

## COMBINATORIAL DIMENSION OF FRACTIONAL CARTESIAN PRODUCTS

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**ABSTRACT.** The combinatorial dimension of a fractional Cartesian product is the optimal value of an associated linear programming problem.

Fractional Cartesian products were introduced in a context of harmonic analysis to fill gaps between Cartesian products of spectral sets [1]. These fractional products subsequently gave rise to the idea of combinatorial dimension, a measurement of interdependencies that plays a natural role in harmonic and stochastic analysis (e.g., [2, 4, 6]). In this note, which is completely elementary and self-contained, we compute the combinatorial dimension of a general fractional Cartesian product by solving an associated linear programming problem.

**Definition 1** [2, 5]. Let  $E_1, \dots, E_n$  be infinite sets, and let  $F \subset E_1 \times \dots \times E_n$ . For an integer  $s > 0$ , let

$$(1) \quad \Psi_F(s) = \max\{|F \cap (A_1 \times \dots \times A_n)| : A_i \subset E_i, |A_i| \leq s, i = 1, \dots, n\}.$$

Define the *combinatorial dimension* of  $F$  to be

$$(2) \quad \dim F = \limsup_{s \rightarrow \infty} (\ln \Psi_F(s) / \ln s).$$

Equivalently, for each real number  $a > 0$ , let

$$(3) \quad d_F(a) = \limsup_{s \rightarrow \infty} (\Psi_F(s) / s^a),$$

and then  $\dim F = \inf\{a : d_F(a) < \infty\}$ .

The parameter  $\dim F$ , for  $F \subset E_1 \times \dots \times E_n$ , can be viewed as a measurement of interdependence between the canonical projections from  $E_1 \times \dots \times E_n$  onto  $E_j$ ,  $j = 1, \dots, n$ , restricted to  $F$ . The idea of *combinatorial dimension* was motivated at the outset by problems in harmonic analysis [1]. Indeed, it turns out that  $\dim F$  can precisely be linked to harmonic-analytic parameters that are naturally associated with spectral sets [2, 3].

Clearly,  $\dim(E_1 \times \dots \times E_n) = n$ , and if  $F \subset E_1 \times \dots \times E_n$  is infinite then  $1 \leq \dim F \leq n$ . For arbitrary  $\alpha \in (1, n)$ , randomly produced examples of  $\alpha$ -dimensional sets can be found in [5]. General fractional Cartesian products, archetypal instances of which appeared in [1, 3], provide explicit examples

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of sets with fractional combinatorial dimension; we know of no other general scheme which produces explicit examples of sets with arbitrary fractional dimension.

For a positive integer  $m$ , let  $[m] = \{1, \dots, m\}$ . Let  $m$  be some fixed positive integer, and let  $X$  be some fixed infinite set. For a subset  $S \subset [m]$ , let  $X^S$  be the set of all functions from  $S$  into  $X$ . In particular, we can identify  $X^{[m]}$  and  $X^m$ . If  $S_1 \subset S_2 \subset [m]$ , then we let  $t \mapsto t|_{S_1}$  be the projection from  $X^{S_2}$  onto  $X^{S_1}$ ; that is, if  $t \in X^{S_2}$  then  $t|_{S_1} \in X^{S_1}$  is such that, for each  $i \in S_1$ ,  $(t|_{S_1})(i) = t(i)$ .

**Definition 2.** Let  $X$  be an infinite set. Let  $\mathcal{U} = \{S_1, \dots, S_n\}$  be a cover of  $[m]$ ; that is,  $[m] = S_1 \cup \dots \cup S_n$ . The *fractional Cartesian product* of  $X$  based on  $\mathcal{U}$  is

$$(4) \quad X^{(\mathcal{U})} = \{(t|_{S_1}, \dots, t|_{S_n}) : t \in X^m\}.$$

Now let  $\mathcal{U} = \{S_1, \dots, S_n\}$  be a cover of  $[m]$ , and consider the problem of computing the combinatorial dimension of  $X^{(\mathcal{U})}$ , where  $X^{(\mathcal{U})}$  is viewed as a subset of the  $n$ -fold Cartesian product  $X^{S_1} \times \dots \times X^{S_n}$ . In the case that  $\mathcal{U}$  consists of  $k$ -element subsets of  $[m]$  ( $k \geq 1$ ) so that  $|\{j : i \in S_j\}| = I \geq 1$ , for all  $i \in [m]$ , it is shown in [3, Corollary 2.6] that  $\dim X^{(\mathcal{U})} = m/k$ . In the general case, we show that the combinatorial dimension of an arbitrary fractional Cartesian product is the optimal value of the following linear programming problem:

$$(5) \quad \begin{aligned} &\text{Maximize } x_1 + \dots + x_m \\ &\text{subject to the constraints that each } x_i \geq 0 \text{ and that} \\ &\quad \sum_{i \in S_j} x_i \leq 1 \quad \text{for each } j \in [n]. \end{aligned}$$

Since  $\mathcal{U}$  is a cover of  $[m]$ , the feasible set is bounded, and thus the optimal value exists. Let  $\alpha = \alpha(\mathcal{U})$  be this optimal value.

