

EVERY NORMAL BAND WITH (REP) AND $(REP)^{op}$ IS AN AMALGAMATION BASE

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ABSTRACT. We shall prove that every normal band with the representation extension property and its dual is an amalgamation base in the class of all semigroups.

1. INTRODUCTION

A semigroup S is called an *amalgamation base* in the class of all semigroups (simply called an amalgamation base), if for any semigroups T_1, T_2 containing S as a subsemigroup the amalgam $[T_1, T_2; S]$ is embedded into a semigroup. A semigroup S has the *representation extension property* (denoted by (REP)) if for every embedding $S \rightarrow T$ of semigroups and every right S -set X , the canonical map: $X \rightarrow X \otimes T^1$ is injective (see [2, 6, 7]). The left-right dual of (REP) is denoted by $(REP)^{op}$. Hall [6] showed that any semigroup which is an amalgamation base always has (REP) and $(REP)^{op}$. The author [9] constructed an example of a monoid which has (REP) and $(REP)^{op}$ but is not an amalgamation base. However, such an example of regular semigroups is still unknown. In this direction, Bulman-Fleming and McDowell [4] determined the structure of normal bands with (REP) and $(REP)^{op}$ and, consequently, showed that every right (left) normal band with (REP) and $(REP)^{op}$ is left (right) absolutely flat (see [3]) and hence is an amalgamation base. The purpose of this paper is to prove the following stronger result.

Main Theorem. *A normal band has both (REP) and $(REP)^{op}$ if and only if it is an amalgamation base.*

Our method is to appeal the criterion for an amalgamation base given in [9], which is a modified version of Renshaw's Theorem [8, Theorem 6.11].

2. PRELIMINARIES

Throughout this paper, let S denote a semigroup and S^1 the semigroup with the adjoined identity 1 whether S has an identity or not. Let $\mathcal{I} [\mathcal{L}, \mathcal{R}]$ denote Green's \mathcal{I} - $[\mathcal{L}$ -, \mathcal{R} -] relation on a semigroup. We often use the notation

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and conventions from Clifford and Preston’s book [5] for semigroup theory. Let $S\text{-Ens}$ ($\text{Ens-}S$, $S\text{-Ens-}S$) denote the category of all left S -sets (right S -sets, S -bisets). Let $X \in \text{Ens-}S$ and $Y \in S\text{-Ens}$. The tensor product over S of X and Y is denoted by $X \otimes_S Y$ (simply, $X \otimes Y$ if there is no confusion). Also, any element of $X \otimes Y$ is written in a form $x \otimes y$ ($x \in X$, $y \in Y$). For brevity, $X \supset Y$ ($X, Y \in S\text{-Ens}$ ($\text{Ens-}S, S\text{-Ens-}S$)) means that Y is a left S - (right S -, S -bi) subset of X .

We will use the following results in the sequel.

Result 1 [9, Theorem 2.1]. *A semigroup S has (REP) if and only if, for each $M \in S\text{-Ens}$ with $M \supset S^1$ and each $X \in \text{Ens-}S$, the map: $X \rightarrow X \otimes M$ ($x \mapsto x \otimes 1$) is injective.*

Result 2 [9, Theorem 2.2]. *A semigroup S is an amalgamation base if and only if for each $X \in \text{Ens-}S$, $Y \in S\text{-Ens}$, and $N \in S\text{-Ens-}S$ with $N \supset S^1$, the map: $X \otimes Y \rightarrow X \otimes N \otimes Y$ ($x \otimes y \mapsto x \otimes 1 \otimes y$) is injective.*

We recall that a normal band satisfies the identity $xyzx = xzyx$ (equivalently, $xyza = xzya$).

For a normal band S , let $S = \bigcup\{S_\lambda : \lambda \in \Lambda\}$ be the semilattice decomposition. In this case each S_λ is a \mathcal{F} -class of S . So by using the partial order \geq on Λ , we define a quasi-order $\geq_{\mathcal{F}}$ on S by $s \geq_{\mathcal{F}} t$ ($s, t \in S$) if and only if $\mathcal{F}_s \geq \mathcal{F}_t$. Then, for convenience, we sometimes write $t \leq_{\mathcal{F}} s$. Also, $s >_{\mathcal{F}} t$ means both $\mathcal{F}_s \geq \mathcal{F}_t$ and $\mathcal{F}_s \neq \mathcal{F}_t$. If necessary, we extend the quasi order $\geq_{\mathcal{F}}$ from S to S^1 . Clearly, $1 >_{\mathcal{F}} s$ in S^1 for all $s \in S$.

