\textbf{\textit{A}}-GENERATORS FOR THE DICKSON ALGEBRA

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\textbf{Abstract.} Let $D_k$ denote the Dickson algebra in $k$ variables over the field of two elements. We study the problem of determining a minimal set of generators for $D_k$ as a module over the Steenrod algebra $\mathcal{A}$. This is easy for $k = 1$ and 2. In this paper we answer this question for $k = 3$ and 4 and give techniques which may help solve the problem for general $k$.

1. Introduction

Let $P_k = H^*(RP^\infty \times \cdots \times RP^\infty; \mathbb{Z}/2)$, $k$-times, i.e. $P_k = \mathbb{Z}/2[x_1, \ldots, x_k]$, where $|x_i| = 1$. The general linear group $G = GL(k, \mathbb{Z}/2)$ acts on $P_k$, and the ring of invariants, $(P_k)^G = D_k$, was described by Dickson. He showed that $D_k \cong \mathbb{Z}/2[Q_{k-1}, Q_{k-2}, \ldots, Q_0]$, again a polynomial algebra in $k$ variables, with $|Q_s| = 2^k - 2^s$ ([1]). Note that $Q_s$ depends on $k$, and when necessary, will be denoted $Q_{k,s}$. An inductive definition of $Q_{k,s}$ is given by

$$Q_{k,s} = (Q_{k-1,s-1})^2 + V_k \cdot Q_{k-1,s},$$

where, by convention, $Q_{k,1} = 1$, $Q_{k,s} = 0$ for $s < 0$ and

$$V_k = \prod_{\lambda_i \in \mathbb{Z}/2} (\lambda_1 x_1 + \cdots + \lambda_{k-1} x_{k-1} + x_k).$$

Since the operations of $GL(k, \mathbb{Z}/2)$ on $P_k$ commute with the action of the Steenrod algebra, $D_k$ is also a module (in fact an algebra) over $\mathcal{A}$. The importance of $D_k$ to topologists was first shown by Madsen [6] who calculated the dual to the Dyer–Lashof operations of length $k$, and by Mùi [2] who related this to $D_k$. Furthermore, the cohomology operations derived from the invariants in $D_k$ are exactly the Steenrod–Milnor ones of length $k$ (see Madsen–Milgram [7], Mùi [3]).

In this paper we investigate the structure of $\mathbb{Z}/2 \otimes_\mathcal{A} D_k$, the minimal set of $\mathcal{A}$-generators for $D_k$. The motivation for this problem is provided by the fact that a calculation of $\mathbb{Z}/2 \otimes_\mathcal{A} D_k$ for all $k$ also calculates the $\mathcal{A}$-annihilated elements of $PH_*(Q_0 S^0)$ which form the bottom line of the $E^2$-term of the unstable Adams spectral sequence which converges to $\pi_*(Q_0 S^0)$ ($\cong \pi_*(S^0)$) (this has been studied by Wellington [10], who computed this in dimensions

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This problem is also related to the study of the Hurewicz map \( \pi_*(Q_0 S^0) \to H_*(Q_0 S^0; \mathbb{Z}/2) \) (see Lannes–Zarati [4], [5]).

The paper is organized as follows. The key definitions and the main results are stated in Section 2. Then we give an outline of the proofs in Section 3. Section 4 deals with how to determine inductively all the allowed sequences of length \( k \) for any \( k \). In Section 5, the fact that the allowed monomials generate \( \mathbb{Z}/2 \otimes_{\mathbb{Z}} D_k \) for \( k \leq 4 \) is showed to be equivalent to the main lemma, whose proof is given in the next three sections. The linear independence of the allowed monomials is proved in Section 9 for \( k \leq 4 \). Finally, we give some remarks and conjectures in Section 10.

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2. Statement of results

Definition 2.1. (i) Let \( u \) be an integer. Define \( s(u) = s_1(u) \) to be the maximal non-negative integer such that \( u + 1 \) is divisible by \( 2^{s(u)} \) but not by \( 2^{s(u)+1} \).

For \( u \neq -1 \), \( s(u) \) is well defined. Convention: \( s(-1) = \infty \).

In other words, if \( u \neq -1 \), \( u = 2^{s(u)} - 1 \mod 2^{s(u)+1} \). That means, for \( u \geq 0 \), \( s(u) \) is the non-negative integer with \( 2^{s(u)} \) being the first missing 2 power in the dyadic expansion of \( u \).

(ii) Suppose \( u \geq 0 \). Define \( s'(u) = s_2(u) \) to be the positive integer with \( 2^{s(u)} \) being the second missing 2 power in the dyadic expansion of \( u \). That means \( s_2(u) = s(u + 2^{s(u)}) \).

Thus, if \( u \) is a non-negative integer, \( u = 2^{s_1(u)} - 1 + 2^{s_1(u)+1} + \ldots + 2^{s(u)-1} + 2^{s(u)+1} t \), where \( t \) is a non-negative integer.

Definition 2.2. \( u \triangleright v \) if \( v < 2^{s(u)} \). Here \( u, v \) are non-negative integers.

Remark 2.3. If \( u \triangleright v \) then \( s(u) \geq s(v) \). The equality happens if and only if \( v = 2^{s(u)} - 1 \). In particular, \( u \triangleright u \) if and only if \( u = 2^{s} - 1 \) for some \( s \).

Definition 2.4. \( u \triangleright\triangleright v \) if \( 2^{s_1(u)} \leq v < 2^{s_2(u)} \) and \( s_1(u) = s_1(v) = s_2(v) - 1 \).

Let \( I = (i_{k-1}, i_{k-2}, \ldots, i_0) \) be a sequence of \( k \) non-negative integers, with corresponding monomial \( Q^I = Q_0^{i_0} \cdots Q_k^{i_{k-1}} \in D_k \). Let \( h(I) = \sum_{j=0}^{k-1} i_j \), the classical degree of \( Q^I \).

Definition 2.5. (i) \( I \) is strongly allowed if \( h(I) \triangleright i_{k-1} \triangleright i_{k-2} \triangleright \cdots \triangleright i_0 \).

(ii) \( I \) is weakly allowed if one or more \( \triangleright \) in the above definition is replaced by \( \triangleright\triangleright \).

(iii) \( I \) is allowed if it is either strongly or weakly allowed.

(iv) \( Q^I \) is (strongly, weakly) allowed if \( I \) is.
Let us consider small values of $k$. If $k = 1$, then by Remark 2.3 $i_0 \triangleright i_0$ if and only if $i_0 = 2^s - 1$. So it is well-known that the strongly allowed monomials form a basis for $\mathbb{Z}/2 \otimes_{\mathcal{S}} D_1$.

If $k = 2$, then it is easy to check that $i_1 + i_0 \triangleright i_1 \triangleright i_0$ if and only if $i_1 = 2^s - 1$ and $i_0 = 0$. Thus the strongly allowed monomials form a basis of $\mathbb{Z}/2 \otimes_{\mathcal{S}} D_2$ (see, e.g., Singer [9]).

Our first theorem is the following one.

**Theorem 2.6.** Let $k = 3$. Then the set of strongly allowed monomials forms a basis for $\mathbb{Z}/2 \otimes_{\mathcal{S}} D_3$.

We note which $I$ are strongly allowed for $k = 3$:

$$(2^s - 1, 0, 0), \quad s \geq 0,$$

$$(2^r - 2^s - 1, 2^s - 1, 1), \quad r > s > 0.$$

The corresponding $Q^I$ is of dimension $2^{s+2} - 4$ and $2^{r+2} + 2^{s+1} - 3$ respectively.

This pattern does not continue, as the set of strongly allowed monomials does not span $D_k$ for $k > 3$. It should also be noted that there is no weakly allowed sequence for $k \leq 3$.

Our second theorem handles the case $k = 4$.

**Theorem 2.7.** Let $k = 4$. Then the set of allowed monomials is a basis for $\mathbb{Z}/2 \otimes_{\mathcal{S}} D_4$.

We show that the only allowed sequences $I$ for $k = 4$ are the following ones:

$$(2^s - 1, 0, 0, 0), \quad s \geq 0,$$

$$(2^r - 2^s - 1, 2^s - 1, 1, 0), \quad r > s > 0,$$

$$(2^t - 2^r - 1, 2^r - 2^s - 1, 2^s - 1, 2), \quad t > r > s > 1,$$

$$(2^t - 2^{r+1} - 2^s - 1, 2^s - 1, 2^r - 1, 2), \quad r > s + 1 > 2.$$

The corresponding $Q^I$ is of dimension $2^{s+3} - 8$, $2^{r+3} + 2^{s+2} - 6$, $2^{t+3} + 2^{r+2} + 2^{s+1} - 4$ and $2^{r+3} + 2^{s+1} - 4$ respectively.

Note that the last sequence with $r > s + 2$ is the only weakly allowed one. It has

$$h(I) \triangleright i_3 \triangleright i_2 \triangleright i_1 \triangleright i_0.$$

We hope to solve this problem for larger $k$ in the future. The proofs of Theorems 2.6 and 2.7 are very long and complicated and involve studying many cases. In Section 3 we give formulæ that are needed and describe the global structure of the proofs. The detailed lemmas are then proved in later sections.

### 3. Outline of the proofs

Many algebraic topologists have studied the Dickson algebra. For background, see the primer by Wilkerson [11].

We will need the action of $\mathcal{S}$ on $D_k$. This is given by the following theorem (see Hu'ng [8]).
Theorem 3.1 (Hai–Hu’ng).

\[ Sq^i(Q_{k,i}) = \begin{cases} 
Q_{k,i}, & i = 2^r, \quad r \leq s, \\
Q_{k,i}Q_{k,r}, & i = 2^r - 2^s, \quad r \leq s < t, \\
Q_{k,r}, & i = 2^k - 2^t, \quad r \leq s < t, \\
Q_{k,s}^2, & i = 2^k - 2^t, \\
0, & \text{otherwise.}
\]

Corollary 3.2. Suppose \( q \) and \( l \) are non-negative integers with \( 0 \leq l < k \). The submodule spanned by all \( Q^l \) with \( i_1 + i_{l-1} + \cdots + i_0 \geq q \) is a \( \mathcal{A} \)-ideal of \( D_k \). Therefore, if \( Sq^2(Q_l) = Q_l^l + \text{other terms} \), then \( j_1 + j_{l-1} + \cdots + j_0 \geq i_1 + i_{l-1} + \cdots + i_0 \) for any \( l < k \).

From the definition of strongly allowed, if \( I \) is not strongly allowed, then it is easy to find a sequence \( K \) and an integer \( a \) so that \( Sq^2(Q^K) = Q^l + \text{other terms} \). We give here a canonical way to do that.

3.3. If \( m \not\in m_{m-1}, s = s(m) \), then

\[ Sq^{2m-1}Q^{i_1, i_2 + m, \ldots, i_{m-1} - 2, \ldots} = Q^l + \text{other terms}; \]

3.4. Also, if \( h(I) \not\in i_{k-1}, s = s(h) \), then

\[ Sq^{2m-1}Q^{i_{k-2} - 2, i_{k-2}, \ldots} = Q^l + \text{other terms}. \]

The trouble comes from the fact that a given pair \((K, a)\) might be chosen by two or more different \( I \)'s. The pair \((K, a)\) is not uniquely determined by \( I \) and the proof shows that there are enough \((K, a)\) for all of the \( I \) when \( k = 3 \). When \( k = 4 \), there are not enough pairs \((K, a)\) and we need to add the weakly allowed generators.

More precisely,

Definition 3.5. Let \( I = (i_{k-1}, \ldots, i_0) \), define

\[ m(I) = \min\{s(h(I)), s(i_{k-2} + \cdots + i_0 - 1), \ldots, s(i_1 + i_0 - k + 2)\}. \]

The idea of the proof for \( k = 3 \) and \( 4 \) is to prove the following result. Let \( I \) be not allowed for \( k = 3 \) or \( 4 \) and \( \mathcal{A} \) the augmentation ideal in \( \mathcal{A} \). Then we find a pair \((K, a)\) such that \( Sq^2(Q^K) = Q^l + \sum Q^l \mod \text{Im} \mathcal{A} \), where either (1) \( m(J) < m(I) \) or (2) \( m(J) = m(I) \) and \( h(J) < h(I) \) or (3) \( J \) is weakly allowed. These arguments show that the set of allowed sequences forms a spanning set of \( Z/2 \otimes \mathcal{A} D_k \) for \( k = 3 \) or \( 4 \).

One step in the proof of linear independence is to prove the following proposition.

Proposition 3.6. Let \( k = 3 \) or \( 4 \). Suppose \( I \) is strongly allowed and

\[ Sq^2Q^K = Q^l + \text{other terms}. \]

Then \( a = 0 \) and \( K = I \).

We note that this is not true for a weakly allowed \( I \).

The following proposition allows us to focus only to the case of \( k = 4 \).

Proposition 3.7. Theorem 2.7 implies Theorem 2.6.
Proof. By means of Theorem 3.1 one easily sees that the map
\[ D_k/(Q_k,0) \to D_{k-1}, \quad Q_{k,s} \mapsto Q_{k-1,s-1} \]
is an isomorphism over the "halving" map
\[ \mathcal{A} \to \mathcal{A}, \quad Sq^{2i} \mapsto Sq^i, \quad Sq^{2i+1} \mapsto 0. \]
The proposition follows from the lists of the allowed sequences for \( k = 3 \) and \( k = 4 \) given in the previous section.

4. Determination of the allowed sequences

In this section we describe an inductive procedure for getting all the allowed sequences of length \( k \) for any \( k \).

Definition 4.1. Suppose \( a = a_0(a) + a_1(a)2 + \cdots + a_n(a)2^n \) is the dyadic expansion of \( a \) with \( a_n(a) = 1 \). Define \( \lfloor a : 2^t \rfloor \) to be the non-negative integer less than \( 2^t \) with \( a \equiv \lfloor a : 2^t \rfloor \mod 2^t \). So \( \lfloor a : 2^t \rfloor = a_0(a) + a_1(a)2 + \cdots + a_{t-1}(a)2^{t-1} \) if \( t \leq n \), and \( \lfloor a : 2^t \rfloor = a \) if \( t > n \).

Theorem 4.2. The only strongly allowed sequences \( I \)'s of length \( k \) are:
1. (1) either \( I = (z_{i-1}, \ldots, z_1, 0) \), where \( V = (z_{i-1}, \ldots, z_1) \) is strongly allowed at length \( k - 1 \).
2. (2) or \( I = (i_{i-1}, \ldots, i_1, i_0) \) satisfying the three conditions:
   (2a) \( z_0 = \lfloor -2^{i_0(0)} \rfloor > 0 \),
   (2b) \( I'' = (i_{i-2}, \ldots, i_1, i_0 - 1) \) is strongly allowed at length \( k - 1 \).
   (2c) \( i_{i-1} = 2' - \lfloor h' : 2' \rfloor - 1 \) for \( t > \nu(i_{i-2}) \), where \( h' = i_{i-2} + \cdots + i_0 \) and \( \nu(i_{i-2}) = \max\{n|2^n \leq i_{i-2}\} \).

To prove the theorem we need a couple of lemmas.

Lemma 4.3. Suppose \( u > v \) with \( u \geq 0 \) and \( t \) is an integer with \( 2^t - 1 \leq v \). Then
\[ u \equiv 2^t - 1 \mod 2^t, \quad 2^t \equiv -1 \mod 2^t. \]

Proof. Suppose the contrary that there exists \( t \) with \( 2^t - 1 \leq v \) and
\[ u \not\equiv 1 + 2 + \cdots + 2^{t-1} \mod 2^t. \]
That means there exists \( s < t \) such that \( \alpha_s(u) = 0 \). So \( s(u) \leq s < t \). It implies \( v \geq 2^t - 1 > 2^{s(u)} - 1 \). Hence \( u \not\equiv v \). This is a contradiction. The lemma is proved.

Lemma 4.4. If \( s(a + b) \geq s(a) \) then \( s(b - 1) \geq s(a) \).
Proof. See the proof of Lemma 5.1.

Lemma 4.5. If \( I = (i_{i-1}, \ldots, i_1, i_0) \) is strongly allowed, then
\[ i_0 = \lfloor k - 2 : 2^{s(i_1)} \rfloor. \]
Proof. Using Lemma 4.3 we have $i_1 \equiv -1 \mod 2^{s(i_1)}$ (obviously), and

$$i_2 \triangleright i_1 \implies i_2 \equiv -1 \mod 2^{s(i_1)}$$

$$\vdots \quad \vdots \quad \vdots$$

$$i_{k-1} \triangleright i_1 \implies i_{k-1} \equiv -1 \mod 2^{s(i_1)}$$

$$h \triangleright i_1 \implies h \equiv -1 \mod 2^{s(i)}$$

so $h = i_{k-1} + \cdots + i_1 + i_0 \equiv -1 \mod 2^{s(i_1)}$.

Thus $i_0 \equiv k - 1 - 1 \equiv k - 2 \mod 2^{s(i_1)}$. Combining this with the fact that $0 \leq i_0 < 2^{s(i_1)}$ (because of $i_1 \triangleright i_0$) we get $i_0 = \lfloor k - 2 : 2^{s(i_1)} \rfloor$. The lemma is proved.

Lemma 4.6. Suppose $I = (i_{k-1}, \ldots, i_0)$ is strongly allowed with $i_0 > 0$. Then so is $I' = (i_{k-2}, \ldots, i_0 - 1)$.

Proof. We need only to show that $(i_{k-2} + \cdots + i_0 - 1) \triangleright i_{k-2}$. Applying Lemma 4.4 with $a = i_{k-1}$, $b = i_{k-2} + \cdots + i_0$, we get $s(b - 1) \geq s(i_{k-1})$. This together with $i_{k-1} \triangleright i_{k-2}$ implies $b - 1 \triangleright i_{k-2}$. The lemma is proved.

Lemma 4.7. Suppose $h'$ and $s$ are non-negative integers.

(a) If there exists $i$ satisfying $i + h' \triangleright i \triangleright 2^s - 1$ then $h'$ is divisible by $2^s$.

(b) If $h'$ is divisible by $2^s$, then the only $i$ satisfying (a) are $i = 2^s - [h': 2^s] - 1$, $t > s$.

Proof. If $i$ satisfies (a), then $s(i) \geq s(2^s - 1) = s$. Since $i + h' \triangleright i$, by Lemma 4.4, it implies $s(h' - 1) \geq s(i) \geq s$. Because $h'$ is divisible by $2^{s(h'-1)}$, it is divisible by $2^s$.

Write $i = 2^s - 1 - u$ for some $t$ and $u$ with $t \geq s(i) \geq s$. We require $u < 2^{t-1}$ so that this expression is unique. Now we want to show $u = [h': 2^s]$. The fact $u < 2^{t-1}$ implies $i \geq 2^{t-1}$. Then we have $h' - u \equiv 0 \mod 2^s$; otherwise $i + h' = 2^s - 1 + (h' - u) \not\triangleright 2^{t-1}$, hence $i + h' \not\triangleright i$ (as $i \geq 2^{t-1}$). Combining $h' - u \equiv 0 \mod 2^s$ and $u < 2^s$ we get $u = [h': 2^s]$, by Definition 4.1.

It is easy to check that $i$ given by (b) satisfies (a). The lemma is proved.

Proof of Theorem 4.2. Obviously, if $i_0 = 0$, then $I$ is strongly allowed if and only if $I' = (i_{k-1}, \ldots, i_1)$ is.

If $i_0 > 0$, suppose $I$ is strongly allowed. By Lemmas 4.5 and 4.6, $I$ satisfies (2a) and (2b). Furthermore, from $h \triangleright i_{k-1} \triangleright i_{k-2}$ it implies $i_{k-1} + h' \triangleright i_{k-1} \triangleright 2^s - 1$, where $s = s(i_{k-2})$. By Lemma 4.7, $i_{k-1} = 2^s - [h': 2^s] - 1$, for $t \geq s$.

