

A HOMOTOPY THEOREM FOR MATROIDS, I

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1. **Introduction.** By a *matroid* on a finite set M we understand a class \mathbf{M} of non-null subsets of M which satisfies the following axioms.

AXIOM I. *No member of \mathbf{M} contains another as a proper subset.*

AXIOM II. *If $(X, Y) \in \mathbf{M}$, $a \in X \cap Y$ and $b \in X - (X \cap Y)$, then there exists $Z \in \mathbf{M}$ such that $b \in Z \subseteq (X \cup Y) - \{a\}$.*

Such systems were introduced by Hassler Whitney [1].

As an example let \mathbf{L} be any class of subsets of M forming a group under mod 2 addition, and let \mathbf{M} be the class of all minimal non-null members of \mathbf{L} . Then it is easily verified that \mathbf{L} satisfies Axiom II and that each non-null member of \mathbf{L} is a sum of non-null members of \mathbf{M} . It follows that \mathbf{M} satisfies both axioms and is thus a matroid. Such a matroid we call *binary*.

In particular \mathbf{M} may be the set of edges of a finite graph G and \mathbf{L} may be the class of 1-cycles mod 2 of G . Then it is found that the members of \mathbf{M} are those sets of edges of G which define circuits. In this case we call \mathbf{M} the *circuit-matroid* of G .

Given a matroid \mathbf{M} let \mathcal{Q} be the class of all unions of members of \mathbf{M} . Then each element of \mathcal{Q} is a subset of M . We partition \mathcal{Q} into disjoint classes $\mathcal{Q}_{-1}, \mathcal{Q}_0, \mathcal{Q}_1, \mathcal{Q}_2, \dots$ according to the following rules.

(i) *The null subset \emptyset of M , considered as an empty union, is the only member of \mathcal{Q}_{-1} .*

(ii) *When \mathcal{Q}_r has been determined for $-1 \leq r \leq k$ we define \mathcal{Q}_{k+1} as the class of all minimal members of*

$$P_k = \mathcal{Q} - \bigcup_{r=-1}^k \mathcal{Q}_r.$$

That is \mathcal{Q}_{k+1} consists of all members of P_k which have no other members of P_k as a subset. The members of \mathcal{Q} are the *flats* of \mathbf{M} . Those belonging to \mathcal{Q}_d are the flats of *dimension d* , or *d -flats*.

At the end of §2 of this paper we interpret the dimensions of the flats of a circuit-matroid in terms of graph theory.

We shall see that the flats of a matroid \mathbf{M} on a set M have some properties resembling those of the elements of a projective geometry. Because of this analogy we refer to the 0-flats, 1-flats and 2-flats of \mathbf{M} as its *points*, *lines* and *planes* respectively. The points are simply the members of the class \mathbf{M} .

We have to recognize one distinction which has no analogue in projective

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geometry. A flat F is *disconnected* if it can be represented as the union of two disjoint non-null subsets F' and F'' of M such that each point X of M satisfying $X \subseteq F$ satisfies also either $X \subseteq F'$ or $X \subseteq F''$. If no such representation is possible then F is *connected*. Thus the points of M and its (-1) -flat are connected.

A *path* in M is a finite sequence $P = (X_1, \dots, X_k)$ of one or more points of M , not necessarily all distinct, such that any two consecutive terms are distinct points of M which are subsets of the same connected line. The first and last terms of P are its *origin* and *terminus* respectively. If they are the same point we call P *re-entrant*. If P has only one term we call it *degenerate*.

If $P = (X_1, \dots, X_k)$ and $P' = (X_k, \dots, X_m)$ are paths of M such that the origin of P' is the terminus of P then we define their *product* PP' as the path $(X_1, \dots, X_k, \dots, X_m)$. Multiplication of paths is clearly associative. It is therefore permissible to write a path $(PQ)R$ or $P(QR)$ simply as PQR .

Suppose we have two paths PR and PQR where Q is either (i) of the form (X, Y, X) or (ii) of the form (X, Y, Z, X) with X, Y and Z subsets of the same plane. Then we say that each of PR and PQR can be derived from the other by an *elementary deformation*. Two paths P_1 and P_2 are *homotopic* if they are identical or if one can be derived from the other by a finite sequence of elementary deformations. Homotopy is clearly an equivalence relation. A path homotopic to a degenerate path is said to be *null-homotopic*.

In this paper we show that every re-entrant path in a matroid is null-homotopic. Actually we prove a more general theorem, as the result just stated is not sufficient for the purposes of Paper II. We first agree to call a subclass C of M *convex* if it has the following property: if two distinct members X and Y of C are subsets of the same line L then every point of M which is a subset of L is a member of C . Given a convex subclass C of M we say that a path P is *off* C if no term of P is a point of C . We then enquire into the condition that a path P off C can be transformed into a degenerate path by a finite sequence of elementary deformations so that all the intermediate paths are off C . In this paper we show how the idea of an elementary deformation must be generalized so as to make this transformation possible for every re-entrant path P off C .

It is hoped that the technique here developed for the study of matroids will be found useful in graph theory when applied to the circuit-matroids of graphs.

2. Flats. Let M be a matroid on a set M . We refer to the elements of M as the *cells* of the matroid. If S and T are subsets of M we use the symbol $S \subset T$ to denote that S is a proper subset of T . We write $\langle S \rangle$ for the union of all the points of M which are subsets of S . If S is a flat of M we denote its dimension by dS .

(2.1) *If S is a flat of M and k is an integer satisfying $-1 \leq k < dS$ then there exists a flat T of M such that $dT = k$ and $T \subset S$.*

Proof. If the theorem fails let k be the greatest integer satisfying -1

$\leq k < dS$ such that no k -flat T of \mathbf{M} satisfies $T \subset S$. Clearly $k > -1$ and therefore $dS \geq 1$.

By the definition of k there exists a $(k+1)$ -flat T' of \mathbf{M} such that $T' \subseteq S$. By the definition of the classes \mathcal{Q}_r we have

$$T' \in \mathcal{Q} - \bigcup_{r=-1}^k \mathcal{Q}_r \subset \mathcal{Q} - \bigcup_{r=-1}^{k-1} \mathcal{Q}_r = P_{k-1}.$$

But T' is not a minimal member of P_{k-1} , since dT' is not k . Hence there exists a minimal member T of P_{k-1} such that $T \subset T'$. But then $dT = k$. Since $T' \subseteq S$ this contradicts the definition of k . The theorem follows.

(2.2) *If S and T are flats of \mathbf{M} such that $S \subset T$, then $dS < dT$.*

Proof. Since T is non-null we have $dT > -1$. If $dS \geq dT$ there is a flat U of \mathbf{M} such that $U \subseteq S$ and $dU = dT$, by (2.1). But then T is not a minimal member of \mathcal{Q}_{dT} , contrary to the definition of this class.

It follows from (2.2) that $\langle M \rangle$ has a greater dimension than any other flat of \mathbf{M} .

It is convenient to say that a flat S is *on* a flat T if either $S \subseteq T$ or $T \subseteq S$. If S and T are distinct we can distinguish between the two cases by comparing dimensions.

(2.3) *If S is a flat of \mathbf{M} and $a \in S$, then $d\langle S - \{a\} \rangle = dS - 1$.*

Proof. If possible choose S and a so that $d\langle S - \{a\} \rangle \neq dS - 1$ and so that dS has the least value consistent with this.

By (2.1) there is a flat T of \mathbf{M} such that $T \subseteq S$ and $dT = dS - 1$. Choose $b \in S - T$. Then $T \subseteq \langle S - \{b\} \rangle \subset S$. Hence $d\langle S - \{b\} \rangle = dS - 1$, by (2.2).

