On the Period Length of Pseudorandom Vector Sequences Generated by Matrix Generators

By Jürgen Eichenauer-Herrmann, Holger Grothe, and Jürgen Lehn

Abstract. In Tahmi [5], Niederreiter [4], Afflerbach and Grothe [1], and Grothe [2] linear recursive congruential matrix generators for generating r -dimensional pseudorandom vectors are analyzed. In particular, conditions are established which ensure that the period length equals $p^r - 1$ for any nonzero starting vector in case of a prime modulus p . For a modulus of the form p^α , $\alpha \geq 2$ and p prime, this paper describes a simple method for constructing matrix generators having the maximal possible period length $(p^r - 1) \cdot p^{\alpha-1}$ for any starting vector which is nonzero modulo p .

1. Introduction and Notation. A linear recursive congruential matrix generator for generating r -dimensional pseudorandom vectors is of the form

$$(1) \quad \vec{x}_{n+1} \equiv A \cdot \vec{x}_n \pmod{m}, \quad \vec{x}_{n+1} \in \mathbf{Z}_m^r, \quad n \geq 0,$$

where the modulus m is a positive integer, $\mathbf{Z}_m = \{0, 1, \dots, m-1\}$, $\vec{x}_0 \in \mathbf{Z}_m^r$, and $A \in \mathbf{Z}_m^{r \times r}$, i.e., A is an $r \times r$ -matrix with elements in $\mathbf{Z}_m$. In the sequel it is assumed that the matrix A is nonsingular modulo m . Then the vector sequence $(\vec{x}_n)_{n \geq 0}$ generated by (1) is purely periodic, and the smallest positive integer $\lambda = \lambda(A, \vec{x}_0, m)$ with $\vec{x}_\lambda = \vec{x}_0$ is called the *period length of the vector sequence* $(\vec{x}_n)_{n \geq 0}$. Analogously, the matrix sequence $(A_n)_{n \geq 0}$ with $A_n \equiv A^n \pmod{m}$, $A_n \in \mathbf{Z}_m^{r \times r}$, is purely periodic, and the smallest positive integer $\lambda = \lambda(A, m)$ for which A_λ equals the identity matrix I is called the *period length of the matrix sequence* $(A_n)_{n \geq 0}$. The following two remarks are immediate consequences of these definitions.

Remark 1. The period length $\lambda(A, \vec{x}_0, m)$ of the vector sequence $(\vec{x}_n)_{n \geq 0}$ divides the period length $\lambda(A, m)$ of the matrix sequence $(A_n)_{n \geq 0}$ for any starting vector $\vec{x}_0 \in \mathbf{Z}_m^r$.

Remark 2. If $A_\nu = I$ for some positive integer ν , then the period length $\lambda(A, m)$ of the matrix sequence $(A_n)_{n \geq 0}$ divides ν .

It is well known (cf. Tahmi [5], Niederreiter [4], and Grothe [2]) that $\lambda(A, \vec{x}_0, p) = \lambda(A, p) = p^r - 1$ for any starting vector $\vec{x}_0 \in \mathbf{Z}_p^r \setminus \{\vec{0}\}$ in case of a prime modulus $m = p$ if the characteristic polynomial of the matrix A is primitive modulo p . In this paper the case of a modulus $m = p^\alpha$, $\alpha \geq 2$, is considered where p is a prime number. It is shown that for $p \geq 3$ or $r \geq 2$ there exist matrix generators (1) with period length $(p^r - 1) \cdot p^{\alpha-1}$ for any starting vector which is nonzero modulo p , and a simple method is described for determining such a generator. Observe that $(p^r - 1) \cdot p^{\alpha-1}$ is the maximal possible period length according to the following technical lemma.

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2. Matrix Generators with Maximal Period Length.

LEMMA. Let $A \in \mathbf{Z}_{p^{\alpha+1}}^{r \times r}$, $\alpha \geq 1$, be a matrix which is nonsingular modulo p , and define matrices $A_n \in \mathbf{Z}_{p^{\alpha+1}}^{r \times r}$ by $A_n \equiv A^n \pmod{p^{\alpha+1}}$, $n \geq 0$. Let $\lambda_\alpha = \lambda(A, p^\alpha)$ and $\lambda_{\alpha+1} = \lambda(A, p^{\alpha+1})$ denote the period lengths of the matrix sequence $(A_n)_{n \geq 0}$ modulo p^α and modulo $p^{\alpha+1}$, respectively. Then

$$\lambda_{\alpha+1} = \begin{cases} \lambda_\alpha & \text{for } A^{\lambda_\alpha} \equiv I \pmod{p^{\alpha+1}}, \\ \lambda_\alpha \cdot p & \text{for } A^{\lambda_\alpha} \not\equiv I \pmod{p^{\alpha+1}}. \end{cases}$$