Result 3 [4, Theorem 1]. *A normal band $S = \bigcup\{S_\lambda : \lambda \in \Lambda\}$ has (REP) and (REP)^{op} if and only if S has the following:*

- (i) $uau = vav$ for any $u, v, a \in S$ with $u \mathcal{F} v, u >_{\mathcal{F}} a$;
- (ii) $|S_\lambda| \leq 2$ for each $\lambda \in \Lambda$; and
- (iii) if $|S_\lambda| = 2$ ($\lambda \in \Lambda$) then $\bigwedge S_\lambda$ does not exist with respect to the natural ordering \geq of S .

3. PROOF OF THE MAIN THEOREM

To prove the main theorem, it suffices to prove the “only if” part. In this section, we let S be a normal band with (REP) and (REP)^{op}. Then we shall show first the preliminary lemmas.

Lemma 1. *Let S be as above, and $a, u, v \in S$. Let $X \in \text{Ens-}S$, $Y \in S\text{-Ens}$, $x, x' \in X$, and $y, y' \in Y$. Then:*

- (i) $xu = x'v$ implies $xuau = x'vav$; and
- (ii) $uy = vy'$ implies $uauy = vavy'$.

Proof. (i) If $uv >_{\mathcal{F}} uva$, then $vuauv = (vu)^2a(uv)^2 = uv(vuauv)vu$ (by Result 3(i)) = $uvavu$, so that $xuau = xuvau = x(uvavu) = x'v(uvavu) = x'v(vuauv) = x'(vuav) = x'vav$. If $uv \mathcal{F} uva$, then $xuau = xuvau = x(uvu) = xu$, and similarly $x'vav = x'v$. Hence (i) holds.

(ii) Similarly. \square

Lemma 2 (cf. [1, Lemma 2]). *Let S, X, Y, x , and y be as above. Suppose that $x \otimes y = x' \otimes y'$ in $X \otimes_S Y$. Then there exist $s_1, \dots, s_n, t_1, \dots, t_n \in S^1$,*

$x_1, \dots, x_n \in X$, and $y_2, \dots, y_n \in Y$ such that

$$\begin{aligned}
 & x = x_1s_1, & s_1y = t_1y_2, \\
 & x_1t_1 = x_2s_2, & s_2y_2 = t_2y_3, \\
 (1) \quad & \vdots & \vdots \\
 & x_{n-1}t_{n-1} = x_ns_n, & s_ny_n = t_ny', \\
 & x_nt_n = x'
 \end{aligned}$$

and

$$\begin{aligned}
 (2) \quad & s_1 \geq_{\mathcal{F}} t_1 \geq_{\mathcal{F}} \dots \geq_{\mathcal{F}} s_i \geq_{\mathcal{F}} t_i \leq_{\mathcal{F}} s_{i+1} \leq_{\mathcal{F}} t_{i+1} \leq_{\mathcal{F}} \dots \leq_{\mathcal{F}} s_n \leq_{\mathcal{F}} t_n \\
 & \text{(or } s_1 \geq_{\mathcal{F}} t_1 \geq_{\mathcal{F}} \dots \geq_{\mathcal{F}} s_i \leq_{\mathcal{F}} t_i \leq_{\mathcal{F}} s_{i+1} \leq_{\mathcal{F}} t_{i+1} \leq_{\mathcal{F}} \dots \leq_{\mathcal{F}} s_n \leq_{\mathcal{F}} t_n)
 \end{aligned}$$

where $\geq_{\mathcal{F}}$ is the quasi order of S^1 .

According to [1], a set of equations (1) is called a scheme of length n over X and Y joining (x, y) to (x', y') . If a scheme satisfies (2), then we say that it is V -formed.

Proof. By [1, Lemma 2], there exists a scheme (1) joining (x, y) to (x', y') . By appropriate substitution of s_i, t_i , we will show that (2) is satisfied. Let us assume in (1) that

$$s_i \in S \quad (1 < i \leq n), \quad t_i \in S \quad (1 \leq i < n).$$

For if $s_i = 1 \quad (1 < i \leq n)$, then $s_{i-1}y_{i-1} = t_{i-1}t_iy_{i+1}$, $x_{i-1}t_{i-1}t_i = x_{i+1}s_{i+1}$; hence, the scheme gets shorter; similarly, if $t_i = 1 \quad (1 \leq i < n)$.

