

## NEWTON'S METHOD AND THE JENKINS-TRAUB ALGORITHM

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ABSTRACT. In this paper we propose to show how a multi-increment version of Newton's method can be used to obtain starting points for the Jenkins-Traub algorithm.

### 1. The Jenkins-Traub algorithm. Let

$$(1.0) \quad P(z) = \prod_{j=1}^j (z - z_j)^{\alpha_j}, \quad \sum \alpha_j = n,$$

where the  $z_j$  are distinct roots of multiplicity  $\alpha_j$ . The single stage algorithm of Jenkins and Traub starts by creating a sequence of monic polynomials  $P^{(\nu)}(z)$  with

$$(1.1) \quad P^{(1)}(z) = \frac{1}{n} \frac{d}{dz} P(z),$$

and for  $\nu \geq 2$

$$(1.2) \quad P^\nu(z) = \left( P(z) - \frac{P(a)}{P^{\nu-1}(a)} P^{\nu-1}(z) \right) \cdot \frac{1}{(z - a)}$$

where  $a$  is an arbitrary complex number. They show that if

$$(1.3) \quad |z_1 - a| < |z_j - a|, \quad j \geq 2,$$

then

$$(1.4) \quad \lim_{\nu \rightarrow \infty} \frac{P(z)}{P^\nu(z)} = z - z_1.$$

However good the probability of an arbitrary complex number satisfying (1.3), it is destined to certain failure when  $P$  has real coefficients,  $a$  is real and the nearest root is nonreal. Jenkins and Traub overcome this difficulty by constructing an alternative sequence which converges to the quadratic factor

$$(1.5) \quad (z - z_1)(z - z_2)$$

under the condition

$$(1.6) \quad |z - z_1| = |z - z_2| < |z - z_j|, \quad j \geq 3.$$

In this paper we show that the same objective can be realized by saving the data from the two previous iterations.

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THEOREM 1. If  $P$  defined by (1.0) satisfies (1.6), then with

$$(1.7) \quad Q_\nu(z) = \frac{P^{\nu-1}(a)P^\nu(z)P^{\nu-2}(z) - P^{\nu-2}(a)(P_{\nu-1}(z))^2}{P_{\nu-1}(a) - P_{\nu-2}(a)}$$

we have

$$(1.8) \quad \lim_{\nu \rightarrow \infty} \frac{[P(z)]^2}{Q_\nu(z)} = (z - z_1)(z - z_2).$$

PROOF. Let

$$(1.9) \quad F^\nu(z) = \left[ F^{\nu-1}(z) - \frac{F^{\nu-1}(a)P(z)}{P(a)} \right] \cdot \frac{1}{(z - a)}$$

with initial condition

$$(1.10) \quad F^{(1)}(z) = \frac{1}{n} \frac{d}{dz} (P(z)).$$

It follows from (1.1), (1.2), (1.9) and (1.10) and the fact that  $P$  is monic that

$$(1.11) \quad P^\nu(z) = -\frac{P(a)}{F_{\nu-1}(a)} F^\nu(z)$$

and, conversely, that

$$(1.2) \quad F^\nu(z) = -\frac{F^{\nu-1}(a)}{P(a)} P^\nu(z).$$

Now, as Jenkins and Traub observe, using the Lagrange interpolation polynomials

$$(1.13) \quad P_k(z) = P(z)/(z - z_k),$$

we have the following expression for the  $F^\nu$ 's:

$$(1.14) \quad F^\nu(z) = \sum_{i=1}^j \frac{\alpha_i P_i(z)}{(z_i - a)^{\nu-1}}, \quad \nu = 1, 2, \dots$$

Now let us define

$$(1.15) \quad \bar{Q}_\nu(z) = F_\nu(z)F_{\nu-2}(z) - [F_{\nu-1}(z)]^2$$

and the corresponding monic polynomial

$$(1.16) \quad Q_\nu(z) = \frac{\bar{Q}_\nu(z)}{\text{coeff of } z^{2n-2} \text{ in } \bar{Q}_\nu}.$$

The function  $\bar{Q}_\nu(z)$  is a symmetric quadratic form with zeros on the diagonal. Hence after substituting (1.14) into (1.15) and symmetrizing, we have

$$(1.17) \quad \bar{Q}_\nu(z) = \sum_{i < j} \frac{\alpha_i \alpha_j P_i(z) P_j(z)}{(a - z_i)^{\nu+1} (a - z_j)^{\nu+1}} (z_i - z_j)^2.$$

