

OPERATIONS IN POLYADIC ALGEBRAS

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Abstract. A new treatment of P. R. Halmos' theory of terms and operations in (locally finite) polyadic algebras (of infinite degree) is given that is considerably simpler than the original one.

Introduction. By a "polyadic algebra" we shall mean a "locally finite polyadic algebra of infinite degree". The theory of terms and operations in polyadic algebras has been developed by Halmos in [4] and [5] (see also [6]). The difficulty of [4] has proved to be a stumbling block for many a student of algebraic logic. The purpose of this paper is to present a considerably simplified treatment of the main results of [4] and [5] concerning terms and operations. There are three reasons that this treatment is easier and shorter than the original one of Halmos. First we use a simpler definition of term due to Léon LeBlanc [7, pp. 19–20]. Second we take advantage of the canonical method for extending a polyadic algebra to one with equality described by Halmos in [5, §7]. Finally full use is made of the representation theorem for polyadic algebras. For the purpose of substituting terms in operations, a representation theorem for the transformation system of terms is proved that links up with the notion of generalized ultrapower (equivalent to that of limit ultrapower) introduced in [1, p. 115].

1. **Preliminaries.** We shall use the notations of Halmos [3]–[5]. The theory of polyadic algebras with equality is assumed to be known up to and including the representation theorem. Let $F(X^I, B)$, where B is a complete Boolean algebra, X is a nonempty set and I is the fixed infinite set of variables, denote the full functional (polyadic) algebra with value algebra B and domain X . This is the set of all functions $p: X^I \rightarrow B$ with finite support under the Boolean pointwise operations and the quantifiers $(\exists J)$ and transformation endomorphisms $S(\alpha)$ defined by the equations

$$(1) \quad [(\exists J)p](x) = \bigvee \{p(y) \mid yJ_*x\},$$

$$(2) \quad [S(\alpha)p](x) = p(x\alpha),$$

where $J \subset I$; $\alpha \in I^I$; $x, y \in X^I$ and yJ_*x means $y_j = x_j$ for all $j \notin J$. The algebra

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$F(X', B)$ has an equality E_0 known as the functional equality and defined by the conditions

$$(3) \quad \begin{aligned} E_0(i, j)(x) &= 1 && \text{if } x_i = x_j, \\ &= 0 && \text{if } x_i \neq x_j. \end{aligned}$$

These conditions imply that

$$(4) \quad \text{if } E_0(i, j)(x) = 1 \text{ then } x_i = x_j.$$

A subalgebra of $F(X', B)$ may have an equality E other than E_0 . If E has property (4) it is said to be reduced. The representation theorem for polyadic algebras proved in [3, 16.9 and 17.1] and [5, 6.7] can be stated thus:

THEOREM 1.1. *Every polyadic algebra A with equality is isomorphic to a subalgebra A' of some $F(X', B)$ with a reduced equality and which is regular, i.e. such that every supremum (1) with $p \in A'$ is attained for some y .*

The case of polyadic algebras generally is reduced to that of polyadic algebras with equality by means of the following theorem proved in [5, §7]:

THEOREM 1.2. *For every polyadic algebra A there is a canonically defined polyadic algebra with equality \bar{A} to which every constant or term of A extends canonically.*

The fact that every constant or term of A extends to \bar{A} is obvious from the construction of \bar{A} . This is so both for Halmos' original definitions of constant and of term (see [6, pp. 172, 178]) and for LeBlanc's definitions ([2, p. 89] and [7, p. 19]) which we shall adopt and which we now give.

Let A be a polyadic algebra. A constant of A is a function c which assigns to each subset K of I a Boolean endomorphism $c(K)$ of A (also denoted $S(c/K)$) in such a way that

- (C1) $c(\phi) = 1_A$, the identity map on A ,
- (C2) $c(K_1 \cup K_2) = c(K_1)c(K_2)$,
- (C3) $S(\alpha)c(K) = c(K)$ if α is identity on $I - K$,
- (C4) $c(K)S(\sigma) = S(\sigma)c(\sigma^{-1}K)$,

for all $\sigma \in I'$ and $K_1, K_2 \subset I$.