Proof. From $A^{\lambda_\alpha} \equiv I \pmod{p^\alpha}$ it follows that $A^{\lambda_\alpha} = I + p^\alpha \cdot B$ for some matrix $B \in \mathbf{Z}^{r \times r}$. Therefore,

$$A^{\lambda_\alpha \cdot p} = (I + p^\alpha \cdot B)^p = I + \binom{p}{1} \cdot p^\alpha \cdot B + \binom{p}{2} \cdot (p^\alpha \cdot B)^2 + \cdots + (p^\alpha \cdot B)^p,$$

which yields $A^{\lambda_\alpha \cdot p} \equiv I \pmod{p^{\alpha+1}}$, i.e., $\lambda_{\alpha+1}$ divides $\lambda_\alpha \cdot p$ according to Remark 2. From $A^{\lambda_{\alpha+1}} \equiv I \pmod{p^{\alpha+1}}$ it follows that $A^{\lambda_{\alpha+1}} \equiv I \pmod{p^\alpha}$, i.e., λ_α divides $\lambda_{\alpha+1}$ according to Remark 2, which proves the lemma. $\square$

The purpose of this paper is to prove the following result.

THEOREM. Let $B \in \mathbf{Z}_{p^\alpha}^{r \times r}$, $\alpha \geq 2$, be a matrix whose characteristic polynomial is primitive modulo p . Then

$$(2) \quad B^{p^r-1} \equiv I + p \cdot C \pmod{p^2}$$

for some matrix $C \in \mathbf{Z}_p^{r \times r}$. Let $D \in \mathbf{Z}_{p^{\alpha-1}}^{r \times r}$ denote an arbitrary matrix with $B \cdot D \equiv D \cdot B \pmod{p}$,

$$(3) \quad \det(D) \not\equiv 0 \pmod{p} \quad \text{for } p \geq 3,$$

and

$$(4) \quad \det(D) \equiv \det(D + I) \equiv 1 \pmod{2} \quad \text{for } p = 2.$$

Define a matrix $A \in \mathbf{Z}_{p^\alpha}^{r \times r}$ by

$$(5) \quad A \equiv B \cdot (I + p \cdot (C - D)) \pmod{p^\alpha}.$$

Then the period length of the vector sequence $(\vec{x}_n)_{n \geq 0}$ generated according to (1) with matrix A and modulus $m = p^\alpha$ is given by

$$\lambda(A, \vec{x}_0, p^\alpha) = (p^r - 1) \cdot p^{\alpha-1}$$

for any starting vector $\vec{x}_0 \in \mathbf{Z}_{p^\alpha}^r$ with $\vec{x}_0 \not\equiv \vec{0} \pmod{p}$.

Proof. The proof is subdivided into four parts (i) to (iv).

(i) Because of $A \equiv B \pmod{p}$ according to (5) it follows that

$$(6) \quad \lambda(A, \vec{x}_0, p) = \lambda(A, p) = p^r - 1$$

for any starting vector $\vec{x}_0 \in \mathbf{Z}_p^r \setminus \{\vec{0}\}$, since the characteristic polynomial of the matrix B is primitive modulo p . In particular, $B^{p^r-1} \equiv I \pmod{p}$ holds. Hence a matrix C with (2) exists. Observe that (2) yields $B \cdot C \equiv C \cdot B \pmod{p}$, which implies that $B \cdot (C - D) \equiv (C - D) \cdot B \pmod{p}$ because of the hypothesis $B \cdot D \equiv D \cdot B \pmod{p}$. Therefore (5) and (2) yield

$$(7) \quad \begin{aligned} A^{p^r-1} &\equiv [B \cdot (I + p \cdot (C - D))]^{p^r-1} \equiv B^{p^r-1} \cdot (I + (p^r - 1) \cdot p \cdot (C - D)) \\ &\equiv (I + p \cdot C) \cdot (I - p \cdot (C - D)) \equiv I + p \cdot D \pmod{p^2}. \end{aligned}$$

If $p = 2$ then it follows from (7) that

$$A^{2^r-1} = I + 2 \cdot D + 4 \cdot E$$

for some matrix $E \in \mathbf{Z}^{r \times r}$ and hence

$$A^{(2^r-1) \cdot 2} = (I + 2 \cdot D + 4 \cdot E)^2 = I + 4 \cdot D + 4 \cdot D^2 + 8 \cdot F$$

for some matrix $F \in \mathbf{Z}^{r \times r}$, i.e.,

$$A^{(2^r-1) \cdot 2} \equiv I + 4 \cdot D \cdot (D + I) \pmod{8}.$$

(ii) Now it is shown by induction that in case of $p \geq 3$,

$$(8) \quad A^{(p^r-1) \cdot p^\nu} \equiv I + p^{\nu+1} \cdot D \pmod{p^{\nu+2}}$$