Next, if t_i, s_{i+1} are incomparable with respect to $\geq_{\mathcal{F}}$, then one can insert new equations into the equations (1) as follows:

$$\begin{aligned}
 x_it_i &= x_{i+1}(s_{i+1}t_is_{i+1}), & (s_{i+1}t_is_{i+1})y_{i+1} &= (s_{i+1}t_is_{i+1})y_{i+1}, \\
 x_{i+1}(s_{i+1}t_is_{i+1}) &= x_{i+1}s_{i+1}, & s_{i+1}y_{i+1} &= t_{i+1}y_{i+2}.
 \end{aligned}$$

(If s_i, t_i are incomparable with respect to $\geq_{\mathcal{F}}$, then

$$s_iy_i = (t_is_it_i)y_{i+1}, \quad x_i(t_is_it_i) = x_i(t_is_it_i), \quad (t_is_it_i)y_{i+1} = t_iy_{i+1}.)$$

By repeating such insertions, we may assume any adjacent two elements of the sequence $s_1, t_1, \dots, s_n, t_n$ are \mathcal{F} -comparable. If scheme (1) is not V -formed, then several of the following four cases may occur. In each case, we will convert a part of the scheme into a V -formed scheme as follows.

Case 1. $s_i <_{\mathcal{F}} t_i \mathcal{F} \dots \mathcal{F} t_{j-1} \mathcal{F} s_j >_{\mathcal{F}} t_j$. Then, by assumption, all $s_i, t_i, \dots, t_{j-1}, s_j, t_j$ are in S . Set

$$\begin{aligned}
 t'_k &= t_k s_i t_k, & s'_{k+1} &= s_{k+1} s_i s_{k+1}, \\
 t''_k &= t_k t_j t_k, & s''_{k+1} &= s_{k+1} t_j s_{k+1}, \\
 t^*_k &= t_k s_i s_j t_k, & s^*_{k+1} &= s_{k+1} s_i s_j t_j s_{k+1} \quad (i \leq k \leq j-1).
 \end{aligned}$$

By Result 3(i), we have

$$t'_k = s'_i, \quad t''_k = s''_i, \quad t^*_k = s^*_i \quad (i \leq k < j, \quad i < l \leq j).$$

From (1) we get

$$\begin{aligned}
 & s_i y_i (= s'_j y_j = s'_j s_j y_j = s'_j s_j t_j s_j y_j) = s^*_j y_j, \\
 x_i s^*_j &= x_i t^*_i, & t^*_i y_j (= t^*_i y_{i+1} = t_i t_j s_i t_i y_{i+1} = t_i t_j t_i y_{i+1}) &= t^*_i y_{i+1}, \\
 x_i t^*_i &= x_j s''_j, & s''_j y_{i+1} (= s''_j y_j = s_j y_j) &= t_j y_{j+1},
 \end{aligned}$$

and $s_i \geq_{\mathcal{F}} s_i^* \mathcal{F} t_i^* \leq_{\mathcal{F}} t_i'' \leq_{\mathcal{F}} t_j$. This is a required scheme.

Case 2. $t_i <_{\mathcal{F}} s_{i+1} \mathcal{F} \cdots \mathcal{F} t_{j-1} \mathcal{F} s_j >_{\mathcal{F}} t_j$. Then by assumption, all $t_i, s_{i+1}, \dots, t_{j-1}, s_j, t_j$ are in S .

