# A First Order Method for Differential Equations of Neutral Type 

By R. N. Castleton and L. J. Grimm*


#### Abstract

A first order method is presented for solution of the initial-value problem for a differential equation of neutral type with implicit delay in the critical case where the time-lag is zero and the method of stepwise integration does not apply. A convergence theorem is proved, and numerical examples are given.


1. Introduction. In this note, we present a first order method for the numerical solution of the initial-value problem (IVP) for a neutral-type functional-differential equation without previous history:

$$
\begin{equation*}
x^{\prime}(t)=f\left(t, x(t), x(g(t, x(t))), x^{\prime}(g(t, x(t)))\right) \tag{1}
\end{equation*}
$$

where $z_{0}$ is a real root of the algebraic equation

$$
\begin{equation*}
z=f\left(a, x_{0}, x_{0}, z\right) \tag{3}
\end{equation*}
$$

Here, $x(t)$ is a scalar function to be determined on some finite interval $[a, b]$. We shall make the following assumptions regarding $f$ and $g$ :
(H1) $f$ and $g$ are continuous and satisfy uniform Lipschitz conditions of the form

$$
\begin{aligned}
\left|f\left(t, x_{1}, y_{1}, z_{1}\right)-f\left(t, x_{2}, y_{2}, z_{2}\right)\right| & \leqq L\left\{\left|x_{1}-x_{2}\right|+\left|y_{1}-y_{2}\right|\right\}+L_{2}\left|z_{1}-z_{2}\right|, \\
\left|g\left(t, x_{1}\right)-g\left(t, x_{2}\right)\right| & \leqq L_{0}\left|x_{1}-x_{2}\right|
\end{aligned}
$$

in their respective domains $E$ and $E^{\prime}$, where

$$
E=\left\{(t, x, y, z): a \leqq t \leqq b,\left|x-x_{0}\right| \leqq c,\left|y-x_{0}\right| \leqq c,|z| \leqq M\right\}
$$

and $E^{\prime}$ is the projection of $E$ in the $(t, x)$ space; $c, M, L, L_{0}, L_{z}$ are constants, with $L_{z}<1, M$ is such that $\sup _{(t, x, y, z) \in E}|f(t, x, y, z)|<M$, and $M(b-a)<c$.
(H2) $a \leqq g(t, x) \leqq t$ for $(t, x) \in E^{\prime}$.
Our hypotheses, together with additional smoothness and growth conditions on $f$ and $g$, ensure the local existence of a solution of the IVP (1)-(2). Furthermore, $x(t)$ is the only solution having a bounded derivative on [a, b]; see [2], [4]. Our result extends a method developed by Feldstein [3] for the equation of retarded type

[^0]$$
x^{\prime}(t)=f(t, x(t), x(g(t)))
$$
to the neutral-type equation with implicit delay (1). Other methods for implicitdelay equations are given in [1].
2. The Algorithm $\mathfrak{A}$. Let $y(t)=x(g(t, x(t))) ; z(t)=x^{\prime}(g(t, x(t)))$. Let $N$ be a positive integer, and let $h=(b-a) / N$. For each nonnegative integer $n \leqq N$, let $t_{n}=a+n h$. Let $[s]$ denote the integer part of $s$. Define the algorithm $\mathfrak{A}$ as follows:
\[

$$
\begin{array}{rlrl}
f_{n} & =f\left(t_{n}, x_{n}, y_{n}, z_{n}\right), & g_{n} & =g\left(t_{n}, x_{n}\right), \\
q(n) & =\left[\left(g_{n}-a\right) / h\right], \quad r(n)=\left(g_{n}-a\right) / h-q(n), \\
y_{0} & =x_{0}, & & \\
z_{n} & =x_{q(n)}+h r(n) f_{q(n)}, \\
z_{n} & =f_{q(n)},  \tag{8}\\
x_{n+1} & =x_{n}+h f_{n} .
\end{array}
$$
\]

