

## FORCING AND MODELS OF ARITHMETIC<sup>1</sup>

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**ABSTRACT.** It is shown that every countable model of Peano arithmetic with finitely many extra predicates (or of ZFC with finitely many extra predicates) is a reduct of a pointwise definable such model.

This note applies the forcing method to a question concerning definability in models of Peano arithmetic.

**THEOREM.** Let  $M = \langle |M|, +, \cdot \rangle$  be a countable model of Peano arithmetic.<sup>2</sup> Then there is a set  $U \subseteq |M|$  such that

- (i)  $\langle M, U \rangle$  satisfies the first order induction schema for formulas containing an extra predicate  $U(x)$ ;
- (ii) every element of  $|M|$  is first order definable in  $\langle M, U \rangle$ .

**PROOF.** A condition is an  $M$ -finite sequence of 0's and 1's, i.e. a mapping  $p: \{b \mid b <^M a\} \rightarrow \{0, 1\}$  such that  $a \in |M|$  and  $p$  is coded by an element of  $|M|$ . We use  $p, q, \dots$  as variables ranging over conditions. A set of conditions is *dense* if every condition is extended by some condition in the set. Let  $\langle a_n \mid n < \omega \rangle$  enumerate the elements of  $|M|$ . Let  $\langle D_n \mid n < \omega \rangle$  enumerate the dense sets of conditions which are first order definable over  $M$  allowing parameters from  $|M|$ . It is safe to assume:

- (\*) the parameters in the first order definition of  $D_n$  are among  $a_0, a_1, \dots, a_{n-1}$ .

Define a sequence of conditions  $\langle p_n \mid n < \omega \rangle$  by  $p_0 = \emptyset$ ;  $p_{2n+1}$  = the  $<^M$ -least condition  $q \supseteq p_{2n}$  such that  $q \in D_n$ ;  $p_{2n+2} = p_{2n+1}$  followed by a string of  $a_n$  0's followed by a 1. Define  $U \subseteq |M|$  by letting  $\bigcup \{p_n \mid n < \omega\}$  be the characteristic function of  $U$ . To prove that  $\langle M, U \rangle$  satisfies first order induction, use the genericity of  $U$ .

[*Details.* Let  $L$  be the first order language with  $+, \cdot, U(x)$ , and constant symbols  $a$  for each  $a \in |M|$ . For  $\theta$  a sentence of  $L$  define the (strong)

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<sup>2</sup> Peano arithmetic is the theory P of Shoenfield, *Mathematical logic*, Addison-Wesley, 1967.

forcing relation  $p \Vdash \theta$  by

$p \Vdash a + b = c$	iff	$a + b = c;$
$p \Vdash a \cdot b = c$	iff	$a \cdot b = c;$
$p \Vdash U(a)$	iff	$p(a) = 1;$
$p \Vdash \theta_1 \vee \theta_2$	iff	$p \Vdash \theta_1$ or $p \Vdash \theta_2;$
$p \Vdash \neg \theta$	iff	$q \Vdash \theta$ for no $q \supseteq p;$
$p \Vdash \exists x \theta(x)$	iff	$p \Vdash \theta(a)$ for some $a \in  M .$

Prove the basic forcing lemmas as usual. It remains to show that  $\emptyset \Vdash \neg \neg (\exists x \theta(x) \rightarrow \exists \text{ least } x \text{ such that } \theta(x)).$

Suppose  $p \Vdash \exists x \theta(x)$ . Then  $p \Vdash \theta(a)$  for some  $a \in |M|$ . Working within  $M$ , define a sequence of conditions  $\langle q_c \mid c < {}^M b + 1 \rangle$  where  $b < {}^M a + 1$  as follows:  $q_0 = p$ ;  $q_{c+1} =$  the  $< {}^M$ -least  $q \supseteq q_c$  such that  $q \Vdash \neg \theta(c)$ ;  $b =$  the  $< {}^M$ -least  $c$  such that  $q_{c+1}$  is undefined. Then  $q_b \supseteq p$  and  $q_b \Vdash \neg \neg (b \text{ is the least } x \text{ such that } \theta(x)).$ ]

On the other hand, using (\*), one easily shows by induction on  $n$  that  $p_{2n+1}$ ,  $p_{2n+2}$ , and  $a_n$  are first-order definable in  $\langle M, U \rangle$ . This completes the proof.

**REMARK.** One can apply the same method to models of set theory to get the following theorem: *Let  $M = \langle |M|, \in^M \rangle$  be a countable model of ZFC, then there is a set  $U \subseteq |M|$  such that*

(i)  $\langle M, U \rangle$  satisfies the first-order replacement schema for formulas containing an extra predicate  $U(x)$ ;

(ii) every element of  $|M|$  is first-order definable in  $\langle M, U \rangle$ .

This is an improvement of the theorem of U. Felgner, *Fund. Math.* **71** (1971), 43–62.

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