

## RECURSIVE DETERMINATION OF THE SUM-OF-DIVISORS FUNCTION

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**ABSTRACT.** A recursive scheme for determination of the sum-of-divisors function is presented. As all of the formulas involve triangular numbers, the scheme is therefore compared for efficiency with another known recursive triangular-number formula for this function.

For each positive integer  $n$ ,  $\sigma(n)$  denotes the sum of the positive divisors of  $n$ . Recently Ewell [1] derived two recursive schemes for computing the values  $\sigma(n)$ , and regarding efficiency of computation compared each of them with the well-known pentagonal-number formula of Euler. In the present discussion we derive another recursive scheme and show that among the four known recursive determinations of  $\sigma$  this one ranks no worse than second with regard to efficiency. We prepare the way for an easy statement of our result by defining an auxiliary function  $\alpha$  as follows. For each positive integer  $n$ , we express  $n$  uniquely as  $n = 2^{b(n)}O(n)$  where  $b(n)$  is a nonnegative integer and  $O(n)$  is odd. Now define  $\alpha(n) = (2^{b(n)+1} - 3)\sigma(O(n))$ .

**THEOREM.** For each positive integer  $m$ ,

$$\begin{aligned} \sigma(2m - 1) - \sum_{k=0} \{ & \alpha(2m - 1 - (2k + 1)(4k + 1)) \\ & + \alpha(2m - 1 - (2k + 1)(4k + 3)) \} \\ + \sum_{k=0} \{ & \sigma(2m - 1 - (2k + 2)(4k + 3)) \\ & + \sigma(2m - 1 - (2k + 2)(4k + 5)) \} \\ = \begin{cases} n(n + 1)/2, & \text{if } 2m - 1 = n(n + 1)/2, \\ 0, & \text{otherwise,} \end{cases} \end{aligned} \tag{1}$$

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and

$$\begin{aligned} & \alpha(2m) - \sum_{k=0}^{\infty} \{ \sigma(2m - (2k+1)(4k+1)) + \sigma(2m - (2k+1)(4k+3)) \} \\ & \quad + \sum_{k=0}^{\infty} \{ \alpha(2m - (2k+2)(4k+3)) + \alpha(2m - (2k+2)(4k+5)) \} \\ & = \begin{cases} -n(n+1)/2, & \text{if } 2m = n(n+1)/2, \\ 0, & \text{otherwise,} \end{cases} \end{aligned} \quad (2)$$

where in both (1) and (2) summation extends as far as the arguments are positive.

PROOF. Our proof depends on the following two identities due respectively to Euler and Gauss.

$$\begin{aligned} & \prod_{n=1}^{\infty} (1+x^n)(1-x^{2n-1}) = 1, \\ & \prod_{n=1}^{\infty} (1-x^{2n})(1+x^n) = \sum_{n=0}^{\infty} x^{n(n+1)/2}. \end{aligned}$$

For proofs see [2, pp. 277–284]. Since  $1-x^{2n} = (1-x^n)(1+x^n)$ , we use Euler's identity to express Gauss's identity in the equivalent form

$$\prod_{n=1}^{\infty} (1-x^n) \prod_{n=1}^{\infty} (1-x^{2n-1})^{-2} = \sum_{n=0}^{\infty} x^{n(n+1)/2}.$$

Briefly, set  $F(x) = \sum x^{n(n+1)/2}$ . Now take the logarithmic derivative of the foregoing identity and multiply the resulting identity by  $x$  to obtain

$$\sum_{n=1}^{\infty} \frac{nx^n}{1-x^n} - 2 \sum_{n=1}^{\infty} \frac{(2n-1)x^{2n-1}}{1-x^{2n-1}} = -\frac{x F'(x)}{F(x)}. \quad (3)$$

It is well known that the first series on the left of (3), a "Lambert" series, generates  $\sigma(n)$ : i.e.,

$$\sum_{n=1}^{\infty} \frac{nx^n}{1-x^n} = \sum_{n=1}^{\infty} \sigma(n)x^n.$$

The second series is perhaps less well known, but straightforward algebraic manipulation shows that

$$\begin{aligned} \sum_{n=1}^{\infty} \frac{(2n-1)x^{2n-1}}{1-x^{2n-1}} &= \sum_{n=1}^{\infty} x^n \sum_{\substack{d|n \\ d \text{ odd}}} d \\ &= \sum_{m=1}^{\infty} \sigma(2m-1)x^{2m-1} + \sum_{m=1}^{\infty} \sigma(O(2m))x^{2m}. \end{aligned}$$

Thus, identity (3) becomes

$$\sum_{m=1}^{\infty} \sigma(2m-1)x^{2m-1} - \sum_{m=1}^{\infty} \alpha(2m)x^{2m} = x F'(x)/F(x).$$