The condition $i_{k-1} = 2^s - [h': 2^s] - 1 \triangleright i_{k-2}$ obviously implies $t \geq s.$

On the other hand, we need to show that if $I$ satisfies (2a), (2b) and (2c), then $I$ is strongly allowed. Indeed, (2a) implies $i_1 \triangleright i_0$, (2b) implies $i_{k-2} \triangleright \cdots \triangleright i_1$. Finally, by Lemma 4.7, (2c) implies $h = i_{k-1} + h' \triangleright i_{k-1} \triangleright i_{k-2}$. The theorem is proved.

Remark 4.8. A strongly allowed sequence $I = (i_{k-1}, \ldots, i_0)$ is called main if $i_0 > 0$. Making use of the theorem, we list here all the main strongly allowed
sequences for $k \leq 4$:

For $k = 1$:
- $(2^s - 1)$, $s > 0$,

For $k = 2$:
- None

For $k = 3$:
- $(2r - 2s - 1, 2s - 1, 1)$, $r > s > 0$,

For $k = 4$:
- $(2r - 2s - 1, 2s - 1, 2s - 1, 2)$, $s > 1$,
- $(2r - 2s - 1, 2r - 2s - 1, 2s - 1, 2)$, $t > r > s > 1$.

Here is a generalization of Lemma 4.5.

**Remark 4.9.** Set $S(I) = (s_k, \ldots, s_1) = (s(h), s(i_{k-1}), \ldots, s(i_1))$ for $I = (i_{k-1}, \ldots, i_1, i_0)$. Then $S$ is injective on the set of all strongly allowed sequences.

Clearly, in $\text{Im } S$, $s_k \geq s_{k-1} \geq \cdots \geq s_1$, but not all such appear.

Now let us turn to investigate the weakly allowed sequences. We will show that they can be given as "deformations" of the strongly allowed sequences.

**Definition 4.10.** Let $J = (j_{k-1}, \ldots, j_0)$. We set $j_k = h(J)$.

(i) By a deformation of $J$ at $k$ we mean

$I = (j_{k-1} - 2^{s(j_k-1)} - 1, j_{k-2}, \ldots, j_0)$ with $2^{s(j_k-1)} - 1 \geq j_{k-2}$.

(ii) By a deformation of $J$ at $m < k$ we mean

$I = (j_{k-1}, \ldots, j_m - 2^s, j_{m+1} + 2^s, \ldots, j_0)$

with $s(j_m-1) \geq s + 2$, $2^s - 1 \geq j_{m-1}$, $j_m > j_{m-1} + 2^s$.

We call $j_m, j_{m-1}$ the initial and the terminal of the deformation at $m$ ($1 \leq m < k$).

**Proposition 4.11.** Every weakly allowed sequence $I$ can be obtained from a strongly allowed sequence $J$ by applying some deformations with the condition that the terminal of any deformation is not the initial of another deformation.

**Proof.** By definition of $\gg$, there do not exist $u, v, w$ such that $u \gg v \gg w$. Therefore, in a weakly allowed sequence $I$ a $\gg$ cannot stand just after another $\gg$.

Suppose $I$ is weakly allowed. Applying the reverse procedure of deformation at any place where a $\gg$ stands, we get a strongly allowed sequence $J$.

**Remark 4.12.** Making use of the proposition for $k \leq 4$ we conclude: For $k \leq 3$, there are no weakly allowed sequences. For $k = 4$, the only weakly allowed sequences are

$I = (2^r - 2^{s+1} - 2^s - 1, 2^s - 1, 2^s - 1, 2)$, $r > s + 2 > 3$.

In this sequence $h \gg i_3 \gg i_2 \gg i_1 \gg i_0$.

Combining this result with Remark 4.8 we get the lists of allowed sequences for $k \leq 4$, which are shown in Section 2.

### 5. The span of the allowed monomials

From now on, we always work at $k = 4$ unless otherwise specified. Let us begin with a generalization of Lemma 4.4.
Lemma 5.1. For any integers \( a, b \),
\[
\min\{s(a), s(b-1)\} = \min\{s(a+b), s(a)\} = \min\{s(a+b), s(b-1)\}.
\]

Proof. If either \( a = -1 \) or \( b = -1 \), then one can directly check the conclusion by using Definition 2.1 and the remark that \( s(\text{even number}) = 0 \).

Suppose \( a \neq -1, b \neq -1 \). Set \( s(a) = s, s(b-1) = r \). So \( a \equiv 2^s - 1 \mod 2^{s+1}, b-1 \equiv 2^{r-1} \mod 2^{r+1} \). Consider the three cases:

Case 1: \( r < s \). \( a + b \equiv 2^s + 2^r - 1 \mod 2^{s+1} \equiv 2^r - 1 \mod 2^{r+1} \). Hence \( s(a+b) = r \).

\[
\min\{s(a), s(b-1)\} = \min\{s(a+b), s(a)\} = \min\{s(a+b), s(b-1)\} = r.
\]

Case 2: \( r > s \). \( a + b \equiv 2^s + 2^r - 1 \mod 2^{s+1} \equiv 2^s - 1 \mod 2^{s+1} \). Hence \( s(a+b) = s \).

\[
\min\{s(a), s(b-1)\} = \min\{s(a+b), s(a)\} = \min\{s(a+b), s(b-1)\} = s.
\]

Case 3: \( r = s \). \( a + b \equiv 2^{r+1} - 1 \mod 2^{r+1} \). So \( s(a+b) > r \).

\[
\min\{s(a), s(b-1)\} = \min\{s(a+b), s(a)\} = \min\{s(a+b), s(b-1)\} = r = s.
\]

The lemma is proved.

Making use of the lemma we get
\[
\min\{s(i_2), s(i_1 + i_0 - 2)\} = \min\{s(i_2), s(i_2 + i_1 + i_0 - 1)\}
= \min\{s(i_2 + i_1 + i_0 - 1), s(i_1 + i_0 - 2)\},
\]
\[
\min\{s(i_3), s(i_2 + i_1 + i_0 - 1)\} = \min\{s(i_3), s(h)\}
= \min\{s(h), s(i_2 + i_1 + i_0 - 1)\}.
\]

Definition 5.2. If \( I \) is of length 4, by \( m(I) \) we mean the common value
\[
m(I) = \min\{s(i_3), s(i_2), s(i_1 + i_0 - 2)\}
= \min\{s(i_3), s(i_2 + i_1 + i_0 - 1), s(i_1 + i_0 - 2)\},
\]
\[
= \min\{s(h), s(i_3), s(i_1 + i_0 - 2)\}
= \min\{s(h), s(i_2 + i_1 + i_0 - 1), s(i_1 + i_0 - 2)\}.
\]

Main Lemma 5.3. If \( I \) is not allowed, then
(a) either \( Q' \in \text{Im} \mathcal{A} + \text{Span}\{Q'|m(J) < m(I)\} \) (the strong form of the main lemma),
(b) or \( Q' \in \text{Im} \mathcal{A} + \text{Span}\{Q'|\text{either } m(J) < m(I) \text{ or } m(J) = m(I) \text{ and } h(J) < h(I)\} \) (the weak form of the main lemma),
(c) or there exists a weakly allowed sequence \( I' \) with \( m(I) = m(I'), h(I) < h(I') \) and
\[
Q' + Q'' \in \text{Im} \mathcal{A} + \text{Span}\{Q'|m(J) < m(I)\}
\]
(the exceptional form of the main lemma).

The main lemma is proved by showing several lemmas. We state them here, but postpone their proofs until the next three sections.
Lemma 5.4. Suppose $I$ satisfies $s(h), s(i_1), s(i_2)$ all $\geq s(i_1), i_1 \not\equiv i_0$. Then the strong form of the main lemma holds for $I$.

Lemma 5.5. Suppose $I$ satisfies either $s(h), s(i_3)$ all $\geq s(i_2)$, $i_3 \not\equiv i_2 \not\equiv i_1 \not\equiv i_0$, or $s(h), s(i_3)$ all $\geq s(i_2) = s(i_1), i_2 \not\equiv i_1 \not\equiv i_0$. Then the strong form of the main lemma holds for $I$.

Lemma 5.6. Suppose $I$ satisfies $s(h) \geq s(i_3) \geq s(i_2), i_3 \not\equiv i_2 \not\equiv i_1 \not\equiv i_0, \text{ or } s(i_1 + i_0 - 2) = 0$. Then either the weak form or the exceptional form of the main lemma holds for $I$.

Lemma 5.7. Suppose $I$ satisfies $s(h) \geq s(i_3) < s(i_2)$. Then the strong form of the main lemma holds for $I$.

Lemma 5.8. Suppose $I$ satisfies $s(h) \geq s(i_3), h \not\equiv i_3 \not\equiv i_2 \not\equiv i_1 \not\equiv i_0, \text{ and } h \not\equiv i_3$. Then the weak form of the main lemma holds for $I$.

Lemma 5.9. Suppose $I$ satisfies $s(h) < s(i_3)$. Then the weak form of the main lemma holds for $I$.

It should be noted that we have to distinguish the strong form of the main lemma in Lemmas 5.4, 5.5 and 5.7 because it will essentially be used in the proofs of Lemmas 5.5, 5.7, which are parts of Lemma 5.6.

According to the results of Section 4, the main lemma is equivalent to these six lemmas, whose proofs will be given in the next three sections. Now we show how the main lemma leads to Theorem 2.7.

Proof of Theorem 2.7. The linear independence of the allowed monomials will be shown in Section 9.

Suppose that the main lemma is true. To prove the allowed monomials generate $\mathbb{Z}/2 \otimes_{\mathcal{A}} D_4$ we need only to show

Lemma 5.10. Let $k = 4$. Suppose $I$ is not allowed, with $m(I) = 0$. Then $Q^l \in \text{Im} \mathcal{A}$.

Proof. For abbreviation, from now on we sometimes let $I$ stand for $Q^l$. So, for instance, the statements like $I \in D_k$ or $Sq^l I \ldots$ make sense.

Recall $m(I) = \min\{s(h(I)), s(i_3), s(i_1 + i_0 - 2)\}$. We will consider the following three cases, with many subcases.

1. $s(h(I)) = 0$.
2. $s(h(I)) > 0$ and $s(i_3) = 0$.
3. $s(h(I)) > 0, s(i_3) > 0$, and $s(i_1 + i_0 - 2) = 0$.

In this proof, we will abbreviate the expression $i \equiv 0 \mod 2^l$ by $i \equiv 0(2^l)$. There are three easy ways to show that $I \in \text{Im} \mathcal{A}$, and we will call $I$ simple if it can be handled in one of the following three ways.

(a) $i_1 \equiv 0(2)$ and $i_0 \geq 1$, because $Sq^1(i_3, i_2, i_1 + 1, i_0 - 1) = I$, or
(b) $i_j \equiv 0(2)$ all $j$, because then $I$ is a square, or
(c) $i_2 \equiv 0(2)$ and $i_1 \equiv 1(2)$, because

$$Sq^2(i_3, i_2 + 1, i_1 - 1, i_0) = I + \left(\binom{i_1 - 1}{2}\right)(i_3, i_2 + 1, i_1 - 3, i_0 + 2)$$

and this last term is of type (a) above.

Case 1: $s(h(I)) = 0$. 

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(a) Assume \( i_3 > 0 \). Then \( S q^8(i_3 - 1, i_2, i_1, i_0) = I + \sum J \). Here \( h(J) < h(I) \) for all such \( J \) and we must show \( m(J) = 0 \). For most of these terms, \( s(j_3) = 0 \). The ones left are

\[
\binom{i_3 - 1}{2}(i_3 - 3, i_2 + 2, i_1, i_0), \quad \binom{i_2}{4}(i_3 - 1, i_2 - 4, i_1 + 4, i_0),
\]

and

\[
\binom{i_1}{8}(i_3 - 1, i_2, i_1 - 8, i_0 + 8)
\]

with \( i_3 \equiv 0(2) \), \( i_1 + i_0 \equiv 1(2) \), \( i_2 \equiv 1(2) \). In this case we get

\[
S q^8(i_3 + 1, i_2 - 1, i_1, i_0) = I + \binom{i_2 - 1}{2}(i_3 + 1, i_2 - 3, i_1 + 2, i_0)
\]

\[
+ \binom{i_1}{4}(i_3 + 1, i_2 - 1, i_1 - 4, i_0 + 4).
\]

Note that \( i_1 \equiv 1(2) \) or \( I \) is simple. Hence both the later terms are simple.

(b) Assume \( i_3 = 0 \) and \( i_2 > 0 \).

(i) If \( i_2 \equiv 1(2) \), then \( i_1 \equiv 1(2) \) or \( I \) is simple.

\[
S q^8(1, i_2 - 1, i_1, i_0) = I + \binom{i_2 - 1}{2}(1, i_2 - 3, i_1 + 2, i_0)
\]

\[
+ (i_2 - 1)\binom{i_1}{2}(1, i_2 - 2, i_1 - 1, i_0 + 2)
\]

\[
+ \binom{i_1}{4}(1, i_2 - 1, i_1 - 4, i_0 + 4)
\]

works as the other terms are 0 or simple.

(ii) If \( i_2 \equiv 0(2) \), then \( i_1 + i_0 \equiv 0(2) \). If \( i_1 \equiv 0(2) \), then \( I \) is simple. If \( i_1 \equiv 1(2) \), then \( I \) is also simple.

(c) Assume \( i_3 = 0 \) and \( i_2 = 0 \). Then \( i_1 \equiv 1(2) \) or \( I \) is simple. Then \( I \) is also simple.

Case 2: \( s(h(I)) > 0, s(i_3) = 0 \).

(a) Assume \( i_2 \equiv 0(2) \). If \( i_1 \equiv 0(2) \), then \( i_0 > 0 \) and \( I \) is simple. If \( i_1 > 0 \), then \( I \) is also simple.

(b) Assume \( i_2 \equiv 1(2) \).

\[
S q^4(i_3 + 1, i_2 - 1, i_1, i_0) = I + \binom{i_2 - 1}{2}(i_3 + 1, i_2 - 3, i_1 + 2, i_0)
\]

and

\[
S q^2(i_3 + 1, i_2 - 2, i_1 + 1, i_0)
\]

\[
= (i_3 + 1, i_2 - 3, i_1 + 2, i_0) + \binom{i_1 + 1}{2}(i_3 + 1, i_2 - 2, i_1 - 1, i_0 + 2).
\]

If \( i_1 \equiv 1(2) \), then \( i_0 + 2 > 0 \) and it is simple. If \( i_1 \equiv 0(2) \), then \( i_0 \equiv 1(2) \). \( I \) is simple.

Case 3: \( s(h(I)) > 0, s(i_3) > 0, s(i_1 + i_0 - 2) = 0 \).
(a) Assume $i_0 > 0$. If $i_1 \equiv 0(2)$, then $I$ is simple. If $i_1 \equiv 1(2)$, then $i_2 \equiv 0(2)$ and $i_1 > 0$, so $I$ is simple.

(b) Assume $i_0 = 0$, $i_1 > 0$. Then $i_2 \equiv 0(2)$.

$$Sq^2(i_3, i_2 + 1, i_1 - 1, 0) = I + \binom{i_1 - 1}{2}(i_3, i_2 + 1, i_1 - 3, 2)$$

which handles all cases but $i_1 \equiv 0(4)$.

(i) $i_2 \not\equiv 4, 6(8), i_1 \equiv 0(4)$

$$Sq^8(i_3, i_2 + 4, i_1 - 4, 0) = I + \text{simple terms.}$$

(ii) $i_2 \equiv 4(8), i_1 \equiv 0(4), i_1 \geq 8$.

$$Sq^4(i_3, i_2 + 3, i_1 - 5, 2) = (i_3, i_2 + 1, i_1 - 3, 2) + \binom{i_1 - 5}{4}(i_3, i_2 + 3, i_1 - 9, 6) + (i_3 - 1, i_2 + 4, i_1 - 5, 2) + \text{simple terms}$$

and the two latter terms are simple.

(iii) $i_2 \equiv 4(8), i_1 = 4$.

$$Sq^4(i_3, i_2 + 2, 2, 0) = I + (i_3 - 1, i_2 + 3, 2, 0)$$

which is in Case 2.

(iv) $i_2 \equiv 6(8), i_3 \equiv 3(4)$. Then $h(I) \equiv 1(4)$.

$$Sq^8(i_3 - 2, i_2, i_1, 0) = I + \binom{i_3 - 2}{4}(i_3 - 6, i_2 + 4, i_1, 0) + \binom{i_2}{8}(i_3 - 2, i_2 - 8, i_1 + 8, 0) + \binom{i_3 - 2}{2}(i_3 - 2, i_2 - 1, i_1 + 2, 0) + \text{simple terms.}$$

All $h(J) < h(I)$ and only need check $m(J) = 0$. In all cases, $s(j_1 + j_0 - 2) = 0$.

(v) $i_2 \equiv 6(8), i_3 \equiv 1(4), i_1 \equiv 4(8)$.

$$Sq^8(i_3 + 2, i_2 - 2, i_1, 0) = I + (i_3 + 2, i_2 - 6, i_1 + 4, 0) + \text{simple terms.}$$

$$Sq^2(i_3 + 2, i_2 - 5, i_1 + 3, 0) = (i_3 + 2, i_2 - 6, i_1 + 4, 0) + (i_3 + 2, i_2 - 5, i_1 + 1, 2).$$

$$Sq^4(i_3 + 2, i_2 - 3, i_1 - 1, 2) = (i_3 + 2, i_2 - 5, i_1 + 1, 2) + (i_3 + 1, i_2 - 2, i_1 - 1, 2) + \text{simple terms,}$$

and $(i_3 + 1, i_2 - 2, i_1 - 1, 2)$ is simple.
(vi) $i_2 \equiv 6(8), i_3 = 5(8), i_1 \equiv 0(8), i_1 > 0$.

$Sq^8(i_3 + 2, i_2 - 2, i_1, 0) = I + (i_3 + 2, i_2 - 6, i_1 + 4, 0)$

$+ \text{simple terms.}$

$Sq^4(i_3 + 2, i_2 - 4, i_1 + 2, 0) = (i_3 + 2, i_2 - 6, i_1 + 4, 0)$

$+ (i_3 + 1, i_2 - 3, i_1 + 2, 0).$

$Sq^{2^{\lfloor h(I) \rfloor}}(i_3 + 1 - 2^r, i_2 - 3, i_1 + 2, 0) = (i_3 + 1, i_2 - 3, i_1 + 2, 0) + \sum J$

with $h(J) < h(I)$ and all have $j_1 + j_0 - 2 \equiv 0(2)$. Here we denote $s = s(h(I))$

and use $i_3 \equiv 5(8)$ to show $i_3 + 1 \geq 2^{\lfloor h(K) \rfloor}$.

(vii) $i_2 \equiv 6(8), i_3 = 1(8), i_1 \equiv 0(8), i_1 > 0$.

$Sq^{16}(i_3 + 4, i_2 - 4, i_1, 0) = I + (i_3 + 6, i_2 - 4, i_1, 0)$

$+ \left(\frac{i_2 - 4}{8}\right)(i_3 + 4, i_2 - 12, i_1 + 8, 0).$

$Sq^{32}(i_3 + 2, i_2 - 4, i_1, 0) = (i_3 + 6, i_2 - 4, i_1, 0) + \sum J$

with $h(J) < h(I)$ and all $m(J) = 0$.

$Sq^6(i_3 + 4, i_2 - 8, i_1 + 4, 0) = (i_3 + 4, i_2 - 12, i_1 + 8, 0)$

$+ (i_3 + 3, i_2 - 9, i_1 + 6, 0) + \text{simple terms.}$

$Sq^4(i_3 + 3, i_2 - 7, i_1 + 4, 0) = (i_3 + 3, i_2 - 9, i_1 + 6, 0)$

$+ \text{simple terms.}$

(c) Assume $i_1 = i_0 = 0$, $h(I) \not> i_3$.

