

FINITE FAMILIES WITH FEW SYMMETRIC DIFFERENCES

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ABSTRACT. We show that $2^{\lceil \log_2(m) \rceil}$ is the least number of symmetric differences that a family of m sets can produce. Furthermore we give two characterizations of the set-theoretic structure of the families for which that lower bound is actually attained.

1. INTRODUCTION

Throughout this paper \mathbf{F} and \mathbf{G} will always denote finite families of sets.

Definition 1.1. • Let $A\Delta B$ denote the *symmetric difference* between the sets A and B , defined as

$$A\Delta B = \{x \mid (x \in A \wedge x \notin B) \vee (x \notin A \wedge x \in B)\}.$$

- For \mathbf{F} and \mathbf{G} families of sets and A a set let

$$\Delta\mathbf{F} = \{A\Delta B \mid A, B \in \mathbf{F}\};$$

$$\bar{\Delta}\mathbf{F} = \text{the closure under } \Delta \text{ of } \mathbf{F};$$

$$\mathbf{F}\Delta\mathbf{G} = \{A\Delta B \mid A \in \mathbf{F} \wedge B \in \mathbf{G}\};$$

$$A\Delta\mathbf{F} = \{A\Delta B \mid B \in \mathbf{F}\};$$

$$A \cap \mathbf{F} = \{A \cap B \mid B \in \mathbf{F}\}.$$

Notice that $\emptyset \in \Delta\mathbf{F}$ for any \mathbf{F} , while $\emptyset \in \mathbf{F}\Delta\mathbf{G}$ if and only if $\mathbf{F} \cap \mathbf{G} \neq \emptyset$.

If $B \neq C$, then $A\Delta B \neq A\Delta C$: therefore the cardinality $|\Delta\mathbf{F}|$ of $\Delta\mathbf{F}$ is always greater than or equal to the cardinality $|\mathbf{F}|$ of \mathbf{F} . Hence m sets produce at least m symmetric differences (and at most $m(m-1)/2 + 1$: this upper bound is attained for every m , e.g. by m pairwise disjoint sets).

Our first result, which will be proved in section 2, sharpens the above lower bound on $|\Delta\mathbf{F}|$ by showing that if $|\mathbf{F}| = m$ then $|\Delta\mathbf{F}| \geq 2^{\lceil \log_2(m) \rceil}$, i.e. that if $|\mathbf{F}| > 2^n$ then $|\Delta\mathbf{F}| \geq 2^{n+1}$. Since a family of subsets of a set with $n+1$ elements can produce at most 2^{n+1} symmetric differences our lower bound on $|\Delta\mathbf{F}|$ is optimal. Our result in particular entails that if $|\Delta\mathbf{F}| = |\mathbf{F}|$, then $|\mathbf{F}|$ is a power of 2 and we will also prove that if $|\mathbf{F}| > 2^n$ and $|\Delta\mathbf{F}| = 2^{n+1}$, then there exists $\mathbf{F}' \supseteq \mathbf{F}$ with $|\mathbf{F}'| = 2^{n+1}$ and $|\Delta\mathbf{F}'| = |\mathbf{F}'|$ (so that $\Delta\mathbf{F}' = \Delta\mathbf{F}$).

This shows that to describe the set-theoretic structure of the families \mathbf{F} with as few as possible symmetric differences, i.e. such that the lower bound on $|\Delta\mathbf{F}|$

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is attained, it suffices to describe the set-theoretic structure of the families \mathbf{F} such that $|\Delta\mathbf{F}| = |\mathbf{F}|$. The main goal of sections 3 and 4 is to shed some light on this set-theoretic structure and this will be accomplished using two different approaches. In section 3 we will focus on the Venn diagram of the family, while in section 4 we will concentrate on the way the elements of the family can be distinguished by way of a tree construction.

2. THE LEAST NUMBER OF SYMMETRIC DIFFERENCES

The power set $\text{Pow}(X)$ of a set X together with the operation Δ is a group with \emptyset as the identity and the property that the inverse of any element is the element itself, namely a 2-group. Groups of this kind, also called Boolean groups, are well-known and easily seen to be abelian, finite whenever finitely generated, and, in that case, to have order a power of 2.

Proposition 2.1. *If \mathbf{F} is finite, then $\bar{\Delta}\mathbf{F}$ is a finite Boolean group with the operation Δ , and $|\bar{\Delta}\mathbf{F}| = 2^n$ for some n .*

Finite Boolean groups provide the appropriate framework for the study of the operation Δ on finite families.

The following proposition collects a couple of elementary properties of cosets of subgroups of Boolean groups that will be useful in the sequel.

Proposition 2.2. *Let G be a Boolean group and H be a subgroup of G .*

- 1) *Let $g, g' \in G$. g and g' are in the same coset of H if and only if $gg' \in H$.*
- 2) *If $g \in G$, then $H \cup (gH)$, i.e. the union of H with the coset of H containing g , is a subgroup of G .*

Proof. 1) g and g' are in the same coset of H if and only if for some $h \in H$ $g' = gh$ if and only if $gg' \in H$.

2) Every subset of G is closed under inverses. By 1) the product of two elements of gH belongs to H , as does the product of two elements of H , while the product of an element of H with an element of gH belongs to gH . Therefore $H \cup (gH)$ is closed under products. □

The above facts about Boolean groups allow us to establish the lower bound on $|\Delta\mathbf{F}|$ stated in the introduction.

Theorem 2.3. *For any n and any family \mathbf{F} with $|\mathbf{F}| > 2^n$ we have $|\Delta\mathbf{F}| \geq 2^{n+1}$. Hence for every m the least number of symmetric differences that m sets can produce is $2^{\lceil \log_2(m) \rceil}$.*

Proof. Let $\mathbf{G} = \bar{\Delta}\mathbf{F}$, which by Proposition 2.1 is a finite Boolean group with the operation Δ . Moreover $|\mathbf{G}| = 2^{n+1+k}$ for some $k \geq 0$. Let h be maximal such that there exists a subgroup $\mathbf{H} \leq \mathbf{G}$ of order 2^h such that $\mathbf{H} \cap \Delta\mathbf{F} = \{\emptyset\}$.

We claim that $h \leq k$. In fact if $h > k$, then \mathbf{H} has at most 2^n cosets and, since $|\mathbf{F}| > 2^n$, one of them contains at least two distinct elements of \mathbf{F} whose symmetric difference, by Proposition 2.2.1, would be a nonempty element of $\mathbf{H} \cap \Delta\mathbf{F}$.

