Table of Factors of $2^{p}-1: \underset{\text { factor }}{p}$-Continued

| 9109 | 9127 | 9137 | 9157 | 9173 | 9181 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 4955297 | 146033 | 1005071 | 3076753 | 495343 | 1377151 |
| 9221 | 9283 | 9311 | 9323 | 9337 | 9341 |
| 719239 | 5941121 | 1415273 | 895009 | 2614361 | 74729 |
| 9343 | 9371 | 9391 | 9397 | 9403 | 9419 |
| 149489 | 18743 | 93911 | 225529 | 2576423 | 18839 |
| 9421 | 9431 | 9461 | 9479 | 9491 | 9497 |
| 8573111 | 679033 | 75689 | 18959 | 531497 | 531833 |
| 9511 | 9521 | 9539 | 9601 | 9613 | 9619 |
| 95111 | 361799 | 19079 | 2285039 | 57679 | 615617 |
| 9791 | 9811 | 9829 | 9833 | 9851 | 9859 |
| 19583 | 7770313 | 707689 | 4562513 | 78809 | 1104209 |
| 9883 | 9949 | 9973 |  |  |  |
| 158129 | 8098487 | 299191 |  |  |  |

## A Computation of Some Bi-Quadratic Class Numbers

## By Harvey Cohn

A fascinating chapter in computational number theory began when Lagrange showed that every positive integer is representable as the sum of at most four perfect squares [1]. Clearly three would not suffice in every case, as $7=2^{2}+1^{2}+1^{2}+1^{2}$ would be an exception; nevertheless, the problem of expressing some positive integer $n$ as the sum of at most three squares soon achieved a very special role. For, Gauss showed that $r(n)$, the number of such representations, is connected in a very simple way with the much studied (intrinsically positive) class number, $h$, of the field generated by $\sqrt{-n}$. Specifically for $n$ square-free (and $n \neq 1,3$ where $h=1$ ),

$$
\begin{equation*}
r(n)=g h \tag{1}
\end{equation*}
$$

where $g=12$ for $n \equiv 1,2,5,6$ and $g=24$ for $n \equiv 3(\bmod 8)$. Thus we could conclude the existence of at least one such representation for the indicated $n$. Gauss and later, Kronecker, reversed the direction of these equations by making large scale tabulations of $h$ from $r(n)$, although, unfortunately, no location for Kronecker's alleged tabulation (for odd $n$ up to 10,000 ) seems to exist in the literature. In tallying the representation $n=x_{1}{ }^{2}+x_{2}{ }^{2}+x_{3}{ }^{2}$ it might be noted that one must count each ordered triple ( $x_{1}, x_{2}, x_{3}$ ) of positive, negative, or zero integers as a separate unit, so that as much as $2^{3} \cdot 3!=48$ could be contributed to $r(n)$ when such a decomposition into squares is expressed as triples.

In more recent times, the representation theory was extended to integers in the field $k$ generated by $\sqrt{5}$, i.e., to the quantities $\mu=(a+b \sqrt{5}) / 2$ where $a$ and $b$ are of the same parity. Here we seek to represent, necessarily, only those integers $\mu$ which are positive together with their conjugate (i.e., totally positive). Thus, e.g., $a>|b \sqrt{5}| \geq 0$. The special surd $\sqrt{5}$ must be used because then, as Götzky

[^0]showed [2], each totally positive integer in $k$ could be expressed as the sum of the squares of at most four integers in $k$. Later, Maass [3] made the more remarkable discovery that at most three squares would always suffice; in fact he arrived at a formula analogous to that of Gauss. Since Maass' formula is the basis of a machine calculation, we avoid irrelevant complexities by making certain further assumptions. First of all $\mu$ is to be free from square integral divisors in $k$ except for powers of $(\sqrt{5}+3) / 2=[(\sqrt{5}+1) / 2]^{2}=\epsilon$. Secondly $a \geq 5 b \geq 0$, since if $5 b>a>\sqrt{5} b$, we can continually replace $\mu$ by $\mu / \epsilon$. Then $R(\mu)$, the number of representations of $\mu$ as the sum of three squares, is linked to an intrinsically positive quantity $H$, namely the class number of the bi-quadratic field generated by $\sqrt{5}$ and $\sqrt{-\mu}$, by means of the following formula (which excludes $\mu=1,(5+\sqrt{5}) / 2$, and 3 where $H=1$ ):
\[

$$
\begin{equation*}
H=R(\mu) / G \tag{2}
\end{equation*}
$$

\]

