

LOGICS WITH GIVEN CENTERS AND STATE SPACES

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ABSTRACT. Let B be a Boolean algebra and let K be a compact convex subset of a locally convex topological linear space. Then there exists a logic with the center Boolean isomorphic to B and with the state space affinely homeomorphic to K .

Introduction. In the quantum logic approach to the foundations of quantum mechanics, one identifies the event structure of a system with an orthomodular partially ordered set L (called usually a logic). The set of states is then represented by the set $\mathfrak{S}(L)$ of all probability measures on L (see [4, 7]). It can be shown that $\mathfrak{S}(L)$ is a compact convex set and conversely, it was proved by F. W. Shultz [6] that any compact convex subset of a locally convex topological linear space is affinely homeomorphic to $\mathfrak{S}(L)$ for a logic L .

The center $C(L)$ of a logic L is the subset of L consisting of all "absolutely compatible" elements. It is known that the center of L is a Boolean algebra (see [1, 4]). Obviously, any Boolean algebra is the center of a logic.

Let us now consider the center and the state space simultaneously. The question is if for any Boolean algebra B and any compact convex subset of a LCTLS there exists a logic L such that $C(L) = B$ and $\mathfrak{S}(L) = K$. We answer the question in the affirmative. In the construction we use, among other tools, the result of Shultz [6] and the technique of R. Greechie [2] for constructing orthomodular posets.

Notions. Results. Let us first review the basic definitions and state some auxiliary propositions.

DEFINITION 1. A logic is a set L endowed with a partial ordering \leq and a unary operation $'$ such that:

- (i) $0, 1 \in L$;
- (ii) $a \leq b \Rightarrow b' \leq a'$ for any $a, b \in L$;
- (iii) $(a')' = a$ for any $a \in L$;
- (iv) $a \vee a' = 1$ for any $a \in L$;
- (v) $\bigvee_{i=1}^n a_n$ exists in L whenever $a_n \in L$, $a_n \leq a'_k$ for $n \neq k$;
- (vi) $b = a \vee (b \wedge a')$ whenever $a, b \in L$, $a \leq b$.

In the sequel, we shall reserve the symbol L for logics. One can prove easily that if $a, b \in L$, $a \leq b'$ then $a \vee b$, $a \wedge b$ exists in L .

DEFINITION 2. Two elements $a, b \in L$ are called compatible if there are three elements $c, d, e \in L$ such that $c \leq d'$, $d \leq e'$, $e \leq c'$ and $a = c \vee d$, $b = c \vee e$.

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DEFINITION 3. An element $a \in L$ is called central if a is compatible to any element of L . We denote by $C(L)$ the set of all central elements of L and call $C(L)$ the center of L .

PROPOSITION 1. *The set $C(L)$ with the operations $'$, \vee , \wedge inherited from L is a Boolean algebra.*

PROOF. See [1, 4].

DEFINITION 4. Let $\{L_\alpha \mid \alpha \in I\}$ be a collection of logics. Denote by $\prod_{\alpha \in I} L_\alpha$ the ordinary Cartesian product of the sets L_α and endow the set $\prod_{\alpha \in I} L_\alpha$ with the relation \leq and the unary operation $'$ as follows. If $k = \{k_\alpha \mid \alpha \in I\} \in \prod_{\alpha \in I} L_\alpha$ and $h = \{h_\alpha \mid \alpha \in I\} \in \prod_{\alpha \in I} L_\alpha$, then $k \leq h$ (resp. $k' = h$) if and only if $k_\alpha \leq h_\alpha$ (resp. $k'_\alpha = h_\alpha$) for any $\alpha \in I$. The set $\prod_{\alpha \in I} L_\alpha$ with the above defined \leq , $'$ is called the product of the collection $\{L_\alpha \mid \alpha \in I\}$.

PROPOSITION 2. *Let $\{L_\alpha \mid \alpha \in I\}$ be a collection of logics. Then $\prod_{\alpha \in I} L_\alpha$ is a logic. If $C(L_\alpha) = \{0, 1\}$ for any $\alpha \in I$ then $C(\prod_{\alpha \in I} L_\alpha)$ is Boolean isomorphic to the Boolean algebra of all subsets of I .*

PROOF. See [3, 5].

DEFINITION 5. A state on a logic L is a mapping $s: L \rightarrow \langle 0, 1 \rangle$ such that:

- (i) $s(1) = 1$;
- (ii) if $a, b \in L$, $a \leq b'$ then $s(a \vee b) = s(a) + s(b)$.

Let us denote by $\mathfrak{S}(L)$ the set of all states on L . By a result of F. W. Shultz [6], any compact convex subset of a LCTLS equals, up to an affine homeomorphism, $\mathfrak{S}(L)$ for a logic L (and vice versa, which is obvious).

DEFINITION 6. A logic L is called poor (resp. rigid) if $\mathfrak{S}(L) = \emptyset$ (resp. $|\mathfrak{S}(L)| = 1$). It is known (see [2, 6]) that there are (finite) examples of poor and rigid logics.

PROPOSITION 3. *Suppose that L is a poor logic. Put $L_\alpha = L$ for any $\alpha \in I$. Then $\prod_{\alpha \in I} L_\alpha$ is also a poor logic.*

PROOF. Take the mapping $f: L \rightarrow \prod_{\alpha \in I} L_\alpha$ such that $f(k) = (k, k, k, \dots)$ for any $k \in L$. If $s \in \mathfrak{S}(\prod_{\alpha \in I} L_\alpha)$ then $sf \in \mathfrak{S}(L)$.