A consequence of the claim is that \mathbf{H} has at least 2^{n+1} cosets. If one of them has empty intersection with $\Delta\mathbf{F}$, then the union of this coset with \mathbf{H} is a subgroup (by Proposition 2.2.2) of order 2^{h+1} which intersects $\Delta\mathbf{F}$ only in \emptyset , contradicting the maximality of h . Hence every coset of \mathbf{H} contains at least one element of $\Delta\mathbf{F}$ and $|\Delta\mathbf{F}| \geq 2^{n+1}$. □

We now study the families for which the lower bound on the cardinality of the family of the symmetric differences is attained, i.e. families \mathbf{F} such that for some n , $|\mathbf{F}| > 2^n$ and $|\Delta\mathbf{F}| = 2^{n+1}$.

Theorem 2.4. *Let \mathbf{F} be a nonempty finite family of sets. The following are equivalent:*

- i) $|\Delta\mathbf{F}|$ is the least number of symmetric differences $|\mathbf{F}|$ sets can produce;
- ii) $|\mathbf{F}| > |\Delta\mathbf{F}|/2$ and $\Delta\mathbf{F}$ is a group under Δ ;
- iii) there exists $\mathbf{F}' \supseteq \mathbf{F}$ such that $|\mathbf{F}'| > |\mathbf{F}'|/2$ and $|\mathbf{F}'| = |\Delta(\mathbf{F}')|$.

Proof. For $|\mathbf{F}| = 1$ the theorem holds since i), ii), and iii) are all true. So let n be such that $2^n < |\mathbf{F}| \leq 2^{n+1}$.

To prove that i) implies ii) notice that if $|\Delta\mathbf{F}|$ is the least number of symmetric differences $|\mathbf{F}|$ sets can produce then, since, as pointed out in the introduction, the lower bound in Theorem 2.3 is optimal, $|\Delta\mathbf{F}| = 2^{n+1}$ and hence $|\mathbf{F}| > |\Delta\mathbf{F}|/2$. Therefore we need only to show that $\Delta\mathbf{F}$ is closed under symmetric differences.

Let $\mathbf{G} = \bar{\Delta}\mathbf{F}$, so that by Proposition 2.1 $|\mathbf{G}| = 2^{n+1+k}$ for some $k \geq 0$. If $k = 0$, then $\mathbf{G} = \Delta\mathbf{F}$, and we are done; so we assume that $k \geq 1$.

We first show by induction on i , $1 \leq i \leq k$, that for every $X \in \mathbf{G} \setminus (\Delta\mathbf{F})$ there exists a subgroup \mathbf{H} of \mathbf{G} of order 2^i such that $X \in \mathbf{H}$ and $\mathbf{H} \cap \Delta\mathbf{F} = \{\emptyset\}$. For $i = 1$ simply take $\mathbf{H} = \{\emptyset, X\}$. Assume the property holds for $i < k$ and let \mathbf{H} be of order 2^i , with $X \in \mathbf{H}$ and $\mathbf{H} \cap \Delta\mathbf{F} = \{\emptyset\}$. \mathbf{H} has $2^{n+1+k-i} > 2^{n+1}$ cosets, and since $|\Delta\mathbf{F}| = 2^{n+1}$ there are cosets of \mathbf{H} which do not contain any element of $\Delta\mathbf{F}$. If $Y\Delta\mathbf{H}$ is one of them, then necessarily $Y\Delta\mathbf{H} \neq \mathbf{H}$, and then $\mathbf{H} \cup (Y\Delta\mathbf{H})$ is (by Proposition 2.2.2) a subgroup of \mathbf{G} of order 2^{i+1} which contains X and no element of $\Delta\mathbf{F}$ but \emptyset .

Suppose now that A, B are such that $A\Delta B \notin \Delta\mathbf{F}$. Let \mathbf{H} be a subgroup of \mathbf{G} of order 2^k such that $A\Delta B \in \mathbf{H}$ and $\mathbf{H} \cap \Delta\mathbf{F} = \{\emptyset\}$. By Proposition 2.2.1 A and B belong to the same coset of \mathbf{H} , and by the proof of Theorem 2.3 if $|\Delta\mathbf{F}| = 2^{n+1}$, then at most one of A, B is in $\Delta\mathbf{F}$. Hence if $A, B \in \Delta\mathbf{F}$, then $A\Delta B \in \Delta\mathbf{F}$ and $\Delta\mathbf{F}$ is closed under Δ .

To prove that ii) implies iii), assume ii) holds, fix any $A \in \mathbf{F}$, and let $\mathbf{F}' = A\Delta(\Delta\mathbf{F})$. $\mathbf{F}' \supseteq \mathbf{F}$ holds because $A' = A\Delta(A\Delta A')$ for every $A' \in \mathbf{F}$, and hence $\Delta\mathbf{F} \subseteq \Delta(\mathbf{F}')$. Since $\Delta\mathbf{F}$ is a group \mathbf{F}' is a coset of $\Delta\mathbf{F}$ within $\bar{\Delta}\mathbf{F}$, so that $|\mathbf{F}'| = |\Delta\mathbf{F}|$, and $\Delta(\mathbf{F}') \subseteq \Delta\mathbf{F}$. Therefore $\Delta(\mathbf{F}') = \Delta\mathbf{F}$ and hence $|\Delta(\mathbf{F}')| = |\mathbf{F}'|$.

iii) implies i) is immediate because iii) implies that $|\mathbf{F}'| = |\Delta(\mathbf{F}')| = 2^{n+1}$ and $\Delta\mathbf{F} \subseteq \Delta(\mathbf{F}')$. Hence $|\Delta\mathbf{F}| = 2^{n+1}$ since by Theorem 2.3, $|\Delta\mathbf{F}| \geq 2^{n+1}$. \square

Remark 2.5. The proof that ii) implies iii) of Theorem 2.4 shows that if \mathbf{F} produces as few as possible symmetric differences, then either $|\bar{\Delta}\mathbf{F}| = 2^{n+1}$ or $|\bar{\Delta}\mathbf{F}| = 2^{n+2}$. Nevertheless $|\bar{\Delta}\mathbf{F}|$ being small is not equivalent to $|\Delta\mathbf{F}|$ being least: $\mathbf{F} = \{\{1\}, \{2\}, \{3\}, \{4\}, \{1, 2, 3, 4\}\}$ is a family of $2^2 + 1$ sets such that $|\bar{\Delta}\mathbf{F}| = 2^4$ but $|\Delta\mathbf{F}| = 2^3 + 3$.

Remark 2.6. In ii) of Theorem 2.4 both conditions are necessary: any family of 4 sets producing 7 symmetric differences shows that $|\mathbf{F}| > |\Delta\mathbf{F}|/2$ alone does not suffice to ensure that $|\Delta\mathbf{F}|$ is least, while $\mathbf{F} = \{\emptyset, \{1\}, \{2\}, \{3\}, \{4\}, \{1, 2, 3, 4\}\}$ is a family of $2^2 + 2$ sets such that $\Delta\mathbf{F}$ is a group of order 2^4 under Δ , and hence shows that the closure of $\Delta\mathbf{F}$ under Δ does not suffice to ensure that $|\Delta\mathbf{F}|$ is least.