Here $G=12$ when $(a, b) \not \equiv(1,3),(1,5),(2,4),(5,1),(5,7)$, or $(6,0) \bmod 8$ (actually, when $\eta^{2}+\mu \equiv 0 \bmod 4$ is unsolvable for $\eta$ in $k$ ). Otherwise, $G=120$ except when $(a, b) \equiv(1,5),(1,11),(5,7),(5,9),(6,0),(9,3),(9,13),(13,1)$, $(13,15)$, or $(14,0) \bmod 16$ (actually, when $\eta^{2}+\mu \equiv 0 \bmod 8$ is solvable for $\eta$ in $k$ ); in these cases, $G=96$.

A tabulation of $R$ and $H$ for 446 selected values of $\mu=[a, b]=(a+b \sqrt{5}) / 2$ was made on the stored program electronic computer, the IBM 650. The values of $a, b$ were selected with the restrictions

$$
100>a \geq 5 b \geq 0
$$

and that $\mu$ have (except trivially for powers of $\epsilon$ ) no square divisors in $k$; or, in terms of ordinary integral arithmetic, the condition is that $d$, the g.c.d. of ( $\frac{1}{2}[a+b], \frac{1}{2}[a-b]$ ), be relatively prime to $5\left(=(\sqrt{5})^{2}\right)$, and both $d$ and $\left(a^{2}-5 b^{2}\right) / 4 d^{2}$ be square free. The machine assembled 446 such couples automatically into the highest four decimal positions of 446 individual ten digit storage locations, in lexicographic order.

The machine next tallied the decompositions $\mu=\xi_{1}{ }^{2}+\xi_{2}{ }^{2}+\xi_{3}{ }^{2}$ where $\xi_{i}=\left[a_{i}, b_{i}\right], a_{i} \equiv b_{i} \bmod 2$. Here the three couples $\xi_{i}$ were scanned in lexicographic order, with the restriction $0 \leq b_{i} \leq 7$, while $a_{i} \geq 0$ and $a^{\prime}<100$, (see below). Thus, taking all sign possibilities, with

$$
\left\{\begin{array}{l}
a^{\prime}=\left[a_{1}^{2}+a_{2}^{2}+a_{3}{ }^{2}+5\left(b_{1}{ }^{2}+b_{2}{ }^{2}+b_{3}{ }^{2}\right)\right] / 2 \\
b^{\prime}= \pm a_{1} b_{1} \pm a_{2} b_{2} \pm a_{3} b_{3}
\end{array}\right.
$$

the couples $a^{\prime}, b^{\prime}$ were constructed and compared with the 446 cases stored in the memory. Whenever a matching entry was located the count was augmented and accumulated in the last six decimal positions of the memory word. It might be appropriate to mention that the IBM 650 has a special "table look-up" operation that searches the memory at high speed for the appropriate entry. Without such an instruction the search would have had to be programmed with a considerable loss of running time.

In the final phase the "words" $a, b, R(\mu)$ were unpacked and the congruences were examined automatically to calculate $H$ and to produce the output consisting of one IBM card per value of $\mu$. (See attached table).