$Sq^{2^{\lfloor h(I) \rfloor}}(i_3 - 2^{\lfloor h(I) \rfloor}, i_2, 0, 0) = I + \sum J.$

All $h(J) < h(I)$ and since $\dim I \equiv 0(4)$, we have that $j_1 \equiv j_0 \equiv 0(2)$ and $m(J) = 0$.

(d) Assume $i_1 = i_0 = 0$, $h(I) \not> i_3$. Then $i_3 \not> i_2$ as $I$ is not allowed. Let $s = s(i_3)$. Hence

$Sq^{2^{\lfloor h(I) \rfloor}}(i_3 + 2^r, i_2 - 2^r, 0, 0) = I + (i_3 + 2^r + 2^{r-1}, i_2 - 2^r, 0, 0)$

$+ \left(\frac{i_2 - 2^r}{2^{r+1}}\right)(i_3 + 2^r, i_2 - 2^r - 2^{r+1}, 2^{r+1}, 0).$

Denote the two extra terms by $J$ and $K$ respectively. Obviously, $m(K) = 0$, $s(h(K)) > 0$, $s(k_3) > 0$, $k_1 > 0$. So $K$ is handled by Case 3(b). Note remarkably that while handling $K$, one still has the inductive procedure for lowering $h(K)$ in Case 1(a) or Case 2(b). In all possibilities, every extra term $L$ has $l_1 + l_0 \geq k_1 + k_0 > 0$ by Corollary 3.2. Therefore, no extra term needs to be handled by Case 2(c) or (d). So the inductive procedure works.
Next, note that \( s(h(J)) = s(h(I) + 2^{s-1}) = s - 1 = s(j_3) \) and \( \alpha_t(j_2) = \alpha_t(i_2 - 2^s) = 0 \) for \( t < s \) as \( h > i_3 \). If \( j_2 \geq 2^{s-1} \), we handle \( J \) by downward induction on \( s \) using \( Sq^{2^t-1}(j_3 + 2^{s-1}, j_2 - 2^{s-1}, 0, 0) = 0 \). Otherwise, if \( j_2 < 2^{s-1} \), combining this with the fact that \( \alpha_t(j_2) = 0 \) for \( t < s \) we get \( j_2 = 0 \). Then we cannot use Case 3(c) to handle \( J \) because of \( h(J) > h(I) \). Let \( t \) be a number with \( \alpha_t(j_3) = 0, \alpha_{t+1}(j_3) = 1 \). Such a \( t \) exists as \( h(J) \not\succ j_3 = h(J) \). Then \( Sq^{2^t}(j_3 - 2^t, 0, 0, 0) = J \).

In order to start the induction, we need the following case.

Case 4: \( h(I) = 1 \). \((1, 0, 0, 0) \) is allowed, \( Sq^4(1, 0, 0, 0) = (0, 1, 0, 0), Sq^6(1, 0, 0, 0) = (0, 0, 1, 0), \) and \( Sq^7(1, 0, 0, 0) = (0, 0, 0, 1) \).

6. Proofs of Lemma 5.4 and Lemma 5.5

We need the following terminology:

**Definition 6.1.** Let \( s \) be a natural number. \( Q^I \) (or \( I \)) is \( s \)-killed if

\[
Q^I \in \text{Im} \mathcal{A} + \text{Span}\{Q^J | m(J) < s\}.
\]

Therefore, the strong form of the main lemma holds for \( I \) if and only if \( I \) is \( m(I) \)-killed.

**Proof of Lemma 5.4.** Set \( s = s(i_1) \). By Lemma 5.1,

\[
s(h), s(i_3) \text{ all } \geq s \Rightarrow s(i_2 + i_1 + i_0 - 1) \geq s \text{ also } s(i_2) \geq s \Rightarrow s(i_1 + i_0 - 2) \geq s.
\]

As a consequence: \( m(I) = \min\{s(i_3), s(i_2), s(i_1 + i_0 - 2)\} \geq s \).

We kill \( Q^I \) by

\[
Sq^{2^t}(Q_3 Q_2 Q_1 Q_0) = (\frac{i_1 + 2^s}{2^s}) Q^I + \sum_J Q^J.
\]

For any \( J \) in the sum we have

(a) \( |h(J) - h(I)| \leq \frac{2^s}{\dim Q_3} = 2^{s-3} < 2^s \),

(b) \( |j_3 - i_3| \leq \frac{2^s}{\dim Q_2 - \dim Q_3} = 2^{s-2} < 2^s \),

(c) \( |j_2 - i_2| \leq \frac{2^s}{\dim Q_1 - \dim Q_2} = 2^{s-1} < 2^s \).

If \( h(J) \not= h(I) \), combine (a) with the fact \( s(h(I)) \geq s \) and we get \( s(h(J)) < s \). So \( m(J) \leq s(h(J)) < s \leq m(I) \). Similarly, if either \( j_3 \not= i_3 \) or \( j_2 \not= i_2 \), then \( m(J) < s \leq m(I) \).

Suppose now \( h(J) = h(I), j_3 = i_3, j_2 = i_2 \). The dimensional information shows

(d) \( 14(j_1 - (i_1 + 2^s)) + 15(j_0 - (i_0 - 2^s)) = 2^s \).

As \( j_0 - (i_0 - 2^s) \geq 0 \), so \( 14(j_1 - (i_1 + 2^s)) + j_0 - (i_0 - 2^s) \leq 2^s \), hence \( 0 \leq [(j_1 + j_0 - 2) - (i_1 + i_0 - 2)] \leq \frac{14}{15} < 2^s \).

If \( (j_1 + j_0 - 2) - (i_1 + i_0 - 2) = 0 \), then \( (j_1 - i_1) = -(j_0 - i_0) = x \) (say).

Now, (d) becomes \( 14(x - 2^s) + 15(-x + 2^s) = 2^s \). So \( x = 0, j_1 = i_1, j_0 = i_0 \).

The corresponding term is \( Q^I \) but not a term of the sum.
Otherwise $0 < (j_1 + j_0 - 2) - (i_1 + i_0 - 2) < 2^s$. This together with $s(i_1 + i_0 - 2) \geq s$ implies $m(J) \leq s(j_1 + j_0 - 2) < s \leq m(I)$.

The lemma follows.

In order to prove Lemma 5.5, we need the following two technical lemmas.

Let us first state a strong version of Lemma 5.4, which can be proved by the same argument.

**Lemma 6.2.** Let $I = (i_3, i_2, i_1, i_0)$. Suppose there exists $s$ such that

$$\begin{aligned} &\left(\frac{i_1 + 2^s}{2^s}\right) \equiv 1 \text{ mod } 2, \quad i_0 \geq 2^s, \quad m(I) \geq s. 
\end{aligned}$$

Then $I$ is $s$-killed.

Next we show:

**Lemma 6.3.** Suppose $I$ satisfies $s(h), s(i_3) \geq s(i_2), i_2 \not\equiv i_1$. Then, for $s = s(i_2)$:

$$Q^I + \left(\frac{i_1 - 2^s}{2^{s+1}}\right)Q^{(i_3, i_2 + 2^s, i_1 - 2^s - 2^{s+1}, i_0 + 2^{s+1})} \in \text{Im}A + \text{Span}\{Q^J | m(J) < m(I)\}.$$  

**Proof.** We kill $Q^I$ by the usual way

$$S_q^{2^{s+1}}(Q_j^{i_3}Q_{2}^{i_2+2^s}Q_1^{i_1-2^s}Q_0^{i_0}) = Q^I + \sum_J Q^J.$$  

Similarly as in the proof of Lemma 5.4, for any $J$ in the sum, we have

$$|h(J) - h(I)| \leq 2^{s+1}/2^3 = 2^{s-2} < 2^s, |j_3 - i_3| \leq 2^{s+1}/2^2 < 2^s.$$  

Thus, if either $h(J) \neq h(I)$ or $j_3 \neq i_3$ then $m(J) < s \leq m(I)$. We need only to consider the case where $h(J) = h(I), j_3 = i_3$. Note that

$$j_1 + j_0 = (i_1 - 2^s) + i_0 + t \quad \text{with } 0 \leq t < 2^s.$$  

(The only term with $t = 2^s$ is $Q^I$ according to the dimensional information.)

If $0 < t < 2^s$: Using Lemma 5.1, one has

$$s(h), s(i_3) \geq s(i_2) \Rightarrow s(i_1 + i_0 - 2) \geq s = s(i_2).$$  

meanwhile $0 < t < 2^s \Rightarrow s(j_1 + j_0 - 2) = s(i_1 - 2^s + i_0 + t - 2) < s.$  

$$\Rightarrow m(J) < s \leq m(I).$$

If $t = 0$: One gets

$$\begin{aligned} &j_1 + j_0 = (i_1 - 2^s) + i_0, \\
&14j_1 + 15j_0 = 14(i_1 - 2^s) + 15 i_0 + 2^{s+1} \quad \text{(dim. information).} 
\end{aligned}$$  

So $j_0 = i_0 + 2^{s+1}, j_1 = i_1 - 2^s - 2^{s+1}, j_2 = h - (j_3 + j_1 + j_0) = i_0 + 2^s$. The coefficient of the resulting term $Q^J$ in the sum is $\binom{2^s}{2^{s+1}}$. The lemma is proved.

**Proof of Lemma 5.5.** Using actually the argument in the proof of Lemma 4.5, one observes that if $s(h), s(i_3), s(i_2) \geq s(i_1), i_1 \not\equiv i_0$, then either $i_0 = 0$ or $i_0 = 2$. Combining this with Remark 4.12, we need only to consider the four cases:
Case A: \( s(h), s(i_3) \) all \( \geq s(i_2) \), \( i_2 \not\equiv i_1 \) and at least one from the following three equalities is broken: \( s(i_2) = s(i_1), \ s_2(i_1) = s(i_1) + 1, \ i_0 = 2 \).

Case B: \( s(h), s(i_3) \) all \( \geq s(i_2) \), \( i_2 \not\equiv i_1 \) \( \Rightarrow \) at least one from the following three equalities is broken: \( s(i_2) = s(i_1), \ s_2(i_1) = s(i_1) + 1, \ i_0 = 2, \ i_2 \not\equiv i_1 \) (or equivalently \( i_2 + 2s(i_2) \not\equiv i_1 \)).

Case C: \( s(h), s(i_3) \) all \( \geq s(i_2) \), \( i_3 \not\equiv i_2 \) \( \Rightarrow \) \( i_1 \) \( \not\equiv \) \( i_0 = 2 \).

Case D: \( s(h), s(i_3) \) all \( \geq s(i_2) \), \( h \not\equiv i_3 \) \( \not\equiv i_2 \) \( \not\equiv i_1 \) \( \not\equiv i_0 = 2 \).

Proof of Lemma 5.5 in Case A. Using Lemma 6.3 it suffices to study the case of \( (i_1 - 2^s) \equiv 1 \mod 2 \), where \( s = s(i_2) \). First, we note that:

\[
s(i_1 - 2^s - 2^{s+1}) = s + 1 \iff i_1 - 2^s - 2^{s+1} \equiv 2^{s+1} - 1 \mod 2^{s+1}
\]

\[
\iff i_1 \equiv 2^s - 1 \mod 2^{s+2} \iff \begin{cases}
s(i_1) = s \\
s_2(i_1) = s + 1.
\end{cases}
\]

(1) If \( s(i_1 - 2^s - 2^{s+1}) \neq s + 1 \), since \( (i_1 - 2^s) \equiv 1 \), so \( r = s(i_1 - 2^s - 2^{s+1}) < s + 1 \). We have \( s(h) \geq s(i_2) = s \geq r \). Since \( s(i_3), s(h) \) all \( \geq s(i_2) \), so \( s(i_1 + i_0 - 2) \geq s(i_2) = s \). Then

\[
s(i_1 - 2^s - 2^{s+1} + i_0 + 2^{s+1} - 2) = s(i_1 + i_0 - 2 - 2^s) \geq s \geq r.
\]

Now we can apply Lemma 6.2 for

\[
Q^T = Q^{(i_3, i_1 + 2^s, i_1 - 2^s - 2^{s+1}, i_0 + 2^{s+1})}
\]

with \( r \) playing the role of \( s \). So we get

\[
Q^T \in \text{Im} \tilde{\mathcal{A}} + \text{Span}\{Q^I|m(J) < r\} \subset \text{Im} \tilde{\mathcal{A}} + \text{Span}\{Q^I|m(J) < m(I) = s\}.
\]

Combining this with Lemma 6.3, it implies \( Q' \) is \( s \)-killed.

(2) Suppose \( s(i_1 - 2^s - 2^{s+1}) = s + 1 \). By the hypothesis of Lemma 5.5 Case A: \( i_0 \neq 2 \).

By Lemma 4.3:

\[
s(i_1) = s \Rightarrow 2^s \mid (i_1 + 1)
\]

\[
\Rightarrow (i_1 + i_0 - 2) \geq 2^s \mid (i_1 + i_0 - 1) \Rightarrow 2^s \mid (i_0 - 2).
\]

(a) If \( i_0 - 2 > 0 \), then \( i_0 > 2^s \). On the other hand: \( s(h), s(i_3) \), \( s(i_1 + i_0 - 2) \) all \( \geq s \). We can apply Lemma 6.2 and \( I \) is \( s \)-killed.

(b) If \( i_0 - 2 < 0 \), then either \( i_0 = 1 \) (so \( s = 0 \) as \( 2^s - 1 = i_0 - 2 \)), or \( i_0 = 0 \) (and \( s = 0 \) or \( 1 \) because of \( 2^s - 2 = i_0 - 2 \)).

The case \( s = m(I) = 0 \) has been considered in the proof of Theorem 2.7 at the end of Section 5. Now it suffices to study the case \( s = 1, i_0 = 0 \).

By hypothesis \( i_0 = 0, \ i_1 = 1 + c, \ i_2 = 1 + b, \ i_3 = 1 + a \). Here \( c \equiv 0 \mod 8, c > 0, b \equiv 0 \mod 4 \) and \( a \equiv 0 \mod 2 \). The following computations are modulo \( \text{Span}\{Q^I|m(J) = 0\} \).

- If \( a \equiv 0 \mod 4 \),

\[
S_q^4(1 + a, 3 + b, 1 + c - 2, 0) = (1 + a, 1 + b, 1 + c, 0) + (1 + a, 3 + b, 1 + c - 6, 4),
\]

\[
S_q^8(3 + a, 1 + b, 1 + c - 6, 4) = (1 + a, 3 + b, 1 + c - 6, 4) + \sum J.
\]

Here every \( J \) in the sum satisfies

\[
j_3 \equiv 3 \mod 4, j_2 \equiv 1 \mod 4, j_1 \equiv 3 \mod 4, j_0 \geq 4.
\]
Such a $J$ satisfies the hypothesis of Case 1 above, hence it is 1–killed.

- If $a \equiv 2 \mod 4$,
  \[ Sq^{16}(1 + a - 2 , 1 + b , 1 + c , 0) = (1 + a , 1 + b , 1 + c , 0) + \sum J. \]

Here any $J$ satisfies either the hypothesis of Lemma 5.4 with $m(J) = 1$ or $J = (1 + a', 1 + b', 1 + c', 0)$ with $a' \equiv b' \equiv 0 \mod 4$, $c' \equiv 0 \mod 8$, $c' > 0$. Thus $J$ is 1–killed by means of either Lemma 5.4 or the previous item.

**Proof of Lemma 5.5 in Cases B, C, D supposing that Lemma 5.7 is true.** (We will prove Lemma 5.7 by a direct computation in the next section.)

To this end, the use of $m(I)$ is not enough. We need one more control variable defined as follows.

**Definition 6.4.** For $I = (i_3, i_2, i_1, i_0)$, we define
\[ m'(I) = \min\{s(i_3), s'(i_2)\}. \]

Set
\[ S(s) = \{I|s(h), s(i_3) \text{ all } \geq s(i_2) = s(i_1) = s > 1, i_0 = 2\}. \]

Note that $m(I) = s$ for any $I \in S(s)$. Making use of $m'(I)$, we get

**Lemma 6.5.** If $I \in S(s)$, and $I$ is in Cases B, C or D, then
\[ I \in \text{Im.}\mathcal{S} + \text{Span}\{J|m(J) < m(I) = s\} \]
\[ \quad + \text{Span}\{J \in S(s)|m'(J) < m'(I), \text{ I in Cases B, C or D}\} \]
\[ \quad + \text{Span}\{J \text{ satisfying Lemma 5.7's hypothesis and } s(j_3) = s\}. \]

**Corollary 6.6.** If Lemma 5.7 is true, then so is Lemma 5.5 in Cases B, C, D.

To prove the above lemma we need the technical lemma:

**Lemma 6.7.** If $H \in S(s)$, then for any $i$ we have
\[ Sq^{2^i}H = \sum J \mod \text{Span}\{L|m(L) < s\}, \]

where
(a) either $J$ satisfies the hypothesis of Lemma 5.7 with $s(j_3) = s$,
(b) or $J \in S(s)$.

**Proof.** This proof looks very much like the beginning part of that of Lemma 7.7.

It is easy to see that $m(H) = s$.

\[ Sq^{2^i}H = \sum J = \sum (Sq^{a_0}Q_3)^{a_0}\ldots(Sq^{a_n}Q_3)^{2^na_n} \]
\[ \cdot(Sq^{b_0}Q_2)^{b_0}\ldots(Sq^{b_s}Q_2)^{2^sb_s} \]
\[ \cdot(Sq^{m_0}Q_1)^{m_0}\ldots(Sq^{m_i}Q_1)^{2^ic_i} \]
\[ \cdot(Sq^{n_0}Q_0)^{2}. \]

Here $a_i, b_i, c_i$ are the coefficients of $2^i$ in the dyadic expansions of $h_3, h_2, h_1$ respectively. Convention: $k_i = 0$ if $a_i = 0$, $l_i = 0$ if $b_i = 0$, $m_i = 0$ if $c_i = 0$.

Consider a $J$ in the sum. We need only to study the case $m(J) \geq s$. That means $s(j_3), s(j_2), s(j_1 + j_0 - 2)$ all $\geq s$. 
Case 1: \( s(j_3) \geq s, s(j_2) > s \). These imply \( k_s = 4 \) and \( s(j_3) = s \). Hence \( m(J) \leq s \leq m(J) \), so \( m(J) = s \leq s(h') \). Here \( h' = h(J) \). Therefore \( J \) satisfies the hypothesis of Lemma 5.7 with \( s(j_3) = s \).

Case 2: \( s(j_3) \geq s, s(j_2) = s, s(j_1) > s \). Then \( m(J) \leq s(j_2) = s \). So \( m(J) = s \). Then \( s(h') \geq m(J) = s \). Since \( s(j_2) < s(j_1) \) so \( j_2 \notin j_1 \), we can apply Case A of Lemma 5.5 to show \( J \) is \( s \)-killed.

Case 3: \( s(j_3) \geq s, s(j_2) = s, s(j_1) = t < s \). Similarly as before \( m(J) = s \). Hence \( s(j_1 + j_0 - 2) \geq s \). This together with \( s(j_1) = t < s \) implies \( j_0 > 2t \). Therefore, \( j_1 \notin j_0 \). We also have \( s(h') \geq m(J) = s > t \). Hence, by Lemma 5.4, \( J \) is \( s \)-killed.

Case 4: \( s(j_3) \geq s, s(j_2) = s(j_1) = s \). We also have \( m(J) = s \). Then \( s(h') \geq s \).

If \( j_1 \notin j_0 \), then we can apply Lemma 5.4 to show \( J \) is \( s \)-killed.