**Theorem 3.**  $\dim X^{(\mathcal{U})} = \alpha(\mathcal{U})$ .

*Proof.* First we prove the inequality  $\alpha(\mathcal{U}) \leq \dim X^{(\mathcal{U})}$ . Let  $(x_1, \dots, x_m)$  be an optimal vector; that is,  $x_1, \dots, x_m$  satisfy the constraints of the linear programming problem, and  $\alpha = x_1 + \dots + x_m$ . Moreover, by letting  $(x_1, \dots, x_m)$  be an extreme point of the feasible set, we can assume that each of  $x_1, \dots, x_m$  is rational. Let  $s$  be an arbitrarily large integer such that  $s^{x_i}$  is an integer for each  $i \in [m]$ . Let  $D_1, \dots, D_m \subset X$  be such that  $|D_i| = s^{x_i}$  for each  $i \in [m]$ , and consider  $D = D_1 \times \dots \times D_m \subset X^m$ . For each  $j \in [n]$ , let  $A_j = \times \{D_i : i \in S_j\}$ . We have  $A_j \subset X^{S_j}$  and obtain from the constraints in (5)

$$(6) \quad |A_j| = \prod_{i \in S_j} |D_i| = \prod_{i \in S_j} s^{x_i} \leq s.$$

Also, it is clear that  $t \mapsto (t|_{S_1}, \dots, t|_{S_n})$  is a bijection from  $D$  onto  $X^{(\mathcal{U})} \cap (A_1 \times \dots \times A_n)$ , so that

$$(7) \quad |X^{(\mathcal{U})} \cap (A_1 \times \dots \times A_n)| = |D| = \prod_{i=1}^m s^{x_i} = s^\alpha.$$

Since  $s$  is an arbitrary integer, (6) and (7) imply that  $\alpha \leq \dim X^{(\mathbb{Z})}$ .

To prove  $\alpha \geq \dim X^{(\mathbb{Z})}$ , we consider the following dual linear programming problem:

$$(8) \quad \begin{aligned} & \text{Minimize } y_1 + \cdots + y_n \\ & \text{subject to the constraints that each } y_j \geq 0 \text{ and that} \\ & \sum_{i \in S_j} y_j \geq 1 \quad \text{for each } i \in [m]. \end{aligned}$$

By the fundamental duality theorem of linear programming, the value of the dual problem exists and equals  $\alpha = \alpha(\mathbb{Z})$ . Let  $(y_1, \dots, y_n)$  be an optimal vector. Let  $A_1, \dots, A_n$  be arbitrary finite subsets of  $X^{S_1}, \dots, X^{S_n}$ , respectively. We will show

$$(9) \quad |X^{(\mathbb{Z})} \cap (A_1 \times \cdots \times A_n)| \leq |A_1|^{y_1} \cdots |A_n|^{y_n}.$$

In the case  $|A_1| \leq s, \dots, |A_n| \leq s$ , (9) of course implies

$$(10) \quad \Psi_{X^{(\mathbb{Z})}}(s) \leq s^{\sum_j y_j} = s^\alpha,$$

establishing  $\alpha \geq \dim X^{(\mathbb{Z})}$ .

We use the notation: if  $S = \{i_1, \dots, i_k\}$  then  $\sum_{t_i \in X: i \in S} f(t_1, \dots, t_m)$  denotes the iterated sum

$$\sum_{t_{i_1} \in X} \cdots \sum_{t_{i_k} \in X} f(t_1, \dots, t_m).$$

In particular, if  $S = \emptyset$  then  $\sum_{t_i \in X: i \in S} f(t_1, \dots, t_m) = f(t_1, \dots, t_m)$ . Observe that

$$(11) \quad |X^{(\mathbb{Z})} \cap (A_1 \times \cdots \times A_n)| = \sum_{t_i \in X: i \in [m]} \prod_{j=1}^n \mathbf{1}_{A_j}(t_i: i \in S_j)$$

( $\mathbf{1}_A$  denotes the indicator function of  $A$ ). To obtain (9), we apply recursively in (11) the following multilinear Hölder inequality: Let  $k > 1$  be an integer and let  $p_1 \in (0, \infty], \dots, p_k \in (0, \infty]$  be such that  $(1/p_1) + \cdots + (1/p_k) \geq 1$ . Let  $f_1, \dots, f_k$  be real-valued functions defined on  $X$ . Then

$$(12) \quad \sum_{t \in X} |f_1(t) \cdots f_k(t)| \leq \|f_1\|_{p_1} \cdots \|f_k\|_{p_k}$$