Tabulation of $R$ and $H$ for $\mu=(a+b \sqrt{5}) / 2$

| $a$ | $b$ | $R$ | $H$ | $a$ | $b$ | $R$ | $H$ | $a$ | $b$ | $R$ | H | $a$ | $b$ | $R$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | 0 | 6 | 1 | 33 | 3 | 240 | 2 | 47 | 9 | 144 | 12 | 59 | 1 | 144 | 12 |
| 4 | 0 | 12 | 1 | 33 | 5 | 288 | 3 | 48 | 2 | 216 | 18 | 59 | 3 | 168 | 14 |
| 5 | 1 | 24 | 1 | 34 | 0 | 96 | 8 | 48 | 4 | 168 | 14 | 59 | 5 | 96 |  |
| 6 | 0 | 32 | 1 | 34 | 2 | 96 | 8 | 48 | 6 | 144 | 12 | 59 | 7 | 96 |  |
| 7 | 1 | 24 | 2 | 34 | 4 | 120 | 1 | 49 | 1 | 144 | 12 | 59 | 9 | 144 | 12 |
| 9 | 1 | 24 | 2 | 34 | 6 | 72 | 6 | 49 | 3 | 240 | 2 | 59 | 11 | 144 | 12 |
| 10 | 2 | 24 | 2 | 35 | 1 | 96 | 8 | 49 | 5 | 288 | 3 | 60 | 2 | 144 | 12 |
| 1 | 1 | 24 | 2 | 35 | 3 | 144 | 12 | 49 | 7 | 96 | 8 | 60 | 6 | 192 | 16 |
| 12 | 0 | 48 | 4 | 35 | 7 | 48 | 4 | 49 | 9 | 120 | 10 | 60 | 8 | 192 | 16 |
| 12 | 2 | 48 | 4 | 36 | 2 | 96 | 8 | 50 | 2 | 144 | 12 | 61 | 1 | 288 | 3 |
| 13 | 1 | 96 | 1 | 36 | 6 | 96 | 8 | 50 | 6 | 144 | 12 | 61 | 3 | 144 | 12 |
| 14 | 0 | 96 | 1 | 37 | 1 | 240 | 2 | 50 | 8 | 96 | 8 | 61 | 5 | 144 | 12 |
| 14 | 2 | 24 | 2 | 37 | 3 | 168 | 14 | 51 | 1 | 144 | 12 | 61 | 7 | 240 | 2 |
| 15 | 1 | 48 | 4 | 37 | 5 | 96 | 8 | 51 | 3 | 96 | 8 | 61 | 9 | 360 |  |
| 15 | 3 | 48 | 4 | 37 | 7 | 288 | 3 | 51 | 5 | 168 | 14 | 61 | 11 | 96 |  |
| 16 | 2 | 24 | 2 | 38 | 0 | 384 | 4 | 51 | 7 | 144 | 12 | 62 | 0 | 576 |  |
| 17 | 1 | 48 | 4 | 38 | 2 | 96 | 8 | 51 | 9 | 96 | 8 | 62 | 2 | 240 | 20 |
| 17 | 3 | 120 | 1 | 38 | 4 | 96 | 8 | 52 | 0 | 144 | 12 | 62 | 4 | 216 | 18 |
| 18 | 2 | 72 | 6 | 38 | 6 | 144 | 12 | 52 | 2 | 144 | 12 | 62 | 6 | 264 | 22 |
| 19 | 1 | 48 | 4 | 39 | 1 | 72 | 6 | 52 | 6 | 192 | 16 | 62 | 8 | 480 | 5 |
| 19 | 3 | 48 | 4 | 39 | 3 | 96 | 8 | 52 | 8 | 120 | 10 | 62 | 10 | 192 | 16 |
| 20 | 2 | 48 | 4 | 39 | 5 | 72 | 6 | 52 | 10 | 192 | 16 | 62 | 12 | 240 | 20 |
| 21 | 1 | 120 | 1 | 39 | 7 | 96 | 8 | 53 | 1 | 360 | 3 | 63 | , | 288 | 24 |
| 21 | 3 | 48 | 4 | 40 | 2 | 96 | 8 | 53 | 3 | 168 | 14 | 63 | , | 192 | 16 |
| 22 | 0 | 192 | 2 | 40 | 4 | 96 | 8 | 53 | 5 | 144 | 12 | 63 | 7 | 240 | 20 |
| 22 | 2 | 72 | 6 | 40 | 6 | 144 | 12 | 53 | 7 | 480 | 5 | 64 | 2 | 120 | 10 |
