ALGORITHMS FOR MATRIX EXTENSION AND ORTHOGONAL WAVELET FILTER BANKS OVER ALGEBRAIC NUMBER FIELDS

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ABSTRACT. As a finite dimensional linear space over the rational number field $\mathbb{Q}$, an algebraic number field is of particular importance and interest in mathematics and engineering. Algorithms using algebraic number fields can be efficiently implemented involving only integer arithmetics. We observe that all known finitely supported orthogonal wavelet low-pass filters in the literature have coefficients coming from an algebraic number field. Therefore, it is of theoretical and practical interest for us to consider orthogonal wavelet filter banks over algebraic number fields. In this paper, we formulate the matrix extension problem over any general subfield of $\mathbb{C}$ (including an algebraic number field as a special case), and we provide step-by-step algorithms to implement our main results. As an application, we obtain a satisfactory algorithm for constructing orthogonal wavelet filter banks over algebraic number fields. Several examples are provided to illustrate the algorithms proposed in this paper.

1. INTRODUCTION AND MOTIVATION

In a digital world, data and signals are often recorded using integers and dyadic rational numbers. For example, a 8-bit grey-scale image has its pixels taking integer values between 0 and 255. The rational number field $\mathbb{Q}$ has many advantages in scientific computing. Using integer arithmetics and having simple hardware implementation, algorithms using the rational number field are much more efficient than using floating point arithmetics, and avoid roundup error in computation. Consequently, it is one of the fascinating topics in applied mathematics and engineering (see [3, 16, 19]) to construct orthogonal wavelet filter banks with rational coefficients. Indeed, a few examples of orthogonal wavelet low-pass filters with rational coefficients have been reported in [1, 2, 3, 5]. Typically, the associated wavelets in such examples have pretty low regularity and vanishing moments. This significantly limits the use of such orthogonal wavelet low-pass filters in applications. The interest of seeking orthogonal wavelet filter banks with rational coefficients has been renewed recently in the paper Mo and Li [16], where tight framelet filter banks with rational coefficients and with two high-pass filters have been considered. Though a necessary and sufficient condition is obtained in [16] for a dyadic tight framelet

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filter bank with two high-pass filters having rational coefficients, almost none of the known examples in the literature can satisfy the necessary and sufficient condition in [16]. Moreover, the associated wavelets in the illustrative examples presented in [16] have only one vanishing moment. Note that a tight framelet filter bank is a generalization of an orthogonal wavelet filter bank by having more than necessary number of high-pass filters. Despite the interesting work in [16], constructing orthogonal wavelet filter banks or tight framelet filter banks with rational coefficients and with high vanishing moments remains as a challenging task.

We are motivated by [16] to examine again all the orthogonal wavelet filter banks available in the literature. We observe two interesting phenomena. First, though there are very few known examples of orthogonal wavelet low-pass filters with rational coefficients, we notice that to our best knowledge all known finitely supported orthogonal wavelet low-pass filters have their coefficients coming from an algebraic number field, often in the form $\mathbb{Q}(\sqrt{t_1}, \dots, \sqrt{t_n})$ for some positive integers $t_1, \dots, t_n$. For example, see [8] and [1, 2, 5, 13] for many examples of orthogonal wavelet low-pass filters with symmetry. We shall provide a natural explanation in Section 3 for this observation. Recall that an algebraic number field $\mathbb{A}$ is a finite field extension of the rational number field $\mathbb{Q}$. More precisely, $\mathbb{A} = \mathbb{Q}(t_1, \dots, t_n)$, where each of $t_1, \dots, t_n$ is a root of some polynomial with integer coefficients. It is well known that an algebraic number field $\mathbb{A}$, when viewed as a linear space over $\mathbb{Q}$, is a finite dimensional vector space over $\mathbb{Q}$. Consequently, the arithmetics over $\mathbb{A}$ can be implemented by combining the integer arithmetics and matrix/vector operations from linear algebra (for example, see [14] for more detail). In other words, algorithms over an algebraic number field $\mathbb{A}$ have the same order of complexity as those over the rational number field $\mathbb{Q}$. Consequently, it is of theoretical and practical interest for us to consider orthogonal wavelet filter banks over an algebraic number field.

Second, we observe that even though the orthogonal wavelet low-pass filters constructed in [1, Examples 1 and 2] and [2, (5.3)] have rational coefficients, the constructed high-pass filters, which are obtained by a brute force calculation using Gröbner bases, do not have rational coefficients; instead, such derived high-pass filters have rational coefficients only up to a multiplicative constant $\sqrt{t}$ for some positive integer t . In other words, according to the evidence from symbolic computation, even if an orthogonal wavelet low-pass filter has rational coefficients, the derived high-pass filters cannot always be expected to have rational coefficients. This motivates us to examine the classical matrix extension problem in [6, 9, 10, 15, 17, 19] and to seek the right formulation for the matrix extension problem and for the construction of orthogonal wavelet filter banks over an algebraic number field.

In the following, we shall state our formulation of the matrix extension problem over a general field including an algebraic number field as special cases. To do so, let us recall some notation and definitions. Let $\mathbb{F}$ denote a general subfield of $\mathbb{C}$ such that

$$(1.1) \quad \bar{x} \in \mathbb{F} \quad \text{if} \quad x \in \mathbb{F},$$

where $\bar{x}$ denotes the complex conjugate of a complex number $x \in \mathbb{C}$. That is, the subfield $\mathbb{F}$ is closed under complex conjugation. For example, $\mathbb{F}$ can be the rational number field $\mathbb{Q}$ or an algebraic number field $\mathbb{A}$ satisfying $\bar{x} \in \mathbb{A}$ for all $x \in \mathbb{A}$. By $\mathbb{F}[z, z^{-1}]$ we denote the set of all Laurent polynomials with coefficients in $\mathbb{F}$. We

shall use the following notation throughout the paper:

$$P^*(z) := P(z)^* := \sum_{k \in \mathbb{Z}} \overline{P_k}^T z^{-k} \quad \text{for} \quad P(z) = \sum_{k \in \mathbb{Z}} P_k z^k,$$

where $P_k \in \mathbb{C}^{r \times s}$ are $r \times s$ matrices of complex numbers. We also define $P_k^* := \overline{P_k}^T$ for a matrix P_k of complex numbers. With the above $\star$ notation, we often work on $P(z)$ with $z \in \mathbb{T} := \{\zeta \in \mathbb{C} \mid |\zeta| = 1\}$.

Now we are ready to formulate the matrix extension problem over a general field $\mathbb{F}$. Let P be an $r \times s$ matrix, with $1 \leq r \leq s$, of Laurent polynomials satisfying

- (i) P is paraunitary: $P(z)P^*(z) = I_r$ for all $z \in \mathbb{C} \setminus \{0\}$;
- (ii) P takes the form $P(z) = Q_0(z)D_0$, where Q_0 is an $r \times s$ matrix of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and D_0 is a diagonal matrix with all the entries of $D_0 D_0^*$ in $\mathbb{F}$.

The matrix extension problem over a general field $\mathbb{F}$ is to find an $s \times s$ extension matrix P_e of Laurent polynomials such that

- (1) $[I_r, \mathbf{0}]P_e = P$; that is, the submatrix of the first r rows of P_e is P ;
- (2) P_e is paraunitary: $P_e(z)P_e^*(z) = I_s$ for all $z \in \mathbb{C} \setminus \{0\}$;
- (3) P_e takes the form $P_e(z) = D_e Q_e(z)D_0$, where Q_e is an $s \times s$ matrix of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and D_e is a diagonal matrix with all the entries of $D_e D_e^*$ in $\mathbb{F}$.

Moreover, if the given matrix P has certain symmetry structure, then it is desirable that the extension matrix P_e also possesses certain symmetry structure. It is also important that the support of the coefficients of P_e is controllable in some way by that of the given matrix P .

We shall explain in Section 3 why the above formulation of the matrix extension problem over a general field $\mathbb{F}$ leads to a satisfactory solution in the setting of wavelet analysis and filter banks.

The classical formulation of the matrix extension problem corresponds to the special case $D_0 = D_e = I_s$, but only for subfields $\mathbb{F}$ of $\mathbb{C}$ satisfying the following property:

$$(1.2) \quad \bar{x} \in \mathbb{F} \quad \text{and} \quad \sqrt{y} \in \mathbb{F} \quad \forall x, y \in \mathbb{F} \quad \text{with} \quad y > 0.$$

For example, in the context of wavelet analysis and filter banks, the matrix extension problem without symmetry for $\mathbb{F} = \mathbb{R}$ or $\mathbb{C}$ was studied in Lawton, Lee, and Shen [15] and Vaidyanathan [19]. The matrix extension problem with symmetry has been considered in Petukhov [17] for $r = 1$ and $\mathbb{F} = \mathbb{R}$, and in Han [6] for $r = 1$ and a general subfield $\mathbb{F}$ of $\mathbb{C}$ satisfying (1.2). More recently, the matrix extension problem with symmetry has been studied in [10] by the authors for any r satisfying $1 \leq r \leq s$ and any subfield $\mathbb{F}$ of $\mathbb{C}$ satisfying the condition in (1.2). Clearly, the condition in (1.2) is satisfied if $\mathbb{F} = \mathbb{R}$ or $\mathbb{C}$ and therefore, [10] generalizes the results in [6, 17] to the most general setting.

However, even if all the coefficients of P are integers, the smallest such field $\mathbb{F}$, containing $\mathbb{Z}$ and satisfying the condition in (1.2), must contain $\sqrt{n}$ for all positive integers n . Therefore, any subfield $\mathbb{F}$ satisfying (1.2) can never be an algebraic number field and it must be an infinite dimensional linear space over $\mathbb{Q}$. Hence, in the setting of wavelet analysis and filter banks, despite the beautiful results and efficient algorithms on the matrix extension problem in [6, 10, 15, 17, 19], even if an orthogonal wavelet low-pass filter has all its coefficients from an algebraic number

field (which is indeed the case for all known finitely supported orthogonal wavelet low-pass filters in the literature, for example, see [8] for systematic construction and numerous examples of orthogonal wavelet low-pass filters with symmetry), the derived orthogonal wavelet high-pass filters cannot be guaranteed to have coefficients in an algebraic number field as well.

For a matrix $P(z) = \sum_{k=m}^n P_k z^k$ of Laurent polynomials such that $P_m \neq \mathbf{0}$ and $P_n \neq \mathbf{0}$ with $m, n \in \mathbb{Z}$, we define the filter support and length of P to be $\text{fsupp}(P) := [m, n]$ and $\text{len}(P) := n - m$, respectively. Throughout the paper, $\mathbf{0}$ always denotes a zero matrix whose size can be determined from the context.

Let P be an $r \times s$ matrix, with $1 \leq r \leq s$, of Laurent polynomials satisfying $P(z)P^*(z) = I_r$ for all $z \in \mathbb{C} \setminus \{0\}$ and P takes the form $P(z) = Q_0(z)D_0$, where Q_0 is an $r \times s$ matrix of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and D_0 is a diagonal matrix with all the entries of $D_0 D_0^*$ in $\mathbb{F}$. To study and formulate the matrix extension problem over a general subfield $\mathbb{F}$, in this paper we shall investigate the structure of P by factorizing the matrix P into a product of elementary matrices, that is, we study a more general problem: the matrix factorization problem over a general subfield $\mathbb{F}$ of $\mathbb{C}$. More precisely, we shall construct a sequence of elementary matrices $A_0, \dots, A_J$ such that

- (1) $PA_0 \cdots A_J = [I_r, 0]$;
- (2) $\text{len}(PA_0 \cdots A_j) < \text{len}(PA_0 \cdots A_{j-1})$ for all $j = 0, \dots, J$;
- (3) $A_j^*(z)A_j(z) = I_s$ for all $j = 0, \dots, J$ and $\text{len}(A_j) \leq 1$;
- (4) Each A_j takes the form $D_{j-1}^* V_j(z) D_j$ with $V_j(z) = V_j D_j(z)$, where all D_j are diagonal matrices such that all entries of $D_j^* D_j$ in $\mathbb{F}$, V_j is a constant invertible matrix in $\mathbb{F}$, and D_j is a diagonal matrix whose diagonal entries are monomial Laurent polynomials.

Consequently, we have the following representation of P : $P = [I_r, \mathbf{0}] A_J^* \cdots A_0^*$, which factorizes the paraunitary matrix P into a product of elementary paraunitary matrices A_j^* . When P has symmetry, the matrix extension problem and the matrix factorization problem over a general subfield $\mathbb{F}$ are even much more complicated; see Section 4 for details. The matrix extension problem and the matrix factorization problem are of importance in engineering and system sciences. For example, they play a central role in engineering for the design and study of paraunitary perfect reconstruction filter banks (see [19, Sections 13 and 14]); they are also indispensable tools in the general theory of linear systems in system science; see [20] and references therein.

Our major contribution of this paper is a proper formulation of the matrix extension problem and the matrix factorization problem with or without symmetry over a general field, and provides a satisfactory solution to them. Though we shall use some interesting ideas from [6, 10, 15, 17, 19], our results in this paper are not simple generalizations of the results in [6, 10, 15, 17, 19] and new ideas are needed in order to satisfactorily solve the matrix extension problem over a general field without the condition in (1.2). More precisely, to avoid the square root requirement of the subfield $\mathbb{F}$ in (1.2), we have to formulate the matrix extension and factorization problems in an appropriate way by restricting ourselves to considering only elementary matrices with additional structures. Our main new idea is a key observation made in Lemma 2.2 which allows us to avoid directly using the square roots of positive elements in $\mathbb{F}$ in our elementary matrices. Instead, we require that elementary matrices should have the additional structure $D_0^* V D_1$ such that V is

an invertible matrix in $\mathbb{F}$ and D_0, D_1 are diagonal matrices whose diagonal entries are square roots of positive elements in $\mathbb{F}$. In this way, the troubling square roots of positive elements in $\mathbb{F}$ appearing in our elementary matrices are restricted in the diagonal matrices D 's in a nested fashion (see Theorem 2.1 for details). On the other hand, all the papers [6, 10, 15, 19] employ the Householder matrices to transform a constant vector into a normalized one with only one nonzero entry. However, using Householder matrices, one has to unavoidably use the square roots of positive elements in $\mathbb{F}$ in elementary matrices in a nonseparable way, which in turn forces one to use a subfield $\mathbb{F}$ with the square root requirement in (1.2). Once we have the key observation made in Lemma 2.2, to establish the results and algorithms in this paper, we shall follow the main ideas made in [6, 10] by appropriately modifying the key procedures in [6, 10] to the new setting and formulation.

The structure of the paper is as follows. In Section 2, we shall state our result on the matrix extension problem without symmetry constraint over a general field and its associated step-by-step algorithm for deriving the extension matrix P_e from P . In Section 3, we will discuss the application of the matrix extension problem to the orthogonal wavelet filter banks without symmetry. We shall explain in Section 3 that our formulation of the matrix extension problem over a general field $\mathbb{F}$ leads to a satisfactory solution in the setting of wavelet analysis and filter banks. In Section 4, we shall present our result on the matrix extension problem with symmetry constraint over a general field and its associated step-by-step algorithm for deriving the extension matrix P_e with symmetry structure from P . In Section 5, we shall discuss the application of our results on the matrix extension problem over a general field to wavelet analysis and filter banks. We shall see that our result leads to a satisfactory solution to the design of symmetric filter banks over an algebraic number field. We shall also provide several examples in Section 6 to demonstrate our algorithms for designing orthogonal wavelet filter banks over an algebraic number field. Proofs of some key lemmas are postponed to Section 7.

2. THE MATRIX EXTENSION PROBLEM WITHOUT SYMMETRY CONSTRAINT

In this section, we shall investigate the matrix extension problem without symmetry constraint over a general subfield $\mathbb{F}$ of the complex number field $\mathbb{C}$. We also present a step-by-step algorithm for finding the extension matrix P_e . To do so, let us introduce some notation first.

Without further mention, $\mathbb{F}$ always denotes a general subfield of $\mathbb{C}$ satisfying (1.1). We have the following result on the matrix extension problem without symmetry constraint over a general subfield $\mathbb{F}$:

Theorem 2.1. *Let P be an $r \times s$ matrix, with $1 \leq r \leq s$, of Laurent polynomials such that P takes the following form:*

$$(2.1) \quad P(z) = Q_0(z)D_0, \quad z \in \mathbb{C} \setminus \{0\},$$

where Q_0 is an $r \times s$ matrix of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and D_0 is an $s \times s$ diagonal matrix with all the entries of $D_0 D_0^*$ in $\mathbb{F}$. Then $P(z)P^*(z) = I_r$ for all $z \in \mathbb{C} \setminus \{0\}$, if and only if, there exists an $s \times s$ matrix P_e of Laurent polynomials such that

- (i) $[I_r, \mathbf{0}]P_e = P$; that is, the submatrix of the first r rows of P_e is P ;
- (ii) P_e is paraunitary: $P_e(z)P_e^*(z) = I_s$ for all $z \in \mathbb{C} \setminus \{0\}$;

(iii) P_e can be represented as a product of $s \times s$ elementary matrices:

$$P_e(z) = D_e P_J(z) P_{J-1}(z) \cdots P_1(z) P_0(z) D_0, \quad z \in \mathbb{C} \setminus \{0\},$$

where $P_j, j = 0, \dots, J$ are $s \times s$ matrices of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ such that $\text{len}(P_j) \leq 1$, and D_e is an $s \times s$ diagonal matrix with all the entries of $D_e D_e^*$ in $\mathbb{F}$;

(iv) the length of P_e is controlled by that of P in the following sense:

$$\text{len}([P_e]_{j,k}) \leq \max_{1 \leq n \leq r} \text{len}([P]_{n,k}) \quad \text{for all } 1 \leq j, k \leq s.$$

We make some remarks about Theorem 2.1. The integer J in item (iii) of Theorem 2.1 always satisfies $J \leq \text{len}(P)$. Each of the elementary matrices $P_j, j = 0, \dots, J-1$ is a product of a constant matrix in $\mathbb{F}$ and a diagonal matrix whose diagonal elements are either 1 or z^{-1} . The matrix P_J is a product of a constant matrix in $\mathbb{F}$ and a diagonal matrix whose diagonal elements are monomials.

Theorem 2.1 shows that all the coefficients of the paraunitary extension matrix P_e are not “far away” from those of $\mathbb{F}$, provided that the given matrix P has the special form as in (2.1). In particular, when $\mathbb{F}$ is an algebraic number field $\mathbb{A}$, the extension matrix P_e is of the form $P_e = D_e Q_e D_0$ with Q_e being a matrix of Laurent polynomials with coefficients in $\mathbb{A}$. Consequently, all the coefficients of P_e are from an algebraic number field $\tilde{\mathbb{A}}$, where $\tilde{\mathbb{A}}$ is the field extension over $\mathbb{A}$ by adding the entries of D_0 and D_e . For the result on matrix extension with symmetry constraint over a general subfield $\mathbb{F}$; see Theorem 4.1 in Section 4 and also see Examples 6.1 – 6.4 in Section 6 for its application.