3. VENN DIAGRAMS

Theorem 2.4 shows that if a family has few symmetric differences then it is contained in some \mathbf{F} satisfying $|\Delta\mathbf{F}| = |\mathbf{F}|$. In this and in the next section we will study families satisfying $|\Delta\mathbf{F}| = |\mathbf{F}|$.

A special case is of course offered by the families \mathbf{F} such that $\mathbf{F} = \Delta\mathbf{F}$, i.e. such that (\mathbf{F}, Δ) is a Boolean group; a condition which is clearly equivalent to $|\mathbf{F}| = |\Delta\mathbf{F}|$ and $\emptyset \in \mathbf{F}$. A typical example of a family \mathbf{F} satisfying the equality $\mathbf{F} = \Delta\mathbf{F}$ is the power set $\text{Pow}(X)$ of any given set X . On the other hand every finite Boolean group G is easily seen to be isomorphic to a power set with the operation of symmetric difference. Indeed if X is a minimal set of generators of G , every product of elements of X equals the product of the elements of X which actually occur in it an odd number of times, so that to every element of G corresponds a unique subset of X . Furthermore the product of two elements in G is exactly the product of the elements in the symmetric difference of their corresponding subsets of X .

Therefore from the algebraic point of view there are no solutions to the equation $\mathbf{F} = \Delta\mathbf{F}$ but the power sets. Moreover, since $|\mathbf{F}| = |\Delta\mathbf{F}|$ implies $\Delta\mathbf{F} = \Delta(\Delta\mathbf{F})$, from the same algebraic point of view the families \mathbf{F} such that $|\mathbf{F}| = |\Delta\mathbf{F}|$ are again just the power sets. But this is far from being true as far as the set-theoretic structure of \mathbf{F} is concerned.

In this section we explore the Venn diagrams that families satisfying $|\mathbf{F}| = |\Delta\mathbf{F}|$ can have. We will give first a characterization of the families satisfying $\mathbf{F} = \Delta\mathbf{F}$, and then show how to characterize the remaining solutions to the equation $|\mathbf{F}| = |\Delta\mathbf{F}|$ by making use of the operation, to be defined below, of forming a *Venn variant* of a family of sets.

Definition 3.1. • Let \mathcal{G}_n be the family of sets of cardinality n of the form

$$\{\{s_1\} \cup A_1, \dots, \{s_n\} \cup A_n\}$$

where the s_i are all distinct and $\{s_1, \dots, s_n\} \cap (A_1 \cup \dots \cup A_n) = \emptyset$.

• Let $\mathcal{D}_n = \{\bar{\Delta}\mathbf{G} \mid \mathbf{G} \in \mathcal{G}_n\}$.

We will use the following set-theoretic notion, which has been introduced in [1].

Definition 3.2. S is a *differentiating set* for the family \mathbf{F} if for every $A, B \in \mathbf{F}$ with $A \neq B$ we have $A \cap S \neq B \cap S$, i.e. $(A \Delta B) \cap S \neq \emptyset$. S is a *minimal differentiating set* for \mathbf{F} if it is a differentiating set and for every $s \in S$, $S \setminus \{s\}$ is not a differentiating set for \mathbf{F} .

In [1] it is shown that if $|\mathbf{F}| = m$, then \mathbf{F} has a differentiating set S of cardinality $m - 1$.

Theorem 3.3. 1) If $\mathbf{F} \in \mathcal{D}_n$, then $|\mathbf{F}| = 2^n$ and $\mathbf{F} = \Delta\mathbf{F}$;
 2) If $\mathbf{F} = \Delta\mathbf{F}$, then for every minimal differentiating set S for \mathbf{F} the map $A \mapsto A \cap S$ is a bijection between \mathbf{F} and $\text{Pow}(S)$;
 3) $\mathbf{F} = \Delta\mathbf{F}$ if and only if for some n , $\mathbf{F} \in \mathcal{D}_n$.

Proof. 1) Let $\mathbf{G} = \{\{s_1\} \cup A_1, \dots, \{s_n\} \cup A_n\} \in \mathcal{G}_n$ and $\mathbf{F} = \bar{\Delta}\mathbf{G} \in \mathcal{D}_n$. Then $\emptyset \in \mathbf{F}$ and obviously $\mathbf{F} = \Delta\mathbf{F}$. Furthermore the fact that the s_i are all distinct and the condition $\{s_1, \dots, s_n\} \cap (A_1 \cup \dots \cup A_n) = \emptyset$ ensure that no element of \mathbf{G} can be generated from the others by means of Δ . Hence they are a minimal set of generators of \mathbf{F} , which therefore has cardinality 2^n .

- 2) Assume that $\mathbf{F} = \Delta\mathbf{F}$ so that \mathbf{F} is closed under Δ . In general for \mathbf{F} closed under Δ , from the identity $A\Delta(A\Delta(B\Delta C)) = B\Delta C$, it follows that for every $A \in \mathbf{F}$ we have $\Delta\mathbf{F} = A\Delta\mathbf{F}$. Let S be a minimal differentiating set for \mathbf{F} . For $X \in \Delta\mathbf{F}$ let $\pi(X) = X \cap S$. Since S is a differentiating set for \mathbf{F} , π is one-to-one. In fact if $(A\Delta B) \cap S = (A\Delta C) \cap S$ then $(B\Delta C) \cap S = \emptyset$; therefore if for $A, B, C \in \mathbf{F}$ with $B \neq C$ we had $\pi(A\Delta B) = \pi(A\Delta C)$, S would fail to be a differentiating set for \mathbf{F} . Furthermore, from the identity $(X \cap S)\Delta(Y \cap S) = (X\Delta Y) \cap S$ and the closure of $\Delta\mathbf{F}$ under Δ , it follows that the range of π is closed under Δ as well. Finally, by the minimality of S , for every $s \in S$ there exist $X \in \Delta\mathbf{F}$ such that $\pi(X) = X \cap S = \{s\}$. Thus every subset of S , being the symmetric difference of the singletons of its elements, is in the range of π . π is therefore a bijection between $\Delta\mathbf{F}$ and $\text{Pow}(S)$.
- 3) The “if” part follows immediately from 1). For the “only if” part let $S = \{s_1, \dots, s_n\}$ be a minimal differentiating set for \mathbf{F} as in the proof of 2). If $X_1 = \{s_1\} \cup A_1, \dots, X_n = \{s_n\} \cup A_n$ are the elements of $\Delta\mathbf{F}$ that intersect S in singletons, we have

$$\Delta\mathbf{F} = \bar{\Delta}\{\{s_1\} \cup A_1, \dots, \{s_n\} \cup A_n\}$$

and $\{s_1, \dots, s_n\} \cap (A_1 \cup \dots \cup A_n) = \emptyset$. Thus $\mathbf{F} \in \mathcal{D}_n$. □

Remark 3.4. \mathcal{D}_n does not exhaust all the families \mathbf{F} with 2^n elements such that $|\mathbf{F}| = |\Delta\mathbf{F}|$; for example the family $\mathbf{F} = \{\{1\}, \{2\}\}$ is such that $\Delta\mathbf{F} = \{\emptyset, \{1, 2\}\}$, so that $|\mathbf{F}| = |\Delta\mathbf{F}|$, although $\mathbf{F} \neq \Delta\mathbf{F}$.