Note that condition (H2) implies $q(n) \leqq n$, thus, the algorithm is well defined. For $n=0, g_{0}=a, q(0)=0$, and $r(0)=0$. Thus, $y_{0}=x_{0}$ and $z_{0}=f\left(a, x_{0}, x_{0}, z_{0}\right)$. Let $u_{0}$, an approximation of the root $z_{0}$, be chosen independently of $h$. It is of interest to note that such an approximation does not destroy the order $h$ convergence of the algorithm. It is of further interest that (6) may be simplified to $y_{n}=x_{a(n)}$. The error bound established in the convergence theorem for this "simplified" algorithm is larger but still of order $h$, as noted following the proof of convergence of the algorithm $\mathfrak{A}$. The second numerical example of Section 4 demonstrates both the algorithm $\mathfrak{A}$ and the simplified algorithm.

If $g_{n}=t_{n}$ for any $n, 1 \leqq n \leqq N$, then $q(n)=n, r(n)=0$, and (7) becomes $z_{n}=$ $f\left(t_{n}, x_{n}, y_{n}, z_{n}\right)$ which has exactly one root $z$ in the interval $[-M, M]$ under the conditions (H1)-(H2) together with the smoothness and growth conditions mentioned in Section 1. We must in general include a procedure for finding this root, and this in turn will affect the error estimate. As before, such an estimate does not destroy the order $h$ convergence of the algorithm. For simplicity, we do not take this into account, since our aim is to show the convergence of the algorithm $\mathfrak{A}$.

Thus, we shall assume in the convergence proof that (7) will not reduce to $z_{n}=$ $f\left(t_{n}, x_{n}, y_{n}, z_{n}\right), n \geqq 1$.

## 3. Convergence.

Theorem. Let $f$ and $g$ satisfy $(\mathrm{H} 1)-(\mathrm{H} 2)$ and suppose, in addition, that there exists a unique solution $x(t)$ of (1)-(2) with $\sup _{|a, b|}\left|x^{\prime \prime}(t)\right| \leqq B$. Then, for each $t_{n} \in$ $[a, b], 0<n \leqq N$,

$$
\left|x_{n}-x\left(t_{n}\right)\right| \leqq h\left\{L_{z}\left|z_{0}-u_{0}\right| e^{s(b-a)}+\frac{B}{2 s}\left(\frac{1}{1}-\frac{L_{z}}{L_{z}}\right)\left(e^{s(b-a)}-1\right)\right\}+O\left(h^{2}\right)
$$

where

$$
s=L\left(1+c_{0}\right)+L_{2} c_{1}
$$

$$
\begin{aligned}
& c_{0}=1+M L_{v} \\
& c_{1}=\left(L\left(2+M L_{\imath}\right)+B L_{v}\right) /\left(1-L_{z}\right)
\end{aligned}
$$

$u_{0}$ is the approximation to $z_{0}$ mentioned above, and $x_{n}$ is given by algorithm $\mathfrak{A}$.
Proof. Let $e_{n}=\left|x_{n}-x\left(t_{n}\right)\right| ; e_{n}^{*}=\left|y_{n}-y\left(t_{n}\right)\right| ; e_{n}^{* *}=\left|z_{n}-z\left(t_{n}\right)\right|$. From (8) and Taylor's formula, we obtain

$$
\begin{equation*}
e_{n+1} \leqq e_{n}+h\left(L\left(e_{n}+e_{n}^{*}\right)+L_{2} e_{n}^{* *}\right)+h^{2} B / 2 \tag{9}
\end{equation*}
$$

Equation (5) implies that $g_{n}=t_{q(n)}+h r(n)$, and hence, in a similar manner, we have (after replacing $n$ by $(n+1)$ )

$$
\begin{align*}
& e_{n+1}^{*} \leqq M L_{0} e_{n+1}+e_{q(n+1)}  \tag{10}\\
&+h r(n+1)\left\{L\left(e_{u(n+1)}+e_{u(n+1)}^{*}+L_{2} e_{u(n+1)}^{* *}\right\}+h^{2} r^{2}(n+1) B / 2\right. \\
& e_{n+1}^{* *} \leqq B L_{u} e_{n+1}+L\left(e_{q(n+1)}+e_{u(n+1)}^{*}\right)+L_{2} e_{u(n+1)}^{* *}+h r(n+1) B . \tag{11}
\end{align*}
$$