For row vectors u and v in $\mathbb{C}^{1 \times r}$, we define $v^* := \bar{v}^T$, $\langle u, v \rangle := uv^*$, and $\|v\| := \sqrt{\langle v, v \rangle}$. To prove Theorem 2.1 we need the following auxiliary result.

Lemma 2.2. *Let $f = gD_0$ be a $1 \times s$ nonzero row vector for some $1 \times s$ row vector g with all the entries in $\mathbb{F}$ and some nonsingular diagonal matrix D_0 with all the entries of $D_0 D_0^*$ in $\mathbb{F}$. Then there exists an $s \times s$ unitary matrix U_f of the form $U_f = D_0^* V_f D_1$ such that $U_f U_f^* = I_s$ and $f U_f = [\|f\|, 0, \dots, 0]^T$, where V_f is a nonsingular matrix with all the entries in $\mathbb{F}$, and where D_1 is some diagonal matrix with all the entries of $D_1 D_1^*$ in $\mathbb{F}$.*

Proof. Let $\tilde{g}_1 := g$. Since $\tilde{g}_1 \neq 0$, we can find $1 \times s$ vectors $g_2, \dots, g_s$ in $\mathbb{F}^{1 \times s}$ such that $\tilde{g}_1, g_2, \dots, g_s$ are linearly independent. For instance, assuming $g = [g_{11}, \dots, g_{1s}]$ with $g_{11} \neq 0$, we can choose $g_j = e_j = [0, \dots, 1, \dots, 0]$, the standard j th unit coordinate vector in $\mathbb{F}^{1 \times s}$, for $j = 2, \dots, s$. Then $\tilde{g}_1, g_2, \dots, g_s$ are linearly independent. Because D_0 is nonsingular, all the vectors $\tilde{g}_1 D_0, g_2 D_0, \dots, g_s D_0$ are linearly independent. Applying the Gram-Schmidt orthogonalization process to the linearly independent vectors $\tilde{g}_1 D_0, g_2 D_0, \dots, g_s D_0$ as follows:

$$\begin{aligned} v_1 &:= gD_0 = \tilde{g}_1 D_0, \\ v_2 &:= g_2 D_0 - \frac{\langle g_2 D_0, \tilde{g}_1 D_0 \rangle}{\langle \tilde{g}_1 D_0, \tilde{g}_1 D_0 \rangle} \tilde{g}_1 D_0 = (g_2 - c_{2,1} \tilde{g}_1) D_0 =: \tilde{g}_2 D_0, \\ &\vdots \\ v_s &:= g_s D_0 - \sum_{j=1}^{s-1} \frac{\langle g_s D_0, \tilde{g}_j D_0 \rangle}{\langle \tilde{g}_j D_0, \tilde{g}_j D_0 \rangle} \tilde{g}_j D_0 = (g_s - c_{s,1} \tilde{g}_1 - \cdots - c_{s,s-1} \tilde{g}_{s-1}) D_0 =: \tilde{g}_s D_0, \end{aligned}$$

where

$$c_{k,j} := \frac{\langle \mathbf{g}_k D_0, \tilde{\mathbf{g}}_j D_0 \rangle}{\langle \tilde{\mathbf{g}}_j D_0, \tilde{\mathbf{g}}_j D_0 \rangle} = \frac{\langle \mathbf{g}_k D_0 D_0^*, \tilde{\mathbf{g}}_j \rangle}{\langle \tilde{\mathbf{g}}_j D_0 D_0^*, \tilde{\mathbf{g}}_j \rangle}, \quad 1 \leq j < k \leq s$$

and

$$\tilde{\mathbf{g}}_k := \mathbf{g}_k - c_{k,1}\tilde{\mathbf{g}}_1 - \cdots - c_{k,k-1}\tilde{\mathbf{g}}_{k-1}, \quad 1 \leq k \leq s.$$

By our assumption, it is not difficult to verify that all $c_{k,j}$ and all entries of $\tilde{\mathbf{g}}_k$ are in $\mathbb{F}$ for all $1 \leq j < k \leq s$. It is also easy to check that $\langle \mathbf{v}_j, \mathbf{v}_k \rangle = \delta(j-k)\|\mathbf{v}_j\|^2$ for all $j, k = 1, \dots, s$, where δ denotes the Dirac sequence such that $\delta(0) = 1$ and $\delta(k) = 0$ for $k \neq 0$.

Finally, define $V_{\mathbf{f}} := [\tilde{\mathbf{g}}_1^*, \dots, \tilde{\mathbf{g}}_s^*]$ and $D_1 := \text{diag}(\frac{\zeta_1}{\|\mathbf{v}_1\|}, \dots, \frac{\zeta_s}{\|\mathbf{v}_s\|})$ for any $\zeta_1, \dots, \zeta_s \in \mathbb{T} := \{\zeta \in \mathbb{C} \mid |\zeta| = 1\}$. It is straightforward to verify that $U_{\mathbf{f}} := D_0^* V_{\mathbf{f}} D_1$ is the desired unitary matrix. $\square$

Setting $\zeta_1 = \cdots = \zeta_s = 1$ in the above proof, we see that $D_1 = \text{diag}(\frac{1}{\|\mathbf{v}_1\|}, \dots, \frac{1}{\|\mathbf{v}_s\|})$. Define $\mathbb{F}_+^{1/2} := \{\sqrt{x} \mid x \in \mathbb{F}, x \geq 0\}$. Then all D_1 's constructed by employing Lemma 2.2 in this paper can be chosen in such a way that all entries of D_1 's are from $\mathbb{F}_+^{1/2}$.

The role of $U_{\mathbf{f}}$ is to normalize the vector $\mathbf{f}$ by reducing $\mathbf{f}$ into a vector having only one nonzero entry. Note that $U_{\mathbf{f}}$ does not affect the zero entries of $\mathbf{f}$ and there exists a permutation matrix E such that

$$(2.2) \quad [U_{\mathbf{f}} E]_{j,:} = ([U_{\mathbf{f}}]_{:,j})^T = \mathbf{e}_j, \quad \text{provided} \quad [\mathbf{f}]_j = 0.$$

Here, $[A]_{j,:}$, $[A]_{:,j}$ denote the j th row, j th column of a matrix A , respectively, and $[\mathbf{f}]_j$ is the j th entry of a vector $\mathbf{f}$. For simplicity, we also define $U_{\mathbf{f}} = I_s$ for $\mathbf{f} = \mathbf{0}$ and $U_{\mathbf{f}} = \emptyset$ for $\mathbf{f} = \emptyset$, where $\emptyset$ is the emptyset. Throughout the paper, *up to a permutation* means there exists a permutation matrix E such that the identity holds *exactly* with the corresponding matrix A being replaced by AE (e.g., see (2.3)).

As a direct consequence of Lemma 2.2, for a general $r \times s$ constant matrix F taking the form $F = G D_0$, we can construct a unitary matrix $U_F = D_0^* V_F D_1$ so that $F U_F = [R, \mathbf{0}]$ and up to a permutation,

$$(2.3) \quad [U_F]_{j,:} = ([U_F]_{:,j})^T = \mathbf{e}_j, \quad \text{provided} \quad [F]_{:,j} = \mathbf{0},$$

which can be summarized as the following corollary.

Corollary 2.3. *Let $F = G D_0$ be an $r \times s$ matrix, where G is an $r \times s$ matrix with all the entries in $\mathbb{F}$ and D_0 is a nonsingular diagonal matrix with all the entries of $D_0 D_0^*$ in $\mathbb{F}$. Then there exists an $s \times s$ unitary matrix U_F of the form $U_F = D_0^* V_F D_1$ such that $U_F U_F^* = I_s$ and $F U_F = [R, \mathbf{0}]$ for some lower triangular matrix R , where V_F is a nonsingular matrix with all the entries of V_F in $\mathbb{F}$ and D_1 is some diagonal matrix with all the entries of $D_1 D_1^*$ in $\mathbb{F}$.*

Proof. The main idea is to apply Lemma 2.2 to F row by row.

Let $\mathbf{f}_1 := [F]_{1,:} = [G]_{1,:} D_0$ be the first row of F . By Lemma 2.2, we can construct a unitary matrix $U_1 = D_0^* V_1 D_r$ satisfying $\mathbf{f}_1 U_1 = [\|\mathbf{f}_1\|, 0, \dots, 0]$ for some nonsingular $s \times s$ matrix $V_1 \in \mathbb{F}^{s \times s}$ and some diagonal matrix D_r with all the entries of $D_r D_r^*$ in $\mathbb{F}$.

Next, let $F_1 = F U_1$. Then $F_1 = G_1 D_r$ with $G_1 = G D_0 D_0^* V_1 \in \mathbb{F}^{r \times s}$. Let $\mathbf{f}_2 := [F_1]_{2,:} = [G_1]_{2,:} D_r$ be the second row of F_1 . Applying Lemma 2.2 to $\tilde{\mathbf{f}}_2 := [\mathbf{f}_2]_{2:s}$, the second row of F_1 ignoring the first entry of $\mathbf{f}_2$, we can, through a simple extension, find an $s \times s$ unitary matrix $U_2 = D_r^* V_2 D_{r-1}$ satisfying $\mathbf{f}_2 U_2 =$

$[[\mathbf{f}_2]_1, \|\tilde{\mathbf{f}}_2\|, 0, \dots, 0]$ for some nonsingular $s \times s$ matrix $V_2 \in \mathbb{F}^{s \times s}$ and some diagonal matrix D_{r-1} with all the entries of $D_{r-1}D_{r-1}^*$ in $\mathbb{F}$.

Repeating the above procedure, we can construct a sequence of $s \times s$ unitary matrices $U_1 = D_0^*V_1D_r$, $U_2 = D_r^*V_2D_{r-1}$, $\dots$, $U_r = D_2^*V_rD_1$ by Lemma 2.2, which reduce the matrix F to a lower triangular matrix. That is $FU_1 \cdots U_r = [R, \mathbf{0}]$ for some lower triangular matrix R . Let $U_F := D_0^*V_FD_1$ with $V_F := V_1D_rD_r^* \cdots V_{r-1}D_2D_2^*V_r$. Then U_F is a desired matrix. $\square$

Let $\mathbf{p}$ be a $1 \times s$ vector of Laurent polynomials taking the form $\mathbf{p}(z) = \mathbf{q}_0(z)D_0$. Suppose $\text{fsupp}(\mathbf{p}) = [\ell, h]$ with $h - \ell > 0$. Then $\mathbf{p}$ can be written as

$$\mathbf{p}(z) = \mathbf{f}_\ell z^\ell + \cdots + \mathbf{f}_h z^h = (\mathbf{g}_\ell z^\ell + \cdots + \mathbf{g}_h z^h)D_0, \quad z \in \mathbb{C} \setminus \{0\}.$$

Suppose that $\mathbf{p}$ is paraunitary, i.e., $\mathbf{p}(z)\mathbf{p}^*(z) = 1$ for all $z \in \mathbb{C} \setminus \{0\}$. For the row vector $\mathbf{f}_\ell$, we can construct a unitary matrix $U_{\mathbf{f}_\ell}$ by Lemma 2.2 such that $U_{\mathbf{f}_\ell}U_{\mathbf{f}_\ell}^* = I_s$ and $\mathbf{f}_\ell U_{\mathbf{f}_\ell} = [|\mathbf{f}_\ell|, 0, \dots, 0]$. Because $\mathbf{p}\mathbf{p}^* = 1$, we have $\langle \mathbf{f}_\ell U_{\mathbf{f}_\ell}, \mathbf{f}_h U_{\mathbf{f}_\ell} \rangle = 0$. Consequently, $\mathbf{f}_h U_{\mathbf{f}_\ell}$ must take the form $\mathbf{f}_h U_{\mathbf{f}_\ell} = [0, *, \dots, *]$, where $*$ denotes some number in $\mathbb{F}$. Using a diagonal matrix $D(z) := \text{diag}(1, z^{-1}, \dots, z^{-1})$, we can reduce the length of $\mathbf{p}$ by 1. Replacing $\mathbf{p}$ by $\mathbf{p}U_{\mathbf{f}_\ell}D$ and repeating this procedure, we can reduce the length of $\mathbf{p}$ step by step to 0. This is the main idea to prove Theorem 2.1 for the case $r = 1$. The same idea, by employing Corollary 2.3 instead of Lemma 2.2, yields the following proof of Theorem 2.1.

Proof of Theorem 2.1. The sufficiency part is trivial. We next prove the necessity part. Suppose that $\mathbf{P}$ takes the form in (2.1) and $\mathbf{P}(z)\mathbf{P}^*(z) = I_r$ for all $z \in \mathbb{C} \setminus \{0\}$. Assume that $\text{fsupp}(\mathbf{P}) = [\ell, h]$ with $h - \ell > 0$. Then $\mathbf{P} = \mathbf{Q}_0 D_0$ takes the form

$$\mathbf{P}(z) = F_\ell z^\ell + \cdots + F_h z^h = (G_\ell z^\ell + \cdots + G_h z^h)D_0, \quad z \in \mathbb{C} \setminus \{0\},$$

where F_j, G_j for $j = \ell, \dots, h$ are $r \times s$ constant matrices and $\mathbf{Q}_0(z) = G_\ell z^\ell + \cdots + G_h z^h$. We now perform the following steps to reduce the length of $\mathbf{P}$ by at least 1.

- (1) Since $F_\ell = G_\ell D_0 \neq \mathbf{0}$, by Corollary 2.3, we can construct a unitary matrix U_{F_ℓ} of the form $U_{F_\ell} = D_0^*V_{F_\ell}D_1$ such that $U_{F_\ell}U_{F_\ell}^* = I_s$ and $F_\ell U_{F_\ell} = [R, \mathbf{0}]$, where V_{F_ℓ} is a nonsingular matrix with all the entries in $\mathbb{F}$, D_1 is a diagonal matrix with all the entries of $D_1 D_1^*$ in $\mathbb{F}$, and R is an $r \times m$ lower triangular matrix with m being the rank of F_ℓ .
- (2) In view of the paraunitary property, $\mathbf{P}(z)\mathbf{P}^*(z) = I_r$ for all $z \in \mathbb{C} \setminus \{0\}$ deduces that $(F_\ell U_{F_\ell})(F_h U_{F_\ell})^* = \mathbf{0}$. Hence, $\mathbf{P}U_{F_\ell}$ must take the following form:

$$\mathbf{P}(z)U_{F_\ell} = [R, \mathbf{0}_R]z^\ell + \cdots + [\mathbf{0}_R, \tilde{R}]z^h, \quad z \in \mathbb{C} \setminus \{0\},$$

where $\mathbf{0}_R$ is the $r \times m$ zero matrix having the same size as R .

- (3) Define $D_0(z) := \text{diag}(\mathbf{1}_m, z^{-1}\mathbf{1}_{s-m})$ and $\tilde{\mathbf{P}}(z) := \mathbf{P}(z)U_{F_\ell}D_0(z)$. Here $\mathbf{1}_m$ denotes the $1 \times m$ row vector $[1, \dots, 1]$. Then $\tilde{\mathbf{P}}(z) = \mathbf{Q}_1(z)D_1$, with $\mathbf{Q}_1(z) := \mathbf{Q}_0(z)D_0D_0^*V_{F_\ell}D_0(z)$ being an $r \times s$ matrix of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$, satisfies $\text{fsupp}(\tilde{\mathbf{P}}) = [\ell, h-1]$. Define $A_0(z) := U_{F_\ell}D_0(z) = D_0^*V_0(z)D_1$ with $V_0(z) := V_{F_\ell}D_0(z)$. Then $A_0(z)$ is paraunitary.

Repeating the above steps, we can construct a sequence of paraunitary matrices $A_j(z) = D_j^*V_j(z)D_{j+1}$, $j = 0, \dots, J$ such that $\mathbf{P}A_0 \cdots A_J = [I_r, \mathbf{0}]$. Define $\mathbf{P}_e := A_J^* \cdots A_0^*$. Then it is easy to show that $\mathbf{P}_e$ takes the form as in Item (iii). Items (i) – (iii) follow directly from the above construction. Item (iv) follows from the property of U_{F_ℓ} in (2.3). $\square$

Based on the constructive proof of Theorem 2.1, we provide an algorithm (see Algorithm 1) for the matrix extension problem without symmetry constraint. In the algorithm and this paper, for an $r \times s$ matrix $P(z) = \sum_{k \in \mathbb{Z}} P_k z^k$ of Laurent polynomials, $\text{coeff}(P, k)$ refers to the coefficient matrix P_k of z^k .

Algorithm 1. Matrix Extension without Symmetry Constraint

- (a) **Input.** An $r \times s$ paraunitary matrix $P = Q_0 D_0$ as in Theorem 2.1.
 - (b) **Output.** A paraunitary extension matrix P_e satisfying Items (i) – (iv) of Theorem 2.1.
 - (c) **Initialization.** $\tilde{P} \leftarrow P$. $P_e \leftarrow I_s$. $D_1 \leftarrow D_0$. $J \leftarrow 0$.
 - (d) **Support Reduction.**
 - 1: **while** $\text{len}(\tilde{P}) \neq 0$ **do**
 - 2: $[\ell, h] \leftarrow \text{fsupp}(\tilde{P})$. $F \leftarrow \text{coeff}(\tilde{P}, \ell)$. F is of the form $F = G D_1$.
 - 3: Construct U_F by Corollary 2.3 such that $U_F U_F^* = I_s$, $U_F = D_1^* V_F D_2$, and $F U_F = [R, \mathbf{0}]$ for an $r \times m$ lower triangular matrix R with m being the rank of F .
 - 4: $D(z) \leftarrow \text{diag}(\mathbf{1}_m, \mathbf{1}_{s-m} z^{-1})$.
 - 5: $\tilde{P} \leftarrow \tilde{P} U_F D$. $P_J \leftarrow (V_F D D_2 D_2^*)^*$. $P_e \leftarrow (U_F D)^* P_e$.
 - 6: $D_1 \leftarrow D_2$. $J \leftarrow J + 1$.
 - 7: **end while**
 - (e) **Finalization.**
 - 8: $F \leftarrow \text{coeff}(\tilde{P}, \ell)$, where $\text{fsupp}(\tilde{P}) = [\ell, \ell]$ since $\text{len}(\tilde{P}) = 0$. Construct U_F by Corollary 2.3 such that $U_F U_F^* = I_s$, $U_F = D_1^* V_F D_2$, and $F U_F = [I_r, \mathbf{0}]$.
 - 9: $D(z) \leftarrow z^{-\ell} I_s$. $P_e \leftarrow (U_F D)^* P_e$. $P_{J+1} \leftarrow (V_F D)^*$. $D_e \leftarrow D_2^*$.
-

3. APPLICATION TO ORTHOGONAL WAVELET FILTER BANKS

In this section, we shall discuss the application of our result on the matrix extension problem without symmetry in Section 2 to $\mathbf{d}$ -orthogonal wavelet filter banks in electronic engineering and wavelet analysis. We shall also explain in this section why our formulation of the matrix extension problem over a general field leads to a satisfactory solution to the construction of orthogonal wavelet filter banks over algebraic number fields.