A term of A is a function t which assigns to each subset K of I a Boolean endomorphism $t(K)$ (also noted $S(t/K)$) such that for some finite subset J of I (which we shall call a support of t)

(T1) the restriction of t to the subsets of $I - J$ is a J -constant of A , i.e. a constant of the algebra A_J obtained from A by fixing the variables in J ; and

(T2) $t(K) = t(\sigma K)S(\sigma)$ for all transformations σ such that $\sigma K \subset K$ and σ is identity on $I - K$.

We shall call a term with support J a J -term. Every variable i induces an $\{i\}$ -term which we will identify with i . One advantage of these conditions is that quantifiers do not appear in them so that they make sense in transformation algebras generally,

and also that they are strictly equations between endomorphisms of A . LeBlanc's concept of constant was shown in [2, p. 97] to be equivalent to that of Halmos. That the same holds for terms will appear presently (see Corollary 2.6 below).

We insert here a simplified proof (proposed by the referee) that LeBlanc's concept of constant is equivalent to that of Halmos. The problem easily reduces to showing that $\exists(j)c(i) = c(i)\exists(j)$ for $i \neq j$. We may assume that our algebra has an equality. By the proof of [3, 10.1] we have

$$(1) \quad \begin{aligned} \exists(j)c(i)p &= \bigvee \{S(k/j)c(i)p : k \neq i\} \\ &= \bigvee \{c(i)S(k/j)p : k \neq i\} \quad \text{by (C4),} \end{aligned}$$

$$(2) \quad c(i)\exists(j)p = c(i) \vee \{S(k/j)p : k \neq i\}.$$

From (1) and (2) we easily infer

$$(3) \quad \exists(j)c(i)p \leq c(i)\exists(j)p.$$

Also, from (1) and (2) we see that it suffices to show that $c(i)$ is completely additive, i.e., that, if $\langle p_t : t \in T \rangle$ is a family of elements for which $\bigvee \{p_t : t \in T\} = r$ exists, then $\bigvee \{c(i)p_t : t \in T\}$ exists and equals $c(i)r$. (That this is so had also been proved by LeBlanc.) The following facts are useful first. The first two are easily established.

(4) If i, j, k are distinct, then $c(j)E(i, j) = c(k)E(i, k)$,

(5) $c(i)c(j)E(i, j) = 1$,

(6) if $i \neq j$ and $c(i)p \wedge q = 0$, then $p \wedge \exists(i)q \wedge c(j)E(i, j) = 0$.

To prove (6), first choose k so that $k \neq i, j$ and $\exists(k)p = p$. Then

$$\begin{aligned} 0 &= c(i)p \wedge q = c(k)S(k/i)p \wedge q = c(k)\exists(i)(p \wedge E(i, k)) \wedge q \\ &\geq \exists(i)c(k)(p \wedge E(i, k)) \wedge q \quad \text{by (3)} \\ &= \exists(i)(p \wedge c(k)E(i, k)) \wedge q \\ &= \exists(i)(p \wedge c(j)E(i, j)) \wedge q \quad \text{by (4).} \end{aligned}$$

The desired conclusion is then immediate.

Now we show that $c(i)$ is completely additive. Obviously $c(i)r$ is an upper bound for $\{c(i)p_t : t \in T\}$. Suppose that q is another upper bound for $\{c(i)p_t : t \in T\}$. Then, for any $t \in T$, $c(i)p_t \wedge q' = 0$, so by (6), with $j \neq i$ chosen arbitrarily,

$$p_t \wedge \exists(i)q' \wedge c(j)E(i, j) = 0 \quad \text{for all } t \in T.$$

Hence $r \wedge \exists(i)q' \wedge c(j)E(i, j) = 0$. It follows that $r \leq \forall(i)q \vee c(j)E(i, j)$. Applying $c(i)$ to both sides and using (5) we get $c(i)r \leq \forall(i)q \leq q$, as desired.

We shall assume known the theory of constants including their relationship to equality. We also assume known the few facts concerning predicates that can be stated prior to the definition of terms, including the substitution of constants in predicates. All our prerequisites can be found in [2, pp. 57–101]. For the remainder of this paper A will denote a fixed abstract polyadic algebra with equality E .