| 22 | 4 | 72 | 6 | 41 | 1 | 72 | 6 | 53 | 9 | 480 | 5 | 64 | 4 | 168 | 1 |
| 23 | 1 | 72 | 6 | 41 | 3 | 288 | 3 | 54 | 2 | 120 | 10 | 64 | 6 | 144 | 12 |
| 24 | 2 | 72 | 6 | 41 | 5 | 120 | 1 | 54 | 4 | 168 | 14 | 64 | 10 | 96 | 8 |
| 24 | 4 | 48 | 4 | 41 | 7 | 96 | 8 | 54 | 6 | 96 | 8 | 64 | 12 | 168 | 14 |
| 25 | 1 | 48 | 4 | 42 | 0 | 192 | 16 | 54 | 8 | 384 | 4 | 65 | 1 | 144 | 12 |
| 25 | 3 | 192 | 2 | 42 | 2 | 168 | 14 | 54 | 10 | 144 | 12 | 65 | 3 | 480 | 4 |
| 26 | 0 | 96 | 8 | 42 | 4 | 360 | 3 | 55 | 1 | 144 | 12 | 65 | 7 | 192 | 16 |
| 6 | 2 | 48 | 4 | 42 | 6 | 144 | 12 | 55 | 3 | 192 | 16 | 65 | 9 | 144 | 12 |
| 6 | 4 | 120 | 1 | 43 |  | 120 | 10 | 55 | 7 | 144 | 12 | 65 | 11 | 384 |  |
| 7 | 1 | 120 | 10 | 43 | 3 | 192 | 16 | 55 | 9 | 144 | 12 | 65 | 13 | 240 | 2 |
| 27 | 3 | 96 | 8 | 43 | 5 | 144 | 12 | 55 | 11 | 96 | 8 | 66 |  | 192 | 16 |
| 27 | 5 | 96 | 8 | 43 | 7 | 144 | 12 | 56 | 2 | 144 | 12 | 66 |  | 144 | 12 |
| 28 | 0 | 96 | 8 | 44 | 0 | 144 | 12 | 56 | 4 | 96 | 8 | 66 | 4 | 360 | 3 |
| 28 | 2 | 96 | 8 | 44 |  | 96 | 8 | 56 | 6 | 120 | 10 | 66 |  | 144 | 12 |
| 29 | 1 | 192 | 2 | 44 | 6 | 96 | 8 | 56 | 10 | 120 | 10 | 66 | 8 | 144 | 12 |
| 29 | 3 | 48 | 4 | 44 | 8 | 72 | 6 | 57 |  | 216 | 18 | 66 | 10 | 192 | 16 |
| 29 | 5 | 72 | 6 | 45 |  | 384 | 4 | 57 | 3 | 576 |  | 66 | 12 | 240 | 2 |
| 30 | 2 | 96 | 8 | 45 | 3 | 96 | 8 | 57 | 5 | 480 | 6 | 67 | , | 240 | 20 |
| 30 | 4 | 96 | 8 | 45 | 7 | 240 | 2 | 57 | 7 | 192 | 16 | 67 | 3 | 192 | 16 |
| 30 | 6 | 48 | 4 | 46 | 0 | 288 | 3 | 57 | 9 | 192 | 16 | 67 | 5 | 216 | 18 |
| 31 | 1 | 48 | 4 | 46 | 2 | 120 | 10 | 57 | 11 | 360 | 3 | 67 | 7 | 168 | 14 |
| 31 | 3 | 72 | 6 | 46 | 4 | 72 | 6 | 58 | 0 | 288 | 24 | 67 | 9 | 312 | 26 |
| 31 | 5 | 48 | 4 | 46 | 8 | 288 | 3 | 58 | 2 | 144 | 12 | 67 | 11 | 216 | 18 |
| 32 | 2 | 72 | 6 | 47 | 1 | 96 | 8 | 58 | 4 | 360 | 3 | 67 | 13 | 240 | 20 |
| 32 | 4 | 120 | 10 | 47 | 3 | 216 | 18 | 58 | 6 | 288 | 24 | 68 | 0 | 288 | 24 |
| 32 | 6 | 120 | 10 | 47 | 5 | 192 | 16 | 58 | 8 | 192 | 16 | 68 | 2 | 192 | 16 |
| 33 | 1 | 144 | 12 | 47 | 7 | 120 | 10 | 58 | 10 | 120 | 10 | 68 | 6 | 288 | 1 |