Let us first recall some definitions and notation. We say that $\mathbf{d}$ is a *dilation factor* if $\mathbf{d}$ is an integer with $|\mathbf{d}| > 1$. For simplicity of presentation, we further assume that $\mathbf{d}$ is positive, while the case of a negative dilation factor can be handled similarly by a slight modification of the statements in this paper.

Recall that in this paper $\mathbb{F}$ always denotes a general subfield of $\mathbb{C}$ satisfying (1.1). A filter $a = \{a(k)\}_{k \in \mathbb{Z}} : \mathbb{Z} \rightarrow \mathbb{F}^{r \times r}$ with multiplicity r is a finitely supported sequence of $r \times r$ matrices on $\mathbb{Z}$. The z -transform or symbol of the filter a is defined to be

$$\mathbf{a}(z) := \sum_{k \in \mathbb{Z}} a(k) z^k, \quad z \in \mathbb{C} \setminus \{0\},$$

which is a matrix of Laurent polynomials with coefficients in $\mathbb{F}^{r \times r}$. Moreover, the *polyphase components* of $\mathbf{a}$ (or a) are defined by

$$(3.1) \quad \mathbf{a}^{[\gamma]}(z) := \sum_{k \in \mathbb{Z}} a(\gamma + \mathbf{d}k) z^k, \quad \gamma \in \mathbb{Z}.$$

We say that $\mathbf{a}$ (or a) is a $\mathbf{d}$ -orthogonal wavelet filter if

$$(3.2) \quad \sum_{\gamma=0}^{\mathbf{d}-1} \mathbf{a}^{[\gamma]}(z) \mathbf{a}^{[\gamma]}(z)^* = \mathbf{d}^{-1} I_r, \quad z \in \mathbb{C} \setminus \{0\}.$$

Let $b_1, \dots, b_{\mathbf{d}-1} : \mathbb{Z} \rightarrow \mathbb{C}^{r \times r}$ be filters with multiplicity r . We say that the set of filters $\{a; b_1, \dots, b_{\mathbf{d}-1}\}$ is a $\mathbf{d}$ -orthogonal wavelet filter bank if the following polyphase matrix

$$(3.3) \quad \mathcal{P}(z) = \sqrt{\mathbf{d}} \begin{bmatrix} \mathbf{a}^{[0]}(z) & \cdots & \mathbf{a}^{[\mathbf{d}-1]}(z) \\ \mathbf{b}_1^{[0]}(z) & \cdots & \mathbf{b}_1^{[\mathbf{d}-1]}(z) \\ \vdots & \ddots & \vdots \\ \mathbf{b}_{\mathbf{d}-1}^{[0]}(z) & \cdots & \mathbf{b}_{\mathbf{d}-1}^{[\mathbf{d}-1]}(z) \end{bmatrix}, \quad z \in \mathbb{C} \setminus \{0\}$$

is paraunitary; that is, $\mathcal{P}(z) \mathcal{P}^*(z) = I_{dr}$ for all $z \in \mathbb{C} \setminus \{0\}$, where each $\mathbf{b}_m^{[\gamma]}$ is a polyphase component of $\mathbf{b}_m$ defined similarly as in (3.1) for $m, \gamma = 0, \dots, \mathbf{d} - 1$, respectively. Without any a priori condition, any wavelet filter bank naturally corresponds to a frequency-based framelet in the distribution space; see Han [7] for more detail.

There are two major tasks in the construction of orthogonal wavelet filter banks. One is to construct a $\mathbf{d}$ -orthogonal wavelet low-pass filter a with some desirable properties, and the other is to derive the associated high-pass filters $b_1, \dots, b_{\mathbf{d}-1}$ from a given $\mathbf{d}$ -orthogonal wavelet low-pass filter a so that $\{a; b_1, \dots, b_{\mathbf{d}-1}\}$ is a $\mathbf{d}$ -orthogonal wavelet filter bank.

The matrix extension problem plays a major role in the second part for deriving the high-pass filters $b_1, \dots, b_{\mathbf{d}-1}$ from a given $\mathbf{d}$ -orthogonal wavelet low-pass filter a . As an application of Theorem 2.1, we can construct the high-pass filters $b_1, \dots, b_{\mathbf{d}-1}$ systematically when the $\mathbf{d}$ -orthogonal wavelet filter a is given. In fact, let $\mathbf{P} := \sqrt{\mathbf{d}}[\mathbf{a}^{[0]}, \dots, \mathbf{a}^{[\mathbf{d}-1]}] =: \sqrt{\mathbf{d}}\mathbf{Q}_0$. Then $\mathbf{P} = \mathbf{Q}_0 \mathbf{D}_0$, with $\mathbf{D}_0 := \sqrt{\mathbf{d}}\mathbf{I}_s$, satisfies the conditions in Theorem 2.1. Thus, we can extend $\mathbf{P}$ to a full paraunitary matrix $\mathbf{P}_e(z) =: \mathcal{P}(z)$ as in (3.3), from which it is not difficult to construct high-pass filters $b_1, \dots, b_{\mathbf{d}-1}$; see (5.11). More importantly, Theorem 2.1 guarantees that the coefficients of the high-pass filters are not “far away” from the field generated by the coefficients of the low-pass filter. This is summarized in Theorem 3.1 as follows. For its proof, see the proof of the more detailed version in Theorem 5.2 of Section 5.

Theorem 3.1. *Let $a : \mathbb{Z} \rightarrow \mathbb{F}^{r \times r}$ be a $\mathbf{d}$ -orthogonal wavelet low-pass filter with multiplicity r . Then there exist high-pass filters $b_1, \dots, b_{\mathbf{d}-1} : \mathbb{Z} \rightarrow \mathbb{C}^{r \times r}$ such that $\{a; b_1, \dots, b_{\mathbf{d}-1}\}$ forms a $\mathbf{d}$ -orthogonal wavelet filter bank and $b_m = D_m \tilde{b}_m$ for $m = 1, \dots, \mathbf{d} - 1$, where $\tilde{b}_m : \mathbb{Z} \rightarrow \mathbb{F}^{r \times r}$ has coefficients in $\mathbb{F}$ and D_m is some diagonal matrix with all the entries of $D_m D_m^*$ in $\mathbb{F}$.*

Algorithm 3 in Section 5 provides the detail for the construction of the high-pass filters $b_1, \dots, b_{\mathbf{d}-1}$ in Theorem 3.1 and see Examples 6.1 – 6.4 in Section 6 for the illustration. In the rest of this section we shall explain why almost all finitely supported $\mathbf{d}$ -orthogonal wavelet filters have coefficients from an algebraic number field, and why our formulation of the matrix extension problem over a general field is satisfactory.

Assume that we are looking for a $\mathbf{d}$ -orthogonal wavelet filter a with multiplicity r such that $\text{fsupp}(a) \subseteq [m, n]$. Regard the entries in the matrices $a(k), k = m, \dots, n$ as

unknowns. Then the orthogonality condition in (3.2) induces a system of quadratic equations with rational coefficients. Finding a Gröbner basis of the system and setting the unconstrained unknowns in the Gröbner basis to be numbers from an algebraic number field (or simply from the rational number field), we see that all the unknowns are roots of polynomials with coefficients from some algebraic number fields. Consequently, by solving a system of quadratic equations with rational coefficients induced by (3.2), all the coefficients of a d -orthogonal wavelet filter are from certain algebraic number field. This more or less explains why all finitely supported d -orthogonal wavelet filters known in the literature have their coefficients from certain algebraic number field. See [1, 2, 3, 4, 5, 6, 8, 12, 13, 18] and Section 6 for many such examples. In particular, [8] provides a systematic construction for d -orthogonal wavelet low-pass filters with linear-phase moments and/or symmetry.

Now we assume that a is a finitely supported d -orthogonal wavelet filter such that all its coefficients are from an algebraic number field $\mathbb{A}$. It is quite appealing and desirable that one can derive a d -orthogonal wavelet filter bank $\{a; b_1, \dots, b_{d-1}\}$ such that all the coefficients of $b_1, \dots, b_{d-1}$ are also from the same algebraic number field $\mathbb{A}$. However, as demonstrated by many examples available in the literature and obtained by brute force calculation using Gröbner bases in symbolic computation, such a requirement is too strong or ideal to be feasible. For the convenience of the reader, we provide a simple example here to illustrate this.

Example 3.2. Let $d = 3$ and a be given by $a(z) = \frac{1}{3}(1 + z + z^2)$. Then all the coefficients of a are in the rational number field $\mathbb{Q}$, and it is easy to show that a is 3-orthogonal. Suppose that b_1 and b_2 are two high-pass filters such that $\{a; b_1, b_2\}$ forms a 3-orthogonal wavelet filter bank. It is very much desired in both theory and application that $\max(\text{len}(b_1), \text{len}(b_2)) \leq \text{len}(a) = 2$ (most d -orthogonal wavelet filter banks in the literature satisfy this property). In our case, we wish to construct high-pass filters b_1 and b_2 with coefficients in $\mathbb{Q}$ and with length no more than 2. In terms of polyphase components, this is equivalent to requiring that the polyphase matrix should take the form

$$\mathcal{P}(z) = \sqrt{d} \begin{bmatrix} a^{[0]}(z) & a^{[1]}(z) & a^{[2]}(z) \\ b_1^{[0]}(z) & b_1^{[1]}(z) & b_1^{[2]}(z) \\ b_2^{[0]}(z) & b_2^{[1]}(z) & b_2^{[2]}(z) \end{bmatrix} = \sqrt{3} \begin{bmatrix} \frac{1}{3} & \frac{1}{3} & \frac{1}{3} \\ t_1 z^m & t_2 z^m & t_3 z^m \\ s_1 z^n & s_2 z^n & s_3 z^n \end{bmatrix},$$

where $t_1, t_2, t_3, s_1, s_2, s_3 \in \mathbb{Q}$ and $m, n \in \mathbb{Z}$. It is straightforward to see that the orthogonality condition $\mathcal{P}(z)\mathcal{P}^*(z) = I_3$ for all $z \in \mathbb{C} \setminus \{0\}$ implies

$$(3.4) \quad t_1 + t_2 + t_3 = 0, \quad t_1^2 + t_2^2 + t_3^2 = \frac{1}{3}.$$

So, we need to find a rational solution $\{t_1, t_2, t_3\}$ to the system of equations in (3.4). Solving (3.4), we have $t_1 = x$, $t_2 = -\frac{x}{2} \pm \frac{\sqrt{6-27x^2}}{6}$, and $t_3 = -\frac{x}{2} \mp \frac{\sqrt{6-27x^2}}{6}$, where x is a free parameter. To have a solution such that $t_1, t_2, t_3 \in \mathbb{Q}$, we must require $x \in \mathbb{Q}$ and $\sqrt{6-27x^2} \in \mathbb{Q}$. Equivalently, we need to find a rational point $(x, y) \in \mathbb{Q}^2$ satisfying the equation $27x^2 + y^2 - 6 = 0$, which defines an algebraic curve. Finding a rational point on an algebraic curve is a major task in computational algebraic geometry. In other words, our question becomes whether or not the algebraic curve $27x^2 + y^2 - 6 = 0$ on $\mathbb{R}^2$ contains a rational point $(x, y) \in \mathbb{Q}^2$. Using the package `algebraiccurves` in `maple`, we find that the algebraic curve $27x^2 + y^2 - 6 = 0$ has genus 0 and is irreducible. Therefore, the algebraic curve has a parametrization. Using `parametrization` in the software `maple` for symbolic computation, from

the parametrization of the algebraic curve $27x^2 + y^2 - 6 = 0$, we conclude that the algebraic curve does not contain any rational point $(x, y) \in \mathbb{Q}^2$. This shows that it is impossible to have a 3-orthogonal wavelet filter bank $\{a; b_1, b_2\}$ such that $\max(\text{len}(b_1), \text{len}(b_2)) \leq \text{len}(a)$ and b_1, b_2 have rational coefficients. In fact, the above argument shows that if b is a 3-orthogonal high-pass filter to a and b has rational coefficients, then we must have $\text{len}(b) \geq 3 > \text{len}(a)$.

On the other hand, if we relax the “rational coefficients” condition to “rational coefficients up to a multiplicative constant”, we have the following solution: $m = n = 0$, $t_1 = \frac{\sqrt{6}}{6}$, $t_2 = 0$, $t_3 = -\frac{\sqrt{6}}{6}$, $s_1 = \frac{\sqrt{2}}{6}$, $s_2 = -\frac{\sqrt{2}}{3}$, $s_3 = \frac{\sqrt{2}}{6}$. The corresponding high-pass filters are given by

$$b_1(z) = \frac{\sqrt{6}}{6}(1 - z^2), \quad b_2(z) = \frac{\sqrt{2}}{6}(1 - 2z + z^2),$$

each of which has coefficients in $\mathbb{Q}$ up to a multiplicative constant and $\text{len}(b_1) = \text{len}(b_2) = \text{len}(a)$.

To appreciate our results and our formulation of the matrix extension problem, in the following let us look at the multilevel wavelet transform. To do so, let us introduce some notation and definitions.

The multilevel wavelet transform is implemented by using the subdivision operator for reconstruction and the transition operator for decomposition. For a filter $a : \mathbb{Z} \rightarrow \mathbb{R}^{r \times r}$ with multiplicity r , the subdivision operator $\mathcal{S}_{a,d} : (l_2(\mathbb{Z}))^{1 \times r} \rightarrow (l_2(\mathbb{Z}))^{1 \times r}$ is defined to be

$$(3.5) \quad [\mathcal{S}_{a,d}v](n) := d \sum_{k \in \mathbb{Z}} v(k)a(n - dk), \quad n \in \mathbb{Z}, \quad v \in (l_2(\mathbb{Z}))^{1 \times r}$$

and the transition operator $\mathcal{T}_{a,d} : (l_2(\mathbb{Z}))^{1 \times r} \rightarrow (l_2(\mathbb{Z}))^{1 \times r}$ is defined to be

$$(3.6) \quad [\mathcal{T}_{a,d}v](n) := d \sum_{k \in \mathbb{Z}} v(k)\overline{a(k - dn)}^T, \quad n \in \mathbb{Z}, \quad v \in (l_2(\mathbb{Z}))^{1 \times r}.$$

For a positive integer J , the J -level discrete wavelet decomposition is given by

$$v_{j-1} = \frac{\sqrt{d}}{d} \mathcal{T}_{a,d}v_j \quad \text{and} \quad w_{j-1;m} = \frac{\sqrt{d}}{d} \mathcal{T}_{b_m,d}v_j, \quad m = 1, \dots, d-1; \quad j = J, \dots, 1,$$

where $v_J : \mathbb{Z} \rightarrow \mathbb{C}^{1 \times r}$ is an input signal in $(l_2(\mathbb{Z}))^{1 \times r}$. Moreover, entries in v_0 are called low-pass wavelet coefficients and entries in $w_{j-1;m}$ are called high-pass wavelet coefficients.

The J -level discrete wavelet reconstruction is given by

$$\hat{v}_j = \frac{\sqrt{d}}{d} \mathcal{S}_{a,d}\hat{v}_{j-1} + \frac{\sqrt{d}}{d} \sum_{m=1}^{d-1} \mathcal{S}_{b_m,d}\hat{w}_{j-1;m}, \quad j = 1, \dots, J.$$

Suppose that $\{a; b_1, \dots, b_{d-1}\}$ is a d -orthogonal wavelet filter bank. Then we can deduce from the identity $\mathcal{P}(z)\mathcal{P}^*(z) = I_r$ for all $z \in \mathbb{C} \setminus \{0\}$ that

$$\mathcal{S}_{a,d}\mathcal{T}_{a,d}v + \sum_{m=1}^{d-1} \mathcal{S}_{b_m,d}\mathcal{T}_{b_m,d}v = dv, \quad v \in (l_2(\mathbb{Z}))^{1 \times r}.$$

Noting that $\|\mathcal{T}_{b_m,d}v\|_{(l_2(\mathbb{Z}))^{1 \times r}}^2 = \langle \mathcal{T}_{b_m,d}v, \mathcal{T}_{b_m,d}v \rangle = \langle \mathcal{S}_{b_m,d}\mathcal{T}_{b_m,d}v, v \rangle$, we now have the following energy preserving equality:

$$\|v_J\|_{(l_2(\mathbb{Z}))^{1 \times r}}^2 = \|v_0\|_{(l_2(\mathbb{Z}))^{1 \times r}}^2 + \sum_{j=1}^J \sum_{m=1}^{d-1} \|w_{j-1;m}\|_{(l_2(\mathbb{Z}))^{1 \times r}}^2.$$

Moreover, if nothing is performed on the wavelet coefficients, i.e., $\check{v}_0 = v_0$ and $\check{w}_{j-1;m} = w_{j-1;m}$ for all $m = 1, \dots, d-1$ and $J = 1, \dots, J$, then we must have the perfect reconstruction property: $\check{v}_J = v_J$, that is, the original input signal v_J can be exactly reconstructed.

We now explain that our formulation of the matrix extension problem over a general field leads to a satisfactory solution to the construction of orthogonal wavelet filter banks.

Suppose that $a : \mathbb{Z} \rightarrow \mathbb{A}^{r \times r}$ is a d -orthogonal wavelet filter with multiplicity r such that all the coefficients of a belong to an algebraic number field $\mathbb{A}$ (for example, $\mathbb{A} = \mathbb{Q}$). By Theorem 3.1, we can construct a d -orthogonal wavelet filter bank $\{a; b_1, \dots, b_{d-1}\}$ such that the polyphase matrix $\mathcal{P}$ in (3.3) is paraunitary and

$$(3.7) \quad b_m = D_m \tilde{b}_m, \quad m = 1, \dots, d-1,$$

where all the filters $\tilde{b}_m : \mathbb{Z} \rightarrow \mathbb{A}^{r \times r}$, $m = 1, \dots, d-1$ with multiplicity r have their coefficients in $\mathbb{A}$, and all D_m are $r \times r$ diagonal matrices with all the entries of $D_m D_m^*$ in $\mathbb{A}$.