Set

$$\begin{aligned} s'_k &= s_k t_i s_k, & t'_k &= t_k t_i t_k, \\ s''_k &= s_k t_j s_k, & t''_k &= t_k t_j t_k, \\ s^*_k &= s_k t_i s_j t_j s_k, & t^*_k &= t_k t_i s_j t_j t_k \quad (i + 1 \leq k \leq j). \end{aligned}$$

By Result 3(i), we have

$$s'_k = t'_i, \quad s''_k = t''_i, \quad s^*_k = t^*_i \quad (i + 1 \leq k \leq j, \quad i + 1 \leq l < j).$$

From (1) we get

$$\begin{aligned} x_i t_i (= x_{i+1} s'_{i+1}) &= x_j s'_j, & s'_j y_{i+1} (= s'_j y_j = s'_j s_j y_j = s'_j s_j t_j s_j y_j) &= s^*_j y_j, \\ x_j s^*_j (= x_{i+1} s^*_{i+1} = x_{i+1} s_{i+1} t_i s_j t_j s_{i+1} = x_{i+1} s_{i+1} s_j t_j s_{i+1} = x_{i+1} s_{i+1} t_j s_{i+1}) &= x_j s''_j, \\ & & s''_j y_j &= t_j y_{j+1} \end{aligned}$$

and $t_i \geq_{\mathcal{F}} s'_j \geq_{\mathcal{F}} s^*_i \leq_{\mathcal{F}} s''_j \leq_{\mathcal{F}} t_j$. We are done.

Case 3. $s_i <_{\mathcal{F}} t_i \mathcal{F} \cdots \mathcal{F} t_{j-1} >_{\mathcal{F}} s_j$. By reversely ordering the equations (1), it is just Case 2.

Case 4. $t_i <_{\mathcal{F}} s_{i+1} \mathcal{F} \cdots \mathcal{F} t_{j-1} >_{\mathcal{F}} s_j$. In a way similar to the above, this is Case 1.

Notice that the subband of S^1 generated by all the s_i, t_i in (1) is finite (of course, it has finitely many \mathcal{F} -classes) and it contains all the elements $s'_i, t'_i, s''_i, t''_i, s^*_i, t^*_i$ occurring in the substitutions above. Thus by finitely repeating those substitutions of parts of the scheme by V -formed one, scheme (1) becomes V -formed. \square

Lemma 3. Let S, X, Y be as above and $x, x' \in XS$ and $y, y' \in SY$.

- (i) If $x \otimes y = x' \otimes y'$ in $X \otimes_S Y$, then $xa \otimes y = x'a \otimes y'$ in $X \otimes_S Y$ for all $a \in S$.
- (ii) If $xs \otimes y = x' \otimes y', x \otimes y = x't \otimes y'$ in $X \otimes_S Y$ for some $s, t \in S$, then $x \otimes y = x' \otimes y'$.

Proof. (i) By Lemma 2, there exist $x_1, \dots, x_n \in X, y_2, \dots, y_n \in Y, s_1, \dots, s_n$, and $t_1, \dots, t_n \in S^1$ such that

$$\begin{aligned} (3) \quad & \begin{aligned} x &= x_1 s_1, & s_1 y &= t_1 y_2, \\ x_1 t_1 &= x_2 s_2, & s_2 y_2 &= t_2 y_3, \\ & \vdots & & \vdots \\ x_{n-1} t_{n-1} &= x_n s_n, & s_n y_n &= t_n y', \\ & x_n t_n &= x'. \end{aligned} \end{aligned}$$

Here we may assume that all s_i, t_i belong to S . For, if $s_1 = 1$, then s_1, t_1 can be replaced by s, st_1 , respectively, where s is any element of S with $xs = x$. Also if $t_n = 1$, then t_n can be also replaced by some element of S . Further if $s_i = 1$ ($2 \leq i$), then as seen in the proof of Lemma 2 the scheme gets shorter; similarly, if $t_i = 1$ ($1 \leq i < n$).

Note next that $efy = efey$ ($efy' = efey'$) for all $e, f \in S$. For, by assumption, we can write $y = hy$ ($h \in S$) and by normality of S , $efy = ef(hy) = (efh)y = (efeh)y = efey$. (Similarly, $efy' = efey'$.)

Thus by using Lemma 1 and the note above, we get

$$xa = x_1(s_1a), \quad (s_1a)y = (s_1as_1)y = (t_2at_2)y_2$$

$$((s_nas_n)y_n = (t_nat_n)y') = (t_na)y', \quad x_n(t_na) = x'a).$$

So, by Lemma 1, we get a scheme joining (xa, y) to $(x'a, y')$ by replacing s_i, t_i by s_ias_i, t_iat_i respectively. Then (i) holds.