Tabulation of $R$ and $H$ for $\mu=(a+b \sqrt{5}) / 2$-Continued

|  | $b$ | $R$ | H | $a$ | $b$ | R | H | $a$ |  | $R$ | H |  |  | $R$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 68 | 8 | 264 | 22 | 77 | 1 | 672 | 7 | 84 | 2 | 240 | 20 | 91 |  | 192 | 16 |
|  | 10 | 192 | 16 | 77 | 3 | 288 | 24 | 84 | 6 | 288 | 24 | 91 | 9 | 240 | 20 |
|  | 1 | 240 | 2 | 77 | 5 | 168 | 14 | 84 | 8 | 168 | 14 | 91 | 11 | 192 | 16 |
|  | 3 | 144 | 12 | 77 | 7 | 480 | 4 | 84 | 10 | 240 | 20 | 91 | 13 | 144 | 12 |
|  | 5 | 240 | 20 | 77 | 9 | 600 | 5 | 84 | 14 | 192 | 16 | 91 | 15 | 216 | 18 |
| 69 | 7 | 480 | 5 | 77 | 13 | 288 | 24 | 85 | 3 | 288 | 24 | 91 | 17 | 168 | 14 |
| 69 | 11 | 192 | 16 | 77 | 15 | 672 | 7 | 85 | 7 | 576 | 6 | 92 | 0 | 480 | 40 |
| 69 | 13 | 144 | 12 | 78 | 0 | 768 | 8 | 85 | 9 | 768 | 8 | 92 | 2 | 336 | 28 |
| 70 | 2 | 144 | 12 | 78 | 2 | 360 | 30 | 85 | 11 | 144 | 12 | 92 | 6 | 240 | 20 |
| 70 | 4 | 192 | 16 | 78 | 4 | 240 | 20 | 85 | 13 | 192 | 16 | 92 | 8 | 312 | 26 |
| 70 | 6 | 192 | 16 | 78 | 6 | 288 | 24 | 85 | 17 | 240 | 2 | 92 | 10 | 192 | 16 |
| 70 | 8 | 576 | 6 | 78 | 8 | 768 | 8 | 86 | 0 | 672 | 7 | 92 | 14 | 384 | 32 |
| 70 | 12 | 192 | 16 | 78 | 10 | 312 | 26 | 86 | 2 | 168 | 14 | 92 | 16 | 192 | 16 |
| 70 | 14 | 240 | 20 | 78 | 12 | 240 | 20 | 86 | 4 | 192 | 16 | 92 | 18 | 384 | 32 |
| 71 | 1 | 216 | 18 | 78 | 14 | 240 | 20 | 86 | 6 | 192 | 16 | 93 | 1 | 1056 | 11 |
| 71 | 3 | 192 | 16 | 79 | , | 192 | 16 | 86 | 8 | 384 |  | 93 | 3 | 384 | 32 |
| 71 | 5 | 120 | 10 | 79 | 3 | 264 | 22 | 86 | 10 | 144 | 12 | 93 | 5 | 360 | 30 |
| 71 | 7 | 144 | 12 | 79 | 5 | 144 | 12 | 86 | 12 | 264 | 22 | 93 | 7 | 720 |  |
| 71 | 9 | 144 | 12 | 79 | 7 | 168 | 14 | 86 | 14 | 144 | 12 | 93 | 9 | 480 |  |
|  | 11 | 168 | 14 | 79 | 9 | 168 | 14 | 86 | 16 | 384 | 4 | 93 | 11 | 408 | 34 |
|  | 13 | 144 | 12 | 79 | 11 | 144 | 12 | 87 | 1 | 384 | 32 | 93 | 13 | 288 | 24 |
| 72 | 2 | 360 | 30 | 79 | 13 | 144 | 12 | 87 | 3 | 288 | 24 | 93 | 15 | 768 |  |
| 72 | 4 | 336 | 28 | 79 | 15 | 144 | 12 | 87 | 5 | 360 | 30 | 93 | 17 | 864 |  |
| 72 | 6 | 192 | 16 | 80 |  | 192 | 16 | 87 | 7 | 432 | 36 | 94 |  | 480 |  |
| 2 | 10 | 216 | 18 | 80 | 4 | 192 | 16 | 87 | 9 | 288 | 24 | 94 | 2 | 40 | 20 |
