

# ON CERTAIN SUBSETS OF FINITE BOOLEAN ALGEBRAS

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1. The boolean algebra  $B_n$ , of finite dimension  $n$ , may be written as the direct union  $B_1 \times B_1 \times \cdots \times B_1$  of  $n$  copies of  $B_1$ . Consequently each element  $u$  of  $B_n$  may be represented by an  $n$ -digit binary number. Let  $G_n$  be the group of those permutations on the elements of  $B_n$  which interchange or invert various of the factors in the direct union expansion. Thus the elements of  $G_n$  permute the components of the  $u$  or interchange 0, 1 in certain components of every  $u$  in  $B_n$ . The order of  $G_n$  is therefore  $2^n n!$ .

Two subsets of  $B_n$  will be called *congruent modulo  $G_n$*  if one is the image of the other under transformation by an element of  $G_n$ . Clearly sets congruent modulo  $G_n$  have the same number of elements. The number of elements in a subset will be called the *order* of the subset. Let  $N_n^{(s)}$  be the number of congruence classes of subsets of order  $s$ . Note that  $N_n^{(s)} = N_n^{(2^n - s)}$ . Pólya [1] has calculated  $N_n^{(s)}$  for  $0 \leq s \leq 2^n$  and  $n = 1, 2, 3, 4$ , and Slepian [2] has found the values of  $N_n = \sum_{s=0}^{2^n} N_n^{(s)}$  for  $n = 5, 6$ . Trivially,  $N_n^{(0)} = N_n^{(1)} = 1$ , all  $n$ ; and it is almost as obvious that  $N_n^{(2)} = n$ , all  $n \geq 1$  (see §2).

In this note an elementary argument is given which yields an explicit expression for  $N_n^{(3)}$  good for all  $n > 1$  (Theorem 2).

2. The procedure for calculating  $N_n^{(3)}$  is based on the notion of the "dimension" of a subset of  $B_n$ . Let  $S$  be a subset of  $B_n$  whose order is at least 2. Let  $\vee S$  and  $\wedge S$  be the lattice-union and lattice-intersection, respectively, of the elements of  $S$ . (The symbols  $\vee$ ,  $\wedge$ , and  $\leq$  will be used for the lattice operations and the ordering relation in  $B_n$ , while  $\cup$ ,  $\cap$ , and  $\subseteq$  will be reserved for their set-theoretical counterparts.) Since the quotient  $\vee S / \wedge S$  is relatively complemented, it is a boolean algebra, say  $B_r$ , if of dimension  $r$ .  $B_r$  may be called the *connected closure* of  $S$ . The relation between a subset  $S$  and its connected closure will be written:  $S < B_r$ . The *dimension* of  $S$  is defined to be the dimension of  $B_r$  if  $S < B_r$ . Thus  $S < B_r$  implies  $\dim S = r$ . If the order of  $S$  is two,  $\dim S$  is simply the usual metric in  $B_n$ .

The technique to be used for counting incongruent sets in  $B_n$  is to determine the number of incongruent sets of given order having maximal dimension in  $B_k$  for each  $k \leq n$ . Thus, for instance,  $N_n^{(2)} = n$ ,

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all  $n \geq 1$ ; for in  $B_k$  all sets of order two and maximal dimension are congruent.

LEMMA. *Suppose  $S$  is a subset of  $B_n$  of order 3 or more. If  $S = \{u\} \cup T$  and  $T < B_r$ , then  $S$  has maximal dimension if and only if  $u \in B'_r$  where  $B'_r$  is the set of the complements of the elements of  $B_r$ .*

PROOF. If  $u \in B'_r$ , then  $u \wedge (\wedge T) \leq u \wedge u' = 0$  and  $u \vee (\vee T) \geq u \vee u' = I$ , i.e.  $S = \{u\} \cup T$  has maximal dimension. But if  $u \notin B'_r$ , and  $v \in B_r$ , then either  $u \wedge v > 0$  or  $u \vee v < I$ , since complements are unique, and in this case  $\dim S$  cannot be maximal.