Otherwise, if \( j_1 \supset j_0 \), this together with \( s(h'), s(j_3) \) all \( \geq s(j_2) = s(j_1) = s \) implies \( j_0 = 2 \) (by an argument similar to that in the proof of Lemma 4.5). That means \( J \in S(s) \). The lemma is proved.

Proof of Lemma 6.5. In the three parts of this proof, we find some \( i \) and \( H \in S(s) \) such that \( Sq^{2^i}H = I + \sum J \). Using Lemma 6.7, it suffices to consider \( J \in S(s) \). We will show that either such a \( J \) satisfies \( m'(J) < m'(I) \) and \( j_2 \notin j_1 \) or it is a linear combination of some sequences with these properties modulo \( \mathfrak{A} \). Note that if \( J \in S(s) \) with \( j_2 \notin j_1 \), then either it is \( s \)-killed by Lemma 5.4 or it is in Cases A, B, C or D of Lemma 5.5. If \( J \) is in Case A, it is also \( s \)-killed as showed above. Then the lemma follows.

(1) \( J \) is in Case B with \( s(i_3) \geq s'(i_2) \): Set \( s' = s'(i_2) \). Let \( a_i, b_i, c_i \) be the coefficients of \( 2^i \) in the dyadic expansions of \( i_3, i_2 + 2^i, i_1 - 2^i \) respectively. Then, we have

\[
Sq^{2^i+1}(i_3, i_2 + 2^i, i_1 - 2^i, 2) = I + \sum J = \sum (Sq^{k_0}Q_3)^{a_0} \cdots (Sq^{k_s}Q_3)^{2^u a_u} \\
\cdot (Sq^{k_0}Q_2)^{b_0} \cdots (Sq^{k_s}Q_2)^{2^u b_u} \\
\cdot (Sq^{m_0}Q_1)^{c_0} \cdots (Sq^{m_s}Q_1)^{2^u c_u} \\
\cdot (Sq^nQ_0)^2.
\]

Take a \( J \) in the sum. By Lemma 6.7 we can suppose \( J \in S(s) \). Compare \( j_2 \) and \( i_2 + 2^i \).

(a) If \( |j_2 - (i_2 + 2^i)| \equiv 2^i \mod 2^i+1 \), then \( j_2 = i_2 \) and \( J = I \) by dimensional information.

(b) If \( |j_2 - (i_2 + 2^i)| \equiv 0 \mod 2^i+1 \), then \( j_2 = i_2 + 2^i \) by dimensional information.

- If \( j_3 = i_3 \), since \( J \in S(s) \), so \( (j_3, j_2, j_0) = (i_3, i_2 + 2^i, i_0) \). It is contradictory to the fact that \( 2^u+1 = \dim Sq^{2^u+1} \) is not divisible by \( \dim Q_1 = 14 \).

- If \( j_3 \neq i_3 \), also by dimensional information we have

\[
j_3 - i_3 \equiv 2^u \mod 2^{u+1}, \quad u < s'.
\]

From \( s(i_3) \geq s'(i_2) = s' \), it implies \( s(j_3) = u < s' \).
That means $m'(J) < m'(I)$. If $j_2 \not\in j_1$ then $J$ has the properties which we want.

Otherwise, $j_2 \triangleright j_1$ hence $j_1 = 2^s - 1$ (as $J \in S(s)$). We have

\[ S_{q_2(i_3) + 2^{s(i_3)}, i_2 + 2^{s' - 2^{s(i_3)}}, 2^s - 1, 2} = J + \sum K. \]

By Lemma 6.7, it suffices to consider the case $K \in S(s)$. If $k_2 \not\in k_1$ then $K$ has the properties we want. Otherwise, $k_2 \triangleright k_1$ hence $k_1 = 2^s - 1$, $k_0 = 2$ (as $K \in S(s)$). Then, by Corollary 3.2, $k_2 \geq i_2 + 2^{s' - 2^{s(i_3)}}$. On the other hand $s(k_3) < s(j_3)$, otherwise $K = J$ by dimensional information. So $k_3 \not\in k_2$. We can use an equality that is similar to (*) to handle $K$ by induction on $s(k_3)$.

Emphasize that in the inductive procedure, we always have

\[ m(K) = m(I) = s, \quad m'(K) < s(k_3) < s' = m'(I). \]

(c) If $|j_2 - (i_3 + 2^{s'})| = 2^{s' + 1}(t < s')$ ($t > s$ since $s(j_2) = s(i_2) = s$), then $\alpha_t(j_2) = 0$ because of $\alpha_t(i_2) = 1$. Hence

\[ s'(j_2) \leq t < s' = \min\{s(i_3), s'(i_2)\}. \]

Let us consider all the possibilities to get $\alpha_t(j_2) = 0$.

- If $Q_2^{i_3}$ moves to $Q_1^{i_3}$ ($k = 3$), then it contradicts the fact that $J \in S(s)$. This possibility does not occur.
- If $Q_2^{i_3}$ moves to $Q_1^{i_3}$ ($k = 2$), then $j_2 \not\in j_1$.
- If it raises the power of $Q_2$ from 2 to 2' then $(i_3 + i_1 + i_0) = 1$. Since $i_0 = 2$, $i_2 = 2^{s' - 2^s - 1}$ mod $2^{s' + 1}$ and $s(i_1) = s$, so $\alpha_t(i_1) = 1$. Then $j_2 \not\in j_1$.
- If $Q_3^{i_3}$ moves to either $Q_2^{i_3}Q_0^{i_3}$ or $Q_2^{i_3}Q_1^{i_3}$ ($l = 11$ or 10), then the situation looks like that in the first two cases. The first possibility does not happen and in the second one, $j_2 \not\in j_1$.
- If $Q_3^{i_3}$ moves to $Q_3^{i_3}$ ($k = 4$), then $s(j_3) \leq t < s'$, $j_2 > 2^s$. If $j_2 \not\in j_1$, then $J$ has the properties we want. Otherwise, if $j_2 \triangleright j_1$ then $j_1 = 2^s - 1$ (as $J \in S(s)$). We handle $J$ by

\[ S_{q_2(i_3) + 2^{s(i_3)}, j_2 - 2^{s(i_3)}}, 2^s - 1, 2} = J + \sum K \]

and by using a similar argument as in the last item of (b).

(2) $J$ is in either Case B or Case C with $s(i_3) < s'(i_2)$: Put $r = s(i_3) = m'(I)$.

(a) Suppose $r = s = s(i_2)$. Using $S_{q_2(i_3), i_2 + 2^s, i_1 - 2^s, 2}$ we get as in the proof of Lemma 6.3:

\[ I = I' := (i_3, i_2 + 2^s, i_1 - 2^s - 2^{s+1}, 2 + 2^{s+1}) \]

modulo $\text{Im} \mathcal{A} + \text{Span}\{J|m(J) < s\}$. It is easy to check that $I'$ satisfies the hypothesis of Lemma 5.7 with $s(i'_3) = s$.

(b) Suppose $r > s = s(i_2)$.

\[ S_{q_2(i_3) + 2^r, i_2 - 2^r, i_1, 2} = I + \sum J. \]

$(i_3 + 2^r, i_2 - 2^r, i_1, 2) \in S(s)$. We apply Lemma 6.7 to it. Take a $J \neq I$ in the sum. We can assume $J \in S(s)$. It suffices to prove that $m'(J) < m'(I)$ and $j_2 \not\in j_1$. By dimensional information: $J \neq I \Leftrightarrow 0 \leq |j_3 - (i_3 + 2^r)| < 2^r$. 

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• If \( j_3 \neq i_3 + 2^r \): \(|j_3 - (i_3 + 2^r)| = 2^u \mod 2^{u+1} \) so \( s(j_3) = u < r \). Then \( m'(J) < m'(I) \).

By Corollary 3.2, \( j_1 \geq i_1 \). Meanwhile, \( s(j_2) = s(i_2) = s, i_2 \not\equiv i_1 \). Therefore, \( j_2 \not\equiv j_1 \).

• If \( j_3 = i_3 + 2^r \): From \( s'(i_2 - 2^r) = r \) it implies \( s'(j_2) \leq r \). If \( j_2 = i_2 - 2^r \), then \( (i_3 + 2^r, i_2 - 2^r, i_0) = (j_3, j_2, j_0) \). As seen in part (1), it contradicts the fact that \( \dim Q_1 \) does not divide \( 2^{r+2} \). Now we consider \( j_2 \neq i_2 - 2^r \). By dimensional information we get

\[ |j_2 - (i_2 - 2^r)| \leq 2^{r+1}. \]

If the equality happens, also by means of dimensional information we have

\[ J = (i_3 + 2^r, i_2 - 2^r - 2^{r+1}, i_1 + 2^{r+1}, 2). \]

Note that \( m(J) = m(I) = s, m'(J) = m'(I) = r \) and \( J \) is in Case B with \( s(j_3) > r = s'(j_2) \). So \( J \) can be controlled by part (1) of this proof.

If \( |j_2 - (i_2 - 2^r)| < 2^{r+1} \leq 2^s \). Put \( |j_2 - (i_2 - 2^r)| \equiv 2^t \mod 2^{t+1}, t < s \).

Using a similar argument to that of (c) in part (1) we can control \( J \).

(3) \( I \) is in Case C with \( s(i_3) \geq s'(i_2) \) or \( I \) is in Case D: First we show that in the both cases \( s(h) < \min\{s(i_3), s'(i_2)\} = s'(i_2) \). Indeed, suppose \( i_1 \equiv (2^s - 1) + 2^u \mod 2^{u+1} \), \( s < u < s' \). In the two cases: \( s(i_3) \geq s'(i_2) > u \). So

\[ h \equiv (2^{s(i_3)} - 1) + (2^{s(i_2)} - 2^s - 1) + [(2^s - 1) + 2^u] + 2 \mod 2^{u+1} \]

\[ \equiv 2^u - 1 \mod 2^{u+1}. \]

Hence \( s(h) = u < s'(i_2) \leq s(i_3) \).

Now we can apply Lemma 6.7 to \( K = (i_3 - 2^u, i_2, i_1, i_0) \in S(s) \).

\[ Sq^{2u+1}(i_3 - 2^u, i_2, i_1, i_0) = I + \sum J. \]

By dimensional information, \( j_3 \equiv i_3 \mod 2^{u+1} \iff J = I \).

Otherwise, \( |j_3 - i_3| \equiv 2^v \mod 2^{u+1} (v \leq u) \). Thus

\[ s(j_3) = v \leq u < s'(i_2) \), hence \( m'(J) < m'(I) \).

Moreover,

\[ j_0 = i_0 = 2 \quad (J \in S(s)) \]

\[ j_1 + j_0 \geq i_1 + i_0 \quad \Rightarrow \quad j_1 \geq i_1 > 2^u. \]

This together with the fact that \( J \in S(s) \) implies \( j_2 \not\equiv j_1 \).

The lemma is completely proved.

7. Proofs of Lemma 5.6 and Lemma 5.7

First, suppose that Lemma 5.7 is true. Basing ourselves on Lemma 7.1, we prove Lemma 5.6 by combining Lemmas 7.2, 7.3, 7.5, 7.7 below. Then, Lemma 5.7 will be proved at the end of the section.

Lemma 7.1. Suppose \( I \) satisfies the hypothesis of Lemma 5.6, which means \( s(h) \geq s(i_3) \geq s(i_2) \), and \( i_3 \not\equiv i_2 \not\equiv i_1 \not\equiv i_0 \). Then

(a) either \( i_0 = i_1 = 0, i_2 \) is even \( > 0, i_3 \) is even,

(b) or \( i_0 = i_1 = 0, i_2 \) is even \( > 0, i_3 \) is odd, \( s(i_2 - 1) \geq s(i_3) \),
Proof. Using actually the argument of Lemma 4.5 we get \( i_0 = [2 : 2^{s(i_1)}] = 0 \) or 2.

(a) If \( i_0 = 0 \), then \( s(i_1) = 0 \) or 1. Considering the reduced sequence \((i_3, i_2, i_1)\) at length 3 and repeating the argument mentioned above we obtain \( i_1 = 0 \) or 1. If \( i_1 = 0 \), and \( s(h) = 0 \), then \( s(i_3) = s(i_2) = 0 \). That means \( i_2, i_3 \) are even. Since \( i_3 \not\equiv i_2 \), so \( i_2 > 0 \).

(b) If \( i_0 = i_1 = 0 \) and \( s(h) > 0 \), then either \( i_2 \) or \( i_3 \) is odd, the other is even. But \( s(i_3) \geq s(i_2) \) so \( i_3 \) is odd and \( i_2 \) is even. By \( s(h) \geq s(i_3) \) we get \( s(i_2 - 1) = s(i_2 + i_1 + i_0 - 1) \geq s(i_3) \).

(c) If \( i_0 = 0, i_1 = 1 \), from \( s(h) \geq s(i_3) \) it implies \( s(i_2 + i_1 - 1) = s(i_2) \geq s(i_3) \), so \( s(i_3) = s(i_2) \).

(d) If \( i_0 = 2 \), then \( s(i_1) > 1 \). The hypothesis \( s(h) \geq s(i_3) \geq s(i_2) \) implies \( s(i_1 + i_0 - 2) = s(i_1) \geq s(i_2) \), so \( s(i_2) = s(i_1) = s \) (say). This equality together with \( i_2 \not\equiv i_1 \) implies \( i_1 = 2^s - 1 \). Set \( s' = s(i_2) \), that means \( i_2 \equiv 2s' - 2^s - 1 \mod 2^{s'+1} \). Since \( s(h) \geq s(i_3) \) so \( s(i_2 + i_1 + i_0 - 1) \geq s(i_3) \), or equivalently
\[
s(2^{s'} - 2^s - 1 + 2^s - 1 + 2 - 1 \mod 2^{s'+1}) = s' \geq s(i_3).
\]

The lemma follows.

Now we prove Lemma 5.6 by showing Lemmas 7.2, 7.3, 7.5, 7.7 below.

Lemma 7.2. Suppose \( I \) belongs to case (a) of Lemma 7.1, that means \( i_0 = i_1 = 0, i_2 \) is even and positive, \( i_3 \) is even. Then \( I \) is absolutely killed. (That means \( I \in \text{Im.} \).

Lemma 7.3. Suppose \( I \) belongs to case (b) of Lemma 7.1, that means \( i_0 = i_1 = 0, i_2 \) is even > 0, \( i_3 \) is odd, \( s(i_2 - 1) \geq s(i_3) \). Then \( I \) is absolutely killed. (That means \( I \in \text{Im.} \).

In the previous two lemmas \( m(I) = 0 \). So they are special cases of Lemma 5.10.

In order to prove the main lemma in case (c) of Lemma 7.1 we need the following technical lemma.

Lemma 7.4. Let \( H = (h_3, h_2, 1, 0) \). If \( s(h_3) = s(h_2) = s \), then for any \( a \)
\[
S^a_q h_3, h_2, 1, 0, = \sum J,
\]
with \( m(J) \leq m(H) = s \) for every \( J \) in the sum.

Proof.
\[
S^a_q h_3, h_2, 1, 0, = \sum J
= \sum (Sq^{a_q} Q_3)^a \cdots (Sq^{a_k} Q_3)^a (SQ^{a_0} Q_2^{b_q} \cdots (SQ^{b_k} Q_2)^{2^k b_p} (SQ^{m} Q_1).
\]

Here \( a_i, b_i \) are the coefficients of \( 2^i \) in the dyadic expansions of \( h_3, h_2 \) respectively, and \( k_i = 0 \) if \( a_i = 0 \), \( l_i = 0 \) if \( b_i = 0 \).
Take a sequence \( J \) in the sum. Define
\[
x = \sum_{k_i = 4, 6, 7} 2^i a_i, \quad y = \sum_{l_j = 12} 2^l b_j + \begin{cases} 1 & \text{if } m = 12 \text{ or } 13 \text{ or } 14, \\ 0 & \text{otherwise.} \end{cases}
\]
So we have \( j_2 + j_1 + j_0 - 1 = h_2 + h_1 + h_0 - 1 + x + y. \)

**Case A:** \( h(J) > h(H) \). We need only to consider \( J \) with \( s(j_2 + j_1 + j_0 - 1) > s = m(H). \)

1. We show \( x \not\equiv 2^s \) mod \( 2^{s+1} \). Indeed, if \( x \equiv 2^s \) mod \( 2^{s+1} \), by definition of \( x \) then \( a_s = 1 \). However, \( a_s = 0 \) since \( s(h_4) = s \). This is a contradiction.
2. Suppose \( x \equiv 0 \) mod \( 2^{s+1} \). One gets
\[
j_2 + j_1 + j_0 - 1 = h_2 + h_1 + h_0 - 1 + y \text{ mod } 2^{s+1}.
\]
Combining this with \( s(j_2 + j_1 + j_0 - 1) > s = s(h_2 + h_1 + h_0 - 1) \) we obtain \( y \equiv 2^s \) mod \( 2^{s+1} \). Recall that \( b_s = 0 \), so \( l_s = 0 \). Hence
\[
\begin{cases} l_j = 12 & \text{for any } j \text{ with } 0 \leq j < s \text{ (note } b_j = 1 \text{ for any } j < s), \\ m = 12 \text{ or } 14. \end{cases}
\]
(Note that \( m \neq 13 \) since \( 2^a \) is even if \( 2^a \geq 13 \).) Consider \( u = \min\{i | k_i = 8\} \) (\( u \) may not exist).

(a) We claim \( u \neq s \). Indeed, from \( a_s = 0 \) it implies \( k_s = 0 \).
(b) If \( u < s \), combining this with the facts that \( x \equiv 0 \) mod \( 2^{s+1} \), \( l_j = 12 \) for \( j < s \), \( m = 12 \) or \( 14 \) we get \( j_3 \equiv 2^u - 1 \) mod \( 2^{u+1} \). So \( m(J) \leq s(j_3) = u < s \).
(c) If \( u > s \) or there does not exist \( u \), combining this with the facts that \( x \equiv 0 \) mod \( 2^{s+1} \), \( l_j = 12 \) for \( j < s \), \( m = 12 \) or \( 14 \) we have \( j_3 \equiv 2^s - 1 \) mod \( 2^{s+1} \). So \( m(J) \leq s(j_3) = s = m(H) \).

3. Suppose \( x \equiv 2^u \) mod \( 2^{u+1} \), \( v < s \). It is easy to see
\[
j_2 + j_1 + j_0 - 1 = h_2 + h_1 + h_0 - 1 + y \text{ mod } 2^u.
\]
As \( s(j_2 + j_1 + j_0 - 1) > s = s(h_2 + h_1 + h_0 - 1) \), so \( y \equiv 0 \) mod \( 2^u \).

(a) There exists \( i < v \) with \( k_i = 8 \). Then \( m(J) \leq s(j_3) = i < v < s \).
(b) Otherwise, if there does not exist \( i < v \) with \( k_i = 8 \), then \( m(J) \leq s(j_3) = v < s \).

**Case B:** \( h(J) = h(H) \). Then \( y = 0 \) by definition.

1. We show \( x \not\equiv 2^s \) mod \( 2^{s+1} \) by the same argument as in (1) of Case A.
2. If \( x \equiv 0 \) mod \( 2^{s+1} \), then the fact that \( j_2 + j_1 + j_0 - 1 = h_2 + h_1 + h_0 - 1 + x \) implies
\[
m(J) \leq s(j_2 + j_1 + j_0 - 1) = s(h_2 + h_1 + h_0 - 1) = s.
\]
3. Suppose \( x \equiv 2^v \) mod \( 2^{v+1} \), \( v < s \). As \( h(J) = h(H) \), so \( j_3 = h_3 - x \).
Hence \( m(J) \leq s(j_3) = v < s \).

The lemma is proved.