| 2 | 2 | 2 | 16 | 80 | 6 | 336 | 28 | 87 | 11 | 264 | 22 | 94 |  | 40 | 20 |
| 2 | 4 | 216 | 18 | 80 | 12 | 144 | 12 | 87 | 15 | 336 | 28 | 94 |  | 264 | 22 |
| 3 | 3 | 672 | 7 | 80 | 14 | 240 | 20 | 87 | 17 | 264 | 22 | 94 | 8 | 480 |  |
| 3 | 5 | 600 | 5 | 81 | 1 | 144 | 12 | 88 |  | 264 | 22 | 94 | 10 | 144 | 2 |
| 3 | 7 | 192 | 16 | 81 | 3 | 480 | 4 | 88 | 4 | 192 | 16 | 94 | 12 | 216 | 18 |
| 73 | 9 | 288 | 24 | 81 | 5 | 480 | 5 | 88 | 6 | 336 | 28 | 94 | 14 | 216 | 18 |
| 3 | 11 | 360 | 3 | 81 | 7 | 216 | 18 | 88 | 10 | 264 | 22 | 94 | 16 | 480 |  |
| 3 | 13 | 576 | 6 | 81 | 11 | 672 | 7 | 88 | 12 | 336 | 28 | 94 | 18 | 240 | 20 |
| 4 |  | 192 | 16 | 81 | 13 | 360 | 3 | 88 | 14 | 288 | 24 | 95 | 1 | 288 | 24 |
| 4 | 2 | 144 | 12 | 81 | 15 | 192 | 16 | 89 |  | 168 | 14 | 95 |  | 240 | 20 |
| 4 | 4 | 240 | 2 | 82 | 0 | 384 | 32 | 89 | 3 | 576 |  | 95 | 7 | 288 | 24 |
| 4 | 6 | 168 | 14 | 82 | 2 | 216 | 18 | 89 | 5 | 360 | 3 | 95 | 9 | 240 | 20 |
| 74 | 8 | 144 | 12 | 82 | 4 | 480 | 4 | 89 | 7 | 192 | 16 | 95 | 11 | 192 | 16 |
| 74 | 10 | 144 | 12 | 82 | 6 | 336 | 28 | 89 | 9 | 240 | 20 | 95 | 13 | 240 | 20 |
| 74 | 12 | 480 | 4 | 82 | 8 | 192 | 16 | 89 | 11 | 480 | 4 | 95 | 17 | 240 | 20 |
| 74 | 14 | 144 | 12 | 82 | 10 | 312 | 26 | 89 | 13 | 576 | 6 | 95 | 19 | 192 | 16 |
| 75 | 1 | 192 | 16 | 82 | 12 | 720 | 6 | 89 | 15 | 168 | 14 | 96 | 4 | 216 | 18 |
| 75 | 3 | 192 | 16 | 82 | 14 | 240 | 20 | 89 | 17 | 168 | 14 | 96 |  | 240 | 20 |
| 75 | 7 | 288 | 24 | 82 | 16 | 240 | 20 | 90 |  | 240 | 20 | 96 | 10 | 312 | 26 |
| 75 | 9 | 192 | 16 | 83 | 1 | 240 | 20 | 90 | 4 | 480 | 4 | 96 | 12 | 240 | 20 |
| 75 | 11 | 240 | 20 | 83 | 3 | 288 | 24 | 90 | 6 | 288 | 24 | 96 | 14 | 240 | 20 |
| 5 | 13 | 192 | 16 | 83 | 5 | 240 | 20 | 90 | 8 | 288 | 24 | 96 | 18 | 192 | 16 |
| 76 | 0 | 144 | 12 | 83 | 7 | 288 | 24 | 90 | 12 | 480 | 4 | 97 | 1 | 288 | 24 |
| 76 |  | 144 | 12 | 83 |  | 360 | 30 | 90 | 14 | 336 | 28 | 97 | 3 | 600 |  |
| 76 | 6 | 192 | 16 | 83 | 11 | 168 | 14 | 90 | 16 | 288 | 24 | 97 | 5 | 768 | 8 |
| 76 | 8 | 192 | 16 | 83 | 13 | 288 | 24 | 91 |  | 168 | 14 | 97 | 7 | 336 | 28 |
| 76 | 10 | 144 | 12 | 83 | 15 | 336 | 28 | 91 | 3 | 240 | 20 | 97 | 9 | 456 | 38 |
|  | 14 | 144 | 12 | 84 | 0 | 192 | 16 | 91 | 5 | 192 | 16 | 97 | 11 | 960 |  |