We now prove that if $|\Delta\mathbf{F}|$ is minimal then the size of any minimal differentiating set for \mathbf{F} is minimal.

Proposition 3.5. *If $|\mathbf{F}| > 2^n$, $|\Delta\mathbf{F}| = 2^{n+1}$ and S is a minimal differentiating set for \mathbf{F} , then $|S| = n + 1$.*

Proof. If S is a minimal differentiating set for \mathbf{F} , then $|S| \geq n + 1$. For every $s \in S$, $\{s\} \in S \cap \Delta\mathbf{F}$. By Theorem 2.4 $\Delta\mathbf{F}$ is closed under Δ , and therefore $S \cap \Delta\mathbf{F}$ is closed under Δ . It follows that $\text{Pow}(S) = S \cap \Delta\mathbf{F}$ and thus $|S| \leq n + 1$. □

Remark 3.6. The converse of Proposition 3.5 is false even in the special case $|\mathbf{F}| = 2^{n+1}$: \mathbf{F} can be such that $|\mathbf{F}| = 2^{n+1}$ and have only minimal differentiating sets of the smallest possible size, i.e. $n + 1$, but fail to satisfy $|\Delta\mathbf{F}| = 2^{n+1}$. An example is provided by $\mathbf{F} = \{\emptyset, \{1\}, \{2\}, \{1, 2, 3\}\}$ which is a family of 2^2 sets generating $2^2 + 3$ symmetric differences, although its only minimal differentiating set is $\{1, 2\}$.

Starting with families in \mathcal{D}_n , new families satisfying the equality $|\mathbf{F}| = |\Delta\mathbf{F}|$ can be obtained by the operation of making what we call Venn variants.

Definition 3.7. • Let $V(\mathbf{F})$ denote the *Venn diagram* of \mathbf{F} , namely the partition induced on $\bigcup \mathbf{F}$ by the equivalence relation $\sim_{\mathbf{F}}$ defined by

$$x \sim_{\mathbf{F}} y \text{ if and only if } \forall A \in \mathbf{F} (x \in A \iff y \in A).$$

- For $v \in V(\mathbf{F})$ we let

$$\mathbf{F}_v = \{A \setminus v \mid A \in \mathbf{F} \wedge v \subseteq A\} \cup \{A \cup v \mid A \in \mathbf{F} \wedge A \cap v = \emptyset\}$$

and call \mathbf{F}_v the *Venn variant* of \mathbf{F} determined by v .

The following proposition is immediate from the definition.

Proposition 3.8. *For all $v \in V(\mathbf{F})$, $|\mathbf{F}_v| = |\mathbf{F}|$ and $\Delta(\mathbf{F}_v) = \Delta\mathbf{F}$, i.e. a Venn variant of \mathbf{F} has the same cardinality and produces the same collection of symmetric differences as \mathbf{F} .*

Definition 3.9. \mathcal{V}_n is the least family which contains \mathcal{D}_n and is closed under formation of Venn variants, namely

- if $\mathbf{F} \in \mathcal{D}_n$, then $\mathbf{F} \in \mathcal{V}_n$;
- if $\mathbf{F} \in \mathcal{V}_n$ and $v \in V(\mathbf{F})$, then $\mathbf{F}_v \in \mathcal{V}_n$.

By Theorem 3.3 and iterated application of Proposition 3.8 we obtain:

Proposition 3.10. *If $\mathbf{F} \in \mathcal{V}_n$, then $|\mathbf{F}| = |\Delta\mathbf{F}| = 2^n$.*

The analysis we will now carry out will show that not even \mathcal{V}_n exhausts the families of cardinality 2^n which have only 2^n symmetric differences.

Definition 3.11. For $v, v' \in V(\mathbf{F})$ we say that v is *opposite* to v' in \mathbf{F} if

$$\{A \in \mathbf{F} \mid v \subseteq A\} = \{A \in \mathbf{F} \mid A \cap v' = \emptyset\}$$

or, equivalently, if

$$\{A \in \mathbf{F} \mid v \subseteq A\} = \mathbf{F} \setminus \{A \in \mathbf{F} \mid v' \subseteq A\}.$$

Proposition 3.12. *Let \mathbf{F} be a family and $v, v', v'' \in V(\mathbf{F})$:*

- 1) v is not opposite to v in \mathbf{F} ;
- 2) if v is opposite to v' in \mathbf{F} , then v' is opposite to v in \mathbf{F} ;
- 3) if v' and v'' are opposite to v in \mathbf{F} , then $v' = v''$;
- 4) if no element of $V(\mathbf{F})$ is opposite to v in \mathbf{F} , then $V(\mathbf{F}_v) = V(\mathbf{F})$, so that for every $u \in V(\mathbf{F})$ \mathbf{F}_{vu} is defined; furthermore $\mathbf{F}_{vv} = \mathbf{F}$;
- 5) if v and v' are opposite in \mathbf{F} , then

$$V(\mathbf{F}_v) = (V(\mathbf{F}) \setminus \{v, v'\}) \cup \{v \cup v'\},$$

so that for every $u \in V(\mathbf{F}) \setminus \{v, v'\}$ \mathbf{F}_{vu} is defined; furthermore $\mathbf{F}_v = \mathbf{F}_{v'}$;

- 6) if $v \in V(\mathbf{F})$, then either v (if v has no opposite in \mathbf{F}) or $v \cup v'$ (if v and v' are opposite in \mathbf{F}) has no opposite in \mathbf{F}_v ;
- 7) for any $v \in V(\mathbf{F})$, if v' and v'' are not opposite in \mathbf{F} then v' and v'' are not opposite in \mathbf{F}_v .

Proof. 1) and 2) are immediate.

- 3) Let $A \in \mathbf{F}$. $v' \subseteq A$ is equivalent (since v and v' are opposite) to $v \cap A = \emptyset$ which is equivalent (since v and v'' are opposite) to $v'' \subseteq A$. Hence

$$\{A \in \mathbf{F} \mid v' \subseteq A\} = \{A \in \mathbf{F} \mid v'' \subseteq A\},$$

which entails $v' = v''$.