2. Functional terms. In this section we relate terms of regular subalgebras of $F(X^I, B)$ having reduced equalities, with ordinary operations on X .

THEOREM 2.1. *Let T be an n -ary operation on X and (i_1, \dots, i_n) a sequence of variables (not necessarily distinct). If, for $K \subset I$, $p \in F(X^I, B)$ and $x \in X^I$, we set*

$$(1) \quad [S(t/K)p](x) = p(y) \quad \text{where } yK_*x \text{ and } y_k = T(x_{i_1}, \dots, x_{i_n}) \text{ for all } k \in K,$$

then t is a term of $F(X^I, B)$ with support $J = \{i_1, \dots, i_n\}$.

The proof of this theorem is a straightforward calculation. We shall leave out all such calculations. This term t will be denoted by $\hat{T}(i_1, \dots, i_n)$ with no meaning attached to \hat{T} alone for the time being. A term of this sort is called a functional term.

For a converse we need to be able to substitute terms in predicates, and two lemmas.

THEOREM 2.2. *Let P be an n -place predicate of A and t_1, \dots, t_n terms of A with the same support J . Assume that i_1, \dots, i_n are distinct variables in $I - J$. Then*

$$(2) \quad P(t_1, \dots, t_n) = t_1(i_1) \cdots t_n(i_n)P(i_1, \dots, i_n)$$

is independent of the choice of i_1, \dots, i_n .

Proof. This equation concerns the algebra A_J (obtained by fixing the variables in J) in which the t_i are constants. The proof in the case where the t_i are constants is carried out in [2, p. 96] assuming $n=2$ for simplicity. Q.E.D.

LEMMA 2.3. *Let A_1 be a polyadic algebra and A_2 be a subalgebra with equality E . Let c_i be a constant of A_1 for $i=1, 2$. If $E(c_1, c_2)=1$, then $c_1(K)p=c_2(k)p$ for all $p \in A_2$ and $K \subset I$, i.e. A_2 is closed under all $c_1(K)$, and the constant of A_2 so determined by restrictions of c_1 is c_2 .*

The proof is the same as that of 8.1 in [5]. Note that here $E(c_1, c_2)=c_1(i)E(i, c_2)$ by definition whereas $c_2(i)E(c_1, i)$ is undefined.

LEMMA 2.4. *Every J -constant extends in at most one way to a J -term.*

Proof. We first prove a sublemma: If t is a J -term, $K \subset I$, p is an element of support L and $k \notin L$, then $t(K)p=t(K \cup \{k\})p$. If $k \in K$ this is obvious. Assume $k \notin K$. Choose σ such that $\sigma k \in K$ and $\sigma i=i$ if $i \neq k$. Then $\sigma(K \cup \{k\})=K$ and σ satisfies the conditions of (T2) with $K \cup \{k\}$ in place of K . We have, as L supports p and $\sigma i=i$ for all $i \in L$, $S(\sigma)p=p$, and the result follows by (T2).

For the proof of the lemma itself, suppose that the J -constant t has been extended to a J -term still denoted by t . Let $K \subset I$ and p be of finite support L . If $k \in K - (L \cup J)$ and σ is defined by the conditions $\sigma(i)=i$ for $i \in I - K$ and $\sigma(K) = \{k\}$, then, by (T2), $t(K)p=t(k)S(\sigma)p$. Therefore, if $K - (L \cup J)$ is nonempty, the term t is determined by the J -constant t . But, by the sublemma, we may assume that $K - (L \cup J)$ is nonempty. Q.E.D.

We will prove only later that every J -constant does extend to a J -term so as to minimize our efforts.

Note that it is clear from the proof that if t is a J -constant of a subalgebra A_2 of an algebra A_1 , then A_2 is closed under the extension of t to a J -term of A_1 if this extension exists.