for $0 \leq \nu \leq \alpha - 2$. Obviously, (7) is equivalent to (8) for $\nu = 0$. If (8) is valid for some ν with $0 \leq \nu \leq \alpha - 3$, then

$$A^{(p^r-1) \cdot p^\nu} = I + p^{\nu+1} \cdot D + p^{\nu+2} \cdot E_\nu$$

for some matrix $E_\nu \in \mathbf{Z}^{r \times r}$ and hence

$$A^{(p^r-1) \cdot p^{\nu+1}} = (I + p^{\nu+1} \cdot (D + p \cdot E_\nu))^p = I + p^{\nu+2} \cdot (D + p \cdot E_\nu) + p^{\nu+3} \cdot F_\nu$$

for some matrix $F_\nu \in \mathbf{Z}^{r \times r}$ because of $p \geq 3$, which yields

$$A^{(p^r-1) \cdot p^{\nu+1}} \equiv I + p^{\nu+2} \cdot D \pmod{p^{\nu+3}}.$$

Therefore (8) holds for $0 \leq \nu \leq \alpha - 2$. It can be similarly proved that in case of $p = 2$,

$$(9) \quad A^{(2^r-1) \cdot 2^\nu} \equiv I + 2^{\nu+1} \cdot D \cdot (D + I) \pmod{2^{\nu+2}}$$

for $1 \leq \nu \leq \alpha - 2$.

(iii) Because of (3), (4), (6), (7), (8) and (9) it follows from the lemma that

$$(10) \quad \lambda(A, p^{\nu+1}) = (p^r - 1) \cdot p^\nu$$

for $0 \leq \nu \leq \alpha - 1$. Note that if $\vec{x}_0 \not\equiv \vec{0} \pmod{p}$, then

$$D \cdot \vec{x}_0 \not\equiv \vec{0} \pmod{p} \quad \text{for } p \geq 3$$

and

$$D \cdot (D + I) \cdot \vec{x}_0 \not\equiv \vec{0} \pmod{p} \quad \text{for } p = 2$$

because of (3) and (4), respectively. Therefore (7), (8) and (9) show that

$$(11) \quad A^{(p^r-1) \cdot p^\nu} \cdot \vec{x}_0 \not\equiv \vec{x}_0 \pmod{p^{\nu+2}}$$

for $\vec{x}_0 \not\equiv \vec{0} \pmod{p}$ and $0 \leq \nu \leq \alpha - 2$.

(iv) Now it is proved by induction that

$$(12) \quad \lambda(A, \vec{x}_0, p^{\nu+1}) = (p^r - 1) \cdot p^\nu$$

for any starting vector $\vec{x}_0 \in \mathbf{Z}_{p^\alpha}^r$ with $\vec{x}_0 \not\equiv \vec{0} \pmod{p}$ and $0 \leq \nu \leq \alpha - 1$. Obviously, (6) is equivalent to (12) for $\nu = 0$. Now assume that (12) is valid for some ν with $0 \leq \nu \leq \alpha - 2$. Then

$$\lambda(A, \vec{x}_0, p^{\nu+2}) = \mu \cdot (p^r - 1) \cdot p^\nu$$

for some integer $\mu \geq 1$. Since

$$\lambda(A, \vec{x}_0, p^{\nu+2}) \neq (p^r - 1) \cdot p^\nu$$

according to (11), it follows that $\mu > 1$. Remark 1 and (10) imply that $\lambda(A, \vec{x}_0, p^{\nu+2})$ divides $(p^r - 1) \cdot p^{\nu+1}$ and hence $\mu = p$, which proves the theorem. $\square$

Observe that there exist primitive polynomials of degree r over the Galois field $\text{GF}(p)$ for every positive integer r and every prime number p . Such a polynomial, and hence a matrix $B \in \mathbb{Z}_p^{r \times r}$ which satisfies the hypothesis of the theorem, can be determined without any effort if p and r are small integers (see, e.g., Knuth [3, p. 28]).

Since the characteristic polynomial of the matrix B is primitive modulo p , it follows that $\det(B) \not\equiv 0 \pmod{p}$ and that $B \cdot \vec{x}_0 \not\equiv \vec{x}_0 \pmod{2}$ for $p = 2$, $r \geq 2$, and $\vec{x}_0 \not\equiv \vec{0} \pmod{2}$. Hence $\det(B+I) \equiv 1 \pmod{2}$ for $p = 2$ and $r \geq 2$. Therefore, the matrix $D \in \mathbb{Z}_{p^{\alpha-1}}^{r \times r}$ with $D \equiv B \pmod{p^{\alpha-1}}$ satisfies the hypothesis of the theorem if $p \geq 3$ or $r \geq 2$.

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Technische Hochschule Darmstadt
 Fachbereich Mathematik
 Schloßgartenstr. 7
 D-6100 Darmstadt, West Germany
 E-mail: x%xbr1dc3i@ddatd21.bitnet

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