Suppose $a \notin \langle S - \{b\} \rangle$. Then $\langle S - \{b\} \rangle \subseteq \langle S - \{a\} \rangle \subset S$. Hence $d\langle S - \{a\} \rangle = dS - 1$, by (2.2). But this is contrary to the choice of S and a .

We deduce that $a \in \langle S - \{b\} \rangle$. Hence there exists $X \in \mathbf{M}$ such that $X \subseteq S$, $a \in X$ and $b \notin X$. Since $b \in S$ there exists $Y \in \mathbf{M}$ such that $Y \subseteq S$ and $b \in Y$. It follows by Axiom II that there exists $Z \in \mathbf{M}$ such that $Z \subseteq S$, $a \notin Z$ and $b \in Z$. ($Z = Y$ if $a \notin Y$). These results imply

$$(2.3a) \quad \langle \langle S - \{a\} \rangle - \{b\} \rangle \subset \langle S - \{a\} \rangle \subset S,$$

$$(2.3b) \quad \langle \langle S - \{b\} \rangle - \{a\} \rangle \subset \langle S - \{b\} \rangle \subset S.$$

We also have

$$(2.3c) \quad \langle \langle S - \{a\} \rangle - \{b\} \rangle = \langle \langle S - \{b\} \rangle - \{a\} \rangle,$$

since each side of this equation represents the union of those points of \mathbf{M} which include neither a nor b .

Since $d\langle S - \{b\} \rangle = dS - 1$ it follows from (2.3b) and the choice of S and a that $d\langle \langle S - \{b\} \rangle - \{a\} \rangle = dS - 2$. Hence $d\langle S - \{a\} \rangle = dS - 1$, by (2.2), (2.3a) and (2.3c). This contradiction establishes the theorem.

(2.4) *Let S and T be flats of \mathbf{M} such that $S \subseteq T$. Then there exists a flat U of \mathbf{M} such that $U \subseteq T$, $\langle U \cap S \rangle = \emptyset$ and $dU = dT - dS - 1$.*

Proof. Write $S_0=S$, $T_0=T$. If possible choose $a_0 \in S_0$ and write $S_1 = \langle S_0 - \{a_0\} \rangle$, $T_1 = \langle T_0 - \{a_0\} \rangle$. Observe that $S_1 \subseteq T_1$. If possible choose $a_1 \in S_1$ and write $S_2 = \langle S_1 - \{a_1\} \rangle$, $T_2 = \langle T_1 - \{a_1\} \rangle$. Then $S_2 \subseteq T_2$. Continue this process until it terminates. By (2.3) this will be with S_k and T_k , where $k = dS + 1$ and $S_k = \emptyset$. Applying (2.3) to the sequence of the T_i we find that $dT_k = dT - k = dT - dS - 1$. We note that $\langle T_k \cap S \rangle \subseteq \langle S - \{a_0, \dots, a_{k-1}\} \rangle = S_k = \emptyset$. Hence the theorem is satisfied with $U = T_k$.

(2.5) *If S and T are any flats of M then $d(S \cup T) + d\langle S \cap T \rangle \geq dS + dT$.*

Proof. Write $S_0=S$. If possible choose $a_0 \in S_0 - (S_0 \cap T)$ and write $S_1 = \langle S_0 - \{a_0\} \rangle$. If possible choose $a_1 \in S_1 - (S_1 \cap T)$ and write $S_2 = \langle S_1 - \{a_1\} \rangle$, and so on. By (2.3) the process terminates with $S_k = \langle S \cap T \rangle$, where $k = dS - d\langle S \cap T \rangle$. We now have $T = S_k \cup T \subseteq S_{k-1} \cup T \subseteq \dots \subseteq S_0 \cup T = S \cup T$. Hence $d(S \cup T) - dT \geq k = dS - d\langle S \cap T \rangle$, by (2.2).

Many "geometrical" results can be deduced from (2.2) and (2.5). For example any two distinct lines L_1 and L_2 on a plane P have a unique common point. To prove this we first use (2.2) to show that $d\langle L_1 \cap L_2 \rangle < dL_1 = 1$ and $L_1 \cup L_2 = P$. Then $d\langle L_1 \cap L_2 \rangle \geq 0$, by (2.5). Hence $d\langle L_1 \cap L_2 \rangle = 0$ and $\langle L_1 \cap L_2 \rangle$ is a single point of M . We can prove in the same way that if P_1 and P_2 are distinct planes on the same 3-flat E of M , then $\langle P_1 \cap P_2 \rangle$ is a line on E . Similarly if P is a plane and L a line on the same 3-flat E , and L is not on P , then $\langle P \cap L \rangle$ is a point on E .

Not all the axioms of projective geometry are valid for matroids. For example two points are not necessarily on a common line. In general matroids are like geometrical figures but not like complete geometries.

Suppose M is the circuit-matroid of a graph G . If $S \in \mathcal{Q}$ we write $G \cdot S$ for the subgraph of G made up of the edges of S and their incident vertices. We see that the flats S of M correspond to those subgraphs $G \cdot S$ in which each edge belongs to some circuit of the subgraph. In virtue of (2.3) $dS + 1$ is the least number of edges which must be removed from $G \cdot S$ in order to destroy all its circuits, that is $dS + 1$ is the rank or first Betti number of $G \cdot S$. The subgraph $G \cdot S$ is nonseparable if and only if the flat S is connected.

3. Connected flats. We begin this section with a study of the line.

(3.1) *Any line L of M is on at least two points. If X and Y are distinct points on L then $L = X \cup Y$. Moreover $X \cap Y$ is non-null if and only if L is connected.*

Proof. Choose $a \in L$. Then $\langle L - \{a\} \rangle$ is a point on L , by (2.3). Choose $b \in \langle L - \{a\} \rangle$. Then $\langle L - \{b\} \rangle$ is a point on L , by (2.3), which is distinct from $\langle L - \{a\} \rangle$.

Let X and Y be distinct points on L . Then $X \subseteq X \cup Y \subseteq L$. Hence $X \cup Y = L$, by (2.2). If $X \cap Y$ is non-null then L is clearly connected. If $X \cap Y$ is null then either L is disconnected or there exists $Z \in M$ such that $Z \subseteq X \cup Y$ and Z meets both X and Y . In the latter case $X \subseteq X \cup Z \subseteq X \cup Y = L$, by Axiom I. This is impossible, by (2.2).

(3.2) *A disconnected line is on just two points, and a connected line is on at least three points.*

Proof. By (3.1) any two distinct points on a disconnected line L are disjoint and have L as their union. Hence L has at most two points, and therefore just two by (3.1).

By (3.1) any connected line L has two distinct points X and Y , and we can find $a \in X \cap Y$. By (2.3) $\langle L - \{a\} \rangle$ is a point on L distinct from X and Y .

We shall need the following general theorems on connected flats.

(3.3) *Let S and T be connected flats of \mathbf{M} such that $S \subset T$. Then there exists a connected $(dS+1)$ -flat U of \mathbf{M} which is on both S and T .*

Proof. Since T is connected we can find $X \in \mathbf{M}$ such that $X \subseteq T$ and X meets both S and $T - S$. Choose such an X so that $S \cup X$ has the least possible number of cells. Clearly $S \cup X$ is a connected flat of \mathbf{M} . Its dimension exceeds dS , by (2.2).

Suppose $d(S \cup X) > dS + 1$. Choose $a \in (S \cup X) - S$. Then $d\langle (S \cup X) - \{a\} \rangle \geq dS + 1$, by (2.3). Hence there exists $Y \in \mathbf{M}$ such that $Y \subseteq (S \cup X) - \{a\}$ and Y meets $(S \cup X) - S$. But $Y \cap S$ is null, by the choice of X . Hence $Y \subset X$, which is impossible by Axiom I. We deduce that $d(S \cup X) = dS + 1$. Hence the theorem is true with $U = S \cup X$.