Let $\tilde{v}_J := v_J$. We modify the J -level discrete wavelet decomposition as follows:

$$\tilde{v}_{j-1} := \mathcal{T}_{a,d} \tilde{v}_j \quad \text{and} \quad \tilde{w}_{j-1;m} := \mathcal{T}_{\tilde{b}_m,d} \tilde{v}_j$$

for $j = J, \dots, 1$ and $m = 1, \dots, d-1$. Then we have the following relations:

$$v_{j-1} = d^{(j-1-J)/2} \tilde{v}_{j-1}, \quad w_{j-1;m} = d^{(j-1-J)/2} \tilde{w}_{j-1;m} D_m^*$$

for $m = 1, \dots, d-1, j = 1, \dots, J$. That is, the original wavelet coefficients v_{j-1} and $w_{j-1;m}$ are simply scaled versions of the new wavelet coefficients $\tilde{v}_{j-1}$ and $\tilde{w}_{j-1;m}$. Suppose that the input signal v_J also has its entries from the algebraic number field $\mathbb{A}$ (in fact, the coefficients of v_J are often rational numbers). Then all the new wavelet coefficients $\tilde{v}_{j-1}$ and $\tilde{w}_{j-1;m}$ can be efficiently computed using integer arithmetics and matrix/vector operations in linear algebra. Also, by the definition of the transition operator in (3.5) and the relation in (3.7), we see that the J -level discrete wavelet reconstruction can be modified by

$$(3.8) \quad \check{v}_j = d^{-1} \mathcal{S}_{a,d} \check{v}_{j-1} + d^{-1} \sum_{m=1}^{d-1} \mathcal{S}_{\tilde{b}_m,d} (\check{w}_{j-1;m} D_m^* D_m), \quad j = 1, \dots, J.$$

Since all the entries of the diagonal matrices $D_m^* D_m$ are from $\mathbb{A}$, the above modified J -level discrete wavelet reconstruction in (3.8) can be also efficiently implemented using integer arithmetics and matrix/vector operations in linear algebra. Moreover, if nothing is performed on the new wavelet coefficients, i.e., $\check{v}_0 = \tilde{v}_0$ and $\check{w}_{j-1;m} = \tilde{w}_{j-1;m}$ for all $m = 1, \dots, d-1$ and $j = 1, \dots, J$, then we must have the perfect reconstruction property: $\check{v}_J = \tilde{v}_J = v_J$; that is, the original input signal v_J can be exactly reconstructed by the modified discrete wavelet transform.

4. THE MATRIX EXTENSION PROBLEM WITH SYMMETRY CONSTRAINT

In this section, we shall investigate the matrix extension problem with symmetry constraint over a general subfield $\mathbb{F}$ of $\mathbb{C}$ satisfying (1.1). The matrix extension problem with symmetry constraint is much more challenging than its counterpart without the symmetry constraint in Section 2. Extra effort is needed to guarantee the symmetry structure of the extension matrix. Let us first recall some necessary notation and definitions related to symmetry of Laurent polynomials.

Let $\mathbf{p}(z) = \sum_{k \in \mathbb{Z}} p_k z^k$, $z \in \mathbb{C} \setminus \{0\}$ be a Laurent polynomial with complex coefficients $p_k \in \mathbb{C}$. We say that $\mathbf{p}$ has *symmetry* if its coefficient sequence $\{p_k\}_{k \in \mathbb{Z}}$ has symmetry; more precisely, there exist $\varepsilon \in \{-1, 1\}$ and $c \in \mathbb{Z}$ such that

$$(4.1) \quad p_{c-k} = \varepsilon p_k \quad \forall k \in \mathbb{Z}.$$

If $\varepsilon = 1$, then $\mathbf{p}$ is symmetric about the point $c/2$; if $\varepsilon = -1$, then $\mathbf{p}$ is antisymmetric about the point $c/2$. Symmetry of a Laurent polynomial can be conveniently expressed using a symmetry operator $\mathbf{S}$ defined by

$$(4.2) \quad \mathbf{S}\mathbf{p}(z) := \frac{\mathbf{p}(z)}{\mathbf{p}(z^{-1})}, \quad z \in \mathbb{C} \setminus \{0\}.$$

When $\mathbf{p}$ is not identically zero, it is evident that (4.1) holds if and only if $\mathbf{S}\mathbf{p}(z) = \varepsilon z^c$. For the zero polynomial, it is very natural that $\mathbf{S}0$ can be assigned any symmetry pattern; that is, for every occurrence of $\mathbf{S}0$ appearing in an identity in this paper, $\mathbf{S}0$ is understood to take an appropriate choice of εz^c for some $\varepsilon \in \{-1, 1\}$ and $c \in \mathbb{Z}$ so that the identity holds. If $\mathbf{P}$ is an $r \times s$ matrix of Laurent polynomials with symmetry, then we can apply the operator $\mathbf{S}$ to each entry of $\mathbf{P}$; that is, $\mathbf{S}\mathbf{P}$ is an $r \times s$ matrix such that $[\mathbf{S}\mathbf{P}]_{j,k} := \mathbf{S}([\mathbf{P}]_{j,k})$ for $1 \leq j \leq r$ and $1 \leq k \leq s$.

For two matrices $\mathbf{P}$ and $\mathbf{Q}$ of Laurent polynomials with symmetry, even though all the entries in $\mathbf{P}$ and $\mathbf{Q}$ have symmetry, their sum $\mathbf{P} + \mathbf{Q}$, difference $\mathbf{P} - \mathbf{Q}$, or product $\mathbf{P}\mathbf{Q}$, if well defined, generally may not have symmetry anymore. This is one of the difficulties for matrix extension with symmetry. In order for $\mathbf{P} \pm \mathbf{Q}$ or $\mathbf{P}\mathbf{Q}$ to possess some symmetry, the symmetry patterns of $\mathbf{P}$ and $\mathbf{Q}$ should be compatible. For example, if $\mathbf{S}\mathbf{P} = \mathbf{S}\mathbf{Q}$; that is, both $\mathbf{P}$ and $\mathbf{Q}$ have the same symmetry pattern, then indeed $\mathbf{P} \pm \mathbf{Q}$ has symmetry and $\mathbf{S}(\mathbf{P} \pm \mathbf{Q}) = \mathbf{S}\mathbf{P} = \mathbf{S}\mathbf{Q}$. In the following, we discuss the compatibility of symmetry patterns of matrices of Laurent polynomials. For an $r \times s$ matrix $\mathbf{P}$ of Laurent polynomials with symmetry, we say that *the symmetry of $\mathbf{P}$ is compatible* or *$\mathbf{P}$ has compatible symmetry*, if

$$(4.3) \quad \mathbf{S}\mathbf{P}(z) = (\mathbf{S}\theta_1)^*(z)\mathbf{S}\theta_2(z),$$

for some $1 \times r$ vector θ_1 and $1 \times s$ vector θ_2 of Laurent polynomials with symmetry. For an $r \times s$ matrix $\mathbf{P}$ and an $s \times t$ matrix $\mathbf{Q}$ of Laurent polynomials with symmetry, we say that $(\mathbf{P}, \mathbf{Q})$ *has mutually compatible symmetry* if

$$(4.4) \quad \mathbf{S}\mathbf{P}(z) = (\mathbf{S}\theta_1)^*(z)\mathbf{S}\theta(z) \quad \text{and} \quad \mathbf{S}\mathbf{Q}(z) = (\mathbf{S}\theta)^*(z)\mathbf{S}\theta_2(z)$$

for some $1 \times r$, $1 \times s$, $1 \times t$ row vectors $\theta_1, \theta, \theta_2$ of Laurent polynomials with symmetry, respectively. If $(\mathbf{P}, \mathbf{Q})$ has mutually compatible symmetry as in (4.4), then it is easy to verify that their product $\mathbf{P}\mathbf{Q}$ has compatible symmetry and in fact $\mathbf{S}(\mathbf{P}\mathbf{Q}) = (\mathbf{S}\theta_1)^*\mathbf{S}\theta_2$. An $s \times s$ matrix $\mathbf{V}$ of Laurent polynomials is *strongly invertible* if $\mathbf{V}^{-1}$ is also a matrix of Laurent polynomials, or equivalently, the determinant of $\mathbf{V}$ is a nonzero monomial.

Our main result in this section is as follows:

Theorem 4.1. *Let $\mathbf{P}$ be an $r \times s$ matrix, with $1 \leq r \leq s$, of Laurent polynomials taking the form*

$$(4.5) \quad \mathbf{P}(z) = \mathbf{Q}_0(z)\mathbf{D}_0, \quad z \in \mathbb{C} \setminus \{0\},$$

where $\mathbf{Q}_0$ is an $r \times s$ matrix of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and $\mathbf{D}_0$ is an $s \times s$ diagonal matrix with all the entries of $\mathbf{D}_0\mathbf{D}_0^*$ in $\mathbb{F}$. Then:

- (1) the symmetry pattern of P is compatible: $SP = (S\theta_1)^*S\theta_2$ for some $1 \times r$ vector θ_1 and some $1 \times s$ vector θ_2 of Laurent polynomials with symmetry;
- (2) P is paraunitary: $P(z)P^*(z) = I_r$ for all $z \in \mathbb{C} \setminus \{0\}$,

if and only if, there exists an $s \times s$ square extension matrix P_e of Laurent polynomials such that

- (i) $[I_r, \mathbf{0}]P_e = P$; that is, the submatrix of the first r rows of P_e is P ;
- (ii) P_e is paraunitary: $P_e(z)P_e^*(z) = I_s$ for all $z \in \mathbb{C} \setminus \{0\}$;
- (iii) the symmetry pattern of P_e is compatible: $SP_e = (S\theta)^*S\theta_2$ for some $1 \times s$ vector θ of Laurent polynomials with symmetry;
- (iv) P_e can be represented as a product of some $s \times s$ elementary matrices:

$$(4.6) \quad P_e(z) = D_e P_{J+1}(z) P_J(z) \cdots P_1(z) P_0(z) D_0, \quad z \in \mathbb{C} \setminus \{0\},$$

where

- (a) D_e is an $s \times s$ diagonal matrix with all the entries of $D_e D_e^*$ in $\mathbb{F}$;
- (b) P_j , $1 \leq j \leq J$ are elementary: P_j is strongly invertible, has coefficients in $\mathbb{F}$, and $\text{fsupp}(P_j) \subseteq [-1, 1]$;
- (c) (P_{j+1}, P_j) is mutually compatible for all $0 \leq j \leq J$;
- (d) $P_0 = U_{S\theta_2}^*$ and $P_{J+1} = \text{diag}(U_{S\theta_1}, I_{s-r})$, where $U_{S\theta_1}$ and $U_{S\theta_2}$ are products of a permutation matrix with a diagonal matrix of monomials, as defined in (4.7);
- (e) the length of P_e is controlled by that of P in the following sense:

$$\text{len}([P_e]_{j,k}) \leq \max_{1 \leq n \leq r} \text{len}([P]_{n,k}) \quad \text{for } 1 \leq j, k \leq s.$$

We make some remarks about Theorem 4.1. The integer J in item (iv) of Theorem 4.1 always satisfies $J \leq \lceil \text{len}(P)/2 \rceil$. Here $\lceil \cdot \rceil$ denotes the ceiling function such that $\lceil x \rceil = n$ if $n - 1 < x \leq n$ for some integer n . Each of the elementary matrices P_j , $j = 1, \dots, J$ is a product of a constant diagonal matrix in $\mathbb{F}$ and a matrix of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ with its length no greater than 2. The matrices P_0 and P_{J+1} are two matrices of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ reducing the symmetry pattern of P into a standard simplified form. For more detail, see subsections below and Algorithm 2.

Theorem 4.1 states that an $r \times s$ paraunitary matrix P having compatible symmetry can be extended into a square paraunitary matrix P_e having a special *cascade structure* as in (4.6) and having compatible symmetry structure. In the rest of this section, we shall prove Theorem 4.1 and provide an algorithm for it (for its illustration, see Examples 6.1 – 6.4 in Section 6). We follow the main idea in the proof of [10, Theorem 1]. The key difference is that we need to take into account the special form of $P = QD_0$ and construct paraunitary matrices having certain form as well. The proof for the sufficiency part of Theorem 4.1 is straightforward. The proof for the necessity part is constructive with three major steps—initialization, support reduction, and finalization—for deriving a desired matrix P_e in Theorem 4.1 from P . The step of initialization normalizes the symmetry pattern of P to a standard form. The step of support reduction is the main body of the proof, producing a sequence of elementary matrices $A_1, \dots, A_J$ that reduce the length of P to 0. The step of finalization generates the desired matrix P_e as in Theorem 4.1. The following subsections provide details for the three major steps and the corresponding algorithm.

4.1. Initialization. Let θ be a $1 \times s$ row vector of Laurent polynomials with symmetry such that $S\theta = [\varepsilon_1 z^{c_1}, \dots, \varepsilon_s z^{c_s}]$ for some $\varepsilon_1, \dots, \varepsilon_s \in \{-1, 1\}$ and $c_1, \dots, c_s \in \mathbb{Z}$. Then the symmetry of every entry in the vector $\theta \text{diag}(z^{-\lceil c_1/2 \rceil}, \dots, z^{-\lceil c_s/2 \rceil})$ belongs to $\{\pm 1, \pm z^{-1}\}$. Moreover, there is a permutation matrix E_θ to regroup these four types of symmetries together so that

$$(4.7) \quad S(\theta U_{S\theta}) = [\mathbf{1}_{s_1}, -\mathbf{1}_{s_2}, z^{-1}\mathbf{1}_{s_3}, -z^{-1}\mathbf{1}_{s_4}],$$

where $U_{S\theta} := \text{diag}(z^{-\lceil c_1/2 \rceil}, \dots, z^{-\lceil c_s/2 \rceil})E_\theta$ and $s_1, \dots, s_4$ are nonnegative integers uniquely determined by $S\theta$ such that $s_1 + \dots + s_4 = s$.

Since P satisfies $SP = (S\theta_1)^* S\theta_2$, the matrix $\tilde{P} := U_{S\theta_1}^* P U_{S\theta_2}$ must have the symmetry pattern

$$(4.8) \quad \tilde{P} = [\mathbf{1}_{r_1}, -\mathbf{1}_{r_2}, z\mathbf{1}_{r_3}, -z\mathbf{1}_{r_4}]^T [\mathbf{1}_{s_1}, -\mathbf{1}_{s_2}, z^{-1}\mathbf{1}_{s_3}, -z^{-1}\mathbf{1}_{s_4}].$$

Note that $U_{S\theta_1}$ and $U_{S\theta_2}$ do not increase the length of P . Moreover,

$$\tilde{P} = U_{S\theta_1}^* P U_{S\theta_2} = U_{S\theta_1}^* Q_0 D_0 U_{S\theta_2} = U_{S\theta_1}^* Q_0 U_{S\theta_2} (U_{S\theta_2}^* D_0 U_{S\theta_2}) =: \tilde{Q} \tilde{D}_0,$$

where $\tilde{D}_0 := U_{S\theta_2}^* D_0 U_{S\theta_2}$ is a diagonal constant matrix with all the entries of $\tilde{D}_0 \tilde{D}_0^*$ in $\mathbb{F}$.

4.2. Support reduction. $\tilde{P}$ takes the following form:

$$(4.9) \quad \tilde{P} = \begin{bmatrix} F_{11} & -F_{21} & G_{31} & -G_{41} \\ -F_{12} & F_{22} & -G_{32} & G_{42} \\ \mathbf{0} & \mathbf{0} & F_{31} & -F_{41} \\ \mathbf{0} & \mathbf{0} & -F_{32} & F_{42} \end{bmatrix} z^{-k} + \begin{bmatrix} F_{51} & -F_{61} & G_{71} & -G_{81} \\ -F_{52} & F_{61} & -G_{72} & G_{82} \\ G_{11} & -G_{21} & F_{71} & -F_{81} \\ -G_{12} & G_{22} & -F_{72} & F_{82} \end{bmatrix} z^{-k+1} \\ + \sum_{n=2-k}^{k-2} \text{coeff}(\tilde{P}, n) z^n + \begin{bmatrix} F_{51} & F_{61} & G_{31} & G_{41} \\ F_{52} & F_{61} & G_{32} & G_{42} \\ G_{51} & G_{61} & F_{71} & F_{81} \\ G_{52} & G_{62} & F_{72} & F_{82} \end{bmatrix} z^{k-1} + \begin{bmatrix} F_{11} & F_{21} & \mathbf{0} & \mathbf{0} \\ F_{12} & F_{22} & \mathbf{0} & \mathbf{0} \\ G_{11} & G_{21} & F_{31} & F_{41} \\ G_{12} & G_{22} & F_{32} & F_{42} \end{bmatrix} z^k$$

with at least one of $\text{coeff}(\tilde{P}, k)$ and $\text{coeff}(\tilde{P}, -k)$ being nonzero, where $k \geq 1$ is an integer, all $F_{j\ell}$'s and $G_{j\ell}$'s are constant matrices in $\mathbb{F}$, and $F_{11}, F_{22}, F_{31}, F_{42}$ are constant matrices of size $r_1 \times s_1, r_2 \times s_2, r_3 \times s_3, r_4 \times s_4$, respectively.

We next construct an $s \times s$ paraunitary matrix $A_{\tilde{P}}$ having the following properties:

- (1) $A_{\tilde{P}} = \tilde{D}_0^* V_{\tilde{P}} \tilde{D}_1$ for some strongly invertible matrix $V_{\tilde{P}}$ of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some diagonal matrix $\tilde{D}_1$ with all the entries of $\tilde{D}_1 \tilde{D}_1^*$ in $\mathbb{F}$;
- (2) $\text{fsupp}(A_{\tilde{P}}) \subseteq [-1, 1]$ and $\text{len}(\tilde{P} A_{\tilde{P}}) = \text{len}(\tilde{P}) - \text{len}(A_{\tilde{P}})$; that is, $A_{\tilde{P}}$ is elementary and reduces the length of P by that of $A_{\tilde{P}}$;
- (3) $A_{\tilde{P}}$ has compatible symmetry and $\tilde{P} A_{\tilde{P}}$ satisfies

$$S(\tilde{P} A_{\tilde{P}}) = [\mathbf{1}_{r_1}, -\mathbf{1}_{r_2}, z\mathbf{1}_{r_3}, -z\mathbf{1}_{r_4}]^T [\mathbf{1}_{s'_1}, -\mathbf{1}_{s'_2}, z^{-1}\mathbf{1}_{s'_3}, -z^{-1}\mathbf{1}_{s'_4}],$$

for some nonnegative integers $s'_1, \dots, s'_4$ such that $s'_1 + \dots + s'_4 = s$;

- (4) up to a permutation, $[A_{\tilde{P}}]_{j,:} = ([A_{\tilde{P}}]_{:,j})^T = \mathbf{e}_j$, provided $[\tilde{P}]_{:,j} = \mathbf{0}$.