We then have two cases to consider:
Case 1. $q(n+1)=n+1$ and $r(n+1)=0$. Under these conditions, (9) is unchanged:

$$
\begin{equation*}
e_{n+1} \leqq e_{n}(1+h L)+e_{n}^{*} h L+e_{n}^{* *} h L_{z}+h^{2} B / 2 . \tag{9a}
\end{equation*}
$$

(10) becomes

$$
\begin{equation*}
e_{n+1}^{*} \leqq e_{n+1}\left(1+M L_{v}\right)=e_{n+1} c_{0} \tag{10a}
\end{equation*}
$$

And (11) becomes

$$
e_{n+1}^{* *} \leqq\left(L+B L_{u}\right) e_{n+1}+L e_{n+1}^{*}+L_{2} e_{n+1}^{* *}
$$

or

$$
\begin{equation*}
e_{n+1}^{* *} \leqq\left(\frac{L+B L_{g}+L\left(1+M L_{g}\right)}{1-L_{z}}\right) e_{n+1}=e_{n+1} c_{1} \tag{11a}
\end{equation*}
$$

Define the partial ordering for vectors: $v_{1}=\left(v_{1}^{1}, \cdots, v_{1}^{k}\right) \leqq v_{2}=\left(v_{2}^{1}, \cdots, v_{2}^{k}\right)$ if $v_{1}^{i} \leqq v_{2}^{i}, i=1, \cdots, k$. Then, in vector form, (9a), (10a), and (11a) become

$$
\left[\begin{array}{l}
e_{n+1} \\
e_{n+1}^{*} \\
e_{n+1}^{* *}
\end{array}\right] \leqq\left[\begin{array}{ccc}
1+h L & h L & h L_{2} \\
(1+h L) c_{0} & h L c_{0} & h L_{2} c_{0} \\
(1+h L) c_{1} & h L c_{1} & h L_{2} c_{1}
\end{array}\right]\left[\begin{array}{c}
e_{n} \\
e_{n}^{*} \\
e_{n}^{* *}
\end{array}\right]+h B\left[\begin{array}{c}
h / 2 \\
h c_{0} / 2 \\
h c_{1} / 2
\end{array}\right]
$$

which is of the form $d_{n+1} \leqq A_{1} d_{n}+b_{1}$.
Case 2. $\quad q(n+1) \leqq n$ and $0 \leqq r(n+1)<1$.
Let

$$
\delta_{n}=\max _{1 \leqq i \leqq n} e_{i}, \quad \delta_{n}^{*}=\max _{1 \leqq i \leqq n} e_{i}^{*}, \quad \delta_{n}^{* *}=\max _{1 \leqq i \leqq n} e_{i}^{* *}
$$

Then, (9) becomes

$$
\begin{equation*}
\delta_{n+1} \leqq \delta_{n}(1+h L)+\delta_{n}^{*} h L+\delta_{n}^{* *} h L_{z}+h^{2} B / 2 \tag{9b}
\end{equation*}
$$

And (10) becomes

$$
\delta_{n+1}^{*} \leqq M L_{0} \delta_{n+1}+\delta_{n}(1+h L)+h L \delta_{n}^{*}+h L_{2} \delta_{n}^{* *}+h^{2} B / 2 .
$$

Using (9b), we have

$$
\delta_{n+1}^{*} \leqq\left(\delta_{n}(1+h L)+\delta_{n}^{*} h L+\delta_{n}^{* *} h L_{z}+h^{2} B / 2\right)\left(1+M L_{\theta}\right)
$$

or

$$
\begin{equation*}
\delta_{n+1}^{*} \leqq \delta_{n}(1+h L) c_{0}+\delta_{n}^{*} h L c_{0}+\delta_{n}^{* *} h L_{2} c_{0}+h^{2} c_{0} B / 2 \tag{10b}
\end{equation*}
$$

Finally, (11) becomes

$$
\delta_{n+1}^{* *} \leqq \delta_{n+1} B L_{v}+\delta_{n} L+\delta_{n}^{*} L+\delta_{n}^{* *} L_{z}+h B .
$$

