

FORCING AND MODELS OF ARITHMETIC¹

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ABSTRACT. It is shown that every countable model of Peano arithmetic with finitely many extra predicates (or of ZFC with finitely many extra predicates) is a reduct of a pointwise definable such model.

This note applies the forcing method to a question concerning definability in models of Peano arithmetic.

THEOREM. Let $M = \langle |M|, +, \cdot \rangle$ be a countable model of Peano arithmetic.² Then there is a set $U \subseteq |M|$ such that

- (i) $\langle M, U \rangle$ satisfies the first order induction schema for formulas containing an extra predicate $U(x)$;
- (ii) every element of $|M|$ is first order definable in $\langle M, U \rangle$.

PROOF. A condition is an M -finite sequence of 0's and 1's, i.e. a mapping $p: \{b \mid b <^M a\} \rightarrow \{0, 1\}$ such that $a \in |M|$ and p is coded by an element of $|M|$. We use p, q, \dots as variables ranging over conditions. A set of conditions is *dense* if every condition is extended by some condition in the set. Let $\langle a_n \mid n < \omega \rangle$ enumerate the elements of $|M|$. Let $\langle D_n \mid n < \omega \rangle$ enumerate the dense sets of conditions which are first order definable over M allowing parameters from $|M|$. It is safe to assume:

- (*) the parameters in the first order definition of D_n are among a_0, a_1, \dots, a_{n-1} .

Define a sequence of conditions $\langle p_n \mid n < \omega \rangle$ by $p_0 = \emptyset$; p_{2n+1} = the $<^M$ -least condition $q \supseteq p_{2n}$ such that $q \in D_n$; $p_{2n+2} = p_{2n+1}$ followed by a string of a_n 0's followed by a 1. Define $U \subseteq |M|$ by letting $\bigcup \{p_n \mid n < \omega\}$ be the characteristic function of U . To prove that $\langle M, U \rangle$ satisfies first order induction, use the genericity of U .

[*Details.* Let L be the first order language with $+, \cdot, U(x)$, and constant symbols a for each $a \in |M|$. For θ a sentence of L define the (strong)

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² Peano arithmetic is the theory P of Shoenfield, *Mathematical logic*, Addison-Wesley, 1967.

forcing relation $p \Vdash \theta$ by

$p \Vdash a + b = c$	iff	$a + b = c;$
$p \Vdash a \cdot b = c$	iff	$a \cdot b = c;$
$p \Vdash U(a)$	iff	$p(a) = 1;$
$p \Vdash \theta_1 \vee \theta_2$	iff	$p \Vdash \theta_1$ or $p \Vdash \theta_2;$
$p \Vdash \neg \theta$	iff	$q \Vdash \theta$ for no $q \supseteq p;$
$p \Vdash \exists x \theta(x)$	iff	$p \Vdash \theta(a)$ for some $a \in M .$

Prove the basic forcing lemmas as usual. It remains to show that $\emptyset \Vdash \neg \neg (\exists x \theta(x) \rightarrow \exists \text{ least } x \text{ such that } \theta(x)).$

Suppose $p \Vdash \exists x \theta(x)$. Then $p \Vdash \theta(a)$ for some $a \in |M|$. Working within M , define a sequence of conditions $\langle q_c \mid c < {}^M b + 1 \rangle$ where $b < {}^M a + 1$ as follows: $q_0 = p$; $q_{c+1} =$ the $< {}^M$ -least $q \supseteq q_c$ such that $q \Vdash \neg \theta(c)$; $b =$ the $< {}^M$ -least c such that q_{c+1} is undefined. Then $q_b \supseteq p$ and $q_b \Vdash \neg \neg (b \text{ is the least } x \text{ such that } \theta(x)).$]

On the other hand, using (*), one easily shows by induction on n that p_{2n+1} , p_{2n+2} , and a_n are first-order definable in $\langle M, U \rangle$. This completes the proof.

REMARK. One can apply the same method to models of set theory to get the following theorem: *Let $M = \langle |M|, \in^M \rangle$ be a countable model of ZFC, then there is a set $U \subseteq |M|$ such that*

(i) $\langle M, U \rangle$ satisfies the first-order replacement schema for formulas containing an extra predicate $U(x)$;

(ii) every element of $|M|$ is first-order definable in $\langle M, U \rangle$.

This is an improvement of the theorem of U. Felgner, *Fund. Math.* **71** (1971), 43–62.