Tabulation of $R$ and $H$ for $\mu=(a+b \sqrt{5}) / 2$-Continued

| $a$ | $b$ | $R$ | $H$ | $a$ | $b$ | $R$ | $H$ | $a$ | $b$ | $R$ | $H$ | $a$ | $b$ | $R$ | $H$ |
| :---: | :---: | :---: | ---: | :---: | ---: | :---: | ---: | :---: | ---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 97 | 13 | 600 | 5 | 98 | 6 | 432 | 36 | 98 | 18 | 360 | 30 | 99 | 13 | 288 | 24 |
| 97 | 15 | 432 | 36 | 98 | 8 | 336 | 28 | 99 | 1 | 192 | 16 | 99 | 15 | 192 | 16 |
| 97 | 17 | 288 | 24 | 98 | 10 | 336 | 28 | 99 | 3 | 192 | 16 | 99 | 17 | 240 | 20 |
| 97 | 19 | 360 | 3 | 98 | 12 | 840 | 7 | 99 | 5 | 288 | 24 | 99 | 19 | 288 | 24 |
| 98 | 2 | 288 | 24 | 98 | 14 | 288 | 24 | 99 | 7 | 264 | 22 |  |  |  |  |
| 98 | 4 | 840 | 7 | 98 | 16 | 240 | 20 | 99 | 11 | 192 | 16 |  |  |  |  |

The computation was monitored for about the first 200 tally operations to make sure the score-keeping was correct in all possible cases. The tallying was, as before, basically a question of seeing that every permutation and change in sign in the triple ( $\xi_{1}, \xi_{2}, \xi_{3}$ ) counted as a unit. The total running time was roughly 2.5 hours. One might remark that the human time involved in computing these class numbers $H$ from basic algebraic concepts would have to be measured in "life-times," not "man-hours."

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## Multiplication Time on The IBM 709

By D. D. Wall

Average multiply time is useful for roughly estimating problem running time for various problems, as well as for roughly comparing different computing machines. Determining average multiply time for the 709 is complicated, however, due to its zero-skipping feature, and requires an investigation of runs of zeros in binary sequences. The particular problem we solve is that of evaluating $R(n, l)=$ total number of runs of length $l$ in all the $2^{n}$ words of $n$ bits each, and $S(n, l)=\sum_{x=l}^{n} R(n, x)=$ number of runs of length $\geq l$ in the $2^{n}$ words of $n$ bits each. The resulting 709 average multiply time is 193 microseconds fixed point, or 170 microseconds normalized floating point, and the purpose of this note is to derive these two numbers.

We make use of a device which we call "differencing modulo 2," which obtains an $n-1$ bit number from a given $n$ bit number by writing 1 or 0 according as the successive bits in the given number exhibit a change or no change. For example, each of the (complementary) 8 bit numbers 11010001 and 00101110 gives the same result 0111001 as its 7 bit difference modulo 2.

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[^0]:    Received 8 January 1958.