Recall that “up to a permutation” in the above item (4) means that there exists a permutation matrix E such that $[A_{\tilde{P}} E]_{j,:} = [A_{\tilde{P}}]_{:,j} = \mathbf{e}_j$ provided $[\tilde{P}]_{:,j} = \mathbf{0}$. Such a paraunitary matrix $A_{\tilde{P}}$ is of the form $A_{\tilde{P}} = B_{[-k,k]} B_{\tilde{P}_1}$. The construction of $B_{[-k,k]}$ consists of two parts: $\{B_1, \dots, B_r\}$ and $B_{(-k,k)}$. The first part $\{B_1, \dots, B_r\}$

is constructed recursively for each of the r rows of $\tilde{P}$ so that $\tilde{P}_0 := \tilde{P}B_1 \cdots B_r$ reduces $\tilde{P}$ to a special form as follows:

$$(4.10) \quad \left[\begin{array}{cccc} \mathbf{0} & \mathbf{0} & \tilde{G}_{31} & -\tilde{G}_{41} \\ \mathbf{0} & \mathbf{0} & -\tilde{G}_{32} & \tilde{G}_{42} \\ \hline \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \end{array} \right] z^{-k} + \cdots + \left[\begin{array}{cccc} \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \\ \hline \tilde{G}_{11} & \tilde{G}_{21} & \mathbf{0} & \mathbf{0} \\ \tilde{G}_{12} & \tilde{G}_{22} & \mathbf{0} & \mathbf{0} \end{array} \right] z^k.$$

If both $\text{coeff}(\tilde{P}_0, -k) \neq \mathbf{0}$ and $\text{coeff}(\tilde{P}_0, k) \neq \mathbf{0}$, then the second part $B_{(-k,k)}$ is further constructed so that $\tilde{P}_1 := \tilde{P}_0 B_{(-k,k)}$ satisfies either $\text{fsupp}(\tilde{P}_1) \subseteq [-k+1, k]$ or $\text{fsupp}(\tilde{P}_1) \subseteq [-k, k-1]$. $B_{\tilde{P}_1}$ is then constructed so that $\text{fsupp}(\tilde{P}_1 B_{\tilde{P}_1}) \subseteq [-k+1, k-1]$.

4.2.1. *Construction of $B_{[-k,k]}$.* The following lemma is needed for the construction of $B_1, \dots, B_r$, and we postpone its proof to Section 7.

Lemma 4.2. *Let $\mathbf{p}$ be a $1 \times s$ row vector of Laurent polynomials with symmetry satisfying $\mathbf{p}(z)\mathbf{p}^*(z) = 1$ for all $z \in \mathbb{C} \setminus \{0\}$ and $\mathbf{p}(z) = \mathbf{q}_0(z)D_0$, $z \in \mathbb{C} \setminus \{0\}$ for some $1 \times s$ vector $\mathbf{q}_0$ of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some $s \times s$ diagonal matrix D_0 with all the entries of $D_0 D_0^*$ in $\mathbb{F}$. Suppose $\text{fsupp}(\mathbf{p}) = [k_1, k_2]$ with $k_2 - k_1 \geq 2$. Then there exists an $s \times s$ paraunitary matrix $B_{\mathbf{p}}$ satisfying all the following properties:*

- (i) $B_{\mathbf{p}} = D_0^* \mathbf{V} D_1$ for some strongly invertible matrix $\mathbf{V}$ of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some diagonal matrix D_1 with all the entries of $D_1 D_1^*$ in $\mathbb{F}$;
- (ii) $\text{fsupp}(B_{\mathbf{p}}) = [-1, 1]$ and $S B_{\mathbf{p}} = (S \mathbf{p})^* S \mathbf{p}$; that is, $B_{\mathbf{p}}$ has compatible symmetry with its coefficients supported inside $[-1, 1]$;
- (iii) $\text{fsupp}(\mathbf{p} B_{\mathbf{p}}) = [k_1 + 1, k_2 - 1]$ and $S(\mathbf{p} B_{\mathbf{p}}) = S(\mathbf{p})$; that is, $B_{\mathbf{p}}$ reduces the length of $\mathbf{p}$ exactly by 2 and preserves the symmetry pattern of $\mathbf{p}$;
- (iv) for any vector $\tilde{\mathbf{p}}$ of Laurent polynomials such that $\tilde{\mathbf{p}}(z)\mathbf{p}^*(z) = 0$ for all $z \in \mathbb{C} \setminus \{0\}$ and $S\tilde{\mathbf{p}}(z) = \epsilon z^{k_0} S \mathbf{p}(z)$ for some $\epsilon \in \{-1, 1\}$ and some $k_0 \in \mathbb{Z}$, then $S(\tilde{\mathbf{p}} B_{\mathbf{p}}) = S\tilde{\mathbf{p}}$ and $\text{fsupp}(\tilde{\mathbf{p}} B_{\mathbf{p}}) \subseteq \text{fsupp}(\tilde{\mathbf{p}})$; that is, $B_{\mathbf{p}}$ keeps the symmetry pattern of $\tilde{\mathbf{p}}$ and does not increase the length of $\tilde{\mathbf{p}}$;
- (v) up to a permutation, $[B_{\mathbf{p}}]_{j,:} = ([B_{\mathbf{p}}]_{:,j})^T = \mathbf{e}_j$, provided $[\mathbf{p}]_j = 0$.

Suppose that $\tilde{P}$ takes the form (4.9) with $k \geq 1$. If $\text{fsupp}(\tilde{P}) = [-k, k-1]$ or $\text{fsupp}(\tilde{P}) = [-k+1, k]$, we simply let $B_{[-k,k]} := I_s$, $\tilde{P}_1 := \tilde{P}$, and continue to construct $B_{\tilde{P}_1}$ as in Section 4.2.2; otherwise, let us construct $B_1, \dots, B_r$ and $B_{(-k,k)}$ as follows. Let $\mathbf{p}_j := [\tilde{P}]_{j,:}$ be the j th row of $\tilde{P}$ and $B_0 := I_s$. Suppose we have constructed B_{j-1} for $j \geq 1$. We first define $\mathbf{q}_j := \mathbf{p}_j B_0 \cdots B_{j-1}$. Then $B_j := B_{\mathbf{q}_j}$ if $\text{fsupp}(\mathbf{p}_j) = \text{fsupp}(\mathbf{q}_j) = [-k, k]$; otherwise $B_j := I_s$, where $B_{\mathbf{q}_j}$ is constructed as in Lemma 4.2. Let $j \leftarrow j+1$ and repeat this process until $j = r$. Note that each B_j is of the form $B_j = D_{j-1}^* \mathbf{V}_j D_j$ for $j = 1, \dots, r$. Hence, $\tilde{P}_0 := \tilde{P} B_1 \cdots B_r$ preserves the form and symmetry pattern of $\tilde{P}$ as $\tilde{P} = \tilde{Q}_0 \tilde{D}_0$. Moreover, by [10, Lemma 2], the support of $B_1 \cdots B_r$ is contained inside $[-1, 1]$.

Thanks to the properties of B_j as in Lemma 4.2, $\tilde{P}_0$ must take the form as in (4.10). If both $\text{coeff}(\tilde{P}_0, -k) \neq \mathbf{0}$ and $\text{coeff}(\tilde{P}_0, k) \neq \mathbf{0}$, then the second part $B_{(-k,k)}$ is further constructed so that $\tilde{P}_1 := \tilde{P}_0 B_{(-k,k)}$ satisfies either $\text{fsupp}(\tilde{P}_1) \subseteq [-k+1, k]$

or $\text{fsupp}(\tilde{P}_1) \subseteq [-k, k-1]$. The following lemma is needed for us to construct $B_{(-k,k)}$. We postpone its proof to Section 7 as well.

Lemma 4.3. *Let (q_1, q_2) be a pair of $1 \times s$ vectors of Laurent polynomials with symmetry satisfying $q_j(z)q_\ell^*(z) = \delta(j-\ell)$ for all $z \in \mathbb{C} \setminus \{0\}$ and for $j, \ell = 1, 2$. Suppose $q_1 = \tilde{q}_1 D_0$ and $q_2 = \tilde{q}_2 D_0$ for some $1 \times s$ vectors $\tilde{q}_1, \tilde{q}_2$ of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some diagonal matrix D_0 with all the entries of $D_0 D_0^*$ in $\mathbb{F}$. Moreover, $\text{fsupp}(q_1) = [-k, k-1]$ and $\text{fsupp}(q_2) = [-k+1, k]$ with $k \geq 1$, and $Sq_2(z) = \varepsilon z Sq_1(z) = \varepsilon z[\mathbf{1}_{s_1}, -\mathbf{1}_{s_2}, z^{-1}\mathbf{1}_{s_3}, -z^{-1}\mathbf{1}_{s_4}]$ for some $\varepsilon \in \{-1, 1\}$ and some nonnegative integers $s_1, \dots, s_4$ such that $s_1 + \dots + s_4 = s$. Then there exists an $s \times s$ paraunitary matrix $B_{(q_1, q_2)}$ satisfying all the following properties:*

- (i) $B_{(q_1, q_2)} = D_0^* V_{(q_1, q_2)} D_1$ for some strongly invertible matrix $V_{(q_1, q_2)}$ of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some diagonal matrix D_1 with all the entries of $D_1 D_1^*$ in $\mathbb{F}$;
- (ii) $SB_{(q_1, q_2)} = [\mathbf{1}_{s_1}, -\mathbf{1}_{s_2}, z\mathbf{1}_{s_3}, -z\mathbf{1}_{s_4}]^T [\mathbf{1}_{s_1}, -\mathbf{1}_{s_2}, z^{-1}\mathbf{1}_{s_3}, -z^{-1}\mathbf{1}_{s_4}]$ and support $\text{fsupp}(B_{(q_1, q_2)}) = [-1, 1]$; that is, $B_{(q_1, q_2)}$ has compatible symmetry with coefficients supported inside $[-1, 1]$;
- (iii) $\text{fsupp}(q_1 B_{(q_1, q_2)}) \subseteq [-k+1, k-1]$ and $\text{fsupp}(q_2 B_{(q_1, q_2)}) \subseteq [-k+1, k-1]$; that is, $B_{(q_1, q_2)}$ reduces the length of $[q_1^T, q_2^T]^T$ by 2. Moreover, $S(q_1 B_{(q_1, q_2)}) = Sq_1$ and $S(q_2 B_{(q_1, q_2)}) = Sq_2$;
- (iv) if both (p, q_1^*) and (p, q_2^*) have mutually compatible symmetry and $pq_1^* = pq_2^* = 0$, then $S(pB_{(q_1, q_2)}) = Sp$ and $\text{fsupp}(pB_{(q_1, q_2)}) \subseteq \text{fsupp}(p)$; that is, $B_{(q_1, q_2)}$ keeps the symmetry pattern of p and does not increase the length of p ;
- (v) up to a permutation, $[B_{(q_1, q_2)}]_{j,:} = ([B_{(q_1, q_2)}]_{:,j})^T = \mathbf{e}_j$, provided $[q_1]_j = [q_2]_j = 0$.

Now, suppose $\tilde{P}_0$ takes the form in (4.10). Let $B_{(-k,k)} := I_s$ and $\tilde{P}_1 := \tilde{P}_0$. Pick any two rows q_1, q_2 of $\tilde{P}_1$ such that $\text{fsupp}(q_1) = [-k, k-1]$ and $\text{fsupp}(q_2) = [-k+1, k]$. Then, the pair (q_1, q_2) satisfies all the conditions in Lemma 4.3. Hence, we can construct a paraunitary matrix $B_{(q_1, q_2)}$ having the properties as in Lemma 4.3. Replace $B_{(-k,k)}$ and $\tilde{P}_1$ by $B_{(-k,k)} B_{(q_1, q_2)}$ and $\tilde{P}_1 B_{(q_1, q_2)}$, respectively. Pick another two rows (q_1, q_2) such that $\text{fsupp}(q_1) = [-k, k-1]$ and $\text{fsupp}(q_2) = [-k+1, k]$ from the new matrix $\tilde{P}_1$ and repeat the above process. Lemma 4.3 guarantees that this process stops in finite steps and there will be no pair (q_1, q_2) of rows in $\tilde{P}_1$ satisfying $\text{fsupp}(q_1) = [-k, k-1]$ and $\text{fsupp}(q_2) = [-k+1, k]$. Then P_1 must take the form as in (4.10) with at least one of $\text{coeff}(\tilde{P}_1, -k)$ and $\text{coeff}(\tilde{P}_1, k)$ being $\mathbf{0}$.

4.2.2. Construction of $B_{\tilde{P}_1}$. $B_{\tilde{P}_1}$ is constructed to handle the case that $\text{fsupp}(\tilde{P}_1) = [-k, k-1]$ or $\text{fsupp}(\tilde{P}_1) = [-k+1, k]$ so that $\text{fsupp}(\tilde{P}_1 B_{\tilde{P}_1}) \subseteq [-k+1, k-1]$.

If $\tilde{P}_1 := \tilde{P}B_{[-k,k]}$ takes the form in (4.10) with $\text{coeff}(\tilde{P}_1, -k) = \text{coeff}(\tilde{P}_1, k) = \mathbf{0}$, then we simply let $B_{\tilde{P}_1} := I_s$; otherwise, one of $\text{coeff}(\tilde{P}_1, -k)$ and $\text{coeff}(\tilde{P}_1, k)$ is nonzero. For this case, $B_{\tilde{P}_1} := \text{diag}(U_1 W_1, I_{s_3+s_4})E$ with U_1 and W_1 being constructed with respect to $\text{coeff}(\tilde{P}_1, k) \neq \mathbf{0}$ or $B_{\tilde{P}_1} := \text{diag}(I_{s_1+s_2}, U_3 W_3)E$ with U_3 and W_3 being constructed with respect to $\text{coeff}(\tilde{P}_1, -k) \neq \mathbf{0}$, where E is a permutation matrix. Note that $\tilde{P}_1$ is still of the form $\tilde{P}_1 = \tilde{Q}_1 \tilde{D}_1$ similar to $\tilde{P}$. The matrices U_1, W_1, U_3, W_3 , and E are constructed as follows.

Let $U_1 := \text{diag}(U_{\tilde{G}_1}, U_{\tilde{G}_2})$ and $U_3 := \text{diag}(U_{\tilde{G}_3}, U_{\tilde{G}_4})$ with

$$\tilde{G}_1 := \begin{bmatrix} \tilde{G}_{11} \\ \tilde{G}_{12} \end{bmatrix}, \tilde{G}_2 := \begin{bmatrix} \tilde{G}_{21} \\ \tilde{G}_{22} \end{bmatrix}, \tilde{G}_3 := \begin{bmatrix} \tilde{G}_{31} \\ \tilde{G}_{32} \end{bmatrix}, \tilde{G}_4 := \begin{bmatrix} \tilde{G}_{41} \\ \tilde{G}_{42} \end{bmatrix},$$

where $U_{\tilde{G}_1}, \dots, U_{\tilde{G}_4}$ are unitary constant matrices constructed as in Corollary 2.3. If $G_1 G_1^* = G_2 G_2^*$, one can show that U_{G_1} and U_{G_2} can be constructed such that $G_1 U_{G_1} = [R, \mathbf{0}]$ and $G_2 U_{G_2} = [R, \mathbf{0}]$.

Let m_1 and m_3 be the ranks of $\tilde{G}_1$ and $\tilde{G}_3$, respectively ($m_1 = 0$ if $\text{coeff}(\tilde{P}_1, k) = \mathbf{0}$, and $m_3 = 0$ if $\text{coeff}(\tilde{P}_1, -k) = \mathbf{0}$). Note that $\tilde{G}_1 \tilde{G}_1^* = \tilde{G}_2 \tilde{G}_2^*$ and $\tilde{G}_3 \tilde{G}_3^* = \tilde{G}_4 \tilde{G}_4^*$ in view of $\tilde{P}_1 \tilde{P}_1^* = I_r$. The matrices W_1 and W_3 are then constructed to be:

$$W_1 := \left[\begin{array}{c|c|c|c} U_1 & & U_2 & \\ \hline & I_{s_1-m_1} & & \\ \hline U_2 & & U_1 & \\ \hline & & & I_{s_2-m_1} \end{array} \right], \quad W_3 := \left[\begin{array}{c|c|c|c} U_3 & & U_4 & \\ \hline & I_{s_3-m_3} & & \\ \hline U_4 & & U_3 & \\ \hline & & & I_{s_4-m_3} \end{array} \right],$$

where $U_1(z) = -U_2(-z) := \frac{1+z^{-1}}{2} I_{m_1}$ and $U_3(z) = U_4(-z) := \frac{1+z}{2} I_{m_3}$.

Let $W_{\tilde{P}_1} := \text{diag}(U_1 W_1, I_{s_3+s_4})$ for the case that $\text{coeff}(\tilde{P}_1, k) \neq \mathbf{0}$ or $W_{\tilde{P}_1} := \text{diag}(I_{s_1+s_2}, U_3 W_3)$ for the case that $\text{coeff}(\tilde{P}_1, -k) \neq \mathbf{0}$. Then $W_{\tilde{P}_1}$ is paraunitary. By the symmetry pattern and paraunitary property of $\tilde{P}_1$, $W_{\tilde{P}_1}$ reduces the coefficient support of $\tilde{P}_1$ to $[-k+1, k-1]$, i.e., $\text{fsupp}(\tilde{P}_1 W_{\tilde{P}_1}) = [-k+1, k-1]$. Moreover, $W_{\tilde{P}_1}$ changes the symmetry pattern of $\tilde{P}_1$ such that $S(\tilde{P}_1 W_{\tilde{P}_1}) = [\mathbf{1}_{r_1}, -\mathbf{1}_{r_2}, z\mathbf{1}_{r_3}, -z\mathbf{1}_{r_4}]^T S\theta_0$ with

$$S\theta_0 = [z^{-1}\mathbf{1}_{m_1}, \mathbf{1}_{s_1-m_1}, -z^{-1}\mathbf{1}_{m_1}, -\mathbf{1}_{s_2-m_1}, \mathbf{1}_{m_3}, z^{-1}\mathbf{1}_{s_3-m_3}, -\mathbf{1}_{m_3}, -z^{-1}\mathbf{1}_{s_4-m_3}].$$

E is then the permutation matrix such that

$$S(\tilde{P}_1 W_{\tilde{P}_1})E = [\mathbf{1}_{r_1}, -\mathbf{1}_{r_2}, z\mathbf{1}_{r_3}, -z\mathbf{1}_{r_4}]^T S\tilde{\theta}_0,$$

with $S\tilde{\theta}_0 = [\mathbf{1}_{s_1-m_1+m_3}, -\mathbf{1}_{s_2-m_1+m_3}, z^{-1}\mathbf{1}_{s_3-m_3+m_1}, -z^{-1}\mathbf{1}_{s_4-m_3+m_1}] = (S\theta_0)E$.