THEOREM 2.5. *Every term t with support $J = \{i_1, \dots, i_n\}$ of a regular subalgebra A_1 of $F(X^I, B)$ with reduced equality E_1 is obtained by restriction from a term $\hat{T}(i_1, \dots, i_n)$ of $F(X^I, B)$ for some $T: X^n \rightarrow X$. Moreover T is uniquely determined if J and the sequence (i_1, \dots, i_n) are specified.*

Proof. If $i \notin J$ we have $(\exists i)E_1(i, t) = 1$ [2, 10.15]. Therefore $(\forall J)(\exists i)E_1(i, t) = 1$. Since $E_1(i, t) \wedge E_1(j, t) \leq E_1(i, j)$, for $j \notin J$, $(\forall J)(\exists ! i)E_1(i, t) = 1$. From this, the regularity of A_1 and the assumption that E_1 is reduced, it follows that, for any $(x_1, \dots, x_n) \in X^n$, there exists a unique $x \in X$ with the property: for all $y \in X^I$ such that $y_{i_1} = x_1, \dots, y_{i_n} = x_n$, we have $E_1(i, t)(y) = 1$ iff $y_i = x$. Set $T(x_1, \dots, x_n) = x$.

By definition we have

$$(3) \quad E_1[t, \hat{T}(i_1, \dots, i_n)] = S[\hat{T}(i_1, \dots, i_n)/i]E_1(t, i) \quad \text{if } i \notin J.$$

Hence, by virtue of a property of the J -constant $\hat{T}(i_1, \dots, i_n)$ of the algebra $F(X^I, B)$ with equality E_0 [2, 10.18],

$$(4) \quad E_1[t, \hat{T}(i_1, \dots, i_n)] = (\exists i)[E_1(i, t) \wedge E_0(i, \hat{T}(i_1, \dots, i_n))].$$

By the definition of T , the right member of (4) is 1. It follows from Lemma 2.3 that the J -constant $\hat{T}(i_1, \dots, i_n)$ induces the J -constant t in A_1 . From this, Lemma 2.4 and the second remark following the proof of Lemma 2.4, it follows that the term $\hat{T}(i_1, \dots, i_n)$ induces the term t in A_1 .

Finally to show the unicity of T , observe that the T that we have defined is the only one that makes the right member of (4) equal to 1. And, if T_1 satisfies the theorem, then $E_1[t, \hat{T}_1(i_1, \dots, i_n)] = 1$ while (4) is valid for any T . Q.E.D.

COROLLARY 2.6. *The J -terms of a polyadic algebra in the foregoing sense are the same as its J -terms in the sense of Halmos [4, §4].*

Proof. It follows directly from Theorems 1.2, 1.1, 2.5 and the easily verifiable fact that the terms $\hat{T}(i_1, \dots, i_n)$ satisfy Halmos' definition, that any J -term in the foregoing sense is one in the sense of Halmos. The converse follows from the same statements except that, in Theorems 1.2 and 2.5 and their proofs, " J -term" is understood in the sense of Halmos. Q.E.D.

3. The transformation system of terms. By a transformation system we understand a set \mathcal{T} together with a map $I^I \times \mathcal{T} \rightarrow \mathcal{T}$ that sends (σ, t) onto σt in such a way that $\delta t = t$ and $\sigma_1(\sigma_2 t) = (\sigma_1 \sigma_2)t$ where δ is the identity map of I onto itself. A subset J of I supports an element t of \mathcal{T} if $\sigma_1 t = \sigma_2 t$ whenever $\sigma_1|_J = \sigma_2|_J$, i.e. σ_1 and σ_2 have the same effect on J . A polyadic algebra is in an obvious sense a

transformation system and this concept of support coincides with that of Halmos [2, 8.10]. The concept of n -ary predicate of \mathcal{F} is the same as for polyadic algebras and the basic properties of predicates of polyadic algebras remain unchanged for locally finite transformation systems generally [2, 8.11 and 8.12]. In this section we endow the set of all terms of the polyadic algebras A with a transformation structure and we shall henceforth reserve the letter “ \mathcal{F} ” to denote that transformation system.