(ii) This is an immediate consequence of (i). \square

Remarks. 1. Lemma 3(i) is false without assumption that $x, x' \in XS$ and $y, y' \in SY$. For instance, let S be a left zero semigroup. Then $1 \otimes a = a \otimes a$ in $S^1 \otimes S$, but $b \otimes a \neq aba \otimes a$.

2. Given a scheme (3) of length n joining (x, y) to (x', y') (not necessarily, $x, x' \in XS$), it is shown, in the proofs of Lemmas 2 and 3, that it is possible to assume all the s_i, t_i except possibly s_1, t_n belong to S and that s_1 is in S if $x \in XS$ (t_n is in S if $x' \in XS$). Under these assumptions, if $x \in X - XS$ and $y \in Y - SY$, then $s_1 = t_1 = 1$ and $n = 1$; that is, $x = x'$ and $y = y'$. Otherwise, one can find $x'' \in XS$ and $y'' \in SY$ such that $x \otimes y = x'' \otimes y''$.

The proof of the “only if” part of the main theorem. We will appeal to Result 2. Let S be a normal band with (REP) and (REP)^{op}. Suppose

$$(4) \quad x \otimes (1 \otimes y) = x' \otimes (1 \otimes y') \quad \text{in } X \otimes (W \otimes Y)$$

where $x, x' \in X, y, y' \in Y, S^1 \subset W, X \in \text{Ens-}S, W \in S\text{-Ens-}S,$ and $Y \in S\text{-Ens}$. Then we shall show that

$$(5) \quad x \otimes y = x' \otimes y' \quad \text{in } X \otimes Y.$$

By the remarks after Lemma 3, we may assume that $x, x' \in XS$ and $y, y' \in SY$.

Here we may assume that W has the following property:

$$(6) \quad aws \in S, \quad a \mathcal{R} b, \quad \text{and } a >_{\mathcal{F}} s \quad (a, b, s \in S, w \in W)$$

implies

$$bws = bsbws \in S.$$

Proof of (6). Let ξ be the congruence on W generated by the relation $(bws, bsbws)$ and $\xi|_S$ the restriction to S of ξ . Then we shall show that $\xi|_S$ is an identity relation on S . For our purpose, it suffices to show that

$$(7) \quad ubwsv = u'bwsv' \quad (u, u', v, v' \in S) \quad \text{implies} \quad ubsbwsv = u'bsbwsv'.$$

If $a = b$, then, by assumption, $bws \in S$ and so, by normality of S , $bws = b^2ws^3 = b(sbw)s$. Hence (7) holds. Then we can assume that $a \neq b$. By Result 3(ii), $\mathcal{L}_a = \mathcal{R}_a$. If $u, u' \geq_{\mathcal{F}} b$, then $ub = b, u'b = b$ and, hence,

$$(ub)sbwsv = bsbwsv = bs(ub)wsv = bs(u'bwsv') = u'bsbwsv'$$

as required. If $u \not\geq_{\mathcal{F}} b$ (or $u' \not\geq_{\mathcal{F}} b$), then, by Result 3(i),

$$bub = b(bub)b = a(bub)a = aua$$

so that, by assumption,

$$ubwsv = (ubub)wsv = u(aua)wsv \in S.$$

Then, by normality of S ,

$$\begin{aligned} ubwsv &= ub(bub)wsv = ub(aua)w(ssv) = ub(aua)wsv \\ &= ub(aua)wsv = ubsbwsv. \end{aligned}$$

Then $u'bwsv' \in S$. Similarly, $u'bwsv' = u'bsbwsv'$. In any case, (7) holds. Therefore, $\xi|_S$ is an identity relation on S . So S can be naturally embedded in W/ξ . Hence $bws = bsbws = (asa)ws \in S$, which proves (6).

Hereafter, by Result 1, we may identify $y \in Y$ with $1 \otimes y \in W \otimes Y$.

By Lemma 2, we obtain a V -formed scheme of length n over X and $W \otimes Y$ joining $(x, 1 \otimes y)$ to $(x', 1 \otimes y')$ as follows:

$$\begin{aligned} (8) \quad & x = x_1 a_1, & a_1(1 \otimes y) &= b_1(w_2 \otimes y_2), \\ & x_1 b_1 = x_2 a_2, & a_2(w_2 \otimes y_2) &= b_2(w_3 \otimes y_3), \\ & & \vdots & \\ & x_{n-1} b_{n-1} = x_n a_n, & a_n(w_n \otimes y_n) &= b_n(1 \otimes y'), \\ & & x_n b_n &= x' \end{aligned}$$

where $x_i \in X$, $w_i \in W$, $y_i \in Y$, and $a_i, b_i \in S^1$.

