TECHNICAL NOTES AND SHORT PAPERS

	Table of Factors of $2^p - 1$ : $\frac{p}{\text{factor}}$ -Continued												
49	9109 55297	1	9127 46033	10	9137 05071	30	9157 76753	4	9173 95343	13	9181 77151		
7	9221 19239	59	9283 41121	14	9311 15273	8	9323 95009	26	9337 14361		9341 74729		
1	9343 49489		9371 18743		9391 93911	2	9397 25529	25	9403 76423		9419 18839		
85	9421 73111	6	9431 79033		9461 75689		9479 18959	5	9491 31497	5	9497 31833		
	9511 95111	3	9521 61799		9539 19079	22	9601 85039		9613 57679	6	9619 15617		
	9791 19583	77	9811 70313	7	9829 07689	45	9833 62513		9851 78809	11	9859 04209		
1	9883 58129	80	9949 98487	2	9973 99191								

# A Computation of Some Bi-Quadratic Class Numbers

# By Harvey Cohn

A fascinating chapter in computational number theory began when Lagrange showed that every positive integer is representable as the sum of at most four perfect squares [1]. Clearly three would not suffice in every case, as  $7 = 2^2 + 1^2 + 1^2 + 1^2$  would be an exception; nevertheless, the problem of expressing some positive integer *n* as the sum of at most *three* squares soon achieved a very special role. For, Gauss showed that r(n), the number of such representations, is connected in a very simple way with the much studied (intrinsically positive) class number, *h*, of the field generated by  $\sqrt{-n}$ . Specifically for *n* square-free (and  $n \neq 1$ , 3 where h = 1),

$$(1) r(n) = gh$$

where g = 12 for  $n \equiv 1, 2, 5, 6$  and g = 24 for  $n \equiv 3 \pmod{8}$ . Thus we could conclude the existence of at least one such representation for the indicated n. Gauss and later, Kronecker, reversed the direction of these equations by making large scale tabulations of h from r(n), although, unfortunately, no location for Kronecker's alleged tabulation (for odd n up to 10,000) seems to exist in the literature. In tallying the representation  $n = x_1^2 + x_2^2 + x_3^2$  it might be noted that one must count each ordered triple  $(x_1, x_2, x_3)$  of positive, negative, or zero integers as a separate unit, so that as much as  $2^3 \cdot 3! = 48$  could be contributed to r(n) when such a decomposition into squares is expressed as triples.

In more recent times, the representation theory was extended to integers in the field k generated by  $\sqrt{5}$ , i.e., to the quantities  $\mu = (a + b\sqrt{5})/2$  where a and b are of the same parity. Here we seek to represent, necessarily, only those integers  $\mu$  which are positive together with their conjugate (i.e., totally positive). Thus, e.g.,  $a > |b\sqrt{5}| \ge 0$ . The special surd  $\sqrt{5}$  must be used because then, as Götzky

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showed [2], each totally positive integer in k could be expressed as the sum of the squares of at most *four* integers in k. Later, Maass [3] made the more remarkable discovery that at most *three* squares would *always* suffice; in fact he arrived at a formula analogous to that of Gauss. Since Maass' formula is the basis of a machine calculation, we avoid irrelevant complexities by making certain further assumptions. First of all  $\mu$  is to be free from square integral divisors in k except for powers of  $(\sqrt{5} + 3)/2 = [(\sqrt{5} + 1)/2]^2 = \epsilon$ . Secondly  $a \ge 5b \ge 0$ , since if  $5b > a > \sqrt{5}b$ , we can continually replace  $\mu$  by  $\mu/\epsilon$ . Then  $R(\mu)$ , the number of representations of  $\mu$  as the sum of three squares, is linked to an intrinsically positive quantity H, namely the class number of the bi-quadratic field generated by  $\sqrt{5}$  and  $\sqrt{-\mu}$ , by means of the following formula (which excludes  $\mu = 1$ ,  $(5 + \sqrt{5})/2$ , and 3 where H = 1):

$$(2) H = R(\mu)/G.$$

Here G = 12 when  $(a, b) \neq (1, 3)$ , (1, 5), (2, 4), (5, 1), (5, 7), or  $(6, 0) \mod 8$ (actually, when  $\eta^2 + \mu \equiv 0 \mod 4$  is unsolvable for  $\eta$  in k). Otherwise, G = 120except when  $(a, b) \equiv (1, 5)$ , (1, 11), (5, 7), (5, 9), (6, 0), (9, 3), (9, 13), (13, 1), (13, 15), or  $(14, 0) \mod 16$  (actually, when  $\eta^2 + \mu \equiv 0 \mod 8$  is solvable for  $\eta$  in k); in these cases, G = 96.