Let t be the J -term of a regular subalgebra A_1 of $F(X^I, B)$ with reduced equality. Then, in the sense of Theorem 2.5, $t = \hat{T}(i_1, \dots, i_n)$ where $J = \{i_1, \dots, i_n\}$, for some unique $T: X^n \rightarrow X$. It is easy to compute that

$$(1) \quad S[\hat{T}(\sigma_{i_1}, \dots, \sigma_{i_n})/K]q = S(\alpha)S(t/\rho(K \cap L))S(\rho)q$$

where $q \in A_1$, L is a finite support of q , ρ is one-one on L , $\rho L \cap J = \emptyset$, $\alpha|J = \sigma|J$ and $\alpha\rho i = i$ for all $i \in L$. This shows that A_1 is closed under the term $\hat{T}(\sigma_{i_1}, \dots, \sigma_{i_n})$. We define σt to be the σJ -term of A_1 induced by $\hat{T}(\sigma_{i_1}, \dots, \sigma_{i_n})$. The representation theorem (Theorem 1.1) makes it possible to transfer this definition of σt to the abstract algebra A . This definition of σt , for t a term of A , is independent of the representation chosen since the right member of (1) is computed in A . Since, in A_1 , $\sigma t = \hat{T}(\sigma_{i_1}, \dots, \sigma_{i_n})$, it is obvious that \mathcal{F} becomes a transformation system.

We define an n -ary operation of A to be an n -ary predicate of \mathcal{F} . If the regular subalgebra A_1 of $F(X^I, B)$ is closed under $\hat{T}(i_1, \dots, i_n)$ for all $i_1, \dots, i_n \in I$, then \hat{T} can be viewed as an operation of A_1 . Moreover, Theorem 2.5 means that any operation Q of A_1 is of the form \hat{T} for a unique T . We set $T = Q^*$.

The set of all functions $u: X^I \rightarrow X$ forms a transformation system if we set $(\sigma u)(x) = u(x\sigma)$ for $x \in X^I$. The set of all such functions with finite support is a locally finite transformation system denoted by $F(X^I, X)$. Any operation $T: X^n \rightarrow X$ induces a predicate \check{T} of $F(X^I, X)$ defined by

$$(2) \quad [\check{T}(i_1, \dots, i_n)](x) = T(x_{i_1}, \dots, x_{i_n})$$

where $x \in X^I$. Any predicate P of $F(X^I, X)$ is of the form \check{T} for a uniquely defined T given by

$$(3) \quad T(a_1, \dots, a_n) = [P(i_1, \dots, i_n)](x)$$

where $(a_1, \dots, a_n) \in X^n$, i_1, \dots, i_n are distinct variables and $x \in X^I$ is such that $x_{i_1} = a_1, \dots, x_{i_n} = a_n$.

Every variable i induces an element $\check{1}_x(i)$ of $F(X^I, X)$, 1_x being the identity function on X .

THEOREM 3.1. *Let $\mathcal{F}(X, I, B)$ be the transformation system of functional terms of $F(X^I, B)$. Then an isomorphism $\phi: \mathcal{F}(X, I, B) \rightarrow F(X^I, X)$ is defined by the equation $\phi[\hat{T}(i_1, \dots, i_n)] = \check{T}(i_1, \dots, i_n)$. We will identify these two systems by this isomorphism.*

Proof. The only not quite trivial part of the statement is that ϕ is well defined and injective. This means that for any two operations $T_1: X^n \rightarrow X$ and $T_2: X^m \rightarrow X$ and any two sequences of variables (i_1, \dots, i_n) and (j_1, \dots, j_m) ,

$$(1) \quad \hat{T}_1(i_1, \dots, i_n) = \hat{T}_2(j_1, \dots, j_m)$$

iff

$$(2) \quad \check{T}_1(i_1, \dots, i_n) = \check{T}_2(j_1, \dots, j_m).$$

Note that both members of (1) are $\{i_1, \dots, i_n, j_1, \dots, j_m\}$ -terms and that, by virtue of [5, 8.1] (see Lemma 2.3), (1) means that

$$(3) \quad E_0[\hat{T}_1(i_1, \dots, i_n), \hat{T}_2(j_1, \dots, j_m)] = 1.$$

Since E_0 is reduced, both (2) and (3) mean that, for all $x \in X^I$,

$$(4) \quad T_1(x_{i_1}, \dots, x_{i_n}) = T_2(x_{j_1}, \dots, x_{j_m})$$

(of course, for (4) to happen, T_1 and T_2 must be very closely related: T_2 must be obtainable from T_1 by relocating the relevant places and by leaving out or adding some irrelevant places). Q.E.D.