- 4) It suffices to show that $\sim_{\mathbf{F}}$ and $\sim_{\mathbf{F}_v}$ are the same equivalence relation. Let $x, y \in \bigcup \mathbf{F}$. If $x \sim_{\mathbf{F}} y$, then either $x, y \in v$ or $x, y \notin v$, from which it follows immediately that $x \sim_{\mathbf{F}_v} y$. Conversely let us assume that $x \sim_{\mathbf{F}_v} y$ so that at most one of x, y is in v . If $x, y \notin v$ let $A \in \mathbf{F}$ be such that $x \in A$ and $y \notin A$; then one of $A \setminus v$ and $A \cup v$ is in \mathbf{F}_v and witnesses that $x \not\sim_{\mathbf{F}_v} y$. Now suppose $x \in v$, and hence $y \notin v$. Since v has no opposite element in \mathbf{F} , there is $B \in \mathbf{F}$ such that either $x, y \in B$ or $x, y \notin B$. In the former case $B \setminus v$ contains y but does not contain x ; in the latter case $B \cup v$ contains x but does not contain y . In both cases we have that $x \not\sim_{\mathbf{F}_v} y$.

$\mathbf{F}_{vv} = \mathbf{F}$ follows immediately from the definitions.

- 5) The proof of 4) shows that if $u \in V(\mathbf{F})$, $u \neq v$ and $u \neq v'$, then u is a $\sim_{\mathbf{F}_v}$ equivalence class, namely an element of $V(\mathbf{F}_v)$. On the other hand if $x \in v$ and $y \in v'$, then $x \sim_{\mathbf{F}_v} y$ so that $v \cup v'$ replaces both v and v' in $V(\mathbf{F}_v)$.
- 6) Let \bar{v} be either v (if v has no opposite in \mathbf{F}) or $v \cup v'$ (if v and v' are opposite in \mathbf{F}). Suppose $u \in V(\mathbf{F}_v)$ is opposite to \bar{v} in \mathbf{F}_v . By 1) $u \neq \bar{v}$ and hence by 4) or 5) $u \in V(\mathbf{F})$. Since $u \neq v$ let $A \in \mathbf{F}$ be such that either $u \subseteq A$ and $v \not\subseteq A$ or $u \not\subseteq A$ and $v \subseteq A$. In the former case $A \cup v \in \mathbf{F}_v$ and contains both u and \bar{v} ; in the latter case $A \setminus v \in \mathbf{F}_v$ and contains neither u nor \bar{v} . In both cases u and \bar{v} are not opposite in \mathbf{F}_v .
- 7) If v is either v' or v'' , then the conclusion follows immediately from 6). Otherwise $v', v'' \in V(\mathbf{F})$ and, since v' and v'' are not opposite in \mathbf{F} , there exists $A \in \mathbf{F}$ such that either $v' \subseteq A$ and $v'' \subseteq A$ or $v' \cap A = \emptyset$ and $v'' \cap A = \emptyset$. Then either $A \cup v$, if $v \cap A = \emptyset$, or $A \setminus v$, if $v \subseteq A$, witnesses that v' and v'' are not opposite in \mathbf{F}_v . \square

As an immediate consequence we have the following:

Corollary 3.13. *If $V(\mathbf{F})$ has no pair of opposite elements, then for every $v \in V(\mathbf{F})$, $V(\mathbf{F}_v)$ has no pair of opposite elements.*

The following proposition relates the absence of opposite elements in the Venn diagram of a family with the Venn diagram of the family of its symmetric differences.

Proposition 3.14. *$V(\mathbf{F})$ has no pair of opposite elements if and only if $V(\mathbf{F}) = V(\Delta\mathbf{F})$.*

Proof. Let $x \sim_{\mathbf{F}}^- y$ stand for $\forall A \in \mathbf{F}(x \in A \iff y \notin A)$. As it is easy to check $x \sim_{\Delta\mathbf{F}} y$ if and only if either $x \sim_{\mathbf{F}} y$ or $x \sim_{\mathbf{F}}^- y$, so that $V(\Delta\mathbf{F})$ is the partition induced on $\bigcup \mathbf{F}$ by the equivalence relation $x \sim_{\mathbf{F}} y \vee x \sim_{\mathbf{F}}^- y$. If $V(\mathbf{F})$ has no pair of opposite elements, $\sim_{\mathbf{F}}^-$ is the empty relation and $\sim_{\Delta\mathbf{F}}$ coincides with $\sim_{\mathbf{F}}$ so that $V(\mathbf{F}) = V(\Delta\mathbf{F})$. Conversely if $V(\mathbf{F}) = V(\Delta\mathbf{F})$, then $\sim_{\mathbf{F}} = \sim_{\Delta\mathbf{F}}$ and $\sim_{\mathbf{F}}^-$ must be empty, which entails that in $V(\mathbf{F})$ there are no pairs of opposite elements. \square

Proposition 3.15. *If $\mathbf{F} \in \mathcal{V}_n$, then $V(\mathbf{F})$ has no pair of opposite elements.*

Proof. If $\mathbf{F} \in \mathcal{D}_n$, then $\mathbf{F} = \Delta\mathbf{F}$; thus $V(\mathbf{F}) = V(\Delta\mathbf{F})$ so that by the previous proposition $V(\mathbf{F})$ has no pair of opposite elements. Since the families in \mathcal{V}_n are obtained from those in \mathcal{D}_n by iterating the operation of Venn variant, by Corollary 3.13 their Venn diagrams have no pairs of opposite elements. \square

Remark 3.16. If $\mathbf{F} = \{\{1\}, \{2\}, \{1, 3\}, \{2, 3\}\}$, then $\Delta\mathbf{F} = \{\emptyset, \{1, 2\}, \{3\}, \{1, 2, 3\}\}$ so that $|\mathbf{F}| = |\Delta\mathbf{F}|$. However $\{1\}$ and $\{2\}$ are elements of $V(\mathbf{F})$ which are opposite in \mathbf{F} , so that $\mathbf{F} \notin \mathcal{V}_2$ by Proposition 3.15.

The following proposition shows that the families in \mathcal{V}_n are precisely those which have the least possible number of symmetric differences and, at the same time, have no pair of opposite elements in their Venn diagram.

Proposition 3.17.

$$\mathcal{V}_n = \{ \mathbf{F} \mid |\mathbf{F}| = |\Delta\mathbf{F}| = 2^n \wedge V(\mathbf{F}) \text{ has no pair of opposite elements} \}.$$

Proof. Propositions 3.10 and 3.15 show that \mathcal{V}_n is included in the set on the right hand side of the equality.