DEFINITION 7. A mapping $f: L_1 \rightarrow L_2$ is called an embedding if f is injective and the following requirements are satisfied.

- (i) $f(1) = 1$;
- (ii) $f(a') = f(a)'$ for any $a \in L_1$;
- (iii) $a \leq b$ if and only if $f(a) \leq f(b)$;
- (iv) if $a \leq b'$ then $f(a \vee b) = f(a) \vee f(b)$.

PROPOSITION 4. *Let K be a compact convex subspace of a LCTLS. Take the logic L_1 constructed in [6, Theorem, p. 321]. Thus $\mathfrak{S}(L_1) = K$ and moreover, $C(L_1) = \{0, 1\}$ and L_1 can be embedded into a poor logic L_2 with $C(L_2) = \{0, 1\}$.*

PROOF. We must assume here that the reader is well acquainted with the paper [6] and with the Greechie representation of logics (see [2]). It follows immediately from

the construction of [6] that $C(L_1) = \{0, 1\}$ (see e.g. the plan of the construction, p. 321). Further, let us consider the Greechie diagram D_1 of L_1 and the Greechie diagram D of a finite poor logic L exhibited in [2]. Let us choose “points” $d_1 \in D_1$, $d_2 \in D$ such that d_1, d_2 belong to exactly one Boolean block of D_1, D . Form a new Greechie diagram D_2 by taking the union $D_1 \cup D$ and then “identifying” the points d_1, d_2 . The diagram D_2 then represents the required logic L_2 .

We are now ready to prove our result.

THEOREM. *Let B be a Boolean algebra and let K be a compact convex subset of a LCTLS. Then there exists a logic L such that $C(L)$ is Boolean isomorphic to B and $\mathfrak{S}(L)$ is affinely homeomorphic to K .*

PROOF. We may suppose that B is a Boolean algebra of subsets of a set A . Take a logic M such that $C(M) = \{0, 1\}$, $\mathfrak{S}(M) = K$ and denote by P the poor extension of M (Proposition 4). Take a point $a \in A$ and write $L_c = P$ if $c \in A - \{a\}$, $L_a = M$. Consider the logic $R = \prod_{d \in A} L_d$. The desired logic L will now be obtained as a sublogic of R . Let us describe the elements of L . An element $r \in R$ belongs to L if and only if there exists a finite partition \mathfrak{P} of A , $\mathfrak{P} = \{A_i \mid i = 1, 2, \dots, n\}$ such that $A_i \in B$ for any i , $1 \leq i \leq n$, and $r_p = r_q$ as soon as $\{p, q\} \subset A_i$ for an index i , $1 \leq i \leq n$. We are to show that L is a logic with $C(L) = B$ and $\mathfrak{S}(L) = K$.

Obviously, $1 \in L$ and if $k \in L$ then $k' \in L$. If $k, h \in L$, $k \geq h$ then $k = h \vee (k \wedge h')$. Indeed, if $\mathfrak{P}, \mathfrak{R}$ are partitions corresponding to k, h then $\mathfrak{P} \cap \mathfrak{R}$ is the partition corresponding to $k' \wedge h$. The rest is obvious. Thus L is a logic.

Further, since $C(L_d) = \{0, 1\}$ for any $d \in A$ then any central element of L must have only the elements 0, 1 for the coordinates. One can check easily that $k = \{k_d \mid d \in A\}$, where any k_d is either 0 or 1, belongs to L if and only if $D = \{d \mid k_d = 1\} \in B$. Consequently, $C(L) = B$.

It remains to prove that $\mathfrak{S}(L) = K$. Since $\mathfrak{S}(M) = K$, it suffices to show that there is an affine homeomorphism $g: \mathfrak{S}(L) \rightarrow \mathfrak{S}(M)$. Assume that $s \in \mathfrak{S}(L)$. For any $m \in M$, denote by k^m the element of L which has m for all its coordinates. Define $g(s)$ such that $g(s)(m) = s(k^m)$. We need to show that g is injective.

Let us suppose that $g(s_1) = g(s_2)$. Take an element $k \in L$ and assume that \mathfrak{P} is the partition corresponding to k . Let A_1 be such a set of \mathfrak{P} that $a \in A_1$. Denote by $h = \{h_d \mid d \in A\}$ the element of L with $h_d = 0$ if $d \in A_1$, $h_d = 1$ otherwise. It follows from Proposition 3 that $s_1(k \wedge h) = s_2(k \wedge h) = 0$. Since $g(s_1) = g(s_2)$, we see, again applying Proposition 3, that $s_1(k) = s_1(k \wedge h') = s_2(k \wedge h') = s_2(k)$. Hence the mapping $g: \mathfrak{S}(L) \rightarrow \mathfrak{S}(M)$ is injective and the proof is complete.

Let us state explicitly the following special corollary.

COROLLARY. *Given a Boolean algebra B , there exists a poor (resp. rigid) logic L such that $C(L) = B$.*

Let us observe in conclusion that a similar method yields an analogous result for σ -complete logics and σ -additive states. Naturally, the center then cannot be arbitrary since there are Boolean σ -algebras without any σ -additive state.

THEOREM. *Let B be a Boolean σ -algebra of subsets of a set and let K be a compact convex subset of a LCTLS. Then there is a σ -complete logic L such that $C(L)$ is Boolean σ -isomorphic to B and the space of σ -additive states on L is affinely homeomorphic to K .*

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