The definition of the image of a term of a polyadic algebra A_1 by a polyadic onto homomorphism $A_1 \rightarrow A_2$ poses no difficulty. This is done exactly as in the case of constants [3, 15.3].

A regular reduced representation of A is a polyadic homomorphism of A onto a regular subalgebra A_1 of some $F(X^I, B)$ with reduced equality E_1 . Theorems 2.5 and 3.1 make it possible to define in an obvious way, for any such representation of A , a companion representation for \mathcal{F} which will be a transformation homomorphism from \mathcal{F} into $F(X^I, X)$.

This generalizes to the case where E_1 is not reduced: for that one must consider the system $F(X^I, Y)$ where Y is the quotient of X by the equivalence relation \sim on X determined by $E_1: a \sim b$ iff $[E_1(i, j)](x) = 1$ for $x \in X^I$ such that $x_i = a$ and $x_j = b$ for two distinct variables i and j .

The natural surjection $X \rightarrow Y$ induces a surjection $X^I \rightarrow Y^I$ which, by composition, yields a transformation monomorphism $F(Y^I, Y) \rightarrow F(X^I, Y)$ and a polyadic monomorphism $F(Y^I, B) \rightarrow F(X^I, B)$. The (regular) representation $A \rightarrow F(X^I, B)$ splits uniquely through a (regular) reduced representation $A \rightarrow F(Y^I, B)$ so that $A \rightarrow F(X^I, B) = A \rightarrow F(Y^I, B) \rightarrow F(X^I, B)$. The companion representation $\mathcal{F} \rightarrow F(Y^I, Y)$ of the reduced representation yields a companion representation $\mathcal{F} \rightarrow F(X^I, Y) [= \mathcal{F} \rightarrow F(Y^I, Y) \rightarrow F(X^I, Y)]$ of the given representation.

The notion of the K -compression of an I -transformation system can be defined just as with polyadic algebras, K being a subset of I . The K -compression $\mathcal{F}^{(K)}$ of \mathcal{F} is the set of J -terms of A such that $J \subset K$. A regular (not necessarily reduced) representation of A into $F(X^I, B)$ will induce a companion representation of $\mathcal{F}^{(K)}$ into $F(X^K, Y)$. In [1, p. 115] generalized ultrapowers of a set Y have been defined as

quotients of sets of the form $F(X^K, Y)$ by (Boolean) ultrafilters of the (K -polyadic) algebra of all functions with finite support from X^K to $2 = \{0, 1\}$.

4. Terms in operations. If Q is an n -ary operation of A and $t_1, \dots, t_n \in \mathcal{T}$, we wish to define the term $t = Q(t_1, \dots, t_n)$. We wish that the definition be such that

$$(1) \quad S(t/K)p = S(t_1/i_1) \cdots S(t_n/i_n)S[Q(i_1, \dots, i_n)/K]p$$

whenever i_1, \dots, i_n are distinct variables none of which is in a common support of p, t_1, \dots, t_n and K is an arbitrary subset of I . To avoid the rather cumbersome problem of showing directly that (1) is an unambiguous definition of a term t we rely on the representation theorem (Theorem 1.1) and on the representation theorem for the system of terms $\mathcal{T}(X, I, B)$ (Theorem 3.1).

Identify momentarily A with a regular subalgebra of some $F(X^I, B)$ with reduced equality. Then Q is essentially \hat{T} for some $T: X^n \rightarrow X$. Before transforming \hat{T} from a function $I^n \rightarrow \mathcal{T}$ into an n -ary operation on \mathcal{T} , we transform \check{T} from a function $I^n \rightarrow F(X^I, X)$ into an n -ary operation on $F(X^I, X)$. For this we let $u_1, \dots, u_n \in F(X^I, X)$ and $x \in X^I$ and set