To prove the reverse inclusion assume that $|\mathbf{F}| = |\Delta\mathbf{F}| = 2^n$ and $V(\mathbf{F})$ has no pair of opposite elements. Let S be a minimal differentiating set for \mathbf{F} . Since S is a differentiating set for \mathbf{F} , the map $\lambda : \mathbf{F} \rightarrow \text{Pow}(S)$ defined by $\lambda(A) = A \cap S$ is one-to-one. Since $|\mathbf{F}| = |\Delta\mathbf{F}|$, $\Delta\mathbf{F} = \Delta(\Delta\mathbf{F})$, and, by Theorem 3.3.2, $|\Delta\mathbf{F}| = |\text{Pow}(S)|$. Therefore λ is actually a bijection. Thus there exists a unique $A_0 \in \mathbf{F}$ such that $A_0 \cap S = \emptyset$. Let $v_1, \dots, v_k \in V(\mathbf{F})$ be such that $A_0 = v_1 \cup \dots \cup v_k$. Since $V(\mathbf{F})$ has no pair of opposite elements $\mathbf{G} = \mathbf{F}_{v_1 \dots v_k}$ is well defined. Clearly $\emptyset \in \mathbf{G}$ and by Proposition 3.8 $|\mathbf{G}| = |\Delta\mathbf{G}|$. As noticed earlier these conditions entail $\mathbf{G} = \Delta\mathbf{G}$. Thus by Theorem 3.3 we have that $\mathbf{G} \in \mathcal{D}_n$. Finally, since by Proposition 3.12.4

$$\begin{aligned} \mathbf{F} &= \mathbf{F}_{v_1 \dots v_k v_k v_{k-1} \dots v_1} \\ &= \mathbf{G}_{v_k \dots v_1}, \end{aligned}$$

we have that $\mathbf{F} \in \mathcal{V}_n$. □

By Proposition 3.12.6 and 7, eliminating pairs of opposite elements is simply a matter of iterating the operation of making Venn variants.

Definition 3.18. If $(v_1, v'_1), \dots, (v_k, v'_k)$ are all the pairs of opposite elements in \mathbf{F} , we let

$$\mathbf{F}^* = \mathbf{F}_{v_1 \dots v_k}.$$

Note that \mathbf{F}^* depends neither on the order in which v_1, \dots, v_k are taken (since, in general, $\mathbf{F}_{vu} = \mathbf{F}_{uv}$ as long as both Venn variants are legal), nor on which element of the pair (v_i, v'_i) is used to make the Venn variant (by Proposition 3.12.5). As an immediate consequence of Proposition 3.12.6 and 7 we have the following:

Proposition 3.19. $V(\mathbf{F}^*)$ has no pair of opposite elements.

We can now state our characterization of the finite families satisfying $|\mathbf{F}| = |\Delta\mathbf{F}|$.

Theorem 3.20. $|\mathbf{F}| = |\Delta\mathbf{F}| = 2^n$ if and only if $\mathbf{F}^* \in \mathcal{V}_n$.

Proof. Since by Proposition 3.8 $|\mathbf{F}| = |\mathbf{F}^*|$ and $\Delta\mathbf{F} = \Delta(\mathbf{F}^*)$, the if part follows from Proposition 3.10, while the only if part follows from Propositions 3.19 and 3.17. □

4. TREES DESCRIBING A FAMILY

In this section we provide a different characterization of the families \mathbf{F} satisfying $|\mathbf{F}| = |\Delta\mathbf{F}|$. This characterization is based on the analysis of how elements of $(\bigcup \mathbf{F}) \setminus (\bigcap \mathbf{F})$ discriminate between the sets in \mathbf{F} .

Definition 4.1. Given a family \mathbf{F} and an element $x \in (\bigcup \mathbf{F}) \setminus (\bigcap \mathbf{F})$, we let

$$\mathbf{F}_x = \{ A \in \mathbf{F} \mid x \in A \} \quad \text{and} \quad \mathbf{F}_{\bar{x}} = \{ A \in \mathbf{F} \mid x \notin A \}.$$

We begin with some simple facts that will turn out to be useful.

Proposition 4.2. For any $x \in (\bigcup \mathbf{F}) \setminus (\bigcap \mathbf{F})$ we have

- 1) $\Delta\mathbf{F} = (\Delta\mathbf{F}_x) \cup (\Delta\mathbf{F}_{\bar{x}}) \cup (\mathbf{F}_x \Delta\mathbf{F}_{\bar{x}})$;
- 2) $[(\Delta\mathbf{F}_x) \cup (\Delta\mathbf{F}_{\bar{x}})] \cap (\mathbf{F}_x \Delta\mathbf{F}_{\bar{x}}) = \emptyset$;
- 3) $|\Delta\mathbf{F}| \geq 2 \max(|\mathbf{F}_x|, |\mathbf{F}_{\bar{x}}|)$.

Proof. 1) is immediate. 2) follows from the fact that for every $A \in (\Delta\mathbf{F}_x) \cup (\Delta\mathbf{F}_{\bar{x}})$ we have $x \notin A$ while for every $B \in \mathbf{F}_x \Delta\mathbf{F}_{\bar{x}}$ we have $x \in B$. 3) is a consequence of 1) and 2) together with $|\mathbf{F} \Delta \mathbf{G}| \geq \max(|\mathbf{F}|, |\mathbf{G}|)$. □

Proposition 4.3. *If $|\mathbf{F}| = |\Delta\mathbf{F}|$ and $x \in (\bigcup\mathbf{F}) \setminus (\bigcap\mathbf{F})$, then $|\Delta\mathbf{F}_x| = |\mathbf{F}_x| = |\Delta\mathbf{F}_{\bar{x}}| = |\mathbf{F}_{\bar{x}}| = |\mathbf{F}|/2$ and $\Delta\mathbf{F}_x = \Delta\mathbf{F}_{\bar{x}}$.*

Proof. $|\mathbf{F}_x| = |\mathbf{F}_{\bar{x}}| = |\mathbf{F}|/2$ follows immediately from the hypothesis and Proposition 4.2.3. Since $|\mathbf{F}_x \Delta\mathbf{F}_{\bar{x}}| \geq |\mathbf{F}|/2$ by Proposition 4.2.2 we must have $\Delta\mathbf{F}_x = \Delta\mathbf{F}_{\bar{x}}$ and $|\Delta\mathbf{F}_x| = |\Delta\mathbf{F}_{\bar{x}}| = |\mathbf{F}|/2$. \square

Definition 4.4. A tree *describing* \mathbf{F} is a binary tree satisfying the following properties:

1. the nodes are pairwise different subsets of \mathbf{F} ,
2. the root is \mathbf{F} ,
3. the leaves are singleton subsets of \mathbf{F} ,
4. the children of any internal node $\mathbf{G} \subseteq \mathbf{F}$ are \mathbf{G}_x and $\mathbf{G}_{\bar{x}}$ for some $x \in (\bigcup\mathbf{G}) \setminus (\bigcap\mathbf{G})$.