(3.4) *Let S be a connected d -flat on a connected $(d+2)$ -flat T of \mathbf{M} . Then there exist distinct connected $(d+1)$ -flats U and V of \mathbf{M} such that $S = \langle U \cap V \rangle$ and $T = U \cup V$.*

Proof. By (3.3) there is a connected $(d+1)$ -flat U which is on both S and T . Choose $a \in U - S$ and write $W = \langle T - \{a\} \rangle$. By (2.3) W is another $(d+1)$ -flat on S and T . By (2.4) there is a line L on T having no point in common with S . It meets U and W in points X and Z respectively, by (2.5). (See Figure I.) By (2.2) we have $S \cup X = U$ and $S \cup Z = W$. Hence Z is not on U and therefore $U \cup Z = T$, by (2.2).

Assume W is not connected. Then $S \cap Z = \emptyset$.

Suppose $U \cap Z = \emptyset$. By the connection of T there exists $Z' \in \mathbf{M}$ such that $Z' \subseteq T$ and Z' meets both U and Z . Then $U \subset U \cup Z' \subset U \cup Z = T$, by Axiom I. This is impossible by (2.2). We deduce that $U \cap Z \neq \emptyset$. A similar argument in which X, S and U replace Z, U and T respectively shows that $X \cap S \neq \emptyset$. Choose $b \in Z \cap U$ and $c \in X \cap S$.

Write $V = \langle T - \{b\} \rangle$. By (2.3) V is a $(d+1)$ -flat. It is on S since $b \in Z$ and $S \cap Z = \emptyset$. By (2.5) it has a common point Y with L . Clearly V is distinct from U and W and therefore Y is distinct from X and Z , since $V = S \cup Y$ by (2.2).

Now $c \in S \cap X \subseteq S \cap L = S \cap (Y \cup Z) = S \cap Y$, by (3.1). Hence V is connected.

If instead W is connected we write $V = W$.

We now have two connected $(d+1)$ -flats U and V of \mathbf{M} each of which is on both S and T . Hence $S \subseteq \langle U \cap V \rangle \subset U \subset U \cup V \subseteq T$, since U and V are

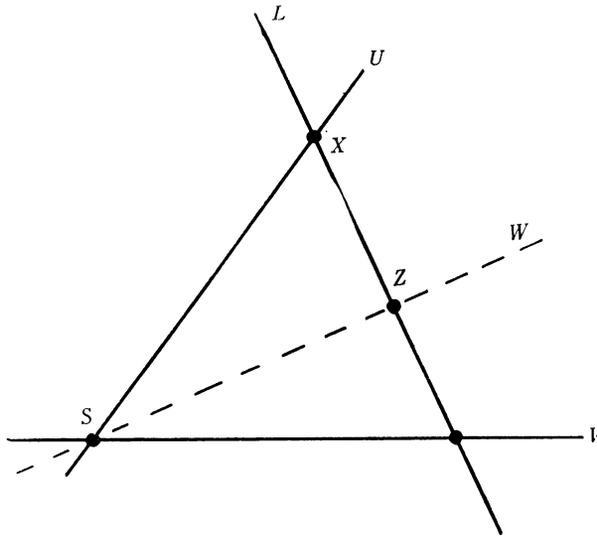


FIG. I

distinct and, by (2.2), neither is a subset of the other. In view of (2.2) this is possible only if $S = \langle U \cap V \rangle$ and $T = U \cup V$.

(3.5) *Let S, T and U be flats of M such that S and T are connected, $S \cup U \subseteq T$ and $\langle S \cap U \rangle = \emptyset$. Then there exists a connected flat R of M such that $S \subseteq R \subseteq T$, $\langle R \cap U \rangle = \emptyset$ and $dR = dT - dU - 1$.*

Proof. If possible choose S, T and U so that the theorem fails and dU has the least value consistent with this. Then $dU > -1$ since otherwise the theorem holds with $T = R$. Let W be a connected flat of M of greatest possible dimension such that $S \subseteq W \subseteq T$ and W does not contain U . Then $dW = dT - 1$ since otherwise, by (3.3) and (3.4), there exist distinct connected $(dW + 1)$ -flats K and L of M on T such that $\langle K \cap L \rangle = W$, and these cannot both contain U . By the choice of S, T and U there is a connected flat R of M such that $S \subseteq R \subseteq W \subseteq T$, $\langle R \cap U \rangle = \emptyset$ and $dR = dW - d\langle U \cap W \rangle - 1$. But then $dR \geq dT - dU - 1$, by (2.2), and therefore $dR = dT - dU - 1$, by (2.2) and (2.5). This contradiction establishes the theorem.

The foregoing results can be applied to circuit-matroids to obtain rather simple theorems about graphs. Thus from (3.5) with $S = \emptyset$ we find that if a nonseparable graph G has rank r and a subgraph $G \cdot U$ has a rank s then there is a nonseparable subgraph $G \cdot R$ of G of rank $r - s$ having no circuit in common with $G \cdot U$.

4. The disconnected line. By a *separation* $\{S_1, S_2\}$ of a disconnected flat S of M we mean a pair of complementary non-null subsets of S such that each $X \in M$ satisfying $X \subseteq S$ satisfies either $X \subseteq S_1$ or $X \subseteq S_2$.

(4.1) *If $\{S_1, S_2\}$ is a separation of a flat S of M and X_1 and X_2 are*

points of \mathbf{M} such that $X_1 \subseteq S_1$ and $X_2 \subseteq S_2$, then $X_1 \cup X_2$ is a disconnected line of \mathbf{M} .

Proof. Suppose Y is a point on $X_1 \cup X_2$ distinct from X_1 and X_2 . Since $Y \subseteq S$ we have $Y \subseteq S_1$ or $Y \subseteq S_2$. Hence $Y \subset X_1$ or $Y \subset X_2$, contrary to Axiom I. Thus the only subsets of $X_1 \cup X_2$ which are flats of \mathbf{M} are \emptyset , X_1 , X_2 and $X_1 \cup X_2$. Hence, by the definition of dimension, $X_1 \cup X_2$ is a line of \mathbf{M} having $\{X_1, X_2\}$ as a separation.

(4.2) *Let L be a disconnected line on a connected d -flat S of \mathbf{M} , where $dS > 1$. Then there exists a connected plane P of \mathbf{M} such that $L \subset P \subseteq S$.*

Proof. Let the two points on L be X and Y . Let P be a connected flat of \mathbf{M} of least possible dimension such that $L \subset P \subseteq S$. Assume $dP > 2$.

Suppose first that there is a disconnected line L' , distinct from L , on X and P . Let its point other than X be Z . By (3.5) there is a connected $(dP - 2)$ -flat U on Y and P having no point in common with L' . By (3.4) there are distinct connected $(dP - 1)$ -flats V and W of \mathbf{M} on P such that $\langle V \cap W \rangle = U$. By (2.2) and (2.5) V and W meet L' in distinct points. Since there are only two points on L' we may suppose X is on V . But then L is on V and the definition of P is contradicted. A similar argument applies if there is a disconnected line distinct from L on Y and P .

In the remaining case we choose $a \in P - L$ and write $R = \langle P - \{a\} \rangle$. Then $L \subseteq R$. Moreover $dR = dP - 1$, by (2.3). By the definition of P the flat R is disconnected. But there is no disconnected line on R , other than L , which is on either X or Y . Hence, by (4.1), the only possible separation of R is $\{X, Y\}$. Accordingly $R = L$ and $dP = 2$, contrary to assumption. From this contradiction we deduce that P is a plane.