4.3. Finalization. In summary, $A_{\tilde{P}} = B_{[-k,k]} B_{\tilde{P}_1}$ reduces the support of $\tilde{P}$ and preserves the compatible symmetry of $\tilde{P}$. The properties (1), (3), and (4) of $A_{\tilde{P}}$ follow directly from the above construction while the support property (2) of $A_{\tilde{P}}$ follows from [10, Lemmas 1 and 2]. Replacing $\tilde{P}$ by $\tilde{P} A_{\tilde{P}}$ and repeating the above construction, we can construct paraunitary matrices $A_1, \dots, A_J$ so that $(U_{S_{\theta_1}}^* P U_{S_{\theta_2}}) A_1 \cdots A_J = [I_r, \mathbf{0}]$. Since each A_j is of the form $A_j = \tilde{D}_{j-1}^* V_j \tilde{D}_j$ for some strongly invertible matrix V_j of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some diagonal matrices $\tilde{D}_{j-1}, \tilde{D}_j$ with all the entries of $\tilde{D}_{j-1} \tilde{D}_{j-1}^*, \tilde{D}_j \tilde{D}_j^*$ in $\mathbb{F}$, it is easy to show that the paraunitary extension matrix $P_e := \text{diag}(U_{S_{\theta_1}}, I_{s-r}) A_J^* \cdots A_1^* U_{S_{\theta_2}}^*$ takes the form as in (4.6) with $P_j := \tilde{D}_j \tilde{D}_j^* V_j^*$ for $j = 1, \dots, J-1$, $P_J := V_J^*$, and $D_e := \text{diag}(U_{S_{\theta_1}}, I_{s-r}) \tilde{D}_J^* \text{diag}(U_{S_{\theta_1}}^*, I_{s-r})$. All the properties of P_e follow directly from the properties of each $A_{\tilde{P}}$.

4.4. Algorithm. Based on the above three steps: initialization, support reduction, and finalization, we present an algorithm (see Algorithm 2) for the matrix extension problem with symmetry constraint.

Algorithm 2. Matrix Extension with Symmetry Constraint

- (a) **Input.** $P = Q_0 D_0$ as in Theorem 4.1 with $SP = (S\theta_1)^* S\theta_2$ for some $1 \times r$ vector θ_1 and $1 \times s$ vector θ_2 of Laurent polynomials with symmetry.
- (b) **Output.** A desired extension matrix P_e satisfying all the properties in Theorem 4.1.
- (c) **Initialization.** $\tilde{P} \leftarrow U_{S\theta_1}^* P U_{S\theta_2}$. $D_1 \leftarrow U_{S\theta_2}^* D_0 U_{S\theta_2}$. $P_0 \leftarrow U_{S\theta_2}^*$. $J \leftarrow 1$. Then $\tilde{P}$ has the symmetry pattern as in (4.8), where all nonnegative integers $r_1, \dots, r_4, s_1, \dots, s_4$ are uniquely determined by SP .
- (d) **Support Reduction.**
- 1: **while** ($\text{len}(\tilde{P}) > 0$) **do**
 - 2: $\tilde{P}_0 \leftarrow \tilde{P}$. $\tilde{P}_1 \leftarrow \tilde{P}$. $[k_1, k_2] \leftarrow \text{fsupp}(\tilde{P})$. $A_J \leftarrow I_s$. $k \leftarrow k_2$.
 - 3: **if** $k_2 = -k_1$ **then**
 - 4: **for** $j = 1$ to r **do**
 - 5: $q \leftarrow [\tilde{P}_0]_{j,:}$ and $p \leftarrow [\tilde{P}]_{j,:}$ //the j th rows of $\tilde{P}_0$ and $\tilde{P}$.
 - 6: $[\ell_1, \ell_2] \leftarrow \text{fsupp}(q)$. $\ell \leftarrow \ell_2 - \ell_1$. $B_j \leftarrow I_s$.
 - 7: **if** $\text{fsupp}(q) = \text{fsupp}(p)$ and $\ell \geq 2$ and $(\ell_1 = k_1 \text{ or } \ell_2 = k_2)$ **then**
 - 8: $B_j \leftarrow B_q =: D_1^* V_q D_2$. // B_q is constructed by Lemma 4.2.
 - 9: $A_J \leftarrow A_J B_j$. $\tilde{P}_0 \leftarrow \tilde{P}_0 B_j$. $D_1 \leftarrow D_2$.
 - 10: **end if**
 - 11: **end for**
 - 12: $\tilde{P}_0$ must take the form in (4.10).
 - 13: $B_{(-k,k)} \leftarrow I_s$. $\tilde{P}_1 \leftarrow \tilde{P}_0$. $j_1 \leftarrow 1$. $j_2 \leftarrow r_3 + r_4 + 1$.
 - 14: **while** $j_1 \leq r_1 + r_2$ and $j_2 \leq r$ **do**
 - 15: $q_1 \leftarrow [\tilde{P}_1]_{j_1,:}$. $q_2 \leftarrow [\tilde{P}_1]_{j_2,:}$.
 - 16: **if** $\text{coeff}(q_1, -k) = 0$ **then** $j_1 \leftarrow j_1 + 1$ **end if**
 - 17: **if** $\text{coeff}(q_2, k) = 0$ **then** $j_2 \leftarrow j_2 + 1$ **end if**
 - 18: **if** $\text{coeff}(q_1, -k) \neq 0$ and $\text{coeff}(q_2, k) \neq 0$ **then**
 - 19: $B_{(q_1, q_2)} =: D_1^* V_{(q_1, q_2)} D_2$. // See Lemma 4.3 for the construction.
 - 20: $B_{(-k,k)} \leftarrow B_{(-k,k)} B_{(q_1, q_2)}$. $\tilde{P}_1 \leftarrow \tilde{P}_1 B_{(q_1, q_2)}$. $A_J \leftarrow A_J B_{(q_1, q_2)}$. $D_1 \leftarrow D_2$.
 - 21: $j_1 \leftarrow j_1 + 1$. $j_2 \leftarrow j_2 + 1$.
 - 22: **end if**
 - 23: **end while** // end inner while loop
 - 24: **end if**
 - 25: $\tilde{P}_1$ takes the form in (4.10) with either $\text{coeff}(\tilde{P}_1, -k) = 0$ or $\text{coeff}(\tilde{P}_1, k) = 0$.
 - 26: $B_{\tilde{P}_1} =: D_1^* V_{\tilde{P}_1} D_2$. // See Section 4.2.2 for the construction.
 - 27: $A_J \leftarrow A_J B_{\tilde{P}_1}$. Then A_J is of the form $A_J = \tilde{D}_1^* V_J \tilde{D}_2$ and

$$\tilde{S}P = [1_{r_1}, -1_{r_2}, z1_{r_3}, -z1_{r_4}]^T [1_{s'_1}, -1_{s'_2}, z^{-1}1_{s'_3}, -z^{-1}1_{s'_4}],$$
 for some nonnegative integers $s'_1, \dots, s'_4$.
 - 28: **for** $j = 1$ to 4 **do** $s_j \leftarrow s'_j$ **end for**
 - 29: $\tilde{P} \leftarrow \tilde{P} A_J$. $P_J \leftarrow \tilde{D}_2 \tilde{D}_1^* V_J^*$. $D_1 \leftarrow D_2$. $J \leftarrow J + 1$.
 - 30: **end while** // end outer while loop
- (e) **Finalization.** $\tilde{P} = \text{diag}(F_1, F_2, F_3, F_4) D_1 = \text{diag}(\tilde{F}_1, \tilde{F}_2, \tilde{F}_3, \tilde{F}_4)$ for some $r_j \times s_j$ constant matrices F_j in $\mathbb{F}$, $j = 1, \dots, 4$. Let $U := \text{diag}(U_{\tilde{F}_1}, U_{\tilde{F}_2}, U_{\tilde{F}_3}, U_{\tilde{F}_4}) = D_1^* V_0 D_2$ so that $\tilde{P}U = [I_r, 0]$. Define $P_J := V_0^*$, $P_{J+1} := \text{diag}(U_{S\theta_1}, I_{s-r})$, and $D_e := P_{J+1} D_2^* P_{J+1}^*$.
-

5. APPLICATION TO ORTHOGONAL WAVELET FILTER BANKS WITH SYMMETRY STRUCTURE

In this section, we shall discuss the application of our results on the matrix extension problem with symmetry in Section 4 to the construction of d -orthogonal wavelet filter banks with symmetry structure in electronic engineering and wavelet analysis.

Symmetry of the filters in a filter bank is a very much desired property in many applications. We say that a low-pass filter a with multiplicity r has symmetry if its symbol $\mathbf{a}$ satisfies

$$(5.1) \quad \mathbf{a}(z) = \text{diag}(\varepsilon_1 z^{dc_1}, \dots, \varepsilon_r z^{dc_r}) \mathbf{a}(z^{-1}) \text{diag}(\varepsilon_1 z^{-c_1}, \dots, \varepsilon_r z^{-c_r})$$

for some $\varepsilon_1, \dots, \varepsilon_r \in \{-1, 1\}$ and some $c_1, \dots, c_r \in \mathbb{R}$ such that

$$(5.2) \quad dc_\ell - c_j \in \mathbb{Z} \quad \forall \ell, j = 1, \dots, r.$$

Under the symmetry condition in (5.1), to apply Theorem 4.1, we first show that there exists a suitable paraunitary matrix $\mathbf{U}$ acting on $\mathbf{P}_a := [\mathbf{a}^{[0]}, \dots, \mathbf{a}^{[d-1]}]$ so that $\mathbf{P} := \sqrt{d}\mathbf{P}_a\mathbf{U}$ has compatible symmetry. Note that $\mathbf{P}_a$ itself may not have any symmetry. More importantly, $\mathbf{P}$ has a special structure $\mathbf{P} = \mathbf{Q}_0 D_0$ and satisfies all the conditions in Theorem 4.1.

Lemma 5.1. *Let $a : \mathbb{Z} \rightarrow \mathbb{F}^{r \times r}$ be a d -orthogonal wavelet filter with multiplicity r whose symbol $\mathbf{a}$ satisfies the symmetry property in (5.1). Let $\mathbf{P}_a := [\mathbf{a}^{[0]}, \dots, \mathbf{a}^{[d-1]}]$, where $\mathbf{a}^{[0]}, \dots, \mathbf{a}^{[d-1]}$ are polyphase components of $\mathbf{a}$. Then there exists a $dr \times dr$ paraunitary matrix $\mathbf{U}$ such that $\mathbf{P} := \sqrt{d}\mathbf{P}_a\mathbf{U}$ satisfies the following properties:*

- (i) $\mathbf{P} = \mathbf{Q}_0 D_0$ for some $r \times dr$ matrix $\mathbf{Q}_0$ of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some $dr \times dr$ diagonal matrix D_0 with all the entries of $D_0 D_0^*$ in $\mathbb{F}$;
- (ii) $\mathbf{P}$ has compatible symmetry; that is $\mathbf{S}\mathbf{P} = (\mathbf{S}\theta_1)^* \mathbf{S}\theta_2$ for some $1 \times r$ vector θ_1 and some $1 \times dr$ vector θ_2 of Laurent polynomials with symmetry.

Proof. From (5.1), we deduce that

$$(5.3) \quad [\mathbf{a}^{[\gamma]}(z)]_{\ell,j} = \varepsilon_\ell \varepsilon_j z^{R_{\ell,j}^\gamma} [\mathbf{a}^{[Q_{\ell,j}^\gamma]}(z^{-1})]_{\ell,j}, \quad \gamma = 0, \dots, d-1; \ell, j = 1, \dots, r,$$

where $\gamma, Q_{\ell,j}^\gamma \in \Gamma := \{0, \dots, d-1\}$ and $R_{\ell,j}^\gamma, Q_{\ell,j}^\gamma$ are uniquely determined by

$$(5.4) \quad dc_\ell - c_j - \gamma = dR_{\ell,j}^\gamma + Q_{\ell,j}^\gamma \quad \text{with} \quad R_{\ell,j}^\gamma \in \mathbb{Z}, \quad Q_{\ell,j}^\gamma \in \Gamma.$$

Since $dc_\ell - c_j \in \mathbb{Z}$ for all $\ell, j = 1, \dots, r$, we have $c_\ell - c_j \in \mathbb{Z}$ for all $\ell, j = 1, \dots, r$ and therefore, $Q_{\ell,j}^\gamma$ is independent of ℓ . Consequently, by (5.3), for every $1 \leq j \leq r$, the j th column of the matrix $\mathbf{a}^{[\gamma]}$ is a flipped version of the j th column of the matrix $\mathbf{a}^{[Q_{\ell,j}^\gamma]}$. Let $\kappa_{j,\gamma} \in \mathbb{Z}$ be the integer such that $\text{len}([\mathbf{a}^{[\gamma]}]_{:,j} + z^{\kappa_{j,\gamma}} [\mathbf{a}^{[Q_{\ell,j}^\gamma]}]_{:,j})$ is the smallest. Define $\tilde{\mathbf{a}} := [\tilde{\mathbf{a}}^{[0]}, \dots, \tilde{\mathbf{a}}^{[d-1]}]$ as follows:

$$(5.5) \quad [\tilde{\mathbf{a}}^{[\gamma]}]_{:,j} := \begin{cases} [\mathbf{a}^{[\gamma]}]_{:,j}, & \gamma = Q_{\ell,j}^\gamma; \\ \frac{1}{\sqrt{2}}([\mathbf{a}^{[\gamma]}]_{:,j} + z^{\kappa_{j,\gamma}} [\mathbf{a}^{[Q_{\ell,j}^\gamma]}]_{:,j}), & \gamma < Q_{\ell,j}^\gamma; \\ \frac{1}{\sqrt{2}}([\mathbf{a}^{[\gamma]}]_{:,j} - z^{\kappa_{j,\gamma}} [\mathbf{a}^{[Q_{\ell,j}^\gamma]}]_{:,j}), & \gamma > Q_{\ell,j}^\gamma, \end{cases}$$

where $[\mathbf{a}^{[\gamma]}]_{:,j}$ denotes the j th column of $\mathbf{a}^{[\gamma]}$. Let $\mathbf{U}$ denote the unique transform matrix corresponding to (5.5). It is easily seen that $\mathbf{U}$ is paraunitary and can be rewritten as a product of a strongly invertible matrix of Laurent polynomials with a diagonal matrix whose diagonal entries are either 1 or $\frac{1}{\sqrt{2}}$.

Define $P := \sqrt{d}P_{\tilde{a}} = \sqrt{d}P_a U$. Then item (i) holds. We now show that P has compatible symmetry. Indeed, by (5.3) and (5.5),

$$(5.6) \quad [\tilde{S}\tilde{a}^{[\gamma]}(z)]_{\ell,j} = \operatorname{sgn}(Q_{\ell,j}^\gamma - \gamma) \varepsilon_\ell \varepsilon_j z^{R_{\ell,j}^\gamma + \kappa_{j,\gamma}},$$

where $\operatorname{sgn}(x) = 1$ for $x \geq 0$ and $\operatorname{sgn}(x) = -1$ for $x < 0$. By (5.4) and noting that $Q_{\ell,j}^\gamma$ is independent of ℓ , we have

$$\frac{[\tilde{S}\tilde{a}^{[\gamma]}(z)]_{\ell,j}}{[\tilde{S}\tilde{a}^{[\gamma]}(z)]_{n,j}} = \varepsilon_\ell \varepsilon_n z^{R_{\ell,j}^\gamma - R_{n,j}^\gamma} = \varepsilon_\ell \varepsilon_n z^{c_\ell - c_n}, \quad \ell, n = 1, \dots, r,$$

which is equivalent to saying that P has compatible symmetry. $\square$

In view of Lemma 5.1 and Theorem 4.1, we have the following result.

Theorem 5.2. *Let $a : \mathbb{Z} \rightarrow \mathbb{F}^{r \times r}$ be a d -orthogonal wavelet filter with multiplicity r whose symbol a satisfies the symmetry property in (5.1). Then there exist high-pass filters $b_1, \dots, b_{d-1} : \mathbb{Z} \rightarrow \mathbb{C}^{r \times r}$ such that $\{a; b_1, \dots, b_{d-1}\}$ forms a d -orthogonal wavelet filter bank with $b_m = D_m \tilde{b}_m$, $m = 1, \dots, d-1$, where $\tilde{b}_m : \mathbb{Z} \rightarrow \mathbb{F}^{r \times r}$ has coefficients in $\mathbb{F}$ and D_m is a diagonal matrix with all the entries of $D_m D_m^*$ in $\mathbb{F}$. Moreover, all filters b_m , $m = 1, \dots, d-1$, have the following symmetry property:*

$$(5.7) \quad b_m(z) = \operatorname{diag}(\varepsilon_1^m z^{dc_1^m}, \dots, \varepsilon_r^m z^{dc_r^m}) b_m(z^{-1}) \operatorname{diag}(\varepsilon_1 z^{-c_1}, \dots, \varepsilon_r z^{-c_r}),$$

where $c_\ell^m := (k_\ell^m - k_\ell) + c_\ell \in \mathbb{R}$, $\varepsilon_\ell^m \in \{-1, 1\}$, and $k_\ell^m, k_\ell \in \mathbb{Z}$, for $\ell, j = 1, \dots, r$ and $m = 1, \dots, d-1$.

Proof. Let $P_a := [a^{[0]}, \dots, a^{[d-1]}]$. By Lemma 5.1, there exists a paraunitary matrix U such that $P := \sqrt{d}P_a U$ has compatible symmetry: $SP = [\varepsilon_1 z^{k_1}, \dots, \varepsilon_r z^{k_r}]^T S \theta$ for some $k_1, \dots, k_r \in \mathbb{Z}$ and some $1 \times dr$ row vector θ of Laurent polynomials with symmetry. Moreover, P is of the form $P = Q_0 D_0$, where Q_0 is a matrix of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and D_0 is a diagonal matrix with all the entries of $D_0 D_0^*$ in $\mathbb{F}$. Because a is a d -orthogonal wavelet filter, we have $P_a P_a^* = d^{-1} I_r$ and consequently P is a paraunitary matrix satisfying $PP^* = I_r$. Now by Theorem 4.1, there exists a paraunitary extension matrix P_e of the form $P_e = D_e Q_e D_0$ having the following compatible symmetry pattern:

$$(5.8) \quad [\varepsilon_1 z^{k_1}, \dots, \varepsilon_r z^{k_r}, \varepsilon_1^1 z^{k_1^1}, \dots, \varepsilon_r^1 z^{k_r^1}, \dots, z^{k_1^{d-1}}, \dots, \varepsilon_r^{d-1} z^{k_r^{d-1}}]^T S \theta(z) := SP_e(z).$$

Rewrite $P_e = (\tilde{a}_m^{[\gamma]})_{0 \leq m, \gamma \leq d-1}$ as a $d \times d$ block matrix with $r \times r$ blocks $\tilde{a}_m^{[\gamma]}$. Since P_e has compatible symmetry as in (5.8), we have $[\tilde{S}\tilde{a}_m^{[\gamma]}]_{\ell,j} = \varepsilon_\ell^m \varepsilon_j z^{k_\ell^m - k_j} [\tilde{S}\tilde{a}_m^{[\gamma]}]_{\ell,j}$, for $\ell = 1, \dots, r$ and $m = 1, \dots, d-1$. By (5.6), we have

$$(5.9) \quad [\tilde{S}\tilde{a}_m^{[\gamma]}(z)]_{\ell,j} = \operatorname{sgn}(Q_{\ell,j}^\gamma - \gamma) \varepsilon_\ell^m \varepsilon_j z^{R_{\ell,j}^\gamma + k_{j,\gamma} + k_\ell^m - k_j}, \quad \ell, j = 1, \dots, r.$$

Let $\mathcal{P} := P_e U^* := \sqrt{d}(\mathbf{b}_m^{[\gamma]})_{0 \leq m, \gamma \leq d-1}$ with $\mathbf{b}_0 := a$. By (5.9) and the definition of U^* in (5.5), we deduce that

$$(5.10) \quad [\mathbf{b}_m^{[\gamma]}(z)]_{\ell,j} = \varepsilon_\ell^m \varepsilon_j z^{R_{\ell,j}^\gamma + k_\ell^m - k_j} \left[\mathbf{b}_m^{[Q_{\ell,j}^\gamma]}(z^{-1}) \right]_{\ell,j}.$$

This implies that $[\mathbf{S}\mathbf{b}_m(z)]_{\ell,j} = \varepsilon_\ell^m \varepsilon_j z^{d(k_\ell^m - k_j + c_\ell) - c_j}$, which is equivalent to (5.7) with $c_\ell^m := k_\ell^m - k_\ell + c_\ell$ for $m = 1, \dots, d-1$ and $\ell = 1, \dots, r$.