The *height* $h(\mathbf{F})$ of a family \mathbf{F} is the height of the highest tree describing \mathbf{F} .

The next theorem gives another characterization of the finite families satisfying $|\mathbf{F}| = |\Delta\mathbf{F}|$.

Theorem 4.5. $|\mathbf{F}| = |\Delta\mathbf{F}|$ if and only if $h(\mathbf{F}) = \log_2(|\mathbf{F}|)$.

Proof. By induction on $|\mathbf{F}|$.

If $|\mathbf{F}| = 1$ the only tree describing \mathbf{F} consists only of the root, and hence the theorem holds.

If $|\mathbf{F}| > 1$ consider any tree describing \mathbf{F} and let $x \in (\bigcup\mathbf{F}) \setminus (\bigcap\mathbf{F})$ be such that \mathbf{F}_x and $\mathbf{F}_{\bar{x}}$ are the children of the root. The subtrees lying above these nodes describe respectively \mathbf{F}_x and $\mathbf{F}_{\bar{x}}$; if $|\mathbf{F}| = |\Delta\mathbf{F}|$ by Proposition 4.3 and the inductive hypothesis we have that their heights are $\log_2(|\mathbf{F}|) - 1$ and hence that the original tree has height $\log_2(|\mathbf{F}|)$.

If $h(\mathbf{F}) = \log_2(|\mathbf{F}|)$, then both \mathbf{F}_x and $\mathbf{F}_{\bar{x}}$ have height less than or equal to $\log_2(|\mathbf{F}|) - 1$. Hence their size is bounded by $2^{\log_2(|\mathbf{F}|-1)} = |\mathbf{F}|/2$. From this we can conclude that $|\mathbf{F}_x| = |\mathbf{F}_{\bar{x}}| = |\mathbf{F}|/2$ and that $h(\mathbf{F}_x) = h(\mathbf{F}_{\bar{x}}) = \log_2(|\mathbf{F}|) - 1$. By inductive hypothesis we have that $|\mathbf{F}_x| = |\Delta\mathbf{F}_x| = |\mathbf{F}_{\bar{x}}| = |\Delta\mathbf{F}_{\bar{x}}| = |\mathbf{F}|/2$.

We begin by showing that $|\mathbf{F}| = |\Delta\mathbf{F}|$ holds if $|\mathbf{F}| = 4$. Let $\mathbf{F} = \{A, B, C, D\}$ and notice that $C\Delta D = A\Delta B$; otherwise $A\Delta B\Delta C\Delta D \neq \emptyset$ and there exists $x \in (\bigcup\mathbf{F}) \setminus (\bigcap\mathbf{F})$ belonging to either exactly one or exactly three elements of \mathbf{F} : in both cases we could construct a tree describing \mathbf{F} of height 3. Therefore we have also $B\Delta C = A\Delta D$ and $B\Delta D = A\Delta C$, so that $\Delta\mathbf{F} = \{A\Delta A, A\Delta B, A\Delta C, A\Delta D\}$ and hence $|\mathbf{F}| = |\Delta\mathbf{F}|$.

We now turn to the general case. Let α be a sequence indexing a node in a tree describing \mathbf{F} , and let β be a sequence obtained from α replacing zero or more characters x by \bar{x} , and zero or more characters \bar{x} by x . For example: $\alpha = x\bar{y}z\bar{w}$ and $\beta = \bar{x}y\bar{z}w$.

Let also denote by $\bar{\gamma}$ the sequence obtained from γ replacing each x by \bar{x} and each \bar{x} by x .

The following equality will be proved by induction on $|\mathbf{F}|$:

$$\mathbf{F}_\alpha \Delta \mathbf{F}_\beta = \mathbf{F}_{\bar{\alpha}} \Delta \mathbf{F}_{\bar{\beta}}.$$

The base case is immediate.

The inductive step is proved by a further induction on $\log_2(|\mathbf{F}|) - |\alpha|$.

The base case corresponds to the case in which $|\alpha| = \log_2(|\mathbf{F}|)$ and is the most complex. In this case, let $\mathbf{F}_\alpha = \{A\}$, $\mathbf{F}_\beta = \{B\}$, $\mathbf{F}_{\bar{\alpha}} = \{A'\}$, and $\mathbf{F}_{\bar{\beta}} = \{B'\}$; we must prove that $A\Delta B = A'\Delta B'$.

If $\alpha = \bar{\beta}$ the result is obvious. Otherwise let x be an element occurring both in α and in β (the argument for the case where only elements of the form \bar{x} occur both in α and in β is analogous). We can assume without loss of generality that $\alpha = x\alpha_1$ and $\beta = x\beta_1$. By inductive hypothesis on the cardinality of the family,

$$\mathbf{F}_\alpha \Delta \mathbf{F}_\beta = \mathbf{F}_{x\alpha_1} \Delta \mathbf{F}_{x\beta_1} = \mathbf{F}_{x\bar{\alpha}_1} \Delta \mathbf{F}_{x\bar{\beta}_1}.$$

Moreover

$$\begin{aligned} \mathbf{F}_{\bar{\alpha}} = \mathbf{F}_{\bar{x}\bar{\alpha}_1} = \{A'\} &\implies \mathbf{F}_{\bar{\alpha}_1} = \{A', A''\} \text{ with } x \in A'', \\ \mathbf{F}_{\bar{\beta}} = \mathbf{F}_{\bar{x}\bar{\beta}_1} = \{B'\} &\implies \mathbf{F}_{\bar{\beta}_1} = \{B', B''\} \text{ with } x \in B''. \end{aligned}$$

From this it follows that $\mathbf{F}_{x\bar{\alpha}_1} = \{A''\}$ and $\mathbf{F}_{x\bar{\beta}_1} = \{B''\}$. Hence $\mathbf{F}_\alpha \Delta \mathbf{F}_\beta = \mathbf{F}_{x\bar{\alpha}_1} \Delta \mathbf{F}_{x\bar{\beta}_1}$ implies that $A\Delta B = A''\Delta B''$ and it suffices to prove $A''\Delta B'' = A'\Delta B'$.