(4.3) *Let L be a disconnected line on a connected plane P of \mathbf{M} . Let X and Y be the two points of L and let Z be any other point on P . Then $X \cup Z$ and $Y \cup Z$ are connected lines. Moreover they are the only lines of \mathbf{M} which are on both Z and P .*

Proof. Any line on Z and P has a common point with the line $X \cup Y$. Hence, by (3.1), the only flats on Z and P which can be lines are $X \cup Z$ and $Y \cup Z$. By (3.4) both these flats must be connected lines.

(4.4) *Let L be a disconnected line on a connected plane P of \mathbf{M} . Then every line on P other than L is connected.*

Proof. Let L' be any such line. By (3.1) it is on a point Z distinct from X and Y . Hence, by (4.3) it is one of the connected lines $X \cup Z$ and $Y \cup Z$.

5. **Convex subclasses.** Convex subclasses of a matroid \mathbf{M} were defined in the Introduction. As an example we may take the class of all points of \mathbf{M} on a given flat. The convexity of this class follows from (3.1).

Consider any path $P = (X_1, X_2, \dots, X_k)$ of \mathbf{M} . We say P is a path from X_1 to X_k . Any two consecutive terms of P have a non-null intersection, by (3.1). Hence the flat $X_1 \cup X_2 \cup \dots \cup X_k$ is connected. We denote this flat by $F(P)$. If S is any flat of \mathbf{M} such that $F(P) \subseteq S$ we say that P is a path on S .

(5.1) *Let \mathbf{C} be any convex subclass of \mathbf{M} . Let S be a non-null connected flat*

of \mathbf{M} and let X and Y be points on S such that $Y \notin \mathbf{C}$. Then there exists a path P from X to Y on S such that no term of P other than the first is a point of \mathbf{C} .

Proof. If possible choose S , X and Y so that the theorem fails and dS has the least value consistent with this. Clearly $dS > 1$. By (3.3) and (3.4) there is a connected $(dS - 2)$ -flat U and two distinct connected $(dS - 1)$ -flats V and W on S such that $X \subseteq U = \langle V \cap W \rangle$. Now Y is not on V or W , for otherwise there would be a path from X to Y on V or W of the kind required. By (3.5) there is a connected line L on S and Y such that $\langle L \cap U \rangle = \emptyset$. This meets V and W in distinct points $Z(V)$ and $Z(W)$ respectively, by (2.2) and (2.5). At least one of these, say $Z(V)$, belongs to $\mathbf{M} - \mathbf{C}$ since \mathbf{C} is convex. By the choice of S , X and Y there is a path Q from X to $Z(V)$ on V such that no term of Q other than the first is a point of \mathbf{C} . Adjoining Y to Q we obtain a path P from X to Y on S of the kind required. This contradiction establishes the theorem.

We now distinguish four kinds of re-entrant paths of \mathbf{M} as *elementary* with respect to a given convex subclass \mathbf{C} of \mathbf{M} . The first kind consists of all paths off \mathbf{C} of the form (X, Y, X) . The second consists of all paths off \mathbf{C} of the form (X, Y, Z, X) such that $d(X \cup Y \cup Z) \leq 2$.

Suppose P is a plane of \mathbf{M} on which there are two distinct points A and B of \mathbf{C} such that each connected line on P is on either A or B . Then any path off \mathbf{C} on P of the form (X, Y, Z, T, X) such that X, Y, Z and T are distinct, the lines $X \cup Y$ and $Z \cup T$ are on A , and the lines $Y \cup Z$ and $T \cup X$ are on B is an elementary re-entrant path of the third kind with respect to \mathbf{C} .

Suppose E is a 3-flat of \mathbf{M} on which there are three points A, B and C such that $A \cup B, B \cup C$ and $C \cup A$ are disconnected lines. Let there be just six connected planes on E , two on each of these disconnected lines. Suppose A, B and C are all in $\mathbf{M} - \mathbf{C}$ but there are two distinct members of \mathbf{C} on each of the six connected planes. Then any path off \mathbf{C} of the form (A, X, B, Y, A) , where X and Y are on distinct connected planes on $A \cup B$ and E , is an elementary re-entrant path of the fourth kind with respect to \mathbf{C} .

In studying the preceding case it is convenient to use the following notation. We write Z_1, Z_2 and Z_3 for A, B and C . We enumerate the six connected planes as P_1, \dots, P_6 in such a way that $\langle P_i \cap P_{i+3} \rangle = Z_j \cup Z_k$, where $1 \leq i \leq 3$ and (i, j, k) is a permutation of $(1, 2, 3)$. In general we write $\langle P_i \cap P_j \rangle = L_{ij}$ for $1 \leq i < j \leq 6$. If $j = i + 3$ then L_{ij} is the disconnected line $(Z_1 \cup Z_2 \cup Z_3) - Z_i$. If $j \neq i + 3$ let k be that integer 1, 2 or 3 which is not congruent to i or j mod 3. Then L_{ij} is on Z_k and it meets P_k and P_{k+3} in two distinct points. It is therefore connected, by (3.2). Clearly it is on no connected plane on E other than P_i and P_j . The 12 lines $L_{ij}, j \neq i + 3$, are the only connected lines on E , for by (3.4) any connected line on E is on two distinct connected planes on E .

We write $\langle P_i \cup P_j \cup P_k \rangle = X_{ijk}$ for $1 \leq i < j < k \leq 6$. Then X_{ijk} is a point of \mathbf{M} , being identical with $\langle L_{ij} \cap P_k \rangle$. If two of the suffices i, j and k are congruent mod 3 then X_{ijk} is one of the points Z_1, Z_2 and Z_3 . The remaining eight points X_{ijk} are all distinct, for on any one of them there can be only three

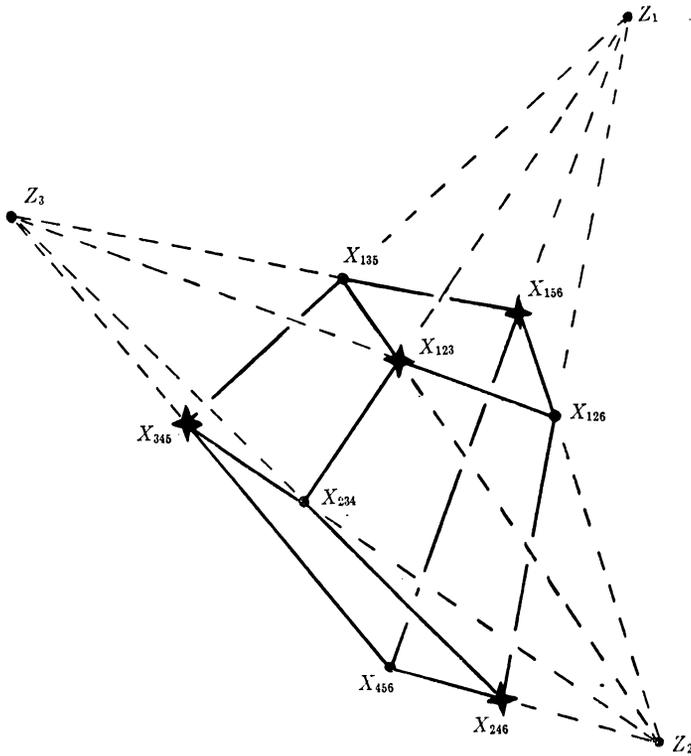


FIG. II

distinct planes such that each is on one of the lines $Z_1 \cup Z_2$, $Z_2 \cup Z_3$ and $Z_3 \cup Z_1$. These eight points, together with Z_1 , Z_2 and Z_3 , are the only points on E . For any point on E is on three distinct connected planes on E , by two applications of (3.4). (See Figure II.)

Consider the plane P_1 . The only points on it are Z_2 , Z_3 , X_{123} , X_{126} , X_{135} and X_{156} . We may adjust the notation so that $X_{123} \in C$. The other point of C on P_1 can have no common connected line with X_{123} and must therefore be X_{156} . We now find that $X_{246} \in C$ since this is the only point on P_2 having no common connected line with X_{123} . Proceeding in this way we find that $X_{ijk} \in C$ if and only if no two of the suffices are congruent mod 3 and the number of suffices less than 4 is odd. In Figures II, III and IV we represent points of C by four-pointed stars.