A tabulation of R and H for 446 selected values of  $\mu = [a, b] = (a + b\sqrt{5})/2$ was made on the stored program electronic computer, the IBM 650. The values of a, b were selected with the restrictions

$$100 > a \ge 5b \ge 0,$$

and that  $\mu$  have (except trivially for powers of  $\epsilon$ ) no square divisors in k; or, in terms of ordinary integral arithmetic, the condition is that d, the g.c.d. of  $(\frac{1}{2}[a+b], \frac{1}{2}[a-b])$ , be relatively prime to  $5(=(\sqrt{5})^2)$ , and both d and  $(a^2 - 5b^2)/4d^2$  be square free. The machine assembled 446 such couples automatically into the highest four decimal positions of 446 individual ten digit storage locations, in lexicographic order.

The machine next tallied the decompositions  $\mu = \xi_1^2 + \xi_2^2 + \xi_3^2$  where  $\xi_i = [a_i, b_i], a_i \equiv b_i \mod 2$ . Here the three couples  $\xi_i$  were scanned in lexicographic order, with the restriction  $0 \le b_i \le 7$ , while  $a_i \ge 0$  and a' < 100, (see below). Thus, taking all sign possibilities, with

$$\begin{cases} a' = [a_1^2 + a_2^2 + a_3^2 + 5(b_1^2 + b_2^2 + b_3^2)]/2, \\ b' = \pm a_1b_1 \pm a_2b_2 \pm a_3b_3, \end{cases}$$

the couples a', b' were constructed and compared with the 446 cases stored in the memory. Whenever a matching entry was located the count was augmented and accumulated in the last six decimal positions of the memory word. It might be appropriate to mention that the IBM 650 has a special "table look-up" operation that searches the memory at high speed for the appropriate entry. Without such an instruction the search would have had to be programmed with a considerable loss of running time.

In the final phase the "words"  $a, b, R(\mu)$  were unpacked and the congruences were examined automatically to calculate H and to produce the output consisting of one IBM card per value of  $\mu$ . (See attached table).

TECHNICAL NOTES AND SHORT PAPERS

# Tabulation of *R* and *H* for $\mu = (a + b\sqrt{5})/2$

$1 \text{ abulation of } I \text{ and } I \text{ for } \mu = (u + b \sqrt{3})/2$															
a	b	R	H	a	b	R	H	a	b	R	H	a	b	R	H
2	0	6	1	33	3	240	2	47	9	144	12	59	1	144	12
4	0	12	1	33	5	288	3	48	2	216	18	59	3	168	14
5	1	24	1	34	0	96	8	48	4	168	14	59	5	96	8
6	0	32	1	34	2	96	8	48	6	144	12	59	7	96	8
7	1	24	2	34	4	120	1	49	1	144	12	59	9	144	12
9	1	24	2	34	6	72	6	49	3	240	2	59	11	144	12
10	2	24	2	35	1	96	8	49	5	288	3	60	2	144	12
11	1	24	2	35	3	144	12	49	7	96	8	60	6	192	16
12	0	48	4	35	7	48	4	49	9	120	10	60	8	192	16
12	2	48	4	36	2	96	8	50	2	144	12	61	1	288	3
13	1	96	1	36	6	96	8	50	6	144	12	61	3	144	12
14	0	96	1	37	1	240	2	50	8	96	8	61	5	144	12
14	2	24	2	37	3	168	14	51	1	144	12	61	7	240	2
15	1	48	4	37	5	96	8	51	3	96	8	61	9	360	3
15	3	48	4	37	7	288	3	51	5	168	14	61	11	96	8
16 17	2	24 48	2	38	0	384	4	51	7	144	12	62	0	576	6
17	1 3	40 120	4	38	2	96	8	51 52	9	96	8	62	2	240	20
18	3 2	72	1 6	38 38	4 6	96 144	8 12	52 52	0 2	144 144	12 12	62 62	4 6	216	18
19	1	48	4	39	1	72	6	52 52	6	192	12	62 62	8	264 480	22 5
19	3	48	4	39	3	96	8	52 52	8	120	10	62	10	480 192	
20	2	48	4	39	5	72	6	52 52	10	192	16	62	12	240	16 20
21	1	120	1	39	7	96	8	53	10	360	3	63	12	288	20 24
21	3	48	4	40	2	96	8	53	3	168	14	63	3	192	16
22	ŏ	192	2	40	$\tilde{4}$	96	8	53	5	144	12	63	7	240	20
$\bar{2}\bar{2}$	Ž	72	6	40	6	144	12	53	7	480	5	64	2	120	10
$\bar{2}\bar{2}$	4	72	Ğ	41	ĭ	72	6	53	9	480	5	64	$\tilde{4}$	168	14
23	1	72	6	41	3	288	3	54	2	120	10	64	6	144	12
24	2	72	6	41	5	120	1	54	4	168	14	64	10	96	8
24	4	48	4	41	7	96	8	54	6	96	8	64	12	168	14
25	1	48	4	42	0	192	16	54	8	384	4	65	1	144	12
25	3	192	2	42	2	168	14	54	10	144	12	65	3	480	4
26	0	96	8	42	4	360	3	55	1	144	12	65	7	192	16
26	2	48	4	42	6	144	12	55	3	192	16	65	9	144	12
26	4	120	1	43	1	120	10	55	7	144	12	65	11	384	4
27	1	120	10	43	3	192	16	55	9	144	12	65	13	240	2
27	3	96	8	43	5	144	12	55	11	96	8	66	0	192	16
27	5	96	8	43	7	144	12	56	2	144	12	66	2	144	12
28	0	96	8	44	0	144	12	56	4	96	8	66	4	360	3
28	2	96	8	44	2	96	8	56	6	120	10	66	6	144	12
29	1 3	192	2	44	6	96 70	8	56	10	120	10	66	8	144	12
29 29	5 5	48 72	4 6	44 45	8 1	72	6 4	57 57	1	216 576	18	66	10	192	16
29 30	2	72 96		45 45	3	384	4 8	57 57	3 5		6	66	12	240	2
30 30	2 4	90 96	8 8	45 45	3 7	96 240	8 2	57	3 7	480 192	4 16	67 67	1 3	240	20
30	6	48	4	45 46	ó	288	$\frac{2}{3}$	57	9	192	16	67	5 5	192 216	16 18
31	ĭ	48	4	46	2	120	10	57	11	360	3	67	7	168	10
31	3	72	6	46	$\frac{2}{4}$	72	6	58	0	288	24	67	9	312	26
31	5	48	4	46	8	288	3	58	ž	144	12	67	11	216	18
32	2	72	6	47	ĭ	96	8	58	4	360	3	67	13	240	20
32	4	120	10	47	$\hat{3}$	216	18	58	6	288	24	68	0	288	20 24
32	6	120	10	47	5	192	16	58	8	192	16	68	ž	192	16
33	1	144	12	47	7	120	10	58	10	120	10	68	6	288	24
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Tabulation of R and H for  $\mu = (a + b\sqrt{5})/2$ —Continued