$$(2) \quad [\check{T}(u_1, \dots, u_n)](x) = T(u_1(x), \dots, u_n(x)).$$

Now define $\hat{T}(t_1, \dots, t_n)$ by

$$(3) \quad \hat{T}(t_1, \dots, t_n) = \phi^{-1}[\check{T}(\phi t_1, \dots, \phi t_n)]$$

where ϕ is the isomorphism from the system of functional terms of $F(X^I, B)$ onto the system $F(X^I, X)$. Note that, a priori, (3) defines a term t of $F(X^I, B)$ under which A might not be closed. That A is closed under t follows from (1) which can be verified for $p \in F(X^I, B)$ by an easy computation. For this, note that, if $J = \{i_1, \dots, i_m\}$ is a common support of t_1, \dots, t_m the term t is $\hat{U}(i_1, \dots, i_m)$ where U is an m -ary operation on X defined by

$$(4) \quad U(x_1, \dots, x_m) = [\check{T}(\phi t_1, \dots, \phi t_n)](y)$$

where $y \in X^I$ is such that $x_k = y_{j_k}$ for $k = 1, \dots, m$ (this definition does not depend on y). That the definition of $Q(t_1, \dots, t_n)$ does not depend on the representation chosen for A also follows from (1).

5. Operations and monovalent predicates. An $(n+1)$ -ary predicate P of A is monovalent in its first place if, for distinct variables k, i, i_1, \dots, i_n ,

$$(1) \quad (\exists i)P(i, i_1, \dots, i_n) = 1.$$

and

$$(2) \quad P(i, i_1, \dots, i_n) \wedge P(k, i_1, \dots, i_n) \rightarrow E(i, k) = 1.$$

The conjunction of the left members of these two equations is written

$$(\exists! i)P(i, i_1, \dots, i_n).$$

The following theorem has first been proved by Halmos [5, 9.2]. The proof given here generalizes that given in [2, 11.5] for the case of constants.

THEOREM 5.1. *The n -ary operations T of A and the $(n+1)$ -ary predicates P of A monovalent in their first place are in one-one correspondence by the following relations:*

$$(3) \quad E[i, T(i_1, \dots, i_n)] = P(i, i_1, \dots, i_n)$$

for arbitrary variables i, i_1, \dots, i_n , and

$$(4) \quad S[T(i_1, \dots, i_n)/K]p = (\exists i)[S(i/K)p \wedge P(i, i_1, \dots, i_n)]$$

where $i \notin L \cup \{i_1, \dots, i_n\}$, L being a support of p .

Proof. Let T be given and define P by (3). We have (1) by [2, 10.15] and (2) by applying $S[T(i_1, \dots, i_n)/j]$ to $E(i, j) \wedge E(k, j) \leq E(i, k)$. We have (4) with $K \cup I - J$, where $J = \{i_1, \dots, i_n\}$, by [2, 10.18]. The case of an arbitrary K follows from this special one: for let $k \notin L \cup \{i_1, \dots, i_n\}$, then as in the proof of Lemma 2.4, with $t = T(i_1, \dots, i_n)$,

$$\begin{aligned} S(t/K)p &= S(t/k)S(k/K)p \\ &= (\exists i)[S(i/k)S(k/K)p \wedge P(i, i_1, \dots, i_n)]. \\ &= (\exists i)[S(i/K)p \wedge P(i, i_1, \dots, i_n)]. \end{aligned}$$

Hence the map $T \mapsto P$ is injective.

Let now P be an $(n+1)$ -ary predicate monovalent in its first place. Assuming that (4) does define an operation T and letting $K = \{i\}$ and $p = E(j, i)$ in (4), we have

$$E[j, T(i_1, \dots, i_n)] = (\exists i)[E(j, i) \wedge P(i, i_1, \dots, i_n)] = P(j, i_1, \dots, i_n).$$

Hence the map $T \mapsto P$ is onto.