To construct a matroid having the structure just described we may use a method based on (2.3). We take M to be a set of six cells in 1-1 correspondence with the planes P_i . Any point X_{ijk} is represented by the set of those cells not corresponding to planes on X_{ijk} .

Suppose we have two paths PQR and PR off C , where Q is an elementary re-entrant path of the k th kind with respect to C . Then we call the process

of deriving one of the paths PQR and PR from the other an *elementary deformation* of the k th kind with respect to \mathbf{C} . We say that two given paths P' and P'' off \mathbf{C} are *homotopic* with respect to \mathbf{C} (written $P' \sim P'' (\mathbf{C})$) if they are identical or if one can be derived from the other by a finite sequence of elementary deformations with respect to \mathbf{C} . Homotopy with respect to \mathbf{C} is an equivalence relation. A path P homotopic to a degenerate path with respect to \mathbf{C} is said to be *null-homotopic* with respect to \mathbf{C} (written $P \sim 0 (\mathbf{C})$).

The null subset of \mathbf{M} is clearly convex. If \mathbf{C} is null we have only elementary deformations of the first and second kind to consider, and homotopy with respect to \mathbf{C} becomes identical with the homotopy defined in the Introduction.

If P is any path of \mathbf{M} we write P^{-1} for the path obtained by taking the terms of P in reverse order.

(5.2) *If P is any path off \mathbf{C} then $PP^{-1} \sim 0 (\mathbf{C})$.*

Proof. If possible choose P so that the theorem fails and P has the least number s of terms consistent with this. If $s > 1$ we can write $P = QR$, where Q and R have each fewer than s terms. Since RR^{-1} and QQ^{-1} can be converted into degenerate paths by elementary deformations we have $PP^{-1} = QRR^{-1}Q^{-1} \sim QQ^{-1} \sim 0 (\mathbf{C})$. If $s = 1$ then PP^{-1} is an elementary re-entrant path of the first kind, and so $PP^{-1} \sim 0 (\mathbf{C})$. The theorem follows.

(5.3) *If PUR and PVR are paths off \mathbf{C} such that $UV^{-1} \sim 0 (\mathbf{C})$, then $PUR \sim PVR (\mathbf{C})$.*

Proof. By (5.2) we have $V^{-1}V = V^{-1}(V^{-1})^{-1} \sim 0 (\mathbf{C})$. Hence $PUR \sim PUV^{-1}VR \sim PVR (\mathbf{C})$.

(5.4) *Let \mathbf{C} be any convex subclass of a matroid \mathbf{M} . Let S be a d -flat of \mathbf{M} on a $(d+1)$ -flat T of \mathbf{M} . Suppose all the points on S and at least one other point on T are members of \mathbf{C} . Then all the points on T are members of \mathbf{C} .*

Proof. Suppose the theorem false. Then we can find points $X \in \mathbf{C}$ and $Y \notin \mathbf{C}$, both on T but not on S . The flat $X \cup Y$ is connected since otherwise $S \subset S \cup X \subset T$, contrary to (2.2). By (5.1) there is a path from X to Y on $X \cup Y$ whose second term, X' say, is not a member of \mathbf{C} . By (2.5) the line $X \cup X'$ has a point X'' in common with S . But $X'' \notin \mathbf{C}$, by the definition of a convex subclass. This is contrary to hypothesis.

6. Proof of the main theorem.

(6.1) *Let \mathbf{C} be any convex subclass of a matroid \mathbf{M} and let P be any re-entrant path of \mathbf{M} off \mathbf{C} . Then $P \sim 0 (\mathbf{C})$.*

Proof. Assume the theorem false. Let P be any re-entrant path off \mathbf{C} which is not null-homotopic with respect to \mathbf{C} , and for which $dF(P)$ has the least value, n say, consistent with this condition. For an arbitrary path Q of \mathbf{M} we call $dF(Q)$ the *dimension* of Q .

By far the most difficult part of the proof is that covered by the following lemma.

LEMMA. *Suppose $n \geq 3$. Let $Q = (W, X, Y, Z, W)$ be a path off \mathbf{C} of dimension*

n such that $W \cup X \cup Y$ and $Y \cup Z \cup W$ are connected planes and $W \cup Y$ is a disconnected line. Then $Q \sim 0$ (\mathbf{C}).

Proof. Write $F_1 = W \cup X \cup Y$ and $F_2 = Y \cup Z \cup W$.

We note that if $Q' = (W, X', Y, Z', W)$ is a path off \mathbf{C} such that X' is on F_1 and Y' on F_2 , then

$$(6.1a) \quad Q' \sim Q(\mathbf{C}).$$

For

$$Q' \sim (W, X', Y)(Y, X, W)(W, X, Y)(Y, Z, W)(W, Z, Y)(Y, Z', W)(\mathbf{C})$$

by (5.2). But $(W, X', Y)(Y, X, W)$ and $(W, Z, Y)(Y, Z', W)$ are re-entrant paths off \mathbf{C} of dimension $< n$ and are therefore null-homotopic with respect to \mathbf{C} . Hence $Q' \sim (W, X, Y)(Y, Z, W) = Q(\mathbf{C})$.

A transversal of dimension $n-1$ is a connected $(n-1)$ -flat of \mathbf{M} which is on $F(Q)$ but not on both W and Y . By (2.2) and (2.5) such a transversal meets each of F_1 and F_2 in a line. These two lines are connected, by (4.4).

A transversal of dimension $n-2$ is a connected $(n-2)$ -flat of \mathbf{M} which is on $F(Q)$ but not on W or Y . By (2.2) and (2.5) the transversal has just one point in common with each of F_1 and F_2 . We call these two points the poles of the transversal.

Let B be any transversal of dimension $n-2$, with poles X' on F_1 and Z' on F_2 . Then B is on two distinct connected $(n-1)$ -flats of \mathbf{M} on $F(Q)$, by (3.4). Using (2.5) we find that each of these is on one, but not both, of W and Y . Hence, by (2.2) they are $B \cup W$ and $B \cup Y$. They are transversals of dimension $n-1$. The flats $X' \cup W$, $X' \cup Y$, $Z' \cup W$ and $Z' \cup Y$ are their connected lines of intersection with F_1 and F_2 . We note that a path (W, X', Y, Z', W) exists.

Assume that Q is not null-homotopic with respect to \mathbf{C} .

Suppose B is a transversal of dimension $n-2$ with poles X' on F_1 and Z' on F_2 . Suppose further that neither X' nor Z' belongs to \mathbf{C} . Then, by (5.1) there is a path R off \mathbf{C} from X' to Z' on B . Now $(W, X')R(Z', W)$ and $(X', Y, Z')R^{-1}$ are paths on the $(n-1)$ -flats $B \cup W$ and $B \cup Y$ respectively. Hence their dimensions are less than n and so they are null-homotopic with respect to \mathbf{C} . Using (6.1a), (5.3) and (5.2) we find $Q \sim (W, X', Y, Z', W) = (W, X')(X', Y, Z')(Z', W) \sim (W, Z')R^{-1}R(Z', W) \sim 0$ (\mathbf{C}). This is contrary to assumption. We deduce that each transversal of dimension $n-2$ has at least one pole in \mathbf{C} .