We are going to prove (5) by induction on the length n of scheme (8).

By the remarks after Lemma 3, we may assume, in (8),

$$\text{all the } a_i, b_i \text{ belong to } S.$$

If $n = 1$, then, obviously, $x \otimes y = x' \otimes y'$. Assuming that (4) implies (5) when $n \leq m$, we proceed to the case where $n = m + 1$. First we may assume

$$(9) \quad a_1 \mathcal{R} b_1 \mathcal{R} a_2 \quad \text{and} \quad b_1 \neq a_2.$$

Proof of (9). If $b_1 >_{\mathcal{F}} a_1$, then we obtain the ascending chain

$$a_1 <_{\mathcal{F}} b_1 \leq_{\mathcal{F}} a_2 \leq_{\mathcal{F}} \cdots \leq_{\mathcal{F}} a_n \leq_{\mathcal{F}} b_n$$

since scheme (8) is V -formed. In this case, regarding scheme (8) as joining (x', y') to (x, y) , we can assume that $a_1 \geq_{\mathcal{F}} b_1$.

Next, if $a_1 >_{\mathcal{F}} b_1$, then $a_1(1 \otimes y) = b_1(w_2 \otimes y_2) = (a_1 b_1 a_1)(1 \otimes y)$. Consequently,

$$\begin{aligned} x &= x_1 a_1, & a_1(1 \otimes y) &= (a_1 b_1 a_1)(1 \otimes y), \\ x_1(a_1 b_1 a_1) &= x(a_1 b_1 a_1), \end{aligned}$$

so that

$$x \otimes y = x(a_1 b_1 a_1) \otimes y,$$

while

$$x(a_1 b_1 a_1) = x_1(a_1 b_1 a_1), \quad (a_1 b_1 a_1)(1 \otimes y) = b_2(w_2 \otimes y_2).$$

Therefore, we may assume that $a_1 \mathcal{F} b_1$.

If $b_1 >_{\mathcal{F}} a_2$, then

$$\begin{aligned} x &= x_1 a_1, & a_1(1 \otimes y) & (= b_1(w_2 \otimes y_2)) = b_1 a_1(1 \otimes y), \\ x_1(b_1 a_1) & (= (x_2 a_2)(b_1 a_1)) & &= (x_1 b_1 a_2)(b_1 a_1) \\ & & &= x_1(a_1 a_2 a_1) a_1 \quad [\text{by Result 3(i)}] = x(a_1 a_2 a_1). \end{aligned}$$

On the other hand,

$$\begin{aligned} x(a_1 a_2 a_1) &= x_1(a_1 a_2 a_1), & (a_1 a_2 a_1)(1 \otimes y) &= (b_1 a_2 b_1)(w_2 \otimes y_2), \\ x_1(b_1 a_2 b_1) &= x_2 a_2. \end{aligned}$$

Hence, we may assume that $a_2 \geq_{\mathcal{F}} b_1$.

If $a_2 >_{\mathcal{F}} b_1$, then, since scheme (8) is V -formed,

$$a_1 <_{\mathcal{F}} b_1 \leq_{\mathcal{F}} a_2 \leq_{\mathcal{F}} \cdots \leq_{\mathcal{F}} a_n \leq_{\mathcal{F}} b_n.$$

In this case, as shown above, we can reduce to the case that $b_n \mathcal{F} a_n \mathcal{F} b_{n-1}$. By renumbering reversely the equations (8), we may assume that $a_1 \mathcal{F} b_1 \mathcal{F} a_2$.

If $\mathcal{F}_{a_1} = \mathcal{L}_{a_1}$, then we can replace $w_2 \otimes y_2$ by $1 \otimes y$ in (8) and scheme (8) gets shorter. On the other hand, if $\mathcal{F}_{a_1} = \mathcal{R}_{a_1}$ and $b_1 = a_2$, then $x = x a_1 = x_1(b_1 a_1) = (x_2 a_2) a_1 = x_2 a_1$. So we can remove $x_1, w_2 \otimes y_2$ from scheme (8). Hence (9) may be assumed.