The proof that (4) does define an operation T begins with an easy modification of a part of the proof of 9.2 in [5]. For this purpose, identify A with a regular subalgebra of some $F(X', B)$ with reduced equality. Then for all $x_1, \dots, x_n \in X$ there is a unique $x \in X$ such that $[P(i, i_1, \dots, i_n)](y) = 1$ whenever $y_i = x$ and $y_{i_k} = x_k$ for $k = 1, \dots, n$. Define $Q: X^n \rightarrow X$ by setting $Q(x_1, \dots, x_n) = x$. Then it is clear that

$$(5) \quad P[\hat{Q}(i_1, \dots, i_n), i_1, \dots, i_n] = 1.$$

Now apply $S[\hat{Q}(i_1, \dots, i_n)/k]$ to

$$P(i, i_1, \dots, i_n) \wedge P(k, i_1, \dots, i_n) \leq E(i, k)$$

to get, using (5),

$$P(i, i_1, \dots, i_n) \leq E[i, \hat{Q}(i_1, \dots, i_n)].$$

On the other hand,

$$P(k, i_1, \dots, i_n) \wedge E(i, k) = P(i, i_1, \dots, i_n) \wedge E(i, k).$$

And hence,

$$P(k, i_1, \dots, i_n) \wedge E(i, k) \leq P(i, i_1, \dots, i_n).$$

Applying $S[\hat{Q}(i_1, \dots, i_n)/k]$ and using (5) again

$$E[i, \hat{Q}(i_1, \dots, i_n)] \leq P(i, i_1, \dots, i_n).$$

Hence

$$(6) \quad E[i, \hat{Q}(i_1, \dots, i_n)] = P(i, i_1, \dots, i_n).$$

Now assume first that A is simple. Then B can be taken to be 2. (This, incidentally, follows from 1.1 since any regular representation f of A into $F(X^I, B)$ yields a regular representation \bar{f} of A into $F(X^I, B/N)$ for any ideal N of B simply by setting $(\bar{f}p)(x) = (fp)(x)/N$ and this representation \bar{f} can be reduced.) Then E is the functional equality and (4) holds with \hat{Q} instead of T in view of (6) as can easily be calculated. When A is not simple (4), with $T = \hat{Q}$, is not so obvious though it is true. It holds if $K \subset I - \{i_1, \dots, i_n\}$ by virtue of the properties of constants. If we knew that every J -constant extends to a J -term then we could easily claim the validity of (4) with $T = \hat{Q}$ by virtue of this, (T2) and the first paragraph of the proof of Lemma 2.4.

The case where A is not simple will follow by semisimplicity. For this we digress a little. By the direct product of a family of polyadic algebras we understand the subalgebra of all elements with finite support in the (possibly not locally finite) cartesian product algebra with the "operations" $(\exists J)$ and $S(\alpha)$ performed componentwise. This direct product, together with the obvious projection epimorphisms, is the product in the category of (equality) algebras with (equality) homomorphisms. A predicate P of such a product $C = \prod \{C_h \mid h \in H\}$ can be identified with a family $\{P_h \mid h \in H\}$ of predicates of the C_k by the equation

$$(7) \quad P(i_1, \dots, i_n) = \{P_h(i_1, \dots, i_n)\}.$$

Similarly, for operations we have the following obvious

LEMMA 5.2. *If $C = \prod \{C_h \mid h \in H\}$ is a direct product of polyadic algebras and if, for all h , T_h is an n -ary operation of C_h , then an n -ary operation T of C is defined by the equation*

$$(8) \quad S[T(i_1, \dots, i_n)/K]p = \{S[T_h(i_1, \dots, i_n)/K]p_h\}$$

where $p = \{p_h\} \in C$.

From the semisimplicity of polyadic algebras [3, 9.3] it follows that any polyadic algebra is isomorphic to a subalgebra of a direct product of simple algebras in the usual manner. The fact that (4) defines an operation in general follows from the above remarks. This completes the proof of Theorem 5.1.

An immediate consequence which could also have been proved earlier (see Lemma 2.4) with some independent labor is the following

COROLLARY 5.3. *Every J -constant extends to a J -term.*

Proof. Let $J = \{j_1, \dots, j_n\}$ and P be the unique predicate such that $P(i, j_1, \dots, j_n) = S(c/k)E(i, k)$ where $i, k \in I - J$. Then P is monovalent in its first place and if T is obtained by 5.1 then $T(j_1, \dots, j_n)$ extends c . Q.E.D.

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