By (3.5) there is a transversal A of dimension $n-1$ which is not on Y . Let its lines of intersection with F_1 and F_2 be L_1 and L_2 respectively. They are connected lines on W . By (3.2) there is a point X' of $\mathbf{M} - \mathbf{C}$ other than W on L_1 . By (3.5) there is a connected $(n-2)$ -flat B of \mathbf{M} which is on A and X' but not on W . Now B is a transversal of dimension $n-2$. Let its pole on L_2 be U_2 . Then $U_2 \in \mathbf{C}$. Similarly there is a transversal B' of dimension $n-2$ on A having a point Z' of $\mathbf{M} - \mathbf{C}$ as its pole on L_2 and a point U_1 of \mathbf{C}

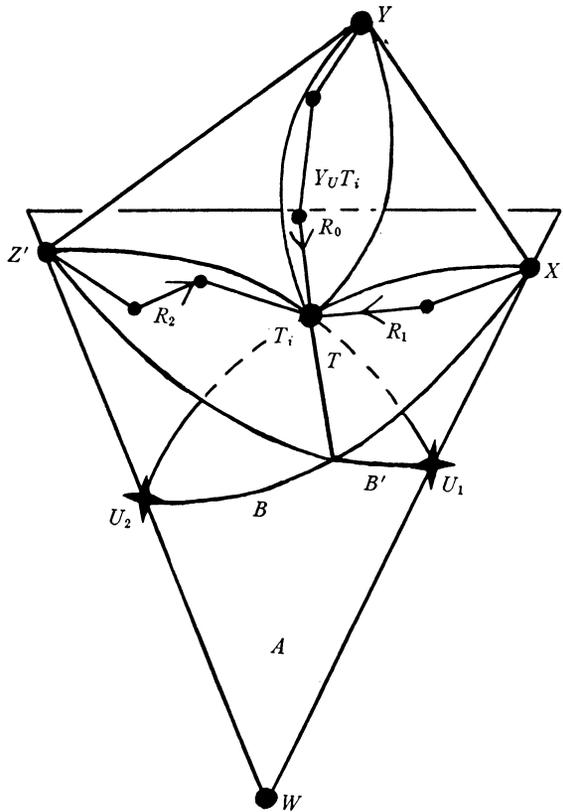


FIG. III

as its pole on L_1 . (See Figure III.) We write $T = \langle B \cap B' \rangle$. By (2.2) and (2.5) T is an $(n-3)$ -flat of \mathbf{M} .

Let \mathbf{S} be the class of all members of $\mathbf{M} - \mathbf{C}$ on T . Since $T \subset B$, $X' \in \mathbf{M} - \mathbf{C}$ and $U_2 \in \mathbf{C}$ it follows by (5.4) that \mathbf{S} is non-null.

Let T_i be any point of \mathbf{S} . Suppose that the flat $Y \cup T_i$ is connected. Then there is a path R_0 from Y to T_i on $Y \cup T_i$ which is off \mathbf{C} , by (5.1). Similarly there is a path R_1 from X' to T_i on B and a path R_2 from Z' to T_i on B' , both R_1 and R_2 being off \mathbf{C} . Now $(X', Y)R_0R_1^{-1}$ is a re-entrant path on the transversal $B \cup Y$ of dimension $n-1$, and $(Y, Z')R_2R_0^{-1}$ is a re-entrant path on the transversal $B' \cup Y$ of dimension $n-1$. Hence both these paths are null-homotopic with respect to \mathbf{C} . Applying (6.1a), (5.3) and (5.2) we find $Q \sim (W, X', Y, Z', W) = (W, X')(X', Y)(Y, Z')(Z', W) \sim (W, X')R_1R_0^{-1}R_0R_2^{-1} \cdot (Z', W) \sim (W, X')R_1R_2^{-1}(Z', W) (\mathbf{C})$. But the last path is on the $(n-1)$ -flat A and is therefore null-homotopic with respect to \mathbf{C} . Hence $Q \sim 0 (\mathbf{C})$, contrary to assumption. We deduce that $Y \cup T_i$ has a separation $\{Y, T_i\}$. That is $Y \cup T_i$ is a disconnected line, and Y and T_i are the two points on it.

We can repeat the above argument with $B \cup Y$ replacing A . Instead of

B' we then obtain a transversal B'' of dimension $n-2$ on $B \cup Y$ with its pole on $U_2 \cup Y$ a member of $\mathbf{M}-\mathbf{C}$ and its pole on $X' \cup Y$ a member of \mathbf{C} . We denote the $(n-3)$ -flat $\langle B \cap B'' \rangle$ by T' . We find that any point T'_i of $\mathbf{M}-\mathbf{C}$ on T' is such that $W \cup T'_i$ is a disconnected line. But, by (2.5), B'' has a point in common with the disconnected line $Y \cup T_i$, and this point can only be T_i . Hence T_i is one of the points of $\mathbf{M}-\mathbf{C}$ on T' .

We conclude that any point T_i of \mathbf{S} is such that $W \cup T_i$ and $Y \cup T_i$ are disconnected lines.

Let E be a connected flat of \mathbf{M} on F_1 and $F(Q)$ which is on some point of \mathbf{S} and has the least dimension consistent with this property. It is clear that either $F(Q)$ or one of its subsets satisfies these conditions. We have

$$(6.1b) \quad n = dF(Q) \geq dE \geq 3,$$

$$(6.1c) \quad d\langle E \cap T \rangle \geq dE - 3,$$

by (2.2) and (2.5). Choose a point N on $\langle E \cap T \rangle$, taking $N \in \mathbf{C}$ if this is possible. By (3.5) there is a connected $(dE-1)$ -flat E' of \mathbf{M} on F_1 and E but not on N . By (2.2) and (2.5) $\langle E' \cap T \rangle$ is a $(d\langle E \cap T \rangle - 1)$ -flat on $\langle E \cap T \rangle$. All the points of \mathbf{M} which are subsets of $\langle E' \cap T \rangle$ belong to \mathbf{C} , by the definition of E . By the choice of N this implies that either $d\langle E' \cap T \rangle = -1$ or $N \in \mathbf{C}$. But in the latter case all the points on $\langle E \cap T \rangle$ belong to \mathbf{C} , by (5.4), contrary to the definition of E . Hence $d\langle E' \cap T \rangle = -1$ and therefore $d\langle E \cap T \rangle = 0$. Hence $dE = 3$, by (6.1b) and (6.1c). Henceforth we use the symbol T_i to denote the single point $\langle E \cap T \rangle$, which must be in \mathbf{S} .

Suppose $n \geq 4$. Then F_2 is not on E . By (3.5) there is a connected $(n-1)$ -flat E'' of \mathbf{M} on F_2 and $F(Q)$ but not on T_i . Write $F_3 = \langle E'' \cap E \rangle$. Then F_3 is a plane on E and $W \cup Y$, by (2.2) and (2.5). By (3.5) there is a connected line L on E and T_i having no common point with $W \cup Y$. Let its common points with F_1 and F_3 be X_1 and X_3 respectively. Neither of these is T_i . We note that $F_3 = W \cup Y \cup X_3$, by (2.2). Now $X_1 \cup W$ and $X_1 \cup Y$ are connected lines by (4.3). Hence $X_1 \cap W$ and $X_1 \cap Y$ are both non-null, by (3.1). But we have shown that $W \cup T_i$ and $Y \cup T_i$ are disconnected lines. Hence $T_i \cap W$ and $T_i \cap Y$ are both null. But $L = X_1 \cup T_i = X_3 \cup T_i$, by (3.1). Hence $X_3 \cap W$ and $X_3 \cap Y$ are both non-null and therefore F_3 is connected.

By (5.1) there is a path R from Y to W on F_3 which is off \mathbf{C} . The re-entrant paths $(W, X, Y)R$ and $(Y, Z, W)R^{-1}$ are on E and E'' respectively and so have dimensions $< n$. Hence they are null-homotopic with respect to \mathbf{C} . Using (5.3) and (5.2) we find $Q = (W, X, Y)(Y, Z, W) \sim R^{-1}R \sim 0$ (\mathbf{C}), contrary to assumption.