Case 1. There exists some $2 \leq i < n$ such that all a_k, b_k ($1 \leq k \leq i$) belong to \mathcal{F}_{a_1} but $a_1 >_{\mathcal{F}} a_{i+1}$. Since $\mathcal{F}_{a_1} = \mathcal{R}_{a_1}$, by multiplying the equations (8) on the left by a_1 from the right, we get

$$x = x a_1 = x_1 a_1 = x_2 a_1 = \cdots = x_i a_1 = x_{i+1} a_{i+1} a_1,$$

so that $x = x_i(a_1 a_{i+1} a_1)$, while, by Result 3(i), $a_k a_{i+1} a_k = b_k a_{i+1} b_k$ for all $1 \leq k \leq i$. So from (8) we obtain a scheme of length $\leq m$ joining $(x, 1 \otimes y)$ to $(x', 1 \otimes y')$ as follows:

$$\begin{aligned} x &= x_i(a_1 a_{i+1} a_1), & (a_1 a_{i+1} a_1)(1 \otimes y) &= (b_i a_{i+1} b_i)(w_{i+1} \otimes y_{i+1}), \\ x_i(b_i a_{i+1} b_i) &= x_{i+1} a_{i+1}, & a_{i+1}(w_{i+1} \otimes y_{i+1}) &= b_{i+1}(w_{i+2} \otimes y_{i+2}), \\ & \vdots & & \vdots \\ x_{n-1} b_{n-1} &= x_n a_n, & a_n(w_n \otimes y_n) &= b_n(1 \otimes y'), \\ x_n b_n &= x'. \end{aligned}$$

By the inductive assumption, $x \otimes y = x' \otimes y'$.

Case 2. There exists some $1 \leq i < n$ such that all a_k, b_k ($1 \leq k \leq i$) belong to \mathcal{F}_{a_1} but $a_1 \mathcal{F} a_{i+1} >_{\mathcal{F}} b_{i+1}$. By applying Lemma 1(ii) to (8) we have

$$\begin{aligned} (a_1 b_{i+1} a_1) y &= (b_1 b_{i+1} b_1)(w_2 \otimes y_2) = \cdots = (b_i b_{i+1} b_i)(w_{i+1} \otimes y_{i+1}) \\ &= (a_{i+1} b_{i+1} a_{i+1})(w_{i+1} \otimes y_{i+1}) = a_{i+1}(w_{i+1} \otimes y_{i+1}) = b_{i+1}(w_{i+2} \otimes y_{i+2}). \end{aligned}$$

Also, $x(a_1 b_{i+1} a_1) = x_{i+1}(a_1 b_{i+1} a_1)$. Then there exists a scheme of length $< n$ over X and $W \otimes Y$ joining $(x(a_1 b_{i+1} a_1), y)$ to (x', y') . Consequently, it follows from the inductive assumption that $x(a_1 b_{i+1} a_1) \otimes y = x' \otimes y'$. We have to prove that $x \otimes y = x \otimes (a_1 b_{i+1} a_1) y$. Since $a_{i+1}(w_{i+1} \otimes y_{i+1}) = a_1(a_1 b_{i+1} a_1 y_{i+1})$, $x_{i+1} a_1 = x$, this case can be reduced to the case for all a_j, b_j ($1 \leq j \leq n$). So we proceed to the next case.

Case 3. All a_i, b_i ($1 \leq i \leq n$) belong to \mathcal{R}_{a_1} . Then

$$(10) \quad x' = x = xs \quad \text{for all } s \in \mathcal{R}_{a_1}.$$

From Lemma 2, it follows that for each $1 \leq i \leq n$, there exists a V -formed scheme of length n_i over W and Y joining $(a_i w_i, y_i)$ to $(b_i w_{i+1}, y_{i+1})$ as follows:

$$(11) \quad \begin{array}{ll} a_i w_i = w_{i1} s_{i1}, & s_{i1} y_i = t_{i1} y_{i2}, \\ w_{i1} t_{i1} = w_{i2} s_{i2}, & s_{i2} y_{i2} = t_{i2} y_{i3}, \\ \vdots & \vdots \\ w_{i \ n_i-1} t_{i \ n_i-1} = w_{in_i} s_{in_i}, & s_{in_i} y_{in_i} = t_{in_i} y_{i+1}, \\ w_{in_i} t_{in_i} = b_i w_{i+1} \end{array}$$

where $w_i (w_1 = 1, w_{n+1} = 1)$, $w_{i1}, \dots, w_{in_i} \in X$, $y_i (y_1 = y, y_{n+1} = y')$, $y_{i2}, \dots, y_{in_i} \in Y$, s_{i1}, \dots, s_{in_i} , and $t_{i1}, \dots, t_{in_i} \in S^1$.