a	b	R	H	a	b	R	H	a	b	R	H	a	b	R	H	
68	8	264	22	77	1	672	7	84	2	240	20	91	7	192	16	
68	10	192	16	77	3	288	24	84	6	288	24	91	9	240	20	
69	1	240	2	77	5	168	14	84	8	168	14	91	11	192	16	
69	3	144	12	77	7	480	4	84	10	240	20	91	13	144	12	
69	5	240	20	77	9	600	5	84	14	192	16	91	15	216	18	
69	7	480	5	77	13	288	24	85	3	288	24	91	17	168	14	
69	11	192 144	16	77	15	672 768	7	85	7	576	6	92	0	480	40	
69 70	13 2	144	12 12	78 78	0 2	360	8 30	85 85	9 11	768 144	8 12	92 92	2 6	336 240	28 20	
70 70	2 4	192	16	78 78	2 4	300 240	30 20	85	13	192	16	92	8	<sup>240</sup> 312	20	
70	6	192	16	78	6	288	$\frac{20}{24}$	85	17	240	2	92	10	192	16	
70	8	576	6	78	8	768	8	86	0	672	7	92	14	384	32	
70	12	192	16	78	10	312	26	86	2	168	14	92	16	192	16	
70	14	240	20	78	12	240	20	86	4	192	16	92	18	384	32	
71	1	216	18	78	14	240	20	86	6	192	16	93	1	1056	11	
71	3	192	16	79	1	192	16	86	8	384	4	93	3	384	32	
71	5	120	10	79	3	264	22	86	10	144	12	93	5	360	30	
71	7	144	12	79	5	144	12	86	12	264	22	93	7	720	6	
71	9	144	12	79	7	168	14	86	14	144	12	93	9	480	4	
71	11	168	14	79	9	168	14	86	16	384	4	93	11	408	34	
71	13	144	12	79 79	11	144 144	12 12	87 87	1 3	384	32	93	13	288	24	
72 72	2 4	360 336	30 28	79	13 15	144	12	87	5	288 360	24 30	93 93	15 17	768 864	8 9	
72	<del>4</del> 6	192	16	80	2	192	16	87	7	432	36	93	0	480	5	
72	10	216	18	80	4	192	16	87	9	288	24	94	2	240	20	
$\frac{72}{72}$	12	192	16	80	6	336	28	87	11	264	22	94	$\tilde{4}$	240	20	
72	14	216	18	80	12	144	12	87	15	336	$\frac{1}{28}$	94	6	264	$\overline{22}$	
73	3	672	7	80	14	240	20	87	17	264	22	94	8	480	5	
73	5	600	5	81	1	144	12	88	2	264	22	94	10	144	12	
73	7	192	16	81	3	480	4	88	4	192	16	94	12	216	18	
73	9	288	24	81	5	480	5	88	6	336	28	94	14	216	18	
73	11	360	3	81	7	216	18	88	10	264	22	94	16	480	5	
73	13	576	6	81	11	672	7 3	88	12	336	28	94	18	240	20	
74	0 2	192 144	16 12	81 81	13 15	360 192	3 16	88 89	14 1	288 168	24 14	95 95	1 3	288 240	24 20	
74 74	2 4	240	2	82	0	384	32	89	3	576	6	95	7	288	20 24	
74	6	168	$14^{2}$	82	2	216	18	89	5	360	3	95	9	240	$\frac{24}{20}$	
74	8	144	12	82	$\overline{4}$	480	4	89	7	192	16	95	11	192	16	
74	10	144	12	82	6	336	28	89	9	240	20	95	13	240	20	
74	12	480	4	82	8	192	16	89	11	480	4	95	17	240	20	
74	14	144	12	82	10	312	26	89	13	576	6	95	19	192	16	
75	1	192	16	82	12	720	6	89	15	168	14	96	4	216	18	
75	3	192	16	82	14	240	20	89	17	168	14	96	6	240	20	
75	7	288	24	82	16	240	20	90	2	240	20	96	10	312	26	
75	9	192	16	83	1	240	20	90	4	480	4	96	12	240	20	
75	11	240	20	83	3 5	288	24	90	6	288 288	24 24	96	14	240	20	
75 76	13	192 144	16 12	83 83	3 7	240 288	20 24	90 90	8 12	288 480	24 4	96 97	18 1	192 288	16 24	
76 76	0 2	144	12	83	9	200 360	$\frac{24}{30}$	90	14	336	28	97	3	200 600	24 5	
76	6	192	16	83	11	168	14	90	16	288	28 24	97	5	768	8	
76	8	192	16	83	13	288	24	91	1	168	14	97	7	336	28	
76	10	144	12	83	15	336	$\overline{28}$	91	3	240	20	97	9	456	38	
76	14	144	12	84	0	192	16		5	192	16	97	11	960	10	
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# TECHNICAL NOTES AND SHORT PAPERS