Now, separating the even and odd triangular numbers  $n(n + 1)/2$  by the least positive residues of  $n \pmod{4}$ , we write

$$\begin{aligned} & \left\{ \sum_{m=1}^{\infty} \sigma(2m - 1)x^{2m-1} \right\} F(x) \\ &= \sum_{m=1}^{\infty} x^{2m-1} \left\{ \sigma(2m - 1) + \sum_{k=0}^{\infty} [\sigma(2m - 1 - (2k + 2)(4k + 3)) \right. \\ & \qquad \qquad \qquad \left. + \sigma(2m - 1 - (2k + 2)(4k + 5))] \right\} \\ &+ \sum_{m=1}^{\infty} x^{2m} \left\{ \sum_{k=0}^{\infty} [\sigma(2m - (2k + 1)(4k + 1)) \right. \\ & \qquad \qquad \qquad \left. + \sigma(2m - (2k + 1)(4k + 3))] \right\}, \end{aligned}$$

and,

$$\begin{aligned} & \left\{ \sum_{m=1}^{\infty} \alpha(2m)x^{2m} \right\} F(x) \\ &= \sum_{m=1}^{\infty} x^{2m-1} \left\{ \sum_{k=0}^{\infty} [\alpha(2m - 1 - (2k + 1)(4k + 1)) \right. \\ & \qquad \qquad \qquad \left. + \alpha(2m - 1 - (2k + 1)(4k + 3))] \right\} \\ &+ \sum_{m=1}^{\infty} x^{2m} \left\{ \alpha(2m) + \sum_{k=0}^{\infty} [\alpha(2m - (2k + 2)(4k + 3)) \right. \\ & \qquad \qquad \qquad \left. + \alpha(2m - (2k + 2)(4k + 5))] \right\}. \end{aligned}$$

We then substitute these last two developments into the identity

$$\begin{aligned} & \left\{ \sum_{m=1}^{\infty} \sigma(2m - 1)x^{2m-1} \right\} F(x) - \left\{ \sum_{m=1}^{\infty} \alpha(2m)x^{2m} \right\} F(x) \\ &= xF'(x) = \sum_{n=1}^{\infty} \frac{n(n + 1)}{2} x^{n(n+1)/2}, \end{aligned}$$

equate coefficients of odd powers  $x^{2m-1}$  to obtain recurrence (1) and equate coefficients of even powers  $x^{2m}$  to obtain recurrence (2).

REMARKS. One of the three recurrences discussed in [1] is

$$\begin{aligned} & \sum_{k=0}^{\infty} (-1)^k (2k + 1)\sigma(n - k(k + 1)/2) \\ &= \begin{cases} (-1)^{j+1} j(j + 1)(2j + 1)/6, & \text{if } n = j(j + 1)/2, \\ 0, & \text{otherwise.} \end{cases} \quad (4) \end{aligned}$$

Like our recurrences (1) and (2) this recurrence involves triangular numbers. Since  $\sigma$  is multiplicative and therefore  $\sigma(n) = (2^{b(n)+1} - 1)\sigma(O(n))$ , we suppose that we are given a large odd number  $n$  and investigate efficiency of computation of  $\sigma(n)$  by recurrences (1) and (4). Theoretically we deduce that each of the recurrences needs about  $\sqrt{2n}$  of the values,  $\sigma(j)$ ,  $1 < j < n$ . But, practically, let us take a not-too-large value of  $n$ , say  $n = 63$ , partially compute  $\sigma(63)$  by each of the recurrences and possibly observe some noteworthy differences.

By recurrence (1),

$$\begin{aligned}\sigma(63) = & (2^2 - 3)\sigma(31) + (2^5 - 3)\sigma(3) + (2^2 - 3)\sigma(9) + (2^3 - 3)\sigma(15) \\ & + (2^2 - 3)\sigma(21) + (2^4 - 3)\sigma(1) - \sigma(57) - \sigma(35) - \sigma(53) - \sigma(27).\end{aligned}$$

By recurrence (4),

$$\begin{aligned}\sigma(63) = & 3\sigma(62) - 5\sigma(60) + 7\sigma(57) - 9\sigma(53) + 11\sigma(48) - 13\sigma(42) \\ & + 15\sigma(35) - 17\sigma(27) + 19\sigma(18) - 21\sigma(8).\end{aligned}$$

Although each recurrence uses 10 lower values, we observe that recurrence (1) uses smaller values  $\sigma(j)$  by separating the binary and odd parts of  $j$ . (This is something that any high-speed computing machine can do easily.) Also, recurrence (1) avoids coefficients such as the  $2k + 1$  of (4).

#### REFERENCES

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