We deduce, using (6.1b), that $n = 3$. This implies $dT = 0$. Hence T is a point of \mathbf{M} , identical with T_i and therefore a member of $\mathbf{M}-\mathbf{C}$. The three flats $W \cup Y$, $Y \cup T$ and $T \cup W$ are disconnected lines. The flat $W \cup Y \cup T$ is not a connected plane of \mathbf{M} by (4.3).

Any plane on $F(Q)$ has a point in common with each of the disconnected

lines $W \cup Y$, $Y \cup T$ and $T \cup W$. It is therefore on one of these lines. Each line of $F(Q)$ is on a plane of $F(Q)$, by (2.2) and (2.3). It follows that each line on $F(Q)$ is on one of the points W , Y and T .

Let F be any transversal of dimension 2. It meets F_1 and F_2 in connected lines L_1 and L_2 respectively. Let the points on L_1 other than W or Y be X_1, \dots, X_k . By (3.4) there is a transversal B_i of dimension 1 on F and X_i for each i . The line B_i must be on T . Hence $B_i = X_i \cup T$, and B_i is uniquely determined for each i . Let X'_i denote the point of intersection of B_i and L_2 . Since $B_i = T \cup X'_i$ for each i the k points X'_1, \dots, X'_k are all distinct. But at most one point on each of the lines L_1 and L_2 belongs to \mathbf{C} , and no transversal of dimension 1 has both its poles in $\mathbf{M} - \mathbf{C}$. Applying (3.2) we deduce that $k = 2$. Moreover we can adjust the notation so that $(X_1, X'_2) \in \mathbf{C}$ and $(X_2, X'_1) \in \mathbf{M} - \mathbf{C}$.

Distinct lines on T and F meet L_1 in distinct points, by (3.1). Hence the only connected lines on T and F are B_1 and B_2 . But each point of F is on two connected lines on F , and one of these is on T . Hence X_1 and X'_2 are the only points of \mathbf{C} on F . We thus prove that each connected plane on $F(Q)$ not on $W \cup Y$ is on just two points of \mathbf{C} .

Any connected line on F_1 is on a transversal of dimension 2, by (3.4). Hence it is on just three points, one of which is in \mathbf{C} . Distinct lines on F_1 and W (or Y) meet a given connected line on F_1 and Y (or W) in distinct points, by (3.1). It follows that on F_1 there are just two connected lines on each of the points W and Y . As each point on F_1 is on two connected lines on F_1 we deduce that F_1 is on just two points of \mathbf{C} . Analogous results hold for F_2 . Two distinct transversals of dimension 2 are both on T and therefore meet F_1 in distinct lines, by (2.2). Accordingly there are just two connected planes on $F(Q)$ and $T \cup W$, and just two on $F(Q)$ and $Y \cup T$ (since $W \cup Y \cup T$ is not a connected plane).

It follows from these results that either Q is an elementary re-entrant path of the fourth kind with respect to \mathbf{C} or there is a third connected plane F_3 on $F(Q)$ and $W \cup Y$. The first alternative must be rejected since it implies $Q \sim 0(\mathbf{C})$.

Let the points of intersection with F_3 of B_1 and B_2 be X''_1 and X''_2 respectively. These are both in $\mathbf{M} - \mathbf{C}$. Each is on two connected lines on F_3 , one on W and the other on Y , by (4.3). We have $(W, X, Y, X''_1, W) = (W, X, Y) \cdot (Y, X''_1, W) \sim 0(\mathbf{C})$. For otherwise we can repeat the first parts of the preceding proof with (W, X, Y, X''_1, W) replacing Q and obtain a contradiction, for the transversal B_2 then has both poles in $\mathbf{M} - \mathbf{C}$. Similarly $(Y, Z, W, X''_2, Y) = (Y, Z, W)(W, X''_2, Y) \sim 0(\mathbf{C})$. Applying (5.3) and (5.2) we find $Q = (W, X, Y)(Y, Z, W) \sim (W, X''_1, Y)(Y, X''_2, W) \sim 0(\mathbf{C})$, contrary to assumption. The lemma follows.

We return to the path P defined at the beginning of this proof. We note that $n = dF(P) \geq 1$ since otherwise P would be trivially null-homotopic with respect to \mathbf{C} . We choose a connected $(n - 1)$ -flat E of \mathbf{M} which is on $F(P)$

and the origin X_0 of P . This choice is possible, by (3.3).

Let $R = (X_0, \dots, X_m, X_0)$ be any re-entrant path with the same origin as P on $F(P)$. We write $u(R)$ for the number of terms of R , counting repetitions, which are not on E . If $u(R) > 0$ we write X_i for the first term of R which is not on E . We then write $v(R) = d(X_{i-1} \cup X_i \cup X_{i+1})$, taking $X_{m+1} = X_0$ if $i = m$. If $u(R) = 0$ we write $v(R) = 0$.

Henceforth we suppose R chosen so as to satisfy the following conditions:

- (i) $R \sim P$ (C),
- (ii) $u(R)$ has the least value consistent with (i),
- (iii) $v(R)$ has the least value consistent with (i) and (ii).

We consider first the case $u(R) > 0$. Then $v(R) > 0$. We may conveniently write R in the form $R_1(X_{i-1}, X_i, X_{i+1})R_2$, noting that R_1 is a path on \dots . We write also $F = X_{i-1} \cup X_i \cup X_{i+1}$.

Suppose $v(R) = 1$. Then F is a connected line. If $X_{i+1} = X_{i-1}$ we have $R \sim R_1R_2$ (C), by an elementary deformation of the first kind. This is impossible since $u(R_1R_2) < u(R)$. If $X_{i+1} \neq X_{i-1}$ then $(X_{i-1}, X_i, X_{i+1}, X_{i-1})$ is an elementary re-entrant path of the second kind with respect to C. Applying (5.3) we find $R \sim R_1(X_{i-1}, X_{i+1})R_2$ (C). This is impossible since

$$u(R_1(X_{i-1}, X_{i+1})R_2) < u(R).$$

Suppose $v(R) = 2$. Then F is a connected plane on $F(Q)$. It meets E in a line L , by (2.2) and (2.5). Let Z be the point of intersection of the lines L and $X_i \cup X_{i+1}$ on F . We discuss first the case $Z \in M - C$. In this case we define Q as the degenerate path (Z) if $Z = X_{i+1}$ and as the path (Z, X_{i+1}) otherwise. Then $(X_i, X_{i+1})Q^{-1}(Z, X_i)$ is an elementary re-entrant path of the first or second kind. Hence $R \sim R_1(X_{i-1}, X_i, Z)QR_2$ (C), by (5.3). If L is connected we have $(X_{i-1}, X_i, Z, X_{i-1}) \sim 0$ (C) and therefore $R \sim R_1(X_{i-1}, Z)QR_2$ (C), by (5.3). If L is not connected it is on a connected plane F' of M on E , by (4.2). We can find a connected line L' on X_{i-1} and F' , and a point T of $M - C$ distinct from X_{i-1} on L' . Then $T \cup X_{i-1}$ and $T \cup Z$ are connected lines, by (4.3). Using the lemma and the definition of n we find $(X_{i-1}, X_i, Z, T, X_{i-1}) \sim 0$ (C). Hence $R \sim R_1(X_{i-1}, T, Z)QR_2$ (C), by (5.3). So whether L is connected or not we have $R \sim R_3QR_2$ (C), where R_3 is on E . This is impossible since $u(R_3QR_2) < u(R)$.