3. Consider now the case  $s=3$  and suppose  $S = \{u_1, u_2, u_3\}$ . Let  $r_i = \dim \{u_j, u_k\}$ , where  $i, j, k$  is some permutation of 1, 2, 3. Clearly  $1 \leq r_i \leq n$ .

THEOREM 1.  $\dim \{u_1, u_2, u_3\} = n$  if and only if  $r_1 + r_2 + r_3 = 2n$ .

PROOF. If  $u_i = x_1^i x_2^i \cdots x_n^i$ ,  $1 \leq i \leq 3$ , are the binary representations of the  $u_i$ , no generality is lost by assuming  $x_i^1 = 1$ ,  $1 \leq i \leq n$ ;  $x_i^2 = 1$ ,  $1 \leq i \leq n - r_3$ ;  $x_i^3 = 0$ ,  $n - r_3 + 1 \leq i \leq n$ . Suppose further that  $x_i^3 = 0$ ,  $1 \leq i \leq k$  and  $n - r_3 + 1 \leq i \leq n - l$ ;  $x_i^3 = 1$ ,  $k + 1 \leq i \leq n - r_3$  and  $n - l + 1 \leq i \leq n$ ; where  $k \leq n - r_3$ ,  $l \leq r_3$ . The lemma implies that  $\dim \{u_1, u_2, u_3\} = n$  if and only if  $k = n - r_3$ . But  $r_1 = k + l$ ,  $r_2 = k + r_3 - l$ , so that  $r_1 + r_2 + r_3 = 2(k + r_3)$ ; from which the theorem follows.

The expressions for the binary components also give the following immediate corollary.

COROLLARY. *If  $\{u_1, u_2, u_3\}$ ,  $\{v_1, v_2, v_3\}$  are subsets of  $B_n$  with maximal dimension, and if*

$$\dim \{u_i, u_j\} = \dim \{v_i, v_j\}, \quad 1 \leq i, j \leq 3,$$

*then  $\{u_1, u_2, u_3\}$  and  $\{v_1, v_2, v_3\}$  are congruent modulo  $G_n$ .*

Theorem 1 and the corollary imply that the number of incongruent sets of order 3 having maximal dimension in  $B_n$  is precisely the number of solutions of the following diophantine system:

$$2n = x + y + z, \quad 1 \leq x \leq y \leq z \leq n.$$

THEOREM 2.

$$N_n^{(3)} = \sum_{k=2}^n \left\{ \left[ \frac{k}{3} \right] + \sum_{r=0}^{\lfloor k/3 \rfloor} \left[ \frac{k - 3r}{2} \right] \right\}.$$

PROOF. For each  $k$ ,  $2 \leq k \leq n$ , it suffices to count the number of solutions,  $N_k^{(3)} - N_{k-1}^{(3)}$ , of

$$2k = x + y + z, \quad 1 \leq x \leq y \leq z \leq k.$$

Suppose  $z = k - r$ , then  $0 \leq r \leq [k/3]$  since  $z = 2k - x - y \geq 2k - 2z$ , or  $3z \geq 2k$ . Now  $y \leq z = k - r$  and  $x = k + r - y \geq k + r - (k - r) = 2r$ , but  $x \leq (k + r)/2$  so  $2r \leq x \leq (k + r)/2$ . For  $r = 0$ , there are  $[k/2]$  possible values for  $x$ ; and if  $r > 0$ , there are  $[(k + r)/2 - (2r - 1)]$ . Hence

$$\begin{aligned} N_k^{(3)} - N_{k-1}^{(3)} &= \left[ \frac{k}{2} \right] + \sum_{r=1}^{[k/3]} \left[ \frac{k - 3r + 2}{2} \right] \\ &= -1 + \sum_{r=0}^{[k/3]} \left[ \frac{k - 3r + 2}{2} \right] \\ &= \left[ \frac{k}{3} \right] + \sum_{r=0}^{[k/3]} \left[ \frac{k - 3r}{2} \right]. \end{aligned}$$

#### REFERENCES

1. G. Pólya, *J. Symbolic Logic*, vol. 5 (1940) pp. 98-103.
2. D. Slepian, *Canadian Journal of Mathematics* vol. 5 (1953) pp. 185-193.

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