{q_i-1} \chi_i(m_1 \delta^{(i)}_1) + m_1 \delta^{(i)}_1 (\alpha_i \gamma^{(i)}_{n} + \beta_i)
\]

\[
= \sum_{\gamma \in F_{q_i}} \chi_i(m_1 \delta^{(i)}_1) + m_1 \delta^{(i)}_1 (\alpha_i \gamma + \beta_i)).
\]

With $\eta_i(\gamma) = \chi_i(m_1 \delta^{(i)}_1) \gamma$ for $\gamma \in F_q$, it follows that

\[
|S_i(h)| = \left| \sum_{\gamma \in F_{q_i}} \eta_i(\gamma + \alpha_i \gamma + \beta_i) \right| = \left| \sum_{\gamma \in F_{q_i}} \eta_i(\gamma + \alpha_i \gamma) \right|
\]

Since $\alpha_i = \beta_i^2 \zeta_i$, one obtains

\[
|S_i(h)| = \left| \sum_{\gamma \in F_{q_i}} \eta_i(\gamma + \beta_i^2 \zeta_i \gamma) \right|
\]

Now, by changing $\beta_i^{-1} \gamma$ into $\gamma'$ in the summation, one sees that

\[
|S_i(h)| = \left| \sum_{\gamma' \in F_{q_i}} \eta_i(\beta_i \gamma' + \beta_i^2 \zeta_i \gamma') \right| = |K_i(\beta_i, \zeta_i)|.
\]

Now, let $0 < t_i \leq \sqrt{\frac{2}{q_i}}$ and let $H_q(t_i)$ be defined as in Lemma 7. Then there exist more than $H_q(t_i)(q_i - 1)$ values of $\beta_i \in F_{q_i}$ with

\[
|S_i(h)| = |K_i(\beta_i, \zeta_i)| \geq t_i q_i^{1/2}
\]

according to Lemma 7. Therefore, one obtains more than $\prod_{i \in \{1, \ldots, r\} \setminus I} H_q(t_i)(q_i - 1)$ values of $(\beta_i)_{i \in \{1, \ldots, r\} \setminus I} \in F_{q_i}^* \times \cdots \times F_{q_i}^*$ with

\[
F_{m^{k}, m}^{*} > \frac{1}{8(\pi + 1)m^{k-2}} \prod_{i \in \{1, \ldots, r\} \setminus I} t_i q_i^{1/2}
\]

\[
= \frac{1}{8(\pi + 1)m^{k-2}} \left( \prod_{i \in \{1, \ldots, r\} \setminus I} t_i \right) m^{-k/2}.
\]

(i) For $k \leq 4$ the choice $I = \emptyset$, i.e., $m_I = 1$, implies that there exist more than $F_{m^{k}, m}^{*} > \frac{1}{8(\pi + 1)} \left( \prod_{i=1}^{r} t_i \right) m^{-k/2}$. 

License or copyright restrictions may apply to redistribution; see https://www.ams.org/journal-terms-of-use
(ii) For \( k \geq 5 \) the choice \( I = \{2, \ldots, r\} \), i.e., \( m_I = m_1 \), implies that there exist more than \( H_{q_I}(t)(q_I - 1) \) values of \( \beta_1 \in E_{q_I^*} \) with
\[
E_{m^*,m}^{(k_s)}(s) \geq \frac{t}{8(\pi + 1)} m_1^{(k-4)/2} m^{-k/2}.
\]

Corollary 1. (i) Let \( k \leq 4 \). Then for fixed underlying ordered bases of \( F_q \) over \( F_p \), for \( 1 \leq i \leq r \) and any \((\zeta_1, \ldots, \zeta_r) \in P_{q_I} \times \cdots \times P_{q_r}\), there exists a value of \((\beta_1, \ldots, \beta_r) \in E_{q_I^*} \times \cdots \times E_{q_r^*}\) such that the discrete star discrepancy \( E_{m^*,m}^{(k_s)} = E_{m^*,m}(v_0, v_1, \ldots, v_{m^*-1}) \) for the sequence of \( k \)-dimensional pseudorandom vectors generated by the compound inverse method with \( \alpha_i = \beta_i^2 \zeta_i \) satisfies
\[
E_{m^*,m}^{(k_s)}(s) \geq \frac{1}{8(\pi + 1)} \left( \prod_{i=1}^r \left( \frac{q_i - 3}{q_i - 1} \right) \right) m^{-k/2}
\]
for all \( s \geq 2 \).

(ii) Let \( k \geq 5 \) and suppose that \( p_i = \min_{1 \leq i \leq r} p_i \). Then for a fixed underlying ordered basis of \( F_q \) over \( F_p \), and any \( \zeta_i \in P_{q_i} \) there exists a value \( \beta_1 \in E_{q_I^*} \) such that the discrete star discrepancy \( E_{m^*,m}^{(k_s)} = E_{m^*,m}(v_0, v_1, \ldots, v_{m^*-1}) \) for the sequence of \( k \)-dimensional pseudorandom vectors generated by the compound inverse method with \( \alpha_i = \beta_i^2 \zeta_i \) satisfies
\[
E_{m^*,m}^{(k_s)}(s) \geq \frac{1}{8(\pi + 1)} \left( \prod_{i=1}^r \left( \frac{q_i - 3}{q_i - 1} \right) \right) m_1^{(k-4)/2} m^{-k/2}
\]
for all \( s \geq 2 \).

For a fixed number \( r \) of prime factors of \( m \), Theorem 3 and Corollary 1 imply that for any underlying ordered basis of \( F_q \) over \( F_p \), and any \( \zeta_i \in P_{q_i} \) there exist compound inverse generators with discrete star discrepancy \( E_{m^*,m}^{(k_s)}(s) \) of the order of magnitude at least \( m^{-k/2} \) for \( k \leq 4 \) and of the order of magnitude at least \( m_1^{(k-4)/2} m^{-k/2} \) for \( k \geq 5 \), where \( m_1 = m/p_1 \) and \( p_i = \min_{1 \leq i \leq r} p_i \). This shows that the order of magnitude of the upper bounds for the discrete discrepancy \( E_{m^*,m}^{(k_s)}(s) \) in Theorem 2 is best possible up to the logarithmic factor. The law of the iterated logarithm for discrepancies [6] implies that the discrepancy of \( M \) independent and uniformly distributed points from \([0, 1]^d\) is almost always of an order of magnitude \( M^{-1/2}(\log \log M)^{1/2} \). A law of the iterated logarithm for discrete discrepancies is not yet known, but it is tempting to conjecture that such a result holds for the discrete discrepancy, too. In that case the lower and upper bounds for the discrete discrepancy in the compound inverse method for pseudorandom vector generation with \( k \leq 4 \) are in good accordance with such a result. On the other hand, for \( k \geq 5 \) and \( r \geq 2 \) the lower bound for the discrete discrepancy is already too large compared with such a result.

Note, that in the cases \( k = 1 \) and \( k = 2 \) similar results for the (star) discrepancy instead of the discrete (star) discrepancy can be established, since the order of the discretization error is small enough.

ACKNOWLEDGMENT

The author would like to thank a referee for valuable comments.
References


Fachbereich Mathematik, AG9, Technische Hochschule Darmstadt, Schloßgartenstraβe 7, D-64289 Darmstadt, Germany