Since the  $P_i$ 's are monic, it follows that the coefficient of  $z^{2n-2}$  is

$$(1.18) \quad \sum_{i < j} \frac{\alpha_i \alpha_j (z_i - z_j)^2}{(a - z_i)^{\nu+1} (a - z_j)^{\nu+1}}.$$

It now is a consequence of (1.3) that

$$(1.19) \quad \lim_{z \rightarrow \infty} \frac{F^2(z)}{Q_v(z)} = (z - z_1)(z - z_2).$$

To complete the proof, we have only to use (1.11) and (1.12) to show that  $Q_v$  has the alternative expression (1.7).

**2. Starting points.** In this section we show that if we take a starting point  $z$  outside a circle containing the roots of a polynomial in its interior and multiply the Newton iteration by  $n$ , we get into the interior of the circle.

**THEOREM 2.** *If the roots of  $P$ , (1.0), satisfy*

$$(2.1) \quad |z_i - \alpha| < R,$$

then with  $|z| > R$

$$(2.2) \quad \beta = z - nf(z)/f'(z)$$

satisfies

$$(2.3) \quad |\beta - \alpha| < R.$$

**PROOF.** We may without loss of generality assume that

$$(2.4) \quad \alpha = 0 \quad \text{and} \quad R = 1.$$

Since

$$(2.5) \quad \frac{f'(z)}{f(z)} = \sum_{i=1}^j \frac{\alpha_i}{z - z_i},$$

with

$$(2.6) \quad g(z) = z - nf(z)/f'(z),$$

we have

$$(2.7) \quad g(z) = \sum_{i=1}^n \frac{\alpha_i z_i}{z - z_i} / \sum_{i=1}^j \frac{\alpha_i}{z - z_i}.$$

For a fixed  $z$ ,  $|z| > 1$ ,  $g(z)$  is an analytic function of the complex variables  $z_i$  in the polydisc

$$(2.8) \quad |z_i| < 1, \quad i = 1, 2, \dots, n.$$

Hence by the maximum principle we may assume that

$$(2.9) \quad |z_i| = 1, \quad i = 1, 2, \dots, n.$$

Now, by the theorem of Gauss and Lucas, the zeros of the derivative  $P'(z)$  are in the convex hull of

$$(2.10) \quad z_1, \dots, z_n$$

and so are in a compact subset of the interior of the unit circle. Hence for each fixed set (2.10) satisfying (2.9),  $g$  is analytic and bounded for  $|z| \geq 1$ . It therefore assumes its maximum for  $|z| = 1$  or at  $\infty$ . But

$$(2.11) \quad \lim_{z \rightarrow \infty} g(z) = \sum_{i=1}^n \frac{\alpha_i z_i}{n}$$

which is a convex combination of the  $z_i$ 's and so has average less than 1. For  $|z| = 1$ , we may write the denominator in (2.7) as

$$(2.12) \quad \sum_{i=1}^j \frac{\alpha_i(\bar{z} - \bar{z}_i)}{|z - z_i|^2} = \frac{1}{z} \sum_{i=1}^n \frac{\alpha_j(1 - \bar{z}z)}{|z - z_i|^2}$$

and the numerator as

$$(2.13) \quad \sum_{i=1}^j \frac{\alpha_i z_i(\bar{z} - \bar{z}_i)}{|z - z_i|^2} = - \sum_{i=1}^n \frac{\alpha_i(1 - \bar{z}z_i)}{|z - z_i|^2}.$$

Hence  $|g(z)| = 1$  for  $|z| = 1$ . This completes the proof.

**3. Automatic deflation.** It follows from (1.4) that

$$(3.1) \quad \lim_{\nu \rightarrow \infty} P^\nu(z) = \frac{P(z)}{z - z_1} = P_1(z),$$

the Lagrange interpolation polynomial corresponding to the factor  $z - z_1$ . Hence if we use  $z_1$  as starting point and compute a root of  $P$ , we may be certain that we do not compute  $z_1$  again when it is simple. If it is multiple, we shall compute  $z_1$  again and by iterating find its multiplicity. Moreover it is clear from (1.14) that if we keep  $a$  a reasonable distance from a multiple root, this method of deflation is more accurate than applying the division algorithm. Since all we ask of the initial approximation is that it be closer to the root it approximates than any other root, this procedure appears to have good prospects.

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