**Lemma 7.5.** Suppose \( I \) belongs to case (c) of Lemma 7.1. That means
\[
i_0 = 0, \ i_1 = 1, \ s(i_2) = s(i_3) = s = m(I), \text{ and } i_3 \not\in i_2.
\]

(a) Further if \( h \not\in i_3 \), then the weak form of the main lemma holds for \( I \).
(b) If \( h \in i_3 \), then the strong form of the main lemma holds for \( I \).
Proof. (a) \( h \not\in i_3 \). Set \( t = s(h) > s \), we kill \( I \) by the usual method
\[
Sq^{2s+1}(i_3 - 2t, i_2, 1, 0) = I + \sum J.
\]
By Lemma 7.4, \( m(J) \leq m(i_3 - 2t, i_2, 1, 0) = s \). By dimensional information \( h(J) < h(I) \).

(b) \( h > i_3 \not\in i_2 \).

This case is treated in the following lemma.

Lemma 7.6. Suppose \( I = (i_3, i_2, i_1, i_0) \) satisfies \( i_3 \not\in i_2 \not\in i_1 = 1, i_0 = 0, s(i_3) = s(i_2) = s, i_3 + 2^s \not\in i_2 \). Then
\[
Q' \in \text{Im} \mathcal{A} + \text{Span}\{Q' | m(J) < m(I) = s\}.
\]

Proof. Note that the hypothesis of the lemma is valid at least when
\( h > i_3 \not\in i_2 \not\in i_1 = 1, i_0 = 0, s(i_3) = s(i_2) = s \).

Suppose \( s < r \). Consider the submodule
\[
R(s, r) = \text{Span}\{Q' | j_3 \not\in j_2 \not\in j_1 = 1, j_0 = 0, s(j_3) = s(j_2) = s, j_3 + 2^s \not\in j_2, s'(j_3) < r\}.
\]

Now we show that if \( I \) is as in the lemma, and \( r = s'(i_3) \), then
\[
Q' \in \text{Im} \mathcal{A} + \text{Span}\{Q' | m(J) < m(I) = s\} + R(s, r - 1).
\]
Furthermore, one could see
\[
R(s, s + 1) \subset \text{Im} \mathcal{A} + \text{Span}\{Q' | m(J) < s\}.
\]
So, the lemma follows.

From \( i_3 + 2^s \not\in i_2 \), it implies \( i_2 \geq 2^r \), where \( r = s'(i_3) \). One gets
\[
Sq^{2s+1}(i_3 + 2^r, i_2 - 2^r, 1, 0) = I + \sum J.
\]
Let \( J \) be a term in the right-hand side.

If \( j_0 > 0 \): Then \( j_0 \) is even (dim information). By Lemma 7.4, it suffices to consider \( m(J) = m(I) \). First we show that \( s(h(J)) \geq s(j_3) \). In fact, if \( h(J) = h(I) \), combine this with \( s(i_3 + 2^r) = s \) and get \( s(j_3) \leq s \). On the other hand, \( s(h(J)) = s(h(I)) > s \). So \( s(h(J)) \geq s(j_3) \). Suppose \( h(J) > h(I) \) and \( m(J) = m(I) \). Going through the proof of Lemma 7.4 we observe that in this case \( s(h(J)) \geq s(j_3) = s = m(I) \).

- If \( j_0 \neq 2 \): It is easy to check that \( J \) is \( s \)-killed by Lemma 5.4, Case A in Lemma 5.5 and Lemma 5.7.

- If \( j_0 = 2 \): Since \( i_0 = 0 \), \( j_0 \) is moved from either \( i_3 + 2^r, i_2 - 2^r \) or \( i_1 = 1 \). Hence

either \( j_3 \not\in j_0 \), or \( j_2 \not\in j_0 \) or \( j_1 \not\in j_0 \).

- If \( j_1 \not\in j_0 \): \( s(j_1) \leq 1 \). We can assume \( s(h), s(j_3), s(j_2) \) all \( \geq 1 \). (Otherwise, \( m(J) = 0 \), then \( J \) is absolutely killed as seen in the proof of Theorem 2.7 at the end of Section 5.) Now we can apply Lemma 5.4 and \( J \) is \( s \)-killed.

- If \( j_1 \not\in j_0, j_2 \not\in j_0 \): \( s(j_2) \leq 1 < s(j_1) \). Again, it suffices to consider \( s(h), s(j_3) \) all \( \geq 1 \). We can apply Case A of Lemma 5.5 and \( J \) is \( s \)-killed.
If \( j_2 \triangleright j_0, j_1 \triangleright j_0, j_3 \not\triangleright j_0 \): \( s(j_3) \leq 1 < s(j_2) \). Again, we need only to consider the case \( s(h) \geq 1 \geq s(j_3) < s(j_2) \). \( J \) is \( s \)-killed by Lemma 5.7.

If \( j_0 = 0 \): According to Lemma 7.4 it suffices to consider \( J \) with \( m(J) = m(I) = s \).

**Case A:** \( h(J) = h(I) = h, m(J) = m(I) = s \).

Since \( s(i_3) = s(i_2 + i_1 + i_0 - 1) = s \), so \( s(h) > s = m(J) \). Hence \( s(j_3) = s(j_2 + j_1 + j_0 - 1) = s \).

1. If \( s(j_2) > s \), then we can apply Lemma 5.7 to show \( J \) is \( s \)-killed.
2. If \( s(j_2) < s \), as \( h(J) = h(I) \) we have

\[
 z = \sum_{i_2=2}^u 2^i b_i \not\equiv 0 \mod 2^s \equiv 2^u \mod 2^{u+1}, \text{ for some } u < s. 
\]

Here we use the notations given in the proof of Lemma 7.4 with \( h_3 = i_3 + 2^s, h_2 = i_2 - 2^s \). Hence \( j_1 = i_1 + z \mod 2^s \). Combining this with the fact that \( s(j_2) = u \) we get \( j_2 \not\triangleright j_1 \). We can apply Lemma 5.5 to show \( J \) is \( s \)-killed.

3. \( s(j_3) = s \).
   a. \( j_1 > 1 \). Combine \( s(j_2 + j_1 + j_0 - 1) = s \), \( s(j_2) = s \), \( j_1 > 1 \), \( j_0 = 0 \) and get \( j_1 > 1 + 2^s \). So \( j_2 \not\triangleright j_1 \). Then \( Q^j \) is \( s \)-killed by Lemma 5.5 again.
   b. \( j_1 = 1 \). We show that either \( J = I \) or \( J \in R(s, r - 1) \).

Because of \( h(J) = h(I), j_0 = i_0 = 0, j_1 = i_1 = 1 \), \( j_3 + j_2 = i_3 + i_2 \).

Recall \( s(i_3) = s(j_2) = s \). We will prove that \( j_3 + 2^s \not\triangleright j_2 \). Indeed \( j_3 = i_3 - \sum_{k=4} a_i 2^i \). Combine this with \( s(j_3) = s(i_3) = s \) and get \( s(j_3 + 2^s) \leq s(i_3 + 2^s) \).

That means \( s'(j_3) \leq s'(i_3) = r \).

On the other hand, \( j_2 \geq i_2 \) by Corollary 3.2, \( i_3 + 2^s \not\triangleright i_2 \). Hence \( j_3 + 2^s \not\triangleright j_2 \).

If \( s'(j_3) = s'(i_3) \), then \( j_3 \equiv i_3 \mod 2^{r+1} \). Note that \( |Sq^{2^{r+1}}| = 4 \cdot 2^r \). By dimensional information we get \( J = I \).

Otherwise, if \( s'(j_3) < s'(i_3) = r \), then \( J \in R(s, r - 1) \).

**Case B:** \( h(J) > h(I), m(J) = m(I) = s \).

Going though the proof of Lemma 7.4, we see that the hypotheses \( h(J) > h(I), m(J) = m(I) = s \) happen only if \( s(j_3) = s \). Consider the following two cases.

1. \( s(j_3) = s, s(j_2 + j_1 + j_0 - 1) > s \).
   a. \( s(j_2) \leq s \). Combine this with \( s(j_2 + j_1 + j_0 - 1) > s, j_0 = 0 \) and get \( j_1 \geq 1 + 2^s \). So \( j_2 \not\triangleright j_1 \). Then \( Q^J \) is \( s \)-killed by Lemma 5.5.
   b. \( s(j_2) > s \). Then \( j_2 \not\triangleright j_1 \). So \( Q^J \) is \( s \)-killed by Lemma 5.7.

2. \( s(j_3) = s, s(j_2 + j_1 + j_0 - 1) = s \). It implies that \( s(h(J)) > s \).
   a. \( s(j_2) < s \). So \( j_1 \geq 1 + 2^{s(j_2)} \). Hence \( j_2 \not\triangleright j_1 \). Meanwhile \( s(h(J)) > s(j_3) \geq s(j_2) \). Then \( Q^J \) is \( s \)-killed by Lemma 5.5.
   b. \( s(j_2) > s \). So \( s(h(J)) > s(j_3) \not< s(j_2) \). Then \( Q^J \) is \( s \)-killed by Lemma 5.7.
   c. \( s(j_3) = s \). If \( j_1 > 1 \), then \( j_1 \geq 1 + 2^{s(j_2)} \). So \( Q^J \) is also \( s \)-killed by Lemma 5.5. Suppose \( j_1 = 1 \). Note that \( j_3 \geq i_3 + 2^r - 2^{r+2}/4 = i_3 \).

If \( j_3 = i_3 \), then \( J = I \) (by dimensional information). This contradicts the fact that \( h(J) > h(I) \).

If \( j_3 > i_3 \), then \( i_3 \leq i_3 + 2^r - 2^{r-1} \). So \( s'(j_3) \leq r - 1 \).

There are two possibilities:

- \( j_3 + 2^s \not\triangleright j_2 \), then \( Q^J \in R(s, r - 1) \).
• \( j_3 + 2^u \gg j_2 \), then \( j_2 < 2^{s(j_3)} \leq 2^{r-1} \). Set

\[
u = \max\{i | \alpha_i(j_3) = 0, \alpha_{i+1}(j_3) = 1\} \geq s'(j_3).
\]

Then \( j_3 - 2^u \geq \alpha_0 \cdot 2^0 + \cdots + 2^u + 0 \cdot 2^{u+1} + j \), where \( j \equiv 0 \mod 2^{u+2} \).

Since \( j_3 + 2^u \gg j_2 \), so \( j_2 < 2^{s(j_3)} \leq 2^u \).

\[S\!q^{2u+3}(Q_3^{j_3-2^u}Q_2^{j_2}Q_1) = Q_3^jS\!q^{2u+3}(Q_3^{j_3-2^u}Q_2^{j_2}Q_1).\]

Therefore, we can reduce to the case \( j = 0 \), i.e., we have

\[j_3 - 2^u = \alpha_0 \cdot 2^0 + \cdots + 2^u.\]

\[S\!q^{2u+3}(Q_3^{j_3-2^u}Q_2^{j_2}Q_1) = (S\!q^8Q_3)^2u(Q_3^{j_3-2^{u+1}}Q_2^{j_2}Q_1)
\]

\[+ (S\!q^7Q_3)^2uS\!q^{2u}(Q_3^{j_3-2^{u+1}}Q_2^{j_2}Q_1)
\]

\[+ (S\!q^6Q_3)^2uS\!q^{2u+1}(Q_3^{j_3-2^{u+1}}Q_2^{j_2}Q_1)
\]

\[+ (S\!q^4Q_3)^2uS\!q^{2u+2}(Q_3^{j_3-2^{u+1}}Q_2^{j_2}Q_1)
\]

\[+ (Q_3)^2uS\!q^{2u+3}(Q_3^{j_3-2^{u+1}}Q_2^{j_2}Q_1).
\]

The first term is obviously \( Q^j \).

Let the second term be \( \sum Q^T \). Since \( S\!q^7Q_3 = Q_0 \), then \( t_0 \geq 2^u \) for any \( T \) in the sum. One can repeat the argument for the previous cases of this lemma to show that \( Q^T \) is s–killed by Lemma 5.4 or Lemma 5.5.

Let the third term be \( \sum Q^T \). Since \( S\!q^6Q_3 = Q_1 \), so \( t_1 \geq 2^u \) for any \( T \) in the sum. Using Lemma 5.5 one can s–kill \( Q^T \) by the same way that is used for the second term.

Let the fourth term be \( \sum Q^T \). Since \( S\!q^4Q_3 = Q_2 \), so \( t_2 \geq 2^u \) for any \( T \) in the sum. One can repeat the argument for the previous cases of this lemma to show that, studying modulo \( \text{Im} \mathcal{A} + \text{Span}\{Q^K | m(K) < s = m(J)\} \), one needs only to consider \( T \) with \( s(t_3) = s, s'(t_3) \leq s'(j_3) - 1 \leq r - 2 \). So \( t_3 + 2^u \neq 2^{s(t_3)} \leq 2^u \). Hence \( t_3 + 2^u \neq t_2 \). Therefore, \( Q^T \) is s–killed modulo \( R(s, r - 1) \) by Lemmas 5.4, 5.5 and 5.7.

Suppose the last term is \( \sum Q^T \). Again, using the argument for the previous cases of this lemma, it suffices to consider \( T \)'s satisfying

\[m(T) = s(t_3) = s(t_2) = s, t_3 + 2^u \gg t_2, t_1 = 1, t_0 = 0.
\]

However, we will arithmetically prove that there is no such a \( T \).

Let us apply Lemma 7.4 to \( S\!q^{2u+3}(Q_3^{j_3-2^{u+1}}Q_2^{j_2}Q_1) \) and use the notations given in its proof \((a_i, b_i, k_i, l_i, m,...)\).

Fix such a \( T \) and take a corresponding sequence \( \{k_i, l_i, m\} \).

(*) First, we show that there exists \( i \) such that \( l_i \neq 0 \). Otherwise, if \( l_i = 0 \) for any \( i \), then combining this with the fact that \( m = 0 \) (because of \( t_1 = 1, t_0 = 0 \)) we claim \( Q^T \) is a term of

\[S\!q^{2u+3}(Q_3^{j_3-2^{u+1}})Q_2^{j_2}Q_1 = 0,
\]

(as \( j_3 - 2^{u+1} < 2^u \)).
Next, we show that one can choose a sequence \( \{k_i, l_i, m\} \) among those corresponding to \( Q^T \) so that, if \( l_x \neq 0 \) then \( k_i = 0 \) for any \( i \leq x \). Suppose the contrary. Let \( v = \min\{i | |i| \leq x, k_i \neq 0\} \). As \( s(t_3) = s(t_2) = s \), we get \( v > s \). Note that \( k_v = 4 \) or \( 8 \) (because of \( t_1 = 1, t_0 = 0 \)). If either \( k_v = 4 \) or \( 8 \), and \( l_v = 0 \) or \( 12 \), then \( \alpha_v(t_3) = \alpha_v(t_3 + 2^s) = 0 \), so \( s(t_3 + 2^s) \leq v \). Meanwhile, \( l_x \neq 0 \) hence \( l_x = 8 \) or \( 12 \) (because of \( t_1 = 1, t_0 = 0 \)). It implies \( t_2 \geq 2^s \geq 2^v \). Thus \( t_3 + 2^s \neq t_2 \). This contradicts the hypothesis on \( T \).

If \( k_v = 4 \) and \( l_v = 8 \), we choose another sequence \( \{k'_i, l'_i, m\} \) which corresponds to \( Q^T \) as well as \( \{k_i, l_i, m\} \) as follows:

\[
k'_i = \begin{cases} k_i & \text{if } i \neq v, \\ 0 & \text{if } i = v, \\ 12 & \text{if } i = v. \end{cases}
\]

The new sequence satisfies the property we state at the beginning of the item.

Finally, set \( j = \min\{i | l_i \neq 0\} \). From \( s(t_3) = s(t_2) = s, t_1 = 1, t_0 = 0 \) it implies that \( j > s, m = 0 \) and if \( k_i \neq 0 \) then \( k_i = 4 \) or \( 8 \), also if \( l_i \neq 0 \) then \( l_i = 8 \) or \( 12 \). They all are divisible by 4. Since \( l_j \neq 0 \) so \( b_j = 1 \) (by definition of \( l_j \)). Furthermore,

\[
2^i l_j = 2^i l_j b_j = 2^{i+3} - \sum_{i > j} 2^i k_i a_i - \sum_{i > j} 2^i l_i b_i \equiv 0 \mod 2^{i+3}.
\]

It implies \( l_j = 8 \) (but not 12).

If \( l_{j+1} = 0 \), then \( 2^i l_j = 2^{i+3} \equiv 2^{i+1} k_{j+1} a_{j+1} \mod 2^{i+4} \). Then \( k_{j+1} = 4 \) (\( a_{j+1} = 1 \)). So \( s(t_3 + 2^s) \leq j + 1, t_2 \geq 2^{j+1} \). Hence \( t_3 + 2^s \neq t_2 \) (contradiction).

If \( l_{j+1} = 8 \), then \( k_{j+1} = 0 \) (see (**)). Thus \( 2^{j+3} \equiv 0 \mod 2^{j+4} \). This cannot happen.

Therefore, \( l_{j+1} = 12 \). So \( b_{j+1} = \alpha_{j+1}(j_2) = 1 \). That means \( j_2 \geq 2^{j+1} \).

Combining this with \( j_3 + 2^s > j_2 \) we get \( a_{j+1} = \cdots = a_{j+1} = 1 \) (recall that \( a_i = \alpha_i(j_3 - 2^{i+1}) \)).

Meanwhile, \( l_j = 8, a_j = 1 \) implies \( \alpha_j(t_3) = 0 \), so \( s(t_3 + 2^s) \leq j \).

On the other hand, \( l_{j+1} = 12 \) implies \( t_2 \geq 2^{j+1} \cdot 2 = 2^{j+2} \). Therefore \( t_3 + 2^s \neq t_2 \) (contradiction).

The fact that \( R(s, s + 1) \subset \text{Im} \alpha + \text{Span}(Q^J | m(J) < s) \) can be read off from the inductive procedure on \( s \).

In the proof of the lemma we need only to use Lemmas 5.4, 5.5 and 5.7, so the strong form of the main lemma holds for \( I \). The lemma, and so Lemma 7.5, is proved.

**Lemma 7.7.** Suppose \( I \) belongs to case (d) of Lemma 7.1. That means

\[
i_0 = 2, \quad i_1 = 2^s - 1 \quad (s > 1), \quad s(i_2) = s(i_1) = s = m(I),
\]

\[
s \leq s(i_3) \leq s'(i_2), \quad s(h) \geq s(i_3).
\]

Then either the strong form or the exceptional form of the main lemma holds for \( I \).

**Proof.** In this proof, one can see how the strong form of the main lemma is needed.
Set $r = s(i_3)$. We kill $J$ by the usual method using $(i_3 + 2', i_2 - 2', i_1, i_0)$. 

$$Sq^{2r+2}(i_3 + 2', i_2 - 2', i_1, i_0) = I + \sum J = \sum (Sq^0 Q_3)^a \cdots (Sq^{k_3} Q_3)^{2^a a} \cdot (Sq^{b_1} Q_2)^{b_1} \cdots (Sq^{b_s} Q_2)^{2^s b_s} \cdot (Sq^{m_1} Q_1)^{c_1} \cdots (Sq^{m_t} Q_1)^{2^t c_t} \cdot (Sq^n Q_0)^2.$$ 

Here $a_i, b_i, c_i$ are the coefficients of $2^i$ in the dyadic expansions of $i_3 + 2', i_2 - 2', i_1$ respectively. Convention: $k_i = 0$ if $a_i = 0$, $l_i = 0$ if $b_i = 0$, $m_i = 0$ if $c_i = 0$. Consider a $J$ in the sum with $h' = h(J)$. We need only to study the case $m(J) \geq m(I) = s$. That means $s(j_3), s(j_2),$ and $s(j_1 + j_0 - 2)$ all $\geq s$.