Noting that $D_e = \operatorname{diag}(I_r, D_1, \dots, D_{d-1})$ for some diagonal matrices D_m with all the entries of $D_m D_m^*$ in $\mathbb{F}$, we complete the proof. $\square$

To end this section, for a $\mathbf{d}$ -orthogonal wavelet low-pass filter $\mathbf{a}$ satisfying (5.1), we provide an algorithm (see Algorithm 3) to construct the associated high-pass filters $\mathbf{b}_1, \dots, \mathbf{b}_{\mathbf{d}-1}$ such that they form a symmetric $\mathbf{d}$ -orthogonal wavelet filter bank.

Algorithm 3. Construction of Wavelet High-Pass Filters with Symmetry Structure

- (a) **Input.** A $\mathbf{d}$ -orthogonal wavelet filter a with symmetry as in (5.1) and coefficients in $\mathbb{F}$.
- (b) **Output.** A symmetric $\mathbf{d}$ -orthogonal wavelet filter bank $\{a, b_1, \dots, b_{\mathbf{d}-1}\}$ such that $\mathcal{P}$ in (3.3) is paraunitary and all filters b_m , $m = 1, \dots, \mathbf{d}-1$, have symmetry pattern satisfying (5.7). Moreover, $b_m = D_m \tilde{b}_m$, $m = 1, \dots, \mathbf{d}-1$, where $\tilde{b}_m : \mathbb{Z} \rightarrow \mathbb{F}^{r \times r}$ has coefficients from $\mathbb{F}$ and D_m is a diagonal matrix with all the entries of $D_m D_m^*$ in $\mathbb{F}$.
- (c) **Initialization.** Let $\mathbf{P}_a := [a^{[0]}, \dots, a^{[\mathbf{d}-1]}]$ and construct $\mathbf{U}$ with respect to (5.5) such that $\mathbf{P} := \sqrt{\mathbf{d}} \mathbf{P}_a \mathbf{U} = \mathbf{Q}_0 D_0$ for some matrix $\mathbf{Q}_0$ of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some diagonal matrix D_0 with all the entries of $D_0 D_0^*$ in $\mathbb{F}$. Moreover, $\mathbf{P}$ has compatible symmetry: $\mathbf{SP}(z) = [\varepsilon_1 z^{k_1}, \dots, \varepsilon_r z^{k_r}]^T \mathbf{S} \theta(z)$ for some $k_1, \dots, k_r \in \mathbb{Z}$ and some $1 \times \mathbf{d}r$ row vector θ of Laurent polynomials with symmetry.
- (d) **Matrix Extension.** Derive $\mathbf{P}_e$ having all the properties as in Theorem 4.1 from $\mathbf{P}$ by Algorithm 2.
- (e) **High-Pass Filters.** Let $\mathcal{P} := \mathbf{P}_e \mathbf{U}^* =: \sqrt{\mathbf{d}} (\mathbf{b}_m^{[\gamma]})_{0 \leq m, \gamma \leq \mathbf{d}-1}$ as in (3.3) with $\mathbf{b}_0 := \mathbf{a}$. Define high-pass filters $b_1, \dots, b_{\mathbf{d}-1}$ through their symbols by

$$(5.11) \quad \mathbf{b}_m(z) := \sum_{\gamma=0}^{\mathbf{d}-1} \mathbf{b}_m^{[\gamma]}(z^{\mathbf{d}}) z^{\gamma}, \quad m = 1, \dots, \mathbf{d}-1.$$

6. ILLUSTRATIVE EXAMPLES

In this section we provide several examples to illustrate our algorithms and results stated in previous sections.

Example 6.1. Let $\mathbf{d} = 3$ and $r = 1$. A 3-orthogonal low-pass wavelet filter a is given by its symbol $\mathbf{a}$ as follows (also see [2]):

$$\mathbf{a}(z) = \left(\frac{z^{-1}+1+z}{3}\right)^2 \left(-\frac{1}{9}z^{-2} - \frac{2}{9}z^{-1} + \frac{5}{3} - \frac{2}{9}z - \frac{1}{9}z^2\right).$$

All the coefficients of $\mathbf{a}$ are rational numbers and the filter a is symmetric about 0: $\mathbf{S}\mathbf{a} = 1$.

Applying Lemma 5.1, we obtain

$$\mathbf{P}_a(z) := [\mathbf{a}^{[0]}(z), \mathbf{a}^{[1]}(z), \mathbf{a}^{[2]}(z)] = \frac{1}{81} [-4z^{-1} + 35 - 4z, 8z^{-1} + 20 - z, -z^{-2} + 20z^{-1} + 8]$$

and $\mathbf{U}$ that symmetrizes $\mathbf{P}_a$ is given by

$$\mathbf{U}(z) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ 0 & \frac{1}{\sqrt{2}}z & -\frac{1}{\sqrt{2}}z \end{bmatrix}.$$

Then $\mathbf{P} := \sqrt{3} \mathbf{P}_a \mathbf{U} = \mathbf{Q}_0 D_0$ satisfies $\mathbf{SP} = [1, 1, -1]$ with $\mathbf{Q}_0$ and D_0 being given by

$$\mathbf{Q}_0(z) = \frac{1}{162} [-8z^{-1} + 70 - 8z, 7z^{-1} + 40 + 7z, 9z^{-1} - 9z], \quad D_0 = \text{diag}(\sqrt{3}, \sqrt{6}, \sqrt{6}).$$

Applying Algorithm 2 to P , we obtain a desired paraunitary matrix $P_e = D_e Q_e D_0$ with $D_e = \text{diag}(\frac{1}{162}, \frac{\sqrt{2}}{324}, \frac{\sqrt{6}}{108})$ and Q_e as follows:

$$Q_e(z) = \begin{bmatrix} -8z^{-1} + 70 - 8z & 7z^{-1} + 40 + 7z & -9z^{-1} + 9z \\ 40z^{-1} + 28 + 40z & -35z^{-1} + 16 - 35z & -45z^{-1} + 45z \\ -8z^{-1} + 8z & 7z^{-1} - 7z & 9z^{-1} + 9z \end{bmatrix}.$$

We have $SP_e = [1, 1, -1]^T [1, 1, -1]$ and $\text{fsupp}([P_e]_{:,j}) \subseteq \text{fsupp}([P]_j)$ for all $1 \leq j \leq 3$. Now, from the polyphase matrix $\mathcal{P} := P_e U^* =: \sqrt{3}(\mathbf{b}_m^{[\gamma]})_{0 \leq m, \gamma \leq 2}$ with $\mathbf{b}_0 := \mathbf{a}$, we derive two high-pass filters $\mathbf{b}_1, \mathbf{b}_2$ as follows:

$$\begin{aligned} \mathbf{b}_1(z) &= \frac{\sqrt{2}}{162}(5z^{-4} + 20z^{-3} - 40z^{-2} + 8z^{-1} + 14 + 8z - 40z^2 + 20z^3 + 5z^4), \\ \mathbf{b}_2(z) &= \frac{\sqrt{6}}{54}(z^{-4} + 4z^{-3} - 8z^{-2} + 8z^2 - 4z^3 - z^4). \end{aligned}$$

The high-pass filter $\mathbf{b}_1$ is symmetric about 0: $S\mathbf{b}_1 = 1$, while the high-pass filter $\mathbf{b}_2$ is antisymmetric about 0: $S\mathbf{b}_2 = -1$. Note that $\text{len}(\mathbf{b}_1) = \text{len}(\mathbf{b}_2) = \text{len}(\mathbf{a}) = 8$.

Example 6.2. Let $d = 5$ and $r = 1$. A 5-orthogonal wavelet low-pass filter a is given by its symbol $\mathbf{a}$ as follows:

$$\mathbf{a}(z) = (\frac{z^{-2} + z^{-1} + 1 + z + z^2}{5})^2 (-\frac{2}{3}z^{-2} + \frac{2}{3}z^{-1} + 1 + \frac{2}{3}z - \frac{2}{3}z^2).$$

All the coefficients of $\mathbf{a}$ are rational numbers and the filter a is symmetric about 0: $S\mathbf{a} = 1$.

Applying Lemma 5.1, we obtain

$$\begin{aligned} P_{\mathbf{a}}(z) &:= [\mathbf{a}^{[0]}(z), \dots, \mathbf{a}^{[4]}(z)] \\ &= \frac{1}{75}[-2z^{-1} + 19 - 2z, z^{-1} + 16 - 2z, 6z^{-1} + 9, 9z^{-1} + 6, -2z^{-2} + 16z^{-1} + 1] \end{aligned}$$

and U that symmetrizes $P_{\mathbf{a}}$ is given by

$$U(z) = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & \frac{1}{\sqrt{2}} & 0 & 0 & \frac{1}{\sqrt{2}} \\ 0 & 0 & \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} & 0 \\ 0 & 0 & \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} & 0 \\ 0 & \frac{1}{\sqrt{2}}z & 0 & 0 & -\frac{1}{\sqrt{2}}z \end{bmatrix}.$$

Then $P := \sqrt{5}P_{\mathbf{a}}U = Q_0D_0$ satisfies $SP = [1, 1, z^{-1}, -z^{-1}, -1]$ with

$$D_0 = \text{diag}(\sqrt{5}, \sqrt{10}, \sqrt{10}, \sqrt{10}, \sqrt{10})$$

and Q_0 being given by

$$Q_0(z) = \frac{1}{150}[-4z^{-1} + 38 - 4z, -z^{-1} + 32 - z, 15z^{-1} + 15, -3z^{-1} + 3, 3z^{-1} - 3z].$$

Applying Algorithm 2 to P , we obtain a desired paraunitary matrix $P_e = D_e Q_e D_0$ with $D_e = \text{diag}(\frac{1}{150}, \frac{1}{300}, \frac{\sqrt{5}}{300}, \frac{\sqrt{5}}{60}, \frac{\sqrt{5}}{300})$ and Q_e as follows:

$$Q_e(z) = \begin{bmatrix} 38 - 4(z^{-1} + z) & 32 - (z^{-1} + z) & 15(z^{-1} + 1) & -3(z^{-1} + 1) & 3z^{-1} - 3z \\ 4(z^{-1} + z) - 88 & z^{-1} + 68 + z & -15z^{-1} - 15 & 3z^{-1} - 3 & -3z^{-1} + 3z \\ -4z^{-1} + 4z & -z^{-1} + z & 15z^{-1} - 15 & -3z^{-1} - 3 & 3(z^{-1} + 1) + 36 \\ -4 - 4z & -1 - z & 6 & 0 & 3 - 3z \\ -4 + 4z & -1 + z & 0 & 42 & 3 + 3z \end{bmatrix}.$$

We have $\mathbf{SP}_e = [1, 1, -1, z, -z]^\top [1, 1, z^{-1}, -z^{-1}, -1]$ and $\text{fsupp}([\mathbf{P}_e]_{:,j}) \subseteq \text{fsupp}([\mathbf{P}]_j)$ for all $1 \leq j \leq 5$. Now, from the polyphase matrix $\mathcal{P} := \mathbf{P}_e \mathbf{U}^* =: \sqrt{5}(\mathbf{b}_m^{[\gamma]})_{0 \leq m, \gamma \leq 4}$ with $\mathbf{b}_0 := \mathbf{a}$, we derive four high-pass filters $\mathbf{b}_1, \dots, \mathbf{b}_4$ as follows:

$$\begin{aligned} \mathbf{b}_1(z) &= \frac{1}{150}(t_1(z^{-1}) + t_1(z)); & \mathbf{b}_2(z) &= \frac{\sqrt{5}}{150}(-t_2(z^{-1}) + t_2(z)); \\ \mathbf{b}_3(z) &= \frac{\sqrt{5}}{30}(-2z^{-1} - 2 + z + 3z^2 + 3z^3 + z^4 - 2z^5 - 2z^6); \\ \mathbf{b}_4(z) &= \frac{\sqrt{5}}{150}(-2z^{-1} - 2 + z + 21z^2 - 21z^3 - z^4 + 2z^5 + 2z^6) \end{aligned}$$

with

$$t_1(z) = -22 + 34z - 9z^2 - 6z^3 - z^4 + 2z^5 + 2z^6; \quad t_2(z) = 18z - 9z^2 - 6z^3 - z^4 + 2z^5 + 2z^6.$$

The high-pass filter $\mathbf{b}_1$ is symmetric about 0: $\mathbf{Sb}_1 = 1$, the high-pass filter $\mathbf{b}_2$ is antisymmetric about 0: $\mathbf{Sb}_2 = -1$, the high-pass filter $\mathbf{b}_3$ is symmetric about $5/2$: $\mathbf{Sb}_3 = z^5$, and the high-pass filter $\mathbf{b}_4$ is antisymmetric about $5/2$: $\mathbf{Sb}_4 = -z^5$. Note that $\text{len}(\mathbf{b}_1) = \text{len}(\mathbf{b}_2) = \text{len}(\mathbf{a}) = 12$ and $\text{len}(\mathbf{b}_3) = \text{len}(\mathbf{b}_4) = 7 < \text{len}(\mathbf{a})$.

Example 6.3. Let $d = 3$ and $r = 1$. A 3-orthogonal low-pass wavelet filter a is given by its symbol $\mathbf{a}$ as follows:

$$\mathbf{a}(z) = \frac{1}{3}\left(\frac{z^{-1}+1+z}{3}\right)^3(- (3+2i\sqrt{3})z^{-1} + 9 + 4i\sqrt{3} - (3+2i\sqrt{3})z).$$

All the coefficients of $\mathbf{a}$ are from the algebraic number field $\mathbb{Q}(i\sqrt{3})$ with $i = \sqrt{-1}$ being the imaginary unit, and the filter a is symmetric about 0: $\mathbf{Sa} = 1$.

Applying Lemma 5.1, we obtain

$$\begin{aligned} \mathbf{P}_a(z) &:= [\mathbf{a}^{[0]}(z), \mathbf{a}^{[1]}(z), \mathbf{a}^{[2]}(z)] \\ &= \frac{1}{81}[-2i\sqrt{3}z^{-1} + (27 + 4i\sqrt{3}) - 2i\sqrt{3}z, (6 - 2i\sqrt{3})z^{-1} + (24 + 4i\sqrt{3}) \\ &\quad - (3 + 2i\sqrt{3})z, -(3 + 2i\sqrt{3})z^{-2} + (24 + 4i\sqrt{3})z^{-1} + (6 - 2i\sqrt{3})] \end{aligned}$$

and $\mathbf{U}$ that symmetrizes $\mathbf{P}_a$ is given by

$$\mathbf{U}(z) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ 0 & \frac{1}{\sqrt{2}}z & -\frac{1}{\sqrt{2}}z \end{bmatrix}.$$

Then $\mathbf{P} := \sqrt{3}\mathbf{P}_a\mathbf{U} = \mathbf{Q}_0\mathbf{D}_0$ satisfies $\mathbf{SP} = [1, 1, -1]$ with $\mathbf{D}_0 = \text{diag}(\sqrt{3}, \sqrt{6}, \sqrt{6})$ and $\mathbf{Q}_0$ being given by

$$\mathbf{Q}_0 = \frac{1}{162}[-4i\sqrt{3}(z^{-1}+z) + (54+8\sqrt{3}i), (3-4i\sqrt{3})(z^{-1}+z) + (48+8i\sqrt{3}), 9(z^{-1}-z)].$$

Applying Algorithm 2 to $\mathbf{P}$, we obtain a desired paraunitary matrix $\mathbf{P}_e = \mathbf{D}_e\mathbf{Q}_e\mathbf{D}_0$ with $\mathbf{D}_e = \text{diag}(\frac{1}{162}, \frac{\sqrt{78}}{4212}, \frac{\sqrt{78}}{1404})$ and $\mathbf{Q}_e$ as follows:

$$\mathbf{Q}_e(z) = \begin{bmatrix} -4i\sqrt{3}(z^{-1}+z) + t_{11} & (3-4i\sqrt{3})(z^{-1}+z) + t_{12} & 9(z^{-1}-z) \\ -4(z^{-1}+z) + t_{21}\sqrt{3} & -(4+i\sqrt{3})(z^{-1}+z) + t_{22} & 3i-\sqrt{3}(z^{-1}-z) \\ -4i\sqrt{3}(z^{-1}-z) & (3-4i\sqrt{3})(z^{-1}-z) & 9(z^{-1}+z) - 36i\sqrt{3} \end{bmatrix},$$

with

$$t_{11} = 54 + 8i\sqrt{3}; \quad t_{12} = 48 + 8i\sqrt{3}; \quad t_{21} = -208 + 36i\sqrt{3}; \quad t_{22} = 116 - 16i\sqrt{3}.$$

We have $\mathbf{SP}_e = [1, 1, -1]^\top [1, 1, -1]$ and $\text{fsupp}([\mathbf{P}_e]_{:,j}) \subseteq \text{fsupp}([\mathbf{P}]_j)$ for all $1 \leq j \leq 3$. Now, from the polyphase matrix $\mathcal{P} := \mathbf{P}_e \mathbf{U}^* =: \sqrt{3}(\mathbf{b}_m^{[\gamma]})_{0 \leq m, \gamma \leq 2}$ with $\mathbf{b}_0 := \mathbf{a}$, we derive two high-pass filters $\mathbf{b}_1, \mathbf{b}_2$ as follows:

$$\begin{aligned} \mathbf{b}_1(z) &= \frac{\sqrt{78}}{2106}(-2 - i\sqrt{3})(z^{-4} + z^4) - 2(z^{-3} + z^3) - (2 + 2i\sqrt{3})(z^{-2} + z^2) \\ &\quad + (58 - 8i\sqrt{3})(z^{-1} + z) - 104 + 18i\sqrt{3}, \\ \mathbf{b}_2(z) &= \frac{\sqrt{78}}{702}(-3 + 2i\sqrt{3})(z^{-4} - z^4) - 2i\sqrt{3}(z^{-3} - z^3) + (6 - 2i\sqrt{3})(z^{-2} - z^2) \\ &\quad + 18i\sqrt{3}(z^{-1} - z) \end{aligned}$$

The high-pass filter $\mathbf{b}_1$ is symmetric about 0: $\mathbf{Sb}_1 = 1$, while the high-pass filter $\mathbf{b}_2$ is antisymmetric about 0: $\mathbf{Sb}_2 = -1$. Note that $\text{len}(\mathbf{b}_1) = \text{len}(\mathbf{b}_2) = \text{len}(\mathbf{a}) = 8$.