We go on to the case $Z \in C$, illustrated in Figure IV. By (3.4) there is a connected line L' other than $X_i \cup X_{i+1}$ on X_{i+1} and F . If L' is on X_{i-1} we have $R \sim R_1(X_{i-1}, X_{i+1})R_2$ (C), using (5.3) with the elementary re-entrant path $(X_{i-1}, X_i, X_{i+1}, X_{i-1})$ of the second kind. This is impossible since $u(R_1(X_{i-1}, X_{i+1})R_2) < u(R)$. Hence L' must meet the lines $X_{i-1} \cup X_i$ and L in distinct points U and V respectively. Since $Z \in C$ and $X_{i-1} \in M - C$ we have $V \in M - C$.

Suppose $U \in M - C$. Using (5.3) with elementary re-entrant paths of the first and second kinds we find

$$R \sim R_1(X_{i-1}, U, X_i, X_{i+1})R_2 \sim R_1(X_{i-1}, V, U, X_{i+1})R_2 \sim R_1(X_{i-1}, V, X_{i+1})R_2(\mathbf{C}).$$

This is impossible since $u(R_1(X_{i-1}, V, X_{i+1})R) < u(R)$.

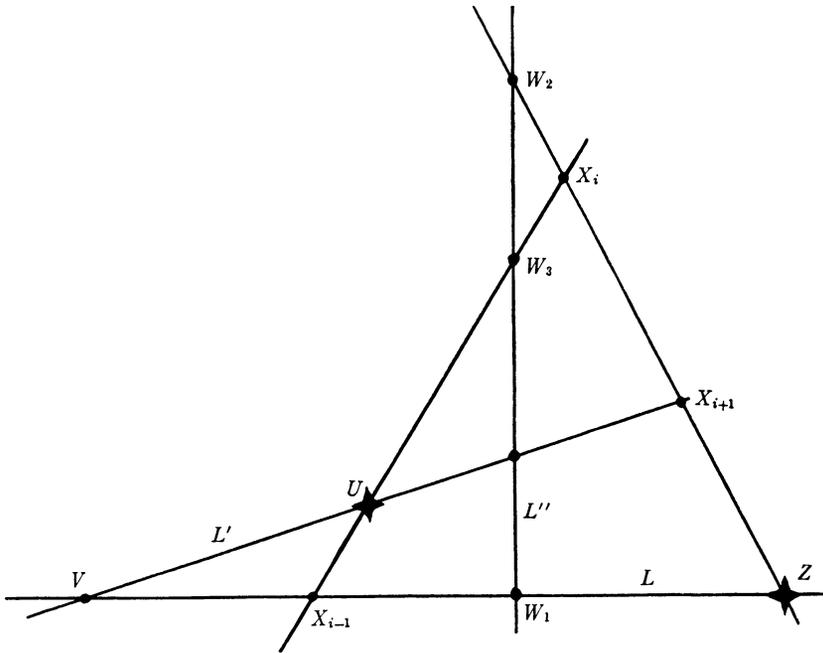


FIG. IV

Suppose $U \in \mathbf{C}$. It may happen that each connected line on F is on either U or Z . Then $(X_{i-1}, X_i, X_{i+1}, V, X_{i-1})$ is an elementary re-entrant path of the third kind with respect to \mathbf{C} . Using (5.3) we have $R \sim R_1(X_{i-1}, V, X_{i+1})R_2(\mathbf{C})$, which is impossible, as before. Hence there is a connected line L'' on F which is not on U or Z . If L'' is on X_{i+1} we can substitute it for L' in the preceding argument and so reduce to the case $U \in \mathbf{M} - \mathbf{C}$. We may therefore suppose L'' is not on X_{i+1} .

If L'' is on X_i it meets L in a point W_1 distinct from X_{i-1} and Z . Writing $R' = R_1(X_{i-1}, W_1, X_i, X_{i+1})R_2$ we have $R' \sim R(\mathbf{C})$, by (5.3). If L'' is not on X_i it meets $X_i \cup X_{i+1}$ in a point W_2 distinct from X_i, X_{i+1} and Z . If L'' is then on X_{i-1} we write $R' = R_1(X_{i-1}, W_2, X_{i+1})R_2$ and have

$$R' \sim R_1(X_{i-1}, W_2, X_i, X_{i+1})R_2 \sim R(\mathbf{C}),$$

by (5.3). If instead L'' is not on X_{i-1} it meets the lines $L, X_i \cup X_{i+1}$ and $X_{i-1} \cup X_i$ in distinct points W_1, W_2 and W_3 respectively of $\mathbf{M} - \mathbf{C}$. We then write $R' = R_1(X_{i-1}, W_1, W_2, X_{i+1})R_2$ and have $R' \sim R_1(X_{i-1}, W_1, W_3, W_2, X_i, X_{i+1})R_2 \sim R_1(X_{i-1}, W_3, X_i, X_{i+1})R_2 \sim R(\mathbf{C})$, by (5.3). For each of these three possibilities we have $R' \sim R(\mathbf{C})$, $u(R') = u(R)$ and $v(R') = v(R) = 2$. Hence we may

replace R by R' in the preceding argument. This reduces the problem to the case $U \in \mathbf{M} - \mathbf{C}$, which we have found to lead to a contradiction.

We now consider the case $v(R) > 2$. By (3.3) there is a connected plane K on $X_{i-1} \cup X_i$ and the connected $v(R)$ -flat F . This plane meets E in a line L . Choose a point T distinct from X_{i-1} on L and if possible in \mathbf{C} . By (3.5) there is a connected $(v(R) - 1)$ -flat F' on $X_i \cup X_{i+1}$ and F but not on T . Now F' is not on X_{i-1} , for otherwise we would have $F \subseteq F'$, contrary to (2.2). Hence F' meets L in a point T' distinct from X_{i-1} and T . It follows that L is connected, by (3.2), and that $T' \in \mathbf{M} - \mathbf{C}$. The flats K and F' intersect in a line L' on X_i and T' . If L' is connected we write

$$R' = R_1(X_{i-1}, T', X_i, X_{i+1})R_2$$

and have $R' \sim R(\mathbf{C})$, by (5.3). If L' is not connected it is on a connected plane K' on F' , by (4.2). K' meets E in a connected line L'' on T' , by (4.4). We can find a point U on L'' distinct from T' and in $\mathbf{M} - \mathbf{C}$. The flat $U \cup X_i$ is a connected line, by (4.3). Using the lemma and the definition of n we find $(T', U, X_i, X_{i-1}, T') \sim 0(\mathbf{C})$. In this case we write

$$R' = R_1(X_{i-1}, T', U, X_i, X_{i+1})R_2.$$

Then by (5.2) and (5.3) we have

$$\begin{aligned} R' &\sim R_1(X_{i-1}, T', U, X_i, X_{i-1}, T', X_{i-1}, X_i, X_{i+1})R_2 \\ &\sim R_1(X_{i-1}, T', X_{i-1}, X_i, X_{i+1})R_2 \sim R(\mathbf{C}). \end{aligned}$$

So whether L' is connected or not we have $R' \sim R(\mathbf{C})$, $u(R') = u(R)$ and $v(R') < v(R)$, which is contrary to the definition of R .

From the above analysis we deduce that $u(R) = 0$. Hence R is on E and has dimension $< n$. Hence $P \sim R \sim 0(\mathbf{C})$, contrary to assumption. The theorem follows.

7. Special cases. With \mathbf{C} null in (6.1) we find that every re-entrant path in a matroid \mathbf{M} is null-homotopic, as stated in the Introduction.

In applying this result to the circuit-matroid \mathbf{M} of a graph G we must remember that a path in \mathbf{M} corresponds to a sequence of circuits of G such that any two consecutive circuits form a nonseparable subgraph of rank 2. It can be shown that such a subgraph is made up of three arcs such that any two have both ends but no other edge or vertex in common. Each of the elementary deformations by which a re-entrant sequence of circuits can be transformed into a sequence with only one member operates within some nonseparable subgraph of rank ≤ 3 .

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