**Case 1:** $s(j_3) \geq s$, $s(j_2) > s$. These imply $k_s = 4$ and $s(j_3) = s$. Hence $m(J) \leq s = m(I) \leq m(J)$, so $m(J) = s \leq s(h')$. We can apply Lemma 5.7 and $J$ is $s$-killed.

**Case 2:** $s(j_3) \geq s$, $s(j_2) = s$, $s(j_1) > s$. Then $m(J) \leq s(j_2) = s = m(I)$. So $m(J) = s$. Then $s(h') \geq m(J) = s$. Since $s(j_2) < s(j_1)$ so $j_2 \not\prec j_1$, we can apply Case A of Lemma 5.5: $J$ is $s$-killed.

**Case 3:** $s(j_3) \geq s$, $s(j_2) = s$, $s(j_1) < s$. Hence $s(j_1 + j_0 - 2) \geq s$. This together with $s(j_1) = t < s$ implies $j_0 > 2'$. Therefore, $j_1 \not\prec j_0$. We also have $s(h') \geq m(J) = s > t$. Hence, by Lemma 5.4, $J$ is $s$-killed.

**Case 4:** $s(j_3) \geq s$, $s(j_2) = s(j_1) = s$. We also have $m(J) = s$. Then $s(h') \geq s$.

If $j_2 \not\prec j_1$ or $j_1 \not\prec j_0$: then we can apply Lemma 5.4 or Lemma 5.5: $J$ is $s$-killed.

If $j_2 \gg j_1 \gg j_0$: we will show $s(h') \geq s(j_3)$. In fact, $s(j_3) \geq s(j_2) = s(j_1) = s \leq s(h')$ and $j_2 \gg j_1 \gg j_0$ imply $j_0 = 2 = i_0, j_1 = 2s - 1 = i_1$ (by the same argument as in the proof of Lemma 4.5). So $j_3 \not\equiv i_3 + 2'$. Otherwise, 

$$2'^{+2} = \dim Q_2 \cdot \{j_2 - (i_2 - 2')\} \equiv 0 \mod 12.$$ 

This cannot happen. We note that $s(j_3) < r$. (The only J with $s(j_3) \geq r$ is I, by dimensional information.)

Suppose the contrary that $s(h') < s(j_3)$ or equivalently $s(j_3) > s'(j_2) := u$. Dimensional information shows 

$$2'^{+2} = \dim Q_3 \{j_3 - (i_3 + 2')\} + \dim Q_2 \{j_2 - (i_2 - 2')\} \equiv 12 \cdot 2^u \mod 2^{u+3}.$$ 

This contradicts $u < r - 1$. So we get $s(h') \geq s(j_3) \geq s(j_2) \geq s(j_1)$.

Assume in addition $j_3 \gg j_2 \gg j_1 \gg j_0$: Recall $s(h') \geq s(j_3)$, and $j_0 \geq i_0 = 2$.

By the method in Lemma 4.5 and Remark 4.9, we get:

(a) either $j_0 = 2, j_1 = 2^s - 1, j_2 = 2^s - 1, j_3 \equiv 2^s - 1 \mod 2^{s+1}$,

(b) or $j_0 = 2, j_1 = 2^s - 1, j_2 = 2^t - 2^s - 1, j_3 \equiv 2^t - 1 \mod 2^{t+1} (t > s)$.

In the both cases, $J$ should be $(i_3 + 2' + 2'^{-1}, i_2 - 2', i_1, i_0)$. Actually case (b) cannot happen. If so, $t = s(i_3) = r - 1$, and $j_2 = 2^t - 2^s - 1 = i_2 - 2'$. Hence, $i_2 = 2^t + 1 + 2^s - 2^s - 1, s'(i_2) = t < s(i_3)$. This contradicts the hypothesis of the lemma: $s(i_3) \leq s'(i_2)$.
In case (a): \( r - 1 = s, \) \( s(i_3) = r = s + 1, \) \( i_3 \equiv 2^{s+1} - 1 \mod 2^{s+2}. \) Thus \( j_3 = i_3 + 2^{s+1} + 2^s \equiv 2^s - 1 \mod 2^{s+2}. \) Since \( j_3 > 2^t - 1, \) we set \( j_3 = 2^t - 2^{s+2} + 2^s - 1 + x, \) where \( x \equiv 0 \mod 2^{t+1}, \) \( t > s + 2. \) We get

\[
i_3 = 2^t - 2^{s+2} - 2^{s+1} - 1 + x, \quad h' = h(J) = 2^t - 2^s - 1 + x.
\]

Note: \( h' \not\geq j_3 \) if and only if \( x = 0. \)

If \( x \neq 0: \) then \( x > 2^t, \) \( \alpha_t(j_3) = 0. \) Let \( v \) be a number \( \geq t \) such that \( \alpha_t(j_3) = 0, \) \( \alpha_{v+1}(j_3) = 1. \) \( (v \) exists as \( x > 2^t. \) \) Hence \( \alpha_v(h) = 0. \)

\[
Sq^{2^{v+3}}(j_3 - 2^v, j_2, j_1, j_0) = J + \delta(j_3 - 2^v - 2^{v+1}, j_2 + 2^{v+1}, j_1, j_0)
\]

\[
+ \text{Span}\{K|m(K) < m(J) = s}\}
\]

Indeed, any change at \( j_2, j_1 \) or \( j_0 \) makes the function \( m(\cdot) \) smaller.

If \( x = 0: \) \( J = I' = (2^t - 2^{s+1} - 2^s - 1, 2^s - 1, 2^s - 1, 2) \) is a weakly allowed sequence (as \( t > s + 2. \)) We observe \( I = (2^t - 2^{t+2} - 2^{s+1} - 1, 2^{s+1} + 2^s - 1, 2^s - 1, 2). \) So \( h(I) = 2^t - 2^{s+1} - 1 < h(I') = 2^t - 2^s - 1, m(I) = m(I') = s. \)

Finally assume that \( j_3 \not\geq j_2 \not\geq j_1 \not\geq j_0: \) Using Lemma 4.5, from these and \( j_0 \geq i_0, \) it implies \( j_0 = 2. \) Now \( J \) satisfies the hypothesis of Lemma 7.7 (see Lemma 7.1). Recall that \( s(j_3) < s(i_3). \) So we can repeat the above procedure to handle \( J \) inductively on \( s(j_3). \)

**Remark.** \( h' \) could be greater than \( h. \) But it does not damage the proof because we are concerned only with the strong form of the main lemma.

In the inductive procedure we might meet a weakly allowed sequence of the form \( J = I' = (2^t - 2^{s+1} - 2^s - 1, 2^s - 1, 2^s - 1, 2) \), when \( x = 0. \) We cannot meet more than one sequence of this form because of dimensional information. That means either the strong form or the exceptional form of the main lemma holds for \( I. \) The lemma is completely proved.

**Example 7.8.** Let \( I = (2^{t-1} - 1, 2^t - 2^s - 1, 2^s - 1, 2). \) The above inductive procedure shows modulo \( \text{Im} \mathcal{A} + \text{Span}\{K|m(K) < s\}: \)

\[
I = (2^t + 2t-2 - 1, 2^t-2 - 2^s - 1, 2^s - 1, 2) = (2^t + 2t-1 + 2t-3 - 1, 2t-2 - 2^s - 1, 2^s - 1, 2) = \ldots = \ldots = (2^t+1 - 2^{s+1} - 2^s - 1, 2^s - 1, 2^s - 1, 2).
\]

The last sequence is weakly allowed.

Now we give a direct proof for Lemma 5.7.

**Proof of Lemma 5.7.** According to Lemma 5.1 and Definition 5.2, \( m(I) = s := s(i_3), \) and

\[
I + \left(\begin{array}{c} i_2 - 2^s \\ 2^{s+1} \end{array}\right) I' + \left(\begin{array}{c} i_1 \\ 2^{s+2} \end{array}\right) I'' \in \text{Im} \mathcal{A} + \text{Span}\{J|m(J) < s\},
\]

where

\[
I' = (i_3 + 2^s, i_2 - 2^s - 2^{s+1}, i_1 + 2^{s+1}, i_0),
\]

\[
I'' = (i_3 + 2^s, i_2 - 2^s, i_1 - 2^{s+2}, i_0 + 2^{s+2}).
\]
It is easy to check that $I''$ is $s$-killed by Lemma 5.4 or Case A of Lemma 5.5, since $i'_0 = i_0 + 2s^2 > 2$. We need only to study $I'$ with $(i_{2s^2+1}) = 1$ or equivalently $\alpha_{s+1}(i_2 - 2s - 2s^2 + 1) = 0$. So $s(i_2 - 2s - 2s^2 + 1) = s$, and

$$s'(i_2 - 2s - 2s^2 + 1) = s + 1.$$ We can concentrate on the case

$$s(i_1 + 2s^2 + 1) = s, \quad s'(i_1 + 2s^2 + 1) = s + 1, \quad i_0 = 2.$$ Otherwise, $I'$ is $s$-killed by Lemma 5.4 or Case A of Lemma 5.5 as before. Hence,

$$i_2 - 2s - 2s^2 + 1 \neq i_1 + 2s^2 + 1.$$ $I'$ turns out to be a very special case in Case B of Lemma 5.5 with $s'(i'_2) = s(i'_2) + 1$. We have

$$Sq^{2s^2 + 1}(i_3 + 2s, i_2 - 2s^2 + 1, i_1 + 2s, 2)$$

$$= I' + (i_3 + 2s, i_2 - 2s^2 + 1, i_1 + 2s - 2s^2 + 1, 2 + 2s^2 + 1)$$

modulo Span{$\{L|m(L) < s\}$. Denote the second term by $J$. We note that

$$s(i_2 - 2s^2 + 1) = s(i_1 + 2s^2 - 2s^2 + 1) = s + 1,$$

and either $s(i_3 + 2s) = s + 1$, hence $s(h(J)) > s + 1$, or $s(h(J)) = s + 1$, hence $s(i_3 + 2s) > s + 1$. From here to the end of this section, all the computations are modulo Span{$\{L|m(L) < s\}$. We always try to show $J = K$ modulo Im.$\mathcal{D}$ + Span{$\{L|m(L) < s\}$, for some $K$ with $s(k_3) = s(k_2) = s(k_1) = s$ in the remaining part of the section.

We need the technical lemma:

Lemma 7.9. Suppose $s(i_3) = s(i_2) = s(i_1) = s$, $s'(i_2) = s + 1$, $i_1 > 2s^2 + 1$. Then $I$ is $s$-killed.

Proof.

$$Sq^{2s^2 + 1}(i_3 + 2s^2 + 1, i_1 - 2s^2 + 1, i_0) = I$$

modulo Span{$\{J|m(J) < s\}$. (Use Lemma 5.4.) The lemma is proved.

Come back to the proof of Lemma 5.7.

Let us consider the two possibilities for $J$ as seen above.

(1) $J = (2s^2 + 1 + a, 2s^2 + 1 + b, 2s^2 + 1 + c, 2 + 2s^2 + 1), \text{ with } a \equiv b \equiv c \equiv 0 \mod 2s^2 + 2.$

$$Sq^{2s^2 + 1}(2s^2 + 1 + a, 2s^2 + 1 + b, 2s^2 + 1 + c, 2)$$

$$= J + (2s^2 + 1 + a, 2s^2 + 1 + b, 2s^2 + 1 + c, 2) \overline{J}.$$
(a) If $\alpha_{s+2}(b) = 0$, then
\[
Sq^{2s+3}(2^{s+1} - 1 + a, 2^{s+2} + 2^s - 1 + b, 2^s - 1 + c, 2) = K + (2^{s+1} + 2^s - 1 + a, 2^{s+2} + 2^s - 1 + b, 2^s - 1 + c, 2).
\]

- If $c \neq 0$: $L$ is $s$–killed by Lemma 7.9.
- If $c = 0$: Note that $\alpha_{s+1}(h(L)) = 0$.
\[
Sq^{2s+4}(2^s - 1 + a, 2^{s+2} + 2^s - 1 + b, 2^s - 1, 2) = L + \alpha_{s+2}(a)(2^s - 1 + a - 2^{s+2}, 2^{s+3} + 2^s - 1 + b, 2^s - 1, 2)
\]

\[+ \{\text{a term, which is } s\text{-killed by Lemma 7.9}\}.
\]
\[
Sq^{2s+3}(2^s - 1 + a - 2^{s+1}, 2^{s+3} - 2^{s+1} + 2^s - 1 + b, 2^s - 1, 2) = M + (2^s - 1 + a - 2^{s+1}, 2^{s+1} + 2^s - 1 + b, 2^{s+2} + 2^s - 1, 2). \]

\[
Sq^{2s+2}(a - 1, 2^{s+1} - 1 + b, 2^{s+2} + 2^s - 1, 2) = N + (a - 1, 2^{s+1} - 1 + b, 2^s - 1, 2 + 2^{s+2}).
\]

The last term is $s$–killed by Lemma 5.4.

(b) If $\alpha_{s+2}(b) = 1$, $K$ can be handled as follows:
\[
Sq^{2s+3}(2^{s+2} - 1 + a, 2^s - 1 + b - 2^{s+1}, 2^{s+2} + 2^s - 1 + c, 2) = K + (2^{s+2} + 2^s - 1 + a, 2^s - 1 + b - 2^{s+1}, 2^{s+2} + 2^s - 1 + c, 2)
\]
\[+ \alpha_{s+3}(c)(2^{s+2} - 1 + a, 2^s - 1 + b - 2^{s+1}, 2^{s+2} + 2^s - 1 + c, 2).
\]

The last term is $s$–killed by Lemma 5.4.

Since $\alpha_{s+2}(b - 2^{s+1}) = 0$, so we get
\[
Sq^{2s+3}(2^{s+2} + 2^s - 1 + a, 2^{s+1} + 2^s - 1 + b, 2^s - 1 + c, 2) = P + \alpha_{s+3}(c)(2^{s+2} + 2^s - 1 + a, 2^{s+1} + 2^s - 1 + b, 2^s - 1 + c - 2^{s+3}, 2 + 2^{s+3}).
\]

This term is $s$–killed by Lemma 5.4.

*Note:* The last equality is based on the fact that
\[
\alpha_{s+1}(2^{s+2} + 2^s - 1 + a) = 0 = \alpha_s(h(Q)),
\]
where
\[
Q = (2^{s+2} + 2^s - 1 + a, 2^{s+1} + 2^s - 1 + b, 2^s - 1 + c, 2).
\]
(2) \( J = (2^t - 1 + a, 2^{s+1} - 1 + b, 2^{s+2} - 1 + c, 2 + 2^{s+1}) \), with \( a \equiv 0 \mod 2^{s+1} \) and \( b \equiv c \equiv 0 \mod 2^{s+2}, \ t > s + 1. \)

\[
S^q 2^{s+1} (2^t - 1 + a, 2^{s+1} - 1 + b, 2^{s+2} - 1 + c, 2) = J + \left( \sum_{k} \frac{2^t - 1 + a, 2^s - 1 + b, 2^{s+2} + 2^t - 1 + c, 2}{K} \right).
\]

Note that \( s(h(K)) = s + 1. \) So

\[
S^q 2^{s+4} (2^t - 2^{s+1} - 1 + a, 2^s - 1 + b, 2^{s+2} + 2^t - 1 + c, 2) = K + \left( \sum_{l_1} \frac{2^t - 2^{s+2} - 2^{s+1} - 1 + a, 2^{s+2} + 2^s - 1 + b, 2^{s+2} + 2^s - 1 + c, 2}{L_1} \right)
\]

\[
+ \alpha_{s+3}(b)(2^t - 2^{s+1} - 1 + a, 2^s - 1 + b - 2^{s+3}, 2^{s+3} + 2^{s+2} + 2^s - 1 + c, 2) \]

\[
+ \alpha_{s+4}(c)(2^t - 2^{s+1} - 1 + a, 2^s - 1 + b, 2^{s+2} + 2^s - 1 + c - 2^{s+4}, 2 + 2^{s+4}).
\]

The last term is \( s \)-killed by Lemma 5.4. 

(a) If \( \alpha_{s+2}(b) = 0, \) first we kill \( L_1: \)

\[
S^q 2^{s+3} (2^t - 2^{s+2} - 1 + a, 2^{s+1} + 2^s - 1 + b, 2^{s+2} + 2^s - 1 + c, 2) = L_1 + \left( \sum_{m_1} \frac{2^t - 2^{s+2} + 2^{s+1} - 1 + a, 2^{s+1} + 2^s - 1 + b, 2^{s+2} + 2^s - 1 + c, 2}{M_1} \right)
\]

\[
+ \alpha_{s+3}(c)(2^t - 2^{s+2} - 1 + a, 2^{s+1} + 2^s - 1 + b, 2^{s+2} + 2^s - 1 + c - 2^{s+3}, 2 + 2^{s+3}).
\]

The last term is \( s \)-killed by Lemma 5.4.

\[
S^q 2^{s+2} (2^t - 2^{s+2} + 2^{s+1} - 1 + a, 2^{s+1} - 1 + b, 2^{s+2} + 2^s - 1 + c, 2) = M_1 + \alpha_{s+2}(c + 2^{s+2})(2^t - 2^{s+2} + 2^{s+1} - 1 + a, 2^{s+1} - 1 + b, 2^{s+2} + 2^s - 1 + c, 2 + 2^{s+2}).
\]

Again, the last term is \( s \)-killed by Lemma 5.4. Now we kill \( L_2: \)

\[
S^q 2^{s+3} (2^t - 2^{s+1} - 1 + a, 2^s - 1 + b - 2^{s+2}, 2^{s+3} + 2^s - 1 + c, 2) = L_2 + \left( \sum_{m_2} \frac{2^t - 2^{s+1} + 2^s - 1 + a, 2^s - 1 + b - 2^{s+2}}{M_2} \right)
\]

\[
+ (2^t - 2^{s+1} - 1 + a, 2^s - 1 + b - 2^{s+2}, 2^s - 1 + c, 2 + 2^{s+3}).
\]

The last two terms are \( s \)-killed by Lemma 7.9 and Lemma 5.4, respectively.
(b) If $\alpha_{s+2}(b) = 1$, we kill $L_1$:

$$S^q(2^t - 2s+2 - 2s+1 - 1 + a, 2s+3 + 2s - 1 + b, 2s - 1 + c, 2)$$

$$= L_1 + (2^t - 2s+2 - 2s+1 + 2^{s+3} - 1 + a, 2s+3 + 2s - 1 + b, 2s - 1 + c, 2)$$

$$+ \alpha_{s+3}(c)(2^t - 2s+2 - 2s+1 - 1 + a, 2s+3 + 2s - 1 + b,$$

$$2s - 1 + c - 2s+3, 2 + 2s+3).$$

The last term is $s$-killed by Lemma 5.4.

$$S^q(2^t - 2s+2 + 2s - 1 + a, 2s+3 + 2s - 1 + b - 2s+1, 2s - 1 + c, 2)$$

$$= M_1 + \alpha_{s+3}(c)(2^t - 2s+2 - 2s - 1 + a, 2s+3 + 2s - 1 + b - 2s+1,$$

$$2s - 1 + c - 2s+3, 2 + 2s+3).$$

The last term is also $s$-killed by Lemma 5.4. Note that the last equality is based on the fact that

$$\alpha_s(h(M_1)) = 0 = \alpha_{s+2}(b - 2s+1).$$

We kill $L_2$:

$$S^q(2^t - 1 + a, 2s - 1 + b - 2s+3 - 2s+1, 2s+3 + 2s+2 + 2s - 1 + c, 2)$$

$$= L_2 + (2^t - 2s+3 - 2s+1 - 1 + a, 2s+3 + 2s - 1 + b - 2s+1,$$

$$2s - 1 + c - 2s+3, 2 + 2s+3).$$

The last term is $s$-killed by Lemma 5.4.

$$S^q(2^t + 2s - 1 + a, 2s - 1 + b - 2s+2 - 2s+1, 2s+3 + 2s - 1 + c, 2)$$

$$= M_2 + (2^t + 2s - 1 + a, 2s - 1 + b - 2s+2,$$

$$2s+3 + 2s - 1 + c, 2 + 2s+3).$$

The last two terms are $s$-killed by Lemma 7.9 and Lemma 5.4, respectively. Lemma 5.7 is completely proved.