Further, enlarging $\delta_{n}$ to $\delta_{n+1}$ and $\delta_{n}^{*}$ to $\delta_{n+1}^{*}$ on the right, and using $1-L_{z}>0$, we find

$$
\delta_{n+1}^{* *} \leqq \delta_{n+1}\left(\frac{L+B L_{o}}{1-L_{z}}\right)+\delta_{n+1}^{*} \frac{L}{1-L_{z}}+\frac{h B}{1-L_{z}} .
$$

Using (9b) and (10b), we have

$$
\delta_{n+1}^{* *} \leqq\left(\frac{L+B L_{o}+L c_{0}}{1-L_{z}}\right)\left(\delta_{n}(1+h L)+\delta_{n}^{*} h L+\delta_{n}^{* *} h L_{2}+\frac{h^{2} B}{2}\right)+\frac{h B}{1-L_{z}}
$$

or

$$
\begin{equation*}
\delta_{n+1}^{* *} \leqq \delta_{n}(1+h L) c_{1}+\delta_{n}^{*} h L c_{1}+\delta_{n}^{* *} h L_{2} c_{1}+\frac{h B}{1-L_{z}}+\frac{h^{2} c_{1} B}{2} \tag{11b}
\end{equation*}
$$

Then, as a vector system, (9b), (10b), and (11b) become

$$
\left[\begin{array}{c}
\delta_{n+1}  \tag{12}\\
\delta_{n+1}^{*} \\
\delta_{n+1}^{* *}
\end{array}\right] \leqq\left[\begin{array}{ccc}
1+h L & h L & h L_{2} \\
(1+h L) c_{0} & h L c_{0} & h L_{2} c_{0} \\
(1+h L) c_{1} & h L c_{1} & h L_{2} c_{1}
\end{array}\right]\left[\begin{array}{c}
\delta_{n} \\
\delta_{n}^{*} \\
\delta_{n}^{* *}
\end{array}\right]+h B\left[\begin{array}{c}
h / 2 \\
h c_{0} / 2 \\
h c_{1} / 2+1 /\left(1-L_{z}\right)
\end{array}\right]
$$

which is of the form $d_{n+1} \leqq A_{2} d_{n}+b_{2}$. Comparing this with the result obtained in Case 1 , we find that $A_{1}$ and $A_{2}$ are identical and that $b_{1} \leqq b_{2}$. Thus, any bound obtained here in Case 2 for $d_{n+1}$ will also bound $d_{n+1}$ in Case 1 .

To complete the proof, we shall use the following lemmas [3] which may be verified by induction:

Lemma 1. Suppose $A$ is a $k \times k$ real matrix and $b$ is a real $k$-vector. Let $\left\{d_{n}\right\}$ $(n=0,1, \cdots)$ satisfy $d_{n+1} \leqq A d_{n}+b$. Then

$$
d_{n+1} \leqq A^{n+1} d_{0}+\left(\sum_{i=1}^{n} A^{i}\right) b
$$

Lemma 2. Let $p=\left(p_{1}, \cdots, p_{k}\right), q=\left(q_{1}, \cdots, q_{k}\right)$. Suppose the $k \times k$ matrix $A$ has the form $A=p^{T} q$. Then

$$
A^{n}=\left(\sum_{i=1}^{k} p_{i} q_{i}\right)^{n-1} A
$$

By Lemma 1,

$$
d_{n+1} \leqq A_{2}^{n+1} d_{0}+\left(\sum_{i=0}^{n} A_{2}^{i}\right) b_{2},
$$

where

$$
d_{0}=\left[\begin{array}{c}
e_{10} \\
e_{11}^{*} \\
e_{11}^{* *}
\end{array}\right]=\left[\begin{array}{c}
0 \\
0 \\
\left|z_{0}-u_{0}\right|
\end{array}\right]
$$

Then, because

$$
A_{2}=\left[\begin{array}{c}
1 \\
c_{1} \\
c_{1}
\end{array}\right]\left(1+h L, h L, h L_{z}\right)
$$

we can make use of Lemma 2 to obtain

$$
A_{2}^{i}=\left(1+h L+h L c_{0}+h L_{2} c_{1}\right)^{i-1} A_{2}=(1+h s)^{i-1} A_{2}
$$