							•	`	•						
a	b	R	H	a	Ь	R	H	a	b	R	H	a	b	R	H
97	13	600	5	98	6	432	36	98	18	360	30	99	13	288	24
														192	
														240	
97	19											99	19	288	24
										264					
98	4	840	7	98	16	240	20	99	11	192	16				

Tabulation of R and H for  $\mu = (a + b\sqrt{5})/2$ —Continued

The computation was monitored for about the first 200 tally operations to make sure the score-keeping was correct in all possible cases. The tallying was, as before, basically a question of seeing that every permutation and change in sign in the triple  $(\xi_1, \xi_2, \xi_3)$  counted as a unit. The total running time was roughly 2.5 hours. One might remark that the human time involved in computing these class numbers H from basic *algebraic* concepts would have to be measured in "life-times," not "man-hours."

The computation was completed 18 April 1958 and was sponsored in part by the National Science Foundation Grant G-4222.

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1. L. E. DICKSON, History of the Theory of Numbers, v. III, G. E. Stechert, New York, 1934 (for references in the first paragraph).

2. F. GÖTZKY, "Über eine zahlentheoretische anwendung von modulfunktionen zweier verändlicher," Mathematische Annalen, v. 100, 1928, p. 411-437.

3. H. MAASS, "Uber die darstellung total positiver zahlen des körpers  $R(\sqrt{5})$  als summe von drei quadraten," Abhandlungen aus dem Mathematischen Seminar der Hansischen Universitat, v. 14, 1941, p. 185–192.

# Multiplication Time on The IBM 709

# By D. D. Wall

Average multiply time is useful for roughly estimating problem running time for various problems, as well as for roughly comparing different computing machines. Determining average multiply time for the 709 is complicated, however, due to its zero-skipping feature, and requires an investigation of runs of zeros in binary sequences. The particular problem we solve is that of evaluating R(n, l) = total number of runs of length l in all the  $2^n$  words of n bits each, and  $S(n, l) = \sum_{x=l}^{n} R(n, x) =$  number of runs of length  $\geq l$  in the  $2^n$  words of nbits each. The resulting 709 average multiply time is 193 microseconds fixed point, or 170 microseconds normalized floating point, and the purpose of this note is to derive these two numbers.

We make use of a device which we call "differencing modulo 2," which obtains an n - 1 bit number from a given n bit number by writing 1 or 0 according as the successive bits in the given number exhibit a change or no change. For example, each of the (complementary) 8 bit numbers 11010001 and 00101110 gives the same result 0111001 as its 7 bit difference modulo 2.

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