8. Proofs of Lemma 5.8 and Lemma 5.9

Lemma 8.1. Suppose $I$ satisfies the hypothesis of Lemma 5.8. That means $s(h) \geq s(i_3)$, $i_3 \triangleright i_2 \triangleright i_1 \triangleright i_0$, $h \not\triangleright i_3$ and $h \not\triangleright i_3$. Then

(a) either $i_0 = i_1 = i_2 = 0$, $m(I) = 0$, $i_3 > 2s+1$,
(b) or $i_0 = 0$, $i_1 = 1$, $i_2 = 2s - 1$, $s(i_3) = s = m(I)$,
(c) or $i_0 = 2$, $i_1 = 2s - 1$, $i_2 = 2s - 2s - 1 + (s' > s + 1)$, $s(i_3) = s'$, $m(I) = s$,
(d) or $i_0 = 2$, $i_1 = 2s - 1 = i_2$, $s(i_3) = \text{either } s \text{ or } s + 1$, $m(I) = s$. 

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Proof. Using the same argument as in the proof of Lemma 4.5, one gets $i_0 = [2: 2^{(i_1)}] = 0$ or 2.

(a) If $i_0 = 0$, consider the reduced sequence $(i_3, i_2, i_1)$ and we get $i_1 = 0$ or 1. A possibility is $i_0 = i_1 = i_2 = 0$. The fact that $i_3 > 2^{(i_3)}$ follows from the hypothesis $h \not\in i_3$.

If $i_0 = i_1 = 0, i_2 > 0$, then $i_3$ and $h$ are odd since $h > i_3 > i_2$. Hence $i_2$ is even, and $s(i_2 - 1) \geq s(i_3)$. So $i_2 \geq 2^{(i_3)}$. This contradicts $i_3 > i_2$.

(b) If $i_0 = 0, i_1 = 1$, since $s(h) \geq s(i_3)$ so $s(i_2) = s(i_2 + i_1 + i_0 - 1) \geq s(i_3)$. Then $s(i_2) = s(i_3) = s$ (say). From $i_3 > i_2$ we have $i_2 = 2^s - 1$.

(c) If $i_0 = 2, s = s(i_1) > 1$. By $s(h) \geq s(i_3) \geq s(i_2)$ we get $s(i_1 + i_0 - 2) = s(i_1) \geq s(i_2) \geq s(i_1)$, thus $s(i_2) = s(i_1) = s$. This together with $i_2 > i_1$ implies $i_1 = 2^s - 1$.

Set $s' = s'(i_2), i_2 \equiv 2^s - 2^s - 1 \mod 2^{s+1}$. Because $s(h) \geq s(i_3)$, we have

\[ s(i_2 + i_1 + i_0 - 1) = s(2^s - 1 \mod 2^{s+1}) = s' \geq s(i_3) \geq s(i_2). \]

If $s' > s + 1$, then the fact $s(h) \geq s(i_3)$ implies $s(i_3) = s'$. So $i_2 = 2^s - 2^s - 1$ as $i_3 > i_2$.

(d) Finally, if $i_0 = 2, i_1 = 2^s - 1, i_2 \equiv 2^s - 2^s - 1 \mod 2^{s+1}$, and $s' = s + 1$, then $s(i_3)$ could be one of the two possible values $s$ and $s + 1$. In the both cases, we have $i_2 = 2^s - 1$ because of $i_3 > i_2$. The lemma is proved.

According to Lemma 8.1, we now prove Lemma 5.8 by showing Lemmas 8.2, 8.3, 8.4 and 8.5 below.

**Lemma 8.2.** Suppose $I$ belongs to case (a) of Lemma 8.1, that means $i_0 = i_1 = i_2 = 0, h \not\in i_3$ or equivalently $i_3 > 2^{(i_3)}$. Then $I$ is absolutely killed, that means $Q^I \in \text{Im}\mathcal{A}$.

Since $m(I) = 0$, the lemma is a special case of Lemma 5.10.

**Lemma 8.3.** Suppose $I$ belongs to case (b) of Lemma 8.1, that means $i_0 = 0, i_1 = 1, i_2 = 2^s - 1, s(i_3) = s = m(I), h \not\in i_3$. Then (the weak form of) the main lemma holds for $I$.

**Proof.** Let $s(h) = r > s$. We kill $I$ by the usual method:

\[ S^r z(i_3 - 2^r, 2^s - 1, 1, 0) = I + \sum J. \]

We use the notations similar as given in the proof of Lemma 7.7 with $a_i, b_i, c_i$ being respectively the coefficients of $2^i$ in the dyadic expansions of $i_3 - 2^r, i_2, i_1$.

By dimensional information, $h(J) < h(I)$ for all the $J$'s. So it suffices to show $m(J) \leq m(I)$ for any $J$. Suppose the contrary that $m(J) > m(I) = s$. Then $s(j_2) > s$. In particular, $\alpha_s(j_2) = 1$ meanwhile $\alpha_s(i_2) = 0$. Since $\alpha_s(i_3 - 2^r) = 0, i_2 = 2^r - 1, i_1 = 1, i_0 = 0$, the only way to fulfil the gap at $2^s$ of $i_2$ (in order to make $j_2$) is

\[ l_0 = \ldots = l_s = m_0 = 12 = |Q^I|. \]

Under this circumstance, $s(j_3) \leq s(i_3 - 2^r) = s$. Thus $m(J) \leq s$. This is a contradiction. The lemma is proved.
Lemma 8.4. Suppose \( I \) belongs to case (c) of Lemma 8.1, that means \( i_0 = 2, i_1 = 2^s - 1, i_2 = 2^{s'} - 2^s - 1 \) \((s' > s + 1)\), \( s(i_3) = s', h \not\in i_3 \). Then (the weak form of) the main lemma holds for \( I \).

Remark. Looking at the list of the allowed sequences for \( k = 4 \), we note that if \( i_0 = 2, i_1 = 2^s - 1, i_2 = 2^{s'} - 2^s - 1 \) \((s' > s + 1)\), \( s(i_3) = s' \), then one does not have \( h \gg i_3 \).

Proof. Let \( r = s(h) \geq s(i_3) = s' \), then \( i_3 = 2^r - 2^{s'} - 1 + x \), with \( x \neq 0 \) but \( x \equiv 0 \mod 2^{r+1} \). We kill \( I \) by the usual method

\[
Sq^{2^{r+3}}(i_3 - 2^r, 2^{s'} - 2^s - 1, 2^s - 1, 2) = I + \sum J.
\]

By dimensional information, \( h(J) < h(I) \) for all the \( J \)'s. We need only to check that \( m(J) \leq m(I) \) for any \( J \). Suppose the contrary that there is a \( J \) with \( m(J) > m(I) = s \).

Let us use the notations given in the proof of Lemma 9.7. In order to get \( J \) with \( m(J) > s \), one has to fulfil the two gaps of \( i_2 \) and \( i_1 \) at \( 2^s \). There are exactly three ways to do that:

1. \( a_5(h_3) = a_5(h_2) = 1 \),
2. \( a_5(t_{i_1}) = 1, a_5(h_2) = a_5(h_3) = 1 \),
3. \( a_5(t_{i_2}) = 1, a_5(h_1) = a_5(h_3) = 1 \).

Similarly as in the end of the proof of Lemma 9.7, one can see that the corresponding coefficients are 0. The lemma follows.

Lemma 8.5. Suppose \( I \) belongs to case (d) of Lemma 8.1. That means \( i_0 = 2, i_1 = 2^s - 1 = i_2, s(i_3) = \) either \( s \) or \( s + 1 \), \( h \not\in i_3 \). Then (the weak form of) the main lemma holds for \( I \).

Proof. If \( s(i_3) = s + 1 \), case (d) of Lemma 8.1 can be considered as a special case of case (c) with \( s' = s'(i_2) = s + 1 \). Actually, we did not use the hypothesis \( s' > s + 1 \) in the proof of Lemma 8.4. That proof still works for this lemma in the both cases \( s(i_3) = \) either \( s \) or \( s + 1 \).

Comment. Suppose \( r > s + 2 \). The sequence \( I = (2^r - 2^{s+1} - 2^s - 1, 2^s - 1, 2^s - 1, 2) \) with \( h \gg i_3 \gg i_2 \gg i_1 \gg i_0 \) is also controlled by Lemma 8.5. However, \( I = I' := (2^r - 2^{s+2} - 2^{s+1} - 1, 2^{s+1} + 2^s - 1, 2^s - 1, 2) \mod \text{Im} \mathcal{A} + \text{Span}\{J|m(J) < s\}, \) with \( m(I) = m(I') = s, h(I') > h(I') \). Note that \( I' \) belongs to case (d) of Lemma 7.1 and satisfies only the exceptional form of the main lemma. Lemma 9.6 will show that \( I = I' \) in \( \mathbb{Z}/2 \otimes_{\mathcal{A}} D_4 \) is an \( \mathcal{A} \)-generator for \( D_4 \).

Finally, we have

Proof of Lemma 5.9. Let \( s = s(h) \). We kill \( Q(I) \) by the usual method:

\[
Sq^{2^{s+3}}(i_3 - 2^s, i_2, i_1, i_0) = I + \sum J.
\]

Obviously, \( h(J) < h(I) \) for every \( J \) in the sum. So it suffices to show that \( m(J) \leq m(I) \) for any \( J \). Consider the two possibilities.
Case 1: $m(I) \geq s$. By the definition of $m(I)$,

$$m(I) = s = \min\{s(h), s(i_3), s(i_2), s(i_1 + i_0 - 2)\}.$$ 

Let $a_i, b_i, c_i, d_i$ be the coefficients of $2^i$ in the 2-adic expansions of $i_3 - 2^s, i_2, i_1, i_0$ respectively. We have

$$S^2e^3(i_3 - 2^s, i_2, i_1, i_0) = \sum (S^k_0 Q_3)^{a_0} \cdots (S^k_0 Q_3)^{2^n a_n}. $$

Let $J$ be a term in the above sum ($J \neq I$), so

$$j_3 = i_3 - 2^s - \sum_i a_i 2^i + y.$$ 

Here the sum runs over all $i$ with $k_i = 4, 6$ or $7$ and $0 < y < 2^s$ by dimensional information. Note that $a_s = a_s(f_a - 2^s) = 0$. Hence $m(J) \leq s_j \leq s = m(I)$.

Case 2: $m := m(I) < s$. Since $s(h) < s(i_3)$, by Lemma 5.1,

$$s(h) = s(i_2 + i_1 + i_0 - 1).$$

Suppose $s(i_2) \neq s(i_1 + i_0 - 2)$ then $s(h) = \min\{s(i_2), s(i_1 + i_0 - 2)\} = m(I)$. This contradicts $m(I) < s$.

So $s(i_2) = s(i_1 + i_0 - 2) = m(I)$, as a consequence $s(h) = m + 1 = s$, where $m = m(I)$. We have $h(J) = h - 2^s + x$ ($0 \leq x < 2^s$).

(a) If $x \neq 0$ mod $2^{m+1}$, then $m(J) \leq s(h(J)) \leq m = m(I)$.

(b) Suppose $x \equiv 0$ mod $2^{m+1}$.

So $x = 0$. Otherwise $x \geq 2^{m+1} = 2^s$. However, we do not use this fact below.

If $\alpha_m(j_2) = 0$, then $m(J) \leq s(j_2) \leq m = m(I)$.

Otherwise, $\alpha_m(j_2) = 1$. Combining this with the facts $s(i_3) = m, x \equiv 0$ mod $2^{m+1}$, we see that $2^m$ is moved from $i_3 - 2^s$ to $j_2$. That means $k_m = 4, \alpha_m(j_3) = 0$. Therefore $m(J) \leq s(j_3) \leq m = m(I)$. The lemma is proved.

9. The Linear Independence of the Allowed Monomials

The purpose of this section is to show the linear independence of the allowed monomials in $D_k = \mathbb{Z}/2 \otimes_{\mathcal{A}} D_k$ for $k = 3$ or 4.

**Definition 9.1.** Suppose $a$ is a non-negative integer. Set

$$\ell(a) = \begin{cases} 
\min\{\min i | \alpha_i(a) \neq 0\} & \text{for } a > 0, \\
+\infty & \text{for } a = 0.
\end{cases}$$

**Proposition 9.2.** Let $k = 3$. Then the allowed monomials are linearly independent in $D_3 = \mathbb{Z}/2 \otimes_{\mathcal{A}} D_3$. 

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Proof. Suppose $I$, $J$ are of length 3 with $I$ allowed, and
\[ Sq^{2^l}Q^l = c \cdot Q^l + \text{other terms}. \]
Then we will show that $c = 0$. The proposition follows.

We are going to describe theoretically the coefficient $c$. Let $a_i$, $b_i$, $c_i$ be the coefficients of $2^i$ in the dyadic expansions of $j_2$, $j_1$, $j_0$ respectively. We have
\[ Sq^{2^l}Q^l = \sum (Sq^{q_0}Q_2)^{a_0} \cdots (Sq^{q_l}Q_2)^{2^l a_l}, \]
\[ \cdot (Sq^{q_0}Q_1)^{b_0} \cdots (Sq^{q_l}Q_1)^{2^l b_l}, \]
\[ \cdot (Sq^{q_0}Q_0)^{c_0} \cdots (Sq^{q_l}Q_0)^{2^l c_l}. \]
Here the sum runs over all the decompositions
\[ 2^l = \sum_i k_i 2^i + \sum_i l_i 2^i + \sum_i m_i 2^i, \]
with $k_i = 0$ if $a_i = 0$, $l_i = 0$ if $b_i = 0$ and $m_i = 0$ if $c_i = 0$.

Pick up a term $Q^l$ in the sum. Define for this term:
- $t_{ij}$ to be the number of $Q_i$, which is sent to $Q_j$ ($i > j$),
- $d_{ij}$ to be the number of $Q_i$, which is sent to $Q_iQ_j$ ($i > l > j$),
- $h_i$ to be the number of $Q_i$, which is sent to $Q_iQ_j$ for any $j < i$.
Obviously one gets
\[ t_2 = t_2 + A_2 + d_{10} + t_{21} - t_{20}, \]
\[ t_1 = t_1 + A_1 - d_{10} + t_{20} - t_{10}, \]
\[ t_0 = t_0 + h_{0} + d_{10} + t_{20} + t_{10}, \]
where all the variables are non-negative integers. They satisfy the Dimensional Equation:
\[ (DE) \quad 2^l = 4A_2 + 6A_1 + 7A_0 + 5d_{10} + 2t_{21} + 3t_{20} + t_{10}. \]
By an easy argument of combination, the coefficient $c$ is determined by $c = \sum_\sigma c(\sigma)$, where
\[ c(\sigma) = \binom{j_0}{h_0} \binom{j_1}{t_{10}} \binom{j_1-t_{10}}{d_{20}} \binom{j_0+j_1-h_0-t_{10}-d_{20}}{h_1} \binom{j_2}{t_{20}}, \]
\[ \cdot \binom{j_2-t_{20}}{t_{21}} \binom{j_0+j_1+j_2-h_0-t_{10}-d_{20}-h_1-t_{20}-t_{21}}{h_2}. \]
Here the sum runs over all the partitions $\sigma$ satisfying (DE).

The list of the allowed sequences is given just after Theorem 2.6. According to this, as $I$ is allowed, $i_0 = 0$ or 1. By Corollary 3.2, $j_0 \leq i_0$. So we need only to consider the three cases.

Case 1: $i_1 = i_0 = 0$. By Corollary 3.2, from $i_0 = 0$ it implies $j_0 = 0$ and $h_0 = d_{10} = t_{20} = t_{10} = 0$. Also by the corollary, $j_1 + j_0 \leq i_1 + i_0 = 0$, hence $j_1 = 0$, $h_1 = t_{21} = 0$. We get $i_2 = j_2 + h_2 = 2^l - 1$. Then
\[ c(\sigma) = \binom{j_0+j_1+j_2}{h_2} = \binom{i_2-h_2}{h_2} = \binom{2^l-1-h_2}{h_2} = 0, \]
as $h_2 \neq 0$ (otherwise, $I = J$, $2^l = 0$, a contradiction).
Case 2: $i_0 = 1$, $j_0 = 0$. Note that $Q_0$ is the only invariant among $Q_2$, $Q_1$, $Q_0$, which is of odd dimension. So $2^i \equiv 1 \pmod{2}$. That means $2^i = 1$. Also by dimensional information, the only solution $J$ in this case is $J = (i_2, i_1 + 1, i_0 - 1)$. We have

$$Sq^1 Q^{(i_2, i_1 + 1, i_0 - 1)} = \binom{i_1 + 1}{1} Q^i = \binom{2^i}{1} Q^i = 0,$$

because of $2^i > 1$ as $i_1 = 2^i - 1 > i_0 = 1$.

Case 3: $i_0 = 1$, $j_0 = 1$. Obviously $h_0 = d_{i_0} = t_{i_2} = tw = 0$, $a_0 = i_0 - h_1 - t_{i_1}$, $J = (J', 0)$ with $J', J''$ are of length 3. By dimensional information, $2^i$ is even. The lemma follows from the proof of Proposition 9.2 and the fact that $D_4/(Q_0) = D_3$ over the “halving” map of the Steenrod algebra, which sends $Sq^{2i}$ to $Sq^i$ and $Sq^{2i+1}$ to 0.

The proposition is proved.

From now on, we deal with the length $k = 4$ only.

Proposition 9.3. The allowed monomials are linearly independent in $D_k = \mathbb{Z}/2 \otimes_{\mathcal{A}} D_k$ for $k = 4$.

According to the list of the allowed sequences for $k = 4$ given just after Theorem 2.7, the proposition is proved by showing Lemmas 9.4 and 9.6 below.

Lemma 9.4. Suppose $I$ is strongly allowed with $j_0 = 0$ and

$$Sq^i Q^i = c \cdot Q^i + \text{other terms}.$$  

Then $c = 0$.

Proof. From $i_0 = 0$ it implies $j_0 = 0$. Then $I = (I', 0), J = (J', 0)$ with $I', J'$ are of length 3. By dimensional information, $2^i$ is even. The lemma follows from the proof of Proposition 9.2 and the fact that $D_4/(Q_0) \cong D_3$ over the “halving” map of the Steenrod algebra, which sends $Sq^{2i}$ to $Sq^i$ and $Sq^{2i+1}$ to 0.

Definition 9.5. Let $M$ be an $\mathcal{A}$–module. Suppose

$$Sq^i(x) = a + \text{other terms},$$

for $x, a \in M$. Then we say $a$ is hit by $Sq^i$.

Suppose given $r > s + 1 > 2$. Set

$$I_n = (2^r - 2^{s+n} - 2^{s+n-1} - 1, 2^{s+n} - 2^s - 1, 2^s - 1, 2),$$

for $1 \leq n < r - s$. Note that it is weakly allowed for $n = 1, r > s + 2$. 

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Lemma 9.6. (a) If I is strongly allowed with \( i_0 = 2 \), then there is no way to hit I by any Steenrod square of positive dimension.

(b) There are exactly two ways to hit \( I_n \) for \( 1 < n < r - s - 1 \), namely

(i) \( Sq^{2s+n-1}(2^r - 2s - 1, 2s+n, 2s+n-1, 2s-1) = I_{n-1} + I_n + \text{other terms} \),

(ii) \( Sq^{2s+n-2}(2^r - 2s+n-1, 2s+n-2s - 1, 2s-1, 2s-1) = I_n + I_{n+1} + \text{other terms} \).