Two results follow from this: $A_{2}^{n+1}=(1+h s)^{n} A_{2} \leqq e^{s(b-a)} A_{2}$, and

$$
\sum_{i=1}^{n} A_{2}^{i}=A_{2} \sum_{i=1}^{n}(1+h s)^{i-1}=\frac{\left((1+h s)^{n}-1\right)}{h s} A_{2} \leqq \frac{1}{h s}(\exp (s(b-a))-1) A_{2} .
$$

Finally,

$$
\begin{aligned}
d_{n+1} & \leqq A_{2}^{n+1} d_{10}+\left(\sum_{i=1}^{n} A_{2}^{i}\right) b_{2} \\
& \leqq h\left\{\left|z_{11}-u_{0}\right| L_{2} e^{s^{s(1-n)}}\left[\begin{array}{c}
1 \\
c_{0} \\
c_{1}
\end{array}\right]\right. \\
& \left.+\frac{B}{2 s}\left(h s+\frac{1+}{1-\frac{L_{2}}{L_{2}}}\right)\left(e^{s(l-a)}-1\right)\left[\begin{array}{c}
1 \\
c_{0} \\
c_{1}
\end{array}\right]+B\left[\begin{array}{c}
\frac{h}{2} \\
\frac{h c_{0}}{2} \\
\frac{h c_{1}}{2}+\frac{1}{1-L_{2}}
\end{array}\right]\right\}
\end{aligned}
$$

which gives

$$
e_{n+1} \leqq \delta_{n+1} \leqq h\left\{\left|z_{0}-u_{0}\right| L_{z} e^{*(b-n)}+\frac{B}{2 s}\left(h s+\frac{1+L_{2}}{1-L_{2}}\right)\left(e^{*(b-a)}-1\right)+\frac{h B}{2}\right\}
$$

and the theorem follows.

For the simplified algorithm, where (6) is replaced by $y_{n}=x_{Q(n)}$ the following bound is possible:

$$
d_{n+1} \leqq h\left\{\left|z_{0}-u_{0}\right| L_{2} e^{s(b-a)}\left[\begin{array}{c}
1 \\
c_{0} \\
c_{1}
\end{array}\right]\right.
$$

$$
\begin{align*}
&+\left(\frac{B}{2 s}\left(h s+\frac{1+L_{z}}{1-L_{z}}\right)+\frac{1}{s}\left(\frac{M L}{1-L_{z}}\right)\right)\left(e^{*(b-a)}-1\right)\left[\begin{array}{c}
1 \\
c_{0} \\
c_{1}
\end{array}\right]  \tag{13}\\
&+B\left[\begin{array}{c}
\frac{h}{2} \\
\frac{h c_{0}}{2} \\
\frac{h c_{1}}{2}+\frac{1}{1-L_{z}}
\end{array}\right]+\left[\begin{array}{c}
0 \\
M \\
\frac{M L}{1-L_{z}}
\end{array}\right]
\end{align*}
$$

and hence

$$
\begin{aligned}
e_{n+1} \leqq & h\left\{\left|z_{0}-u_{0}\right| L_{2} e^{s(b-a)}\right. \\
& \left.+\left(\frac{B}{2 s}\left(h s+\frac{1+L_{z}}{1-L_{z}}\right)+\frac{1}{s}\left(\frac{M L}{1-L_{z}}\right)\right)\left(e^{s(b-a)}-1\right)+\frac{h B}{2}\right\} .
\end{aligned}
$$

Table I. $\quad x_{n}(h)$ denotes the value of $x_{n}$ for step size $h$.