Set $s'_{ij} = s_{ij} a_1 s_{ij}$ and $t'_{ij} = t_{ij} a_1 t_{ij}$.

Subcase 3.1. There exist some of all the s'_{ij}, t'_{ij} , which are under a_1 with respect to $\geq_{\mathcal{F}}$. Then we shall show that there exist $u, v \in S$ such that $a_1 > u$, $a_1 > v$ and $x \otimes y = xu \otimes y$, $x' \otimes y = x'v \otimes y'$. Suppose first that all s'_{pq}, t'_{pq} ($2 \leq i, 1 \leq p \leq i-1, 1 \leq q \leq n_p$), s'_{iq}, t_{iq} ($1 \leq q \leq r-1$) belong to \mathcal{R}_{a_1} , but $a_1 >_{\mathcal{F}} s'_{ir}$.

Set $u = a_1 s'_{ir} a_1$. Since $eu = u$ for all $e \in S$ with $e \geq_{\mathcal{F}} a_1$, it follows from (11) that

$$a_p w_p u = b_p w_{p+1} u \quad (1 \leq p \leq i-1), \quad a_i w_i u = w_{ir-1} u.$$

By applying (6) to the equations just above, we obtain $u = w_{ir-1} u$, so that $u = w_{ir-1} s_{ir-1} u = w_{ir} s'_{r1} u$. By Result 3(iii), u is not the greatest lower bound of \mathcal{R}_{a_1} , since by Result 3(ii) and (9) $|\mathcal{R}_{a_1}| = 2$. So there exists $u' \in S$ such that u' is a lower bound of \mathcal{R}_{a_1} but $u' \neq uu'$. In the same way as above, $u' = u' a_1 = w_{ir} s'_{ir} u'$. Hence $u' = uu'$, which is a contradiction. Thus it must hold that all s'_{pq}, t'_{pq} ($2 \leq i, 1 \leq p \leq i-1, 1 \leq q \leq n_p$), s'_{iq}, t'_{iq} ($1 \leq q \leq r-1$), and s'_{ir} belong to \mathcal{R}_{a_1} , but $a_1 >_{\mathcal{F}} t'_{ir}$.

Then, by Result 3(i), $a_k t'_{ij} a_k = b_k t'_{ij} b_k$ ($1 \leq k \leq n$), say a^* . From equations (11) on the right, we have $a^* y = a^* y_{k+1}$ ($1 \leq k \leq i$), which, together with (10), yields $x \otimes y = x \otimes a^* y$. By the same way as above, we can find $b^* \in S$ satisfying that $b_n > b^*$ and $x' \otimes y' = x' \otimes b^* y'$, as required.

Moreover, by multiplying the right side of (8) by a^*, b^* , respectively that $a^* y = a^* y'$ and $b^* y = b^* y'$. Hence, $x \otimes y = x' \otimes a^* y'$ and $x' \otimes y' = x \otimes b^* y$. By Lemma 3(ii), we conclude that $x \otimes y = x' \otimes y'$.

Subcase 3.2. All the s'_{ij}, t'_{ij} belong to \mathcal{R}_{a_1} . By applying Lemma 1 to (11), we obtain schemes joining (x, y_i) to (x, y_{i+1}) as follows:

$$(12) \quad \begin{array}{ll} x = x s'_{i1}, & s'_{i1} y_i = t'_{i1} y_{i2}, \\ x t'_{i1} = x s'_{i2}, & s'_{i2} y_{i2} = t'_{i2} y_{i3}, \\ \vdots & \vdots \\ x t'_{i \ n_i-1} = x s'_{in_i}, & s'_{in_i} y_{in_i} = t'_{in_i} y_{i+1}, \\ x t'_{in_i} = x. \end{array}$$

From (10) and (12), it follows that $x \otimes y = x' \otimes y'$. This completes the proof of the main theorem.

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