To this end consider the family $\mathbf{G} = \{A', B', A'', B''\}$. We show that $h(\mathbf{G}) = \log_2(|\mathbf{G}|) = 2$: let \mathbf{H} be the first common ancestor of $\mathbf{F}_{\bar{\alpha}_1} = \{A', A''\}$ and $\mathbf{F}_{\bar{\beta}_1} = \{B', B''\}$, and let \mathbf{H}_y and $\mathbf{H}_{\bar{y}}$ be the children of \mathbf{H} . Since y discriminates between $\{A', A''\}$ and $\{B', B''\}$, either y appears in $\bar{\alpha}_1$ or it appears in $\bar{\beta}_1$. Assuming, without loss of generality, that the former is the case, we have that for some ξ , $\mathbf{F}_{\bar{\alpha}} = \mathbf{F}_{\xi\bar{x}y}$ and hence $\mathbf{F}_\xi = \{A', A'', B', B''\} = \mathbf{G}$. If we had a tree describing \mathbf{G} of height greater than two, we could graft such a tree in place of \mathbf{F}_ξ in a tree describing \mathbf{F} and produce a tree describing \mathbf{F} of height greater than $\log_2(|\mathbf{F}|)$, contradicting the hypothesis.

Hence $h(\mathbf{G}) = 2$ and by the case $|\mathbf{F}| = 4$ considered above, $A''\Delta B'' = A'\Delta B'$, which concludes the base case.

For the inductive step pick $x \in \bigcup(\mathbf{F}_\alpha \Delta \mathbf{F}_\beta) \setminus \bigcap(\mathbf{F}_\alpha \Delta \mathbf{F}_\beta)$ and notice that

$$\begin{aligned} \mathbf{F}_\alpha \Delta \mathbf{F}_\beta &= (\mathbf{F}_\alpha \Delta \mathbf{F}_\beta)_x \cup (\mathbf{F}_\alpha \Delta \mathbf{F}_\beta)_{\bar{x}} \\ &= [(\mathbf{F}_{\alpha x} \Delta \mathbf{F}_{\beta \bar{x}}) \cup (\mathbf{F}_{\alpha \bar{x}} \Delta \mathbf{F}_{\beta x})] \cup [(\mathbf{F}_{\alpha x} \Delta \mathbf{F}_{\beta x}) \cup (\mathbf{F}_{\alpha \bar{x}} \Delta \mathbf{F}_{\beta \bar{x}})] \\ &= [(\mathbf{F}_{\bar{\alpha} \bar{x}} \Delta \mathbf{F}_{\bar{\beta} x}) \cup (\mathbf{F}_{\bar{\alpha} x} \Delta \mathbf{F}_{\bar{\beta} \bar{x}})] \cup [(\mathbf{F}_{\bar{\alpha} \bar{x}} \Delta \mathbf{F}_{\bar{\beta} \bar{x}}) \cup (\mathbf{F}_{\bar{\alpha} x} \Delta \mathbf{F}_{\bar{\beta} x})] \\ &= (\mathbf{F}_{\bar{\alpha}} \Delta \mathbf{F}_{\bar{\beta}})_x \cup (\mathbf{F}_{\bar{\alpha}} \Delta \mathbf{F}_{\bar{\beta}})_{\bar{x}} \\ &= \mathbf{F}_{\bar{\alpha}} \Delta \mathbf{F}_{\bar{\beta}} \end{aligned}$$

where the third equality has been obtained by induction hypothesis.

If $\alpha = \beta = x$, the equality we just proved shows that $\Delta \mathbf{F}_x = \Delta \mathbf{F}_{\bar{x}}$. Since we already have $|\Delta \mathbf{F}_x| = |\mathbf{F}|/2$, Proposition 4.2.1 entails that our thesis $|\mathbf{F}| = |\Delta(\mathbf{F})|$ will follow from

$$|\mathbf{F}_x \Delta \mathbf{F}_{\bar{x}}| = |\mathbf{F}|/2.$$

Since $|\Delta \mathbf{F}_x| = |\mathbf{F}_x|$ and $|\Delta \mathbf{F}_{\bar{x}}| = |\mathbf{F}_{\bar{x}}|$ for any $A \in \mathbf{F}_x$ and $B \in \mathbf{F}_{\bar{x}}$, we have $\Delta \mathbf{F}_x = \{A\Delta A' \mid A' \in \mathbf{F}_x\}$ and $\Delta \mathbf{F}_{\bar{x}} = \{B\Delta B' \mid B' \in \mathbf{F}_{\bar{x}}\}$. Fix $B \in \mathbf{F}_{\bar{x}}$ and consider the function $\varphi : \mathbf{F}_x \rightarrow \mathbf{F}_x \Delta \mathbf{F}_{\bar{x}}$ defined by $\varphi(A') = A'\Delta B$.

Clearly φ is injective. φ is also surjective: for any $A'\Delta B' \in \mathbf{F}_x \Delta \mathbf{F}_{\bar{x}}$, since $B\Delta B' \in \Delta \mathbf{F}_{\bar{x}} = \Delta \mathbf{F}_x$ there exists $A'' \in \mathbf{F}_x$ such that $B\Delta B' = A'\Delta A''$. Hence

$$A'\Delta B' = A''\Delta B = \varphi(A'').$$

Hence φ is a bijection between \mathbf{F}_x and $\mathbf{F}_x \Delta \mathbf{F}_{\bar{x}}$: we have $|\mathbf{F}_x \Delta \mathbf{F}_{\bar{x}}| = |\mathbf{F}_x| = |\mathbf{F}|/2$ and the proof is complete. \square

Remark 4.6. The generalization of Theorem 4.5 stating that for any family \mathbf{F} , $|\Delta\mathbf{F}| = 2^{\lceil \log_2(|\mathbf{F}|) \rceil}$ if and only if $h(\mathbf{F}) = \lceil \log_2(|\mathbf{F}|) \rceil$ is false. The “only if” direction follows easily by Theorems 2.4 and 4.5, but the “if” direction does not hold. A counterexample is provided by $\mathbf{F} = \{\emptyset, \{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3, 4\}\}$, which is a family of $2^2 + 1$ sets generating $2^3 + 3$ symmetric differences, although $h(\mathbf{F}) = \lceil \log_2(|\mathbf{F}|) \rceil = 3$. Notice also that this family has only minimal differentiating sets of size 3 (see Remark 3.6).

The result in this section can be interpreted in the following, playful, way. A family of sets \mathbf{F} is given and players I and II have full knowledge of \mathbf{F} . Player I picks $X \in \mathbf{F}$. Player II has to discover X , by asking, one after another, questions of the form “does a belong to X ?”. Player I has to give correct yes/no answers. Player II asks only questions whose answer he/she cannot recover from his/her knowledge of \mathbf{F} and from the answers to the previous questions. When \mathbf{F} is finite, player II will always discover X after asking at most $|\mathbf{F}| - 1$ questions. In general, the number of questions II has to ask depends on X as well as on the sequence of the questions asked. For \mathbf{F} finite, Theorem 4.5 says that the number of questions player II has to ask depends neither on X nor on the sequence of the questions asked if and only if $|\mathbf{F}| = |\Delta\mathbf{F}|$. The results in Section 3 then give a way of constructing plays of that sort.

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