| $t_{n}$ | $x\left(t_{n}\right)$ | $x_{n}\left(2^{-4}\right)$ | $x_{n}\left(2^{-6}\right)$ | $x_{n}\left(2^{-8}\right)$ | $x_{n}\left(2^{-10}\right)$ |
| :--- | ---: | :--- | ---: | :--- | ---: |
| 0 | 0 | 0 | 0 | 0 | 0 |
| .0625 | .0039 | 0 | .0029 | .0034 | .0039 |
| .1250 | .0158 | .0078 | .0138 | .0153 | .0157 |
| .1875 | .0360 | .0238 | .0329 | .0352 | .0358 |
| .2500 | .0653 | .0484 | .0610 | .0642 | .0650 |
| .3125 | .1048 | .0825 | .0990 | .1032 | .1044 |
| .3750 | .1562 | .1275 | .1485 | .1541 | .1556 |
| .4375 | .2224 | .1853 | .2119 | .2196 | .2217 |
| .5000 | .3078 | .2593 | .2942 | .3043 | .3069 |
| .5625 | .4206 | .3547 | .4026 | .4159 | .4194 |
| .6250 | .5771 | .4856 | .5518 | .5705 | .5754 |
| .6875 | .8185 | .6707 | .7778 | .8080 | .8159 |
| .7500 | 1.3244 | .9860 | 1.2205 | 1.2968 | 1.3174 |

Table II. $\quad x_{n}^{(1)}(h)$ denotes the value of $x_{n}$ for step size $h$ by algorithm $\mathfrak{A} ; x_{n}^{(2)}(h)$ denotes the value of $x_{n}$ for step size $h$ by the simplified algorithm.

| $t_{n}$ | $x\left(t_{n}\right)$ | $x_{n}^{(1)}\left(2^{-2}\right)$ | $x_{n}^{(2)}\left(2^{-2}\right)$ | $x_{n}^{(1)}\left(2^{-4}\right)$ | $x_{n}^{(2)}\left(2^{-4}\right)$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| .25 | .2474 | .2500 | .2500 | .2483 | .2478 |
| .50 | .4794 | .4930 | .4892 | .4838 | .4759 |
| .75 | .6816 | .7180 | .6866 | .6942 | .6739 |
| 1.00 | .8414 | .9228 | .8569 | .8697 | .8273 |
| $t_{n}$ | $x\left(t_{n}\right)$ | $x_{n}^{(1)}\left(2^{-8}\right)$ | $x_{n}^{(2)}\left(2^{-8}\right)$ | $x_{n}^{(1)}\left(2^{-12}\right)$ | $x_{n}^{(2)}\left(2^{-12}\right)$ |
| .25 | .2474 | .2475 | .2471 | .2474 | .2474 |
| .50 | .4794 | .4797 | .4787 | .4794 | .4794 |
| .75 | .6816 | .6825 | .6802 | .6817 | .6815 |
| 1.00 | .8414 | .8435 | .8390 | .8416 | .8413 |

4. Examples. (a) We solve the IVP

$$
x^{\prime}(t)=\frac{-4 t x^{2}(t)}{4+\log ^{2} \cos t}+\tan 2 t+\frac{1}{2} \tan ^{-1} z
$$

$\left(z_{0}=0, x_{0}=0, z=x^{\prime}(g(t, x(t))) \equiv x^{\prime}\left(t x^{2}(t) /\left(1+x^{2}(t)\right)\right)\right)$ on the interval $[0, .75]$. The existence and uniqueness of the solution is guaranteed by the results of [2] mentioned earlier. The only solution is $x(t)=-\frac{1}{2} \log \cos 2 t$.

The results of the computation by algorithm $\mathfrak{A}$ are given in Table I.
(b) Consider the IVP

$$
x^{\prime}(t)=\cos t(1+y)+x z-\sin \left(t\left(1+\sin ^{2} t\right)\right)
$$

with $y=x\left(t x^{2}(t)\right), z=x^{\prime}\left(t x^{2}(t)\right), z_{0}=1, x_{0}=0$, on the interval [ 0,1$]$. As in example (a), existence and uniqueness of the solution are guaranteed by the results of [2]. Here, the solution is $x(t)=\sin t$.

The results of the computation by the algorithm $\mathfrak{A}$ and by the simplified algorithm are given in Table II.

## Computation Division <br> Lawrence Livermore Laboratory <br> Livermore, California 94551 <br> Department of Mathematics <br> University of Missouri, Rolla <br> Rolla, Missouri 65401

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