There is exactly one way to hit \( I_n \) for \( n = 1 \) or \( n = r - s - 1 \). It is (ii) for \( n = 1 \) and (i) for \( n = r - s - 1 \).

In order to prove Lemma 9.6, we need

Lemma 9.7. Suppose either I is allowed with \( i_0 = 2 \) or \( I = I_n \) (\( 1 < n < r - s \)). If

\[ Sq^2 Q^j = Q^j + \text{other terms} \]

then \( j_0 = i_0 = 2 \), \( j_1 = i_1 = 2^s - 1 \).

Now suppose that Lemma 9.7 has been proved. Then we show Lemma 9.6 as follows.

Proof of Lemma 9.6 (supposing Lemma 9.7 is proved).

We prove part (b) only. Part (a) can similarly be showed.

Suppose \( I = I_n \) and \( Sq^2 Q^j = c \cdot Q^j + \text{other terms} \). If \( c = 1 \), then according to Lemma 9.7, \( j_0 = i_0 = 2 \), \( j_1 = i_1 = 2^s - 1 \). Define \( h_i \), \( t_{ij} \), \( d_{ij} \) as in the proof of Proposition 9.2 and have:

\[
\begin{align*}
(i_0) & \quad j_3 = j_3 + h_3 + d_3 + d_{30} + d_1 - t_{32} - t_{31} - t_{30} \\
(ii) & \quad j_2 = j_2 + h_2 + d_2 - d_{32} - t_{32} + t_{21} - t_{20} \\
(iii) & \quad j_1 = j_1 + h_1 + d_1 - d_{31} - t_{31} + t_{21} - t_{10} \\
(iv) & \quad j_0 = j_0 + h_0 + d_0 - t_{30} + t_{20} + t_{20} - t_{10}.
\end{align*}
\]

However, we do not need to use this equation system in the general form. In fact, when \( j_0 = i_0 = 2 \), \( j_1 = i_1 = 2^s - 1 \), we get \( h_0 = h_1 = t_{10} = t_{20} = t_{30} = t_{31} = d_0 = d_{30} = d_{31} = 0 \), and \( i_3 = j_3 + h_3 - t_{32} \), \( i_2 = j_2 + h_2 + t_{32} \).

Now the Dimensional Equation becomes

\[ (DE) \quad 2^r = 12 h_2 + 8 h_3 + 4 t_{32} \]

We also have \( c = \sum \sigma c(\sigma) \), with

\[ c(\sigma) = \frac{(2 + i_1 + i_2 - h_2 - t_{32}) (i_3 - h_3 + t_{32})}{h_2 t_{32}} \left(\frac{2 + i_1 + i_2 + i_3 - 2 h_2 - h_3 - t_{32}}{h_3}\right). \]

As \( c = 1 \), there exists at least one \( \sigma \) with \( c(\sigma) \neq 0 \). Fix such a \( \sigma \).

Since \( (2^r+1+i_2-h_2-t_{32}) \neq (\frac{2^r+1+i_2-h_2-t_{32}}{h_2} \neq 0 \), so either \( h_2 = 0 \) or \( \ell(h_2) \neq \ell(t_{32}) \).

Combine this with (DE) and get either \( h_3 = 0 \) or \( \ell(h_3) < \min\{\ell(h_2), \ell(t_{32})\} \).

(1) Suppose \( h_3 = 0 \). Combining \( (i_3-h_3+t_{32}) = (2^r-2^{s+n+1}-t_{32}) \neq 0 \) with the fact that \( t_{32} = t_2 = 2^{s+n} - 2^s - 1 \), we get either \( t_{32} = 0 \) or \( t_{32} = 2^{s+n-1} \).

• If \( t_{32} = 0 \), then \( h_3 = t_{32} = 0 \), \( 2^r = 12 h_2 \). This is a contradiction.
Suppose \( t_{32} = 2^{i+n-1} \). Note that \( \frac{(2^{i+i_1+i_2-h_2-h_{32}})}{h_2} = \frac{(2^{i+n-1-h_2})}{h_2} \neq 0 \) implies \( h_2 < t_{32} = 2^{i+n-1} \). Combining this with (DE) \( 2^t = 12h_2 + 4t_{32} \), we get \( h_2 = 0 \).

The solution \( h_2 = h_3 = 0 \), \( t_{32} = 2^{i+n-1} \) gives rise to equation (i), provided \( n > 1 \).

\( 2 \) If \( h_3 \neq 0 \), then as showed before \( \ell(h_3) < \min\{\ell(h_2), \ell(t_{32})\} \).

Combine this with \( \frac{(2^{i+i_1+i_2-h_2-h_{32}})}{h_3} = \frac{(2^{i+n-1-1}-2h_2-h_{32})}{h_3} \neq 0 \) and get \( \ell(h_3) = s + n - 1 \).

If either \( h_2 > 0 \) or \( t_{32} > 0 \), then by using \( \ell(h_3) < \min\{\ell(h_2), \ell(t_{32})\} \) one obtains \( h_2 \geq 2^{i+n} > i_2 \) or \( t_{32} \geq 2^{i+n} > i_2 \) respectively. They both are contradictions.

Thus \( h_2 = t_{32} = 0 \). Combining (DE) and \( \ell(h_3) = s + n - 1 \) one gets \( h_3 = 2^{i+n-1} \). This solution gives rise to equation (ii), provided that \( n < r - s - 1 \).

The lemma is proved.

Proof of Lemma 9.7.

Step 1: To prove \( j_0 = i_0 = 2 \).

We start with a very simple but useful observation that if \( I \) is as in the lemma, then \( s(i_3) \geq s = s(i_2) = s(i_1) > 1 \). Consider the two cases:

\( 1 \) \( i_0 = 2, \ j_0 = 1 \). Again, by dimensional information we get \( 2^t = \dim Q^l - \dim Q^l \equiv 1 \mod 2 \). That means, \( 2^t = 1, \ t = 0 \). The unique solution is \( J = (i_3, i_2, i_1 + 1, i_0 - 1) \). However,

\[
S q^l Q^l = \begin{pmatrix} i_1 + 1 \\ i_2 + 1 \\ i_1 + 1 + 1 \\ i_0 - 1 \end{pmatrix} Q^l = \begin{pmatrix} 2^i \\ 1 \end{pmatrix} Q^l = 0,
\]

because of \( 2^i - 1 \geq 2 \) or equivalently \( 2^i > 3 \). This is a contradiction.

\( 2 \) \( i_0 = 2, \ j_0 = 0 \). First, we claim that \( \alpha_0(j_1) = \alpha_0(i_1) = 1 \).

Suppose the contrary that \( \alpha_0(j_1) = 0 \). It implies \( \alpha_0(t_{10}) = \alpha_0(d_{20}) = \alpha_0(d_{30}) = 0 \). We observe \( \alpha_0(h_1) = 0 \) as \( \alpha_0(j_0 + j_1) = \alpha_0(0 + 2) = 0 \) or equivalently \( \ell(a_{20}) = 0 \). We get

\[
1 = \alpha_0(j_1) = \alpha_0(j_1) + \alpha_0(h_1) + \alpha_0(d_{31}) + \alpha_0(t_{21}) + \alpha_0(t_{31}) \mod 2.
\]

If \( \alpha_0(t_{21}) = 1 \), then \( \alpha_0(h_2) = 0 \). The only way to hit the new gap of \( j_2 \) at \( 2^0 \) is \( \alpha_0(j_3) = \alpha_0(t_{32}) = 1 \). However, it makes \( \alpha_0(i_3) = 0 \). It contradicts the fact that \( s(i_3) > s > 1 \). Therefore, \( \alpha_0(t_{21}) = 0 \). Thus \( \alpha_0(d_{31}) + \alpha_0(t_{31}) = 1 \).

\( a \) If \( \alpha_0(t_{31}) = 1 \), then \( \alpha_0(d_{31}) = \alpha_0(t_{21}) = 0 \). The only way to hit the new gap of \( h_3 \) at \( 2^0 \) is \( \alpha_0(h_3) = 1 \), \( \alpha_0(j_3) = 1 \). Now we claim that \( \alpha_1(j_1) = 0 \neq \alpha_1(i_1) = 1 \). Otherwise, if \( \alpha_1(j_1) = 1 \), one easily gets

\[
\begin{align*}
\dim Q^l_0 - \dim Q^l_1 & \equiv 2(2^a - 1) \equiv 2^2 - 2 \mod 2^3, \\
\dim Q^l_1 - \dim Q^l_2 & \equiv 2^4 - 2 \mod 2^3, \\
\dim Q^l_2 - \dim Q^l_3 & \equiv 0 \mod 2^3, \\
\dim Q^l_3 & \equiv 0 \mod 2^3.
\end{align*}
\]

So \( 2^l = \dim Q^l - \dim Q^l \equiv 2^5 + 2^4 - 2^2 \mod 2^3 \). Hence \( 2^l = 2^2 = 4 \). On the other hand, \( 2^l \geq 6\alpha_0(t_{31}) = 6 \). This is a contradiction.
Now we have $\alpha_1(j_1) = 0$. It implies $\alpha_1(d_{20}) = \alpha_1(d_{30}) = \alpha_1(t_{10}) = 0$. Besides, since $j_0 = 0$, so $\alpha_1(h_0) = 0$. Then

$$1 = \alpha_1(i_0) - \alpha_1(j_0) = 0 + 0 + 0 + \alpha_1(t_{20}) + \alpha_1(t_{30}) + \alpha_1(d_{30}).$$

Hence, either $\alpha_1(t_{20}) = 1$ or $\alpha_1(t_{30}) = 1$ or $\alpha_1(d_{30}) = 1$.

- $\alpha_1(t_{20}) = 1$: The only way to hit the gap $2^1$ of $j_1$ is $\alpha_1(i_3) = \alpha_1(j_3) = 1$. It makes a gap of $i_3$ at $2^1$, a contradiction.
- $\alpha_1(t_{30}) = 1$: The only way to hit the gap $2^1$ of $j_1$ is $\alpha_1(j_2) = \alpha_1(t_{21}) = 1$. It makes a gap of $i_2$ at $2^1$, a contradiction.
- $\alpha_1(d_{30}) = 1$: The only ways to hit the gaps $2^1$ of $j_2$ and $j_1$ are respectively $\alpha_1(t_{32}) = 1$, $\alpha_1(t_{31}) = 1$. However, each from these equalities damages the other, a contradiction.

(b) If $\alpha_0(d_{31}) = 1$, then $\alpha_0(t_{21}) = \alpha_0(t_{31}) = 0$. The only possibility to hit the new gap of $j_2$ at $2^0$ is $\alpha_0(j_3) = \alpha_0(t_{32}) = 1$.

Again, we claim that $\alpha_1(j_1) = 0$. Essentially, we can repeat the argument in (a) to get a contradiction for this case.

We have just showed that $\alpha_0(j_1) = \alpha_0(i_1) = 1$. Combine this with $i_0 = 2$, $j_0 = 0$ and get $2^t = \dim Q^I - \dim Q^J = 2 \bmod 4$. So $2^t = 2$. The unique solution is $J = (i_3, i_2, i_1 + 2, 0)$. However,

$$Sq^2Q^J = \binom{i_1 + 2}{2}Q^I = \binom{2^s + 1}{2}Q^I = 0,$$

as $2^s - 1 > 2$ or equivalently $2^s > 3$.

Step 2: To prove that $j_1 = i_1 = 2^t - 1$.

Suppose the contrary that $j_1 \neq i_1$, that means $j_1 < i_1$ (as $j_1 + j_0 \leq i_1 + i_0$ by Corollary 3.2 and $j_0 = i_0$). Set $a = s(j_1) < s$. (Indeed, if $a = s$, then $j_1 \geq 2^t - 1 = i_1$, hence $j_1 = i_1$.)

Now we claim that $s(j_2) \geq a$. Suppose the contrary $s(j_2) = u < a$. Let us compare $s(j_3)$ and $s(j_2)$.

If $s(j_3) < s(j_2)$, then the only way to hit the gap $2^{s(j_3)}$ of $j_3$ is to increase $h_3$ by $2^{s(j_3)}$. However, the coefficient should be a multiple of

$$\binom{2^{s(j_3)} - 1 + 2^{s(j_3)} - 1 + 2^{s(j_1)} - 1 + 2}{2^{s(j_3)}} = \binom{2^{s(j_3)} - 1 + 2^u + 2^a}{2^{s(j_3)}} = 0,$$

as $s(j_3) < u < a$, a contradiction.

If $s(j_3) = s(j_2) = u$, then the only way to hit the gaps $2^u$ of $j_3$ and $j_2$ is to increase $h_2$, $h_3$ both by $2^u$. The coefficient should be a multiple of

$$\binom{2^u - 1 + 2^a - 1 + 2}{2^u} \binom{2^u - 1 + 2^u - 1 + 2^a - 1 + 2 - 2^u}{2^u} = \binom{2^u + 2^a}{2^u} \binom{2^u - 1 + 2^a}{2^u} = 0,$$

as $u < a$, a contradiction.
We have just showed that if \( s(j_2) < s(j_1) \), then \( s(j_3) > s(j_2) \). Recall that 
\[ s(i_3) \geq s = s(i_2) = s(i_1) > a. \]
Consider the differences of dimensions:

\[
\begin{align*}
\dim Q_1^j - \dim Q_1^i &\equiv 2^{a+1} \mod 2^{a+2}, \\
\dim Q_2^j - \dim Q_2^i &\equiv 2^{a+2} \mod 2^{a+3}, \\
\dim Q_3^j - \dim Q_3^i &\equiv 0 \mod 2^{a+3}.
\end{align*}
\]

However, \( 2' = \dim Q^j - \dim Q^i \), so \( 2^{a+2} = 2^{a+1} \), or equivalently \( u = a - 1 \).

(1) If \( \alpha_{a+1}(j_1) = \alpha_{a+1}(i_1) = 1 \), we show that \( \alpha_a(j_2) \neq \alpha_a(i_2) = 1 \). Otherwise, if \( \alpha_a(j_2) = \alpha_a(i_2) = 1 \), then

\[
2' = \dim Q^j - \dim Q^i \equiv 2^{a(4^2 - 2)} + 2^{a+1}(2^{a+2} - 2^2) \mod 2^{a+3} \equiv 2^{a+4} + 2^{a+2} \mod 2^{a+3}.
\]

It implies \( 2' = 2^{a+2} \). On the other hand \( 2' > \dim Q_1^j - \dim Q_1^i \geq 2^a(2^{a+2} - 2) > 2^{a+2} \). This is a contradiction.

Since \( i_0 = j_0 \), so \( d_2^{d_2} = d_3^{d_3} = 0 \), and \( i_2 = j_2 + h_2 - d_2^{d_2} + t_{32} - t_{21} - t_{20} \). At level \( 2a^-1 \) we have \( \alpha_{a-1}(d_2^{1}) = \alpha_{a-1}(t_{21}) = \alpha_{a-1}(t_{20}) = 0 \), as \( \alpha_{a-1}(j_2) = 0 \).

Hence \( 1 = \alpha_{a-1}(i_2) = \alpha_{a-1}(h_2) + \alpha_{a-1}(t_{32}) \). There are two possibilities:

(A) \( \alpha_{a-1}(h_2) = 0 \), \( \alpha_{a-1}(t_{32}) = 1 \) it requires \( \alpha_{a-1}(h_3) = 1 \),

(B) \( \alpha_{a-1}(h_2) = 1 \), \( \alpha_{a-1}(t_{32}) = 0 \).

At level \( 2a \), there are three possibilities to hit the two gaps of \( j_1 \) and \( j_2 \):

(a) \( \alpha_j(t_{32}) = 1 \), \( \alpha_j(h_1) = 1 \), \( \alpha_j(h_2) = 0 \), \( \alpha_j(h_3) = 1 \),

(b) \( \alpha_j(t_{31}) = 1 \), \( \alpha_j(h_1) = 0 \), \( \alpha_j(h_2) = 1 \), \( \alpha_j(h_3) = 1 \),

(c) \( \alpha_j(t_{31}) = 0 \), \( \alpha_j(t_{32}) = 0 \), \( \alpha_j(h_1) = 1 \), \( \alpha_j(h_2) = 1 \).

(\( \alpha_a(j_3) \) is either 0 or 1.) Now we have to compute the coefficient \( c \) in the
combined six cases: (Aa) (Ab) (Ac) (Ba) (Bb) (Be). In each case, \( c \) is a multiple
of the following number:

(Aa) \( \left(2^2+1 \right)^{2^2+1} \left(2^2-1+2^2-1+2^2-1+2^2-1+2^2-1\right) = \left(2^2+1 \right)^{2^2-1+2^2-1} = 0. \)

(Bb) \( \left(2^2+1 \right)^{2^2+1} \left(2^2-1+2^2-1+2^2-1+2^2-1+2^2-1\right) = \left(2^2+1 \right)^{2^2-1+2^2-1} = 0. \)

(2) If \( \alpha_{a+1}(j_1) = 0 \neq \alpha_{a+1}(i_1) = 1 \), then similarly as in (1), by computing
the difference of dimensions \( 2a + 3 \), one gets \( \alpha_a(j_2) = \alpha_a(i_2) = 1 \). Also
by an argument similar to that in (1), one can see the coefficients of all the
possibilities are 0.

In any case, the hypothesis \( s(j_2) < s(j_1) \) leads to a contradiction. Therefore,
if \( s(j_1) = a < s \), then \( s(j_2) \geq a \).

Next, we claim that \( s(j_3) \geq a \). This can be showed by the same argument as
in the proof that if \( s(j_2) < s(j_1) \), then \( s(j_3) > s(j_2) \).

By dimensional information \( 2' = \dim Q^j - \dim Q^i \equiv 2^a(2^{a+2} - 2) \mod 2^{a+3} \equiv 2^a \mod 2^{a+2} \). So \( 2' = 2^{a+1} \). On the other hand, \( 2' > \dim Q_1^j - \dim Q_1^i \geq 2^a(2^{a+2} - 2) \). This contradiction finishes the proof of Lemma 9.7.
10. Final remarks

Conjecture on gaps 10.1. Suppose $Sq^rQ^j = Q^j +$ other terms. One gets

(a) $m(J) \geq m(I)$.

(b) If additionally

$$j_0 = i_0, \ldots, j_{n-1} = i_{n-1} \quad (2 \leq n < k),$$

then

$$\min \{s(h(J)), s(j_{k-1}), \ldots, s(j_n)\} \geq \min \{s(h(I)), s(i_{k-1}), \ldots, s(i_n)\}.$$  

Lemma 7.4 is a special case of this conjecture.

If the conjecture is true, then we can simplify the proofs of some lemmas, for instance, Lemmas 7.4, 9.6 and 9.7. Furthermore, it would allow us to show that the allowed monomials are linearly independent in $D_k = \mathbb{Z}/2 \otimes_{\mathfrak{A}} D_k$ for any $k$.

In general, the allowed monomials do not generate $D_k$ for $k > 4$.

Example 10.2. If $k = 12$, the sequence

$$I = (2^{s+4} - 1, 2^{s+1} + 2^s - 1, 2^{s+2} + 2^s - 1, 2^s - 1, \ldots, 2^s - 1, 10)$$

(with $s > 3$)

is not allowed as $2^{s+1} + 2^s - 1 \neq 2^{s+2} + 2^s - 1$ and also $2^{s+1} + 2^s - 1 \neq 2^{s+2} + 2^s - 1$. However, a complicated computation shows that, $I$ does not belong to the subspace spanned by the allowed sequences. It should be considered as a new generator.

Conjecture 10.3. The allowed monomials form a basis for $D_k = \mathbb{Z}/2 \otimes_{\mathfrak{A}} D_k$ for $k < 12$.

References


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