( $\|\cdot\|_p$  denotes the usual  $l^p$ -norm). For  $q = 0, \dots, m$ , define

$$(13) \quad F(q) = \sum_{t_i \in X: i \in [m] \sim [q]} \prod_{j=1}^n \left( \sum_{t_i \in X: i \in [q] \cap S_j} \mathbf{1}_{A_j}(t_i: i \in S_j) \right)^{y_j}.$$

Observe that  $F(0) = |X^{(\mathbb{Z})} \cap (A_1 \times \cdots \times A_n)|$  and that  $F(m) = |A_1|^{y_1} \cdots |A_n|^{y_n}$ , so (9) is just  $F(0) \leq F(m)$ . We let  $q = 0, \dots, m-1$  and proceed to show  $F(q) \leq F(q+1)$ , thus establishing (9). First we rewrite  $F(q)$  as

$$(14) \quad \begin{aligned} F(q) = & \sum_{t_i \in X: i \in [m] \sim [q+1]} \prod_{\{j: q+1 \notin S_j\}} \left( \sum_{t_i \in X: i \in [q] \cap S_j} \mathbf{1}_{A_j}(t_i: i \in S_j) \right)^{y_j} \\ & \cdot \sum_{t_{q+1} \in X} \prod_{\{j: q+1 \in S_j\}} \left( \sum_{t_i \in X: i \in [q] \cap S_j} \mathbf{1}_{A_j}(t_i: i \in S_j) \right)^{y_j}. \end{aligned}$$

Given the constraints in the dual linear programming problem (8) ( $\sum_{q+1 \in S_j} y_j \geq 1$ ), we apply (12) in (14) to the summation over  $t_{q+1}$  with exponents  $1/y_j$  and functions  $(\sum_{t_j \in X: i \in [q] \cap S_j} \mathbf{1}_{A_j}(t_i: i \in S_j))^{y_j}$ , where  $q+1 \in S_j$ . We thus obtain

$$\begin{aligned}
 (15) \quad F(q) &\leq \sum_{t_i \in X: i \in [m] \sim [q+1]} \prod_{\{j: q+1 \notin S_j\}} \left( \sum_{t_i \in X: i \in [q] \cap S_j} \mathbf{1}_{A_j}(t_i: i \in S_j) \right)^{y_j} \\
 &\cdot \prod_{\{j: q+1 \in S_j\}} \left( \sum_{t_{q+1} \in X} \sum_{t_i \in X: i \in [q] \cap S_j} \mathbf{1}_{A_j}(t_i: i \in S_j) \right)^{y_j} \\
 &= \sum_{t_i \in X: i \in [m] \sim [q+1]} \prod_{j=1}^n \left( \sum_{t_i: i \in [q+1] \cap S_j} \mathbf{1}_{A_j}(t_i: i \in S_j) \right)^{y_j} \\
 &= F(q+1). \quad \square
 \end{aligned}$$

*Remarks.* 1. Since the optimal value of the linear programming problem in (5) is a rational number, we deduce that, for every positive integer  $m$  and every cover  $\mathcal{U}$  of  $[m]$ , the combinatorial dimension of  $X^{(\mathcal{U})}$  is rational.

2.  $F \subset E_1 \times \cdots \times E_n$  was said in [5] to be a  $\gamma$ -Cartesian product if

$$(16) \quad 0 < \liminf_{s \rightarrow \infty} (\Psi_F(s)/s^\gamma) \leq \limsup_{s \rightarrow \infty} (\Psi_F(s)/s^\gamma) < \infty.$$

Clearly, the combinatorial dimension of a  $\gamma$ -Cartesian product is  $\gamma$ . However, the statement in (16) is stronger than the statement  $\dim F = \gamma$ ; this can be illustrated by the random constructions in [5]. In the framework of the present paper, following an elementary computation based on (6), (7), and (9), we observe that, for  $\alpha = \alpha(\mathcal{U})$ ,

$$\lim_{s \rightarrow \infty} \Psi_{X^{(\mathcal{U})}}(s)/s^\alpha = 1.$$

3. Corollary 2.6 in [3] can be directly deduced from our main theorem. Suppose that a cover  $\mathcal{U} = \{S_1, \dots, S_n\}$  of  $[m]$  consists of  $k$ -element subsets of  $[m]$  ( $k \geq 1$ ) so that  $|\{j: i \in S_j\}| = I \geq 1$  for all  $i \in [m]$ . Then, by a standard argument, the optimal value of the linear programming problem (5) is  $\geq m/k$ , and that of its dual (8) is  $\leq n/I = m/k$ . Therefore, by Theorem 3 and the fundamental duality theorem,  $\dim X^{(\mathcal{U})} = m/k$ .

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