Example 6.4. Let $\mathbf{d} = 2$ and $r = 2$. A 2-orthogonal wavelet low-pass filter $\mathbf{a}$ with multiplicity 2 in [4] is given by

$$\mathbf{a}(z) = \frac{1}{40} \begin{bmatrix} 12(z^{-1} + 1) & 16\sqrt{2}z^{-1} \\ -\sqrt{2}(z^{-1} - 9 - 9z + z^2) & -2(3z^{-1} - 10 + 3z) \end{bmatrix}.$$

The low-pass filter $\mathbf{a}$ satisfies the symmetry property in (5.1) with $c_1 = -1, c_2 = 0$ and $\varepsilon_1 = \varepsilon_2 = 1$. Note that the coefficients of $\mathbf{a}$ are from the algebraic number field $\mathbb{F} = \mathbb{Q}(\sqrt{2})$.

Applying Lemma 5.1, we obtain $\mathbf{P}_a := [\mathbf{a}^{[0]}, \mathbf{a}^{[1]}]$ and $\mathbf{U}$ as follows:

$$\mathbf{P}_a(z) = \frac{1}{20} \begin{bmatrix} 6 & 0 & 6z^{-1} & 8\sqrt{2}z^{-1} \\ \frac{1}{\sqrt{2}}(9 - z) & 10 & \frac{1}{\sqrt{2}}(9 - z^{-1}) & -3(1 + z^{-1}) \end{bmatrix}$$

and

$$\mathbf{U}(z) = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & \sqrt{2} & 0 & 0 \\ z & 0 & -z & 0 \\ 0 & 0 & 0 & \sqrt{2}z \end{bmatrix}.$$

Then $\mathbf{P} := \sqrt{2}\mathbf{P}_a\mathbf{U}$ satisfies $\mathbf{SP} = [1, z]^\top [1, z^{-1}, -1, 1] = \mathbf{Q}_0 D_0$ with

$$\mathbf{Q}_0(z) = \frac{\sqrt{2}}{20} \begin{bmatrix} 6\sqrt{2} & 0 & 0 & 8\sqrt{2} \\ 4(1 + z) & 10 & 5(1 - z) & -3(1 + z) \end{bmatrix}, \quad D_0 = I_4.$$

We would like to point out that though the filter coefficients of $\mathbf{Q}_0$ are in $\mathbb{Q}(\sqrt{2})$, they are essentially in $\mathbb{Q}$ because $\mathbf{Q}_0$ can be written as

$$\mathbf{Q}_0 = \frac{1}{20} \begin{bmatrix} 1 & 0 \\ 0 & \sqrt{2} \end{bmatrix} \tilde{\mathbf{Q}}_0 \text{ with } \tilde{\mathbf{Q}}_0(z) = \begin{bmatrix} 12 & 0 & 0 & 16 \\ 4(1 + z) & 10 & 5(1 - z) & -3(1 + z) \end{bmatrix},$$

and our algorithms are essentially applied to the part $\tilde{\mathbf{Q}}_0$.

Applying Algorithm 2 to $\mathbf{P}$, we obtain a desired paraunitary matrix $\mathbf{P}_e = D_e \mathbf{Q}_e$ with D_e and $\mathbf{Q}_e$ being given as follows:

$$D_e = \text{diag}(2, \sqrt{2}, \sqrt{2}, 2), \quad \mathbf{Q}_e(z) = \frac{1}{20} \begin{bmatrix} 6 & 0 & 0 & 8 \\ 4(1 + z) & 10 & 5(1 - z) & -3(1 + z) \\ 4(1 + z) & -10 & 5(1 - z) & -3(1 + z) \\ 4(1 - z) & 0 & 5(1 + z) & -3(1 - z) \end{bmatrix}.$$

We have $\mathbf{SP}_e = [1, z, z, -z]^\top [1, z^{-1}, -1, 1]$ and $\text{fsupp}([\mathbf{P}_e]_{:,j}) \subseteq \text{fsupp}([\mathbf{P}]_{:,j})$ for all $1 \leq j \leq 4$.

Now, from the polyphase matrix $\mathcal{P} := \mathbf{P}_e \mathbf{U}^* =: \sqrt{2}(\mathbf{b}_m^{[\gamma]})_{0 \leq m, \gamma \leq 1}$ with $\mathbf{b}_0 := \mathbf{a}$, we derive a high-pass filter $\mathbf{b}_1$ as follows:

$$\mathbf{b}_1(z) = \frac{1}{40} \begin{bmatrix} -\sqrt{2}(z^{-1} - 9 - 9z + z^2) & -2(3z^{-1} + 10 + 3z) \\ 2(-z^{-1} + 9 - 9z + z^2) & -6\sqrt{2}(z^{-1} - z) \end{bmatrix}.$$

Then the high-pass filter $\mathbf{b}_1$ satisfies (5.7) with $c_1^1 = c_2^1 = 0$ and $\varepsilon_1^1 = 1, \varepsilon_2^1 = -1$.

Our algorithms are applicable to many other examples with $d \geq 2$ and $r \in \mathbb{N}$; for example, see [10, Examples 2 and 3] and many examples in [5, 8].

7. PROOFS OF LEMMAS 4.2 AND 4.3

We now prove Lemmas 4.2 and 4.3 stated in Section 4. To do so, we need an auxiliary result. The following lemma shows that for a 1×4 vector $\mathbf{p}$ of Laurent polynomials with symmetry such that $\mathbf{p}(z)\mathbf{p}^*(z) = 1$ for all $z \in \mathbb{C} \setminus \{0\}$, $\mathbf{S}\mathbf{p} = [1, -1, z^{-1}, -z^{-1}]$, and $\mathbf{p}(z) = \mathbf{q}_0(z)D_0$, we can construct a 4×4 paraunitary matrix with a special structure.

Lemma 7.1. *Let $\mathbf{p}$ be a 1×4 row vector of Laurent polynomials with symmetry satisfying $\mathbf{p}(z)\mathbf{p}^*(z) = 1$ for all $z \in \mathbb{C} \setminus \{0\}$, $\mathbf{S}\mathbf{p}(z) = [1, -1, z^{-1}, -z^{-1}]$, and $\mathbf{p}(z) = \mathbf{q}_0(z)D_0$ for some 1×4 vector $\mathbf{q}_0$ of Laurent polynomials in $\mathbb{F}[z, z^{-1}]$ and some diagonal matrix $D_0 = \text{diag}(d_1, d_1, d_3, d_4)$, where d_1, d_3 , and d_4 are positive numbers with d_1^2, d_3^2 , and d_4^2 being numbers in $\mathbb{F}$. Suppose $\text{fsupp}(\mathbf{p}) = [-k, k]$ with a positive integer $k \geq 1$, $\text{coeff}(\mathbf{q}, -k) = [f_1, -f_1, g_1, -g_2]$, and $\text{coeff}(\mathbf{q}, -k+1) = [f_3, -f_4, g_3, -g_4]$ for some numbers $f_1, f_3, f_4, g_1, g_2, g_3, g_4$ in $\mathbb{F}$ with $f_1 > 0$ and $g_1, g_2 \geq 0$, i.e., $\mathbf{q}_0^T$ takes the form:*

$$\mathbf{q}_0(z)^T = \begin{bmatrix} f_1 \\ -f_1 \\ g_1 \\ -g_2 \end{bmatrix} z^{-k} + \begin{bmatrix} f_3 \\ -f_4 \\ g_3 \\ -g_4 \end{bmatrix} z^{-k+1} + \cdots + \begin{bmatrix} f_3 \\ f_4 \\ g_1 \\ g_2 \end{bmatrix} z^{k-1} + \begin{bmatrix} f_1 \\ f_1 \\ 0 \\ 0 \end{bmatrix} z^k.$$

Define a matrix $\mathbf{U}_0(z) := D_0^* \mathbf{V}_0(z) D_1$ with $D_1 = \frac{1}{c} \text{diag}(1, 1, \frac{d_3}{d_1}, \frac{d_4}{d_1})$ and $\mathbf{V}_0$ being given by

(7.1)

$$\mathbf{V}_0(z) := \begin{bmatrix} f_1(z + z^{-1}) + \bar{f} & f_1(z - z^{-1}) & g_1(1 + z^{-1}) & g_2(1 - z^{-1}) \\ -f_1(z - z^{-1}) & -f_1(z + z^{-1}) + \bar{f} & -g_1(1 - z^{-1}) & -g_2(1 + z^{-1}) \\ g_1(1 + z) & -g_1(1 - z) & -\frac{d_1^2}{d_3^2}(2f_1 + f) & 0 \\ g_2(1 - z) & -g_2(1 + z) & 0 & \frac{d_1^2}{d_4^2}(2f_1 - f) \end{bmatrix},$$

where $f := f_3 - f_4$ and $c := \sqrt{4d_1^2 f_1^2 + 2d_3^2 g_1^2 + 2d_4^2 g_2^2 + d_1^2 |f|^2}$. Then $\mathbf{U}_0$ is paraunitary and has the following properties:

- (i) $\text{fsupp}(\mathbf{U}_0) = [-1, 1]$ and $\mathbf{S}\mathbf{U}_0(z) = [1, -1, z, -z]^T [1, -1, z^{-1}, -z^{-1}]$; that is, $\mathbf{U}_0$ has compatible symmetry with filter support on $[-1, 1]$;
- (ii) $\text{fsupp}(\mathbf{p}\mathbf{U}_0) = [-k+1, k-1]$ and $\mathbf{S}(\mathbf{q}\mathbf{U}_0) = \mathbf{S}(\mathbf{p})$; that is, $\mathbf{U}_0$ reduces the length of $\mathbf{p}$ exactly by 2 and preserves the symmetry pattern of $\mathbf{p}$;
- (iii) for any vector $\tilde{\mathbf{p}}$ of Laurent polynomials such that $\tilde{\mathbf{p}}\mathbf{p}^* = 0$ and $\mathbf{S}\tilde{\mathbf{p}}(z) = \epsilon z^{k_0} \mathbf{S}\mathbf{p}(z)$ for some $\epsilon \in \{-1, 1\}$ and some $k_0 \in \mathbb{Z}$, we have $\mathbf{S}(\tilde{\mathbf{p}}\mathbf{U}_0) = \mathbf{S}\tilde{\mathbf{p}}$ and $\text{fsupp}(\tilde{\mathbf{p}}\mathbf{U}_0) \subseteq \text{fsupp}(\tilde{\mathbf{p}})$; that is, $\mathbf{U}_0$ keeps the symmetry pattern of $\tilde{\mathbf{p}}$ and does not increase the length of $\tilde{\mathbf{p}}$.

Proof. Note that $\mathbf{p}$ is of the form:

$$\mathbf{p}(z)^\top = D_0^\top \left(\begin{bmatrix} f_1 \\ -f_1 \\ g_1 \\ -g_2 \end{bmatrix} z^{-k} + \begin{bmatrix} f_3 \\ -f_4 \\ g_3 \\ -g_4 \end{bmatrix} z^{-k+1} + \cdots + \begin{bmatrix} f_3 \\ f_4 \\ g_1 \\ g_2 \end{bmatrix} z^{k-1} + \begin{bmatrix} f_1 \\ f_1 \\ 0 \\ 0 \end{bmatrix} z^k \right).$$

For such a vector $\mathbf{p}$, by [6], a paraunitary matrix, that reduces its length by 2 and preserves its symmetry pattern, is given by

$$(7.2) \quad U_0(z) = \frac{1}{c} \begin{bmatrix} d_1 f_1(z + z^{-1}) + d_1 \bar{f} & d_1 f_1(z - z^{-1}) & d_3 g_1(1 + z^{-1}) & d_4 g_2(1 - z^{-1}) \\ -d_1 f_1(z - z^{-1}) & -d_1 f_1(z + z^{-1}) + d_1 \bar{f} & -d_3 g_1(1 - z^{-1}) & -d_4 g_2(1 + z^{-1}) \\ d_3 g_1(1 + z) & -d_3 g_1(1 - z) & -2d_1 f_1 - d_1 f & 0 \\ d_4 g_2(1 - z) & -d_4 g_2(1 + z) & 0 & 2d_1 f_1 - d_1 f \end{bmatrix}$$

with $c = \sqrt{4d_1^2 f_1^2 + 2d_3^2 g_1^2 + 2d_4^2 g_2^2 + d_1^2 |f|^2}$ and $f = f_3 - f_4$. It is easy to check that U_0 can be rewritten as in the lemma. The properties of U_0 in items (i) – (iii) can be verified directly by computations. $\square$

We remark that the above result holds even when $[\mathbf{p}]_3$ or $[\mathbf{p}]_4$ is empty, or both $[\mathbf{p}]_3$ and $[\mathbf{p}]_4$ are empty. In such a case, we obtain either a 2×2 or a 3×3 paraunitary matrix by simply deleting the corresponding row(s) and column(s) in (7.2).

Thanks to Lemmas 2.2 and 7.1, we now prove Lemma 4.2 as follows.

Proof of Lemma 4.2. The paraunitary matrix $B_{\mathbf{p}}$ is constructed according to the following steps: first reduce the problem size to a 1×4 vector of Laurent polynomials, then construct a 4×4 paraunitary matrix according to Lemma 7.1, and finally extend it in a simple way to an $s \times s$ paraunitary matrix.

(1) Let $\mathbf{p}_0 := \mathbf{p}U_{\mathbf{Sp}}$, where $U_{\mathbf{Sp}}$ is a matrix defined in (4.7). Then

$$\mathbf{Sp}_0 = [\mathbf{1}_{s_1}, -\mathbf{1}_{s_2}, z^{-1}\mathbf{1}_{s_3}, -z^{-1}\mathbf{1}_{s_4}]$$

with $s_1, \dots, s_4$ being nonnegative integers determined by $\mathbf{Sp}$. $\mathbf{p}_0$ must take one of the following forms:

$$(7.3) \quad \mathbf{p}_0 = \left([\mathbf{f}_1, -\mathbf{f}_2, \mathbf{g}_1, -\mathbf{g}_2]z^{\ell_1} + [\mathbf{f}_3, -\mathbf{f}_4, \mathbf{g}_3, -\mathbf{g}_4]z^{\ell_1+1} + \sum_{\ell=\ell_1+2}^{\ell_2-2} \text{coeff}(\mathbf{p}_0, \ell)z^\ell + [\mathbf{f}_3, \mathbf{f}_4, \mathbf{g}_1, \mathbf{g}_2]z^{\ell_2-1} + [\mathbf{f}_1, \mathbf{f}_2, \mathbf{0}, \mathbf{0}]z^{\ell_2} \right) \tilde{D}_0;$$

$$(7.4) \quad \mathbf{p}_0 = \left([\mathbf{0}, \mathbf{0}, \mathbf{f}_1, -\mathbf{f}_2]z^{\ell_1} + [\mathbf{g}_1, -\mathbf{g}_2, \mathbf{f}_3, -\mathbf{f}_4]z^{\ell_1+1} + \sum_{\ell=\ell_1+2}^{\ell_2-2} \text{coeff}(\mathbf{p}_0, \ell)z^\ell + [\mathbf{g}_3, \mathbf{g}_4, \mathbf{f}_3, \mathbf{f}_4]z^{\ell_2-1} + [\mathbf{g}_1, \mathbf{g}_2, \mathbf{f}_1, \mathbf{f}_2]z^{\ell_2} \right) \tilde{D}_0.$$

If $\mathbf{p}_0$ is of the form (7.3), we simply let $D_{\mathbf{p}} := U_{\mathbf{Sp}}$. If $\mathbf{p}_0$ takes the form in (7.4), we further construct a permutation matrix E such that $[\mathbf{g}_1, \mathbf{g}_2, \mathbf{f}_1, \mathbf{f}_2]E = [\mathbf{f}_1, \mathbf{f}_2, \mathbf{g}_1, \mathbf{g}_2]$ and define $D_{\mathbf{p}}(z) := U_{\mathbf{Sp}} E \text{diag}(I_{s-s_g}, z^{-1}I_{s_g})$, where $1 \times s_g$ is the size of the row vector $[\mathbf{g}_1, \mathbf{g}_2]$. In this way, $\mathbf{p}_0 := \mathbf{p}D_{\mathbf{p}}$ always takes the form in (7.3) with $\mathbf{f}_1 \neq \mathbf{0}$. Moreover, the diagonal matrix $\tilde{D}_0$ is given by $\tilde{D}_0 = D_{\mathbf{p}}^* D_0 D_{\mathbf{p}} =: \text{diag}(\tilde{D}_{0,1}, \tilde{D}_{0,2}, \tilde{D}_{0,3}, \tilde{D}_{0,4})$, where $\tilde{D}_{0,1}, \tilde{D}_{0,2}, \tilde{D}_{0,3}, \tilde{D}_{0,4}$ are of the same sizes with respect to the vectors $\mathbf{f}_1, \mathbf{f}_2, \mathbf{g}_1, \mathbf{g}_2$. Define $\tilde{\mathbf{f}}_1 := \mathbf{f}_1 \tilde{D}_{0,1}$, $\tilde{\mathbf{f}}_2 := \mathbf{f}_2 \tilde{D}_{0,2}$, $\tilde{\mathbf{g}}_1 := \mathbf{g}_1 \tilde{D}_{0,3}$, and $\tilde{\mathbf{g}}_2 := \mathbf{g}_2 \tilde{D}_{0,4}$.

- (2) Applying Lemma 2.2 to the vectors $\tilde{\mathbf{f}}_1, \tilde{\mathbf{f}}_2, \tilde{\mathbf{g}}_1, \tilde{\mathbf{g}}_2$, we can construct unitary constant matrices $U_{\tilde{\mathbf{f}}_1} = \tilde{D}_{0,1}^* V_{\tilde{\mathbf{f}}_1} \tilde{D}_{1,1}$, $U_{\tilde{\mathbf{f}}_2} = \tilde{D}_{0,2}^* V_{\tilde{\mathbf{f}}_2} \tilde{D}_{1,2}$, $U_{\tilde{\mathbf{g}}_1} = \tilde{D}_{0,3}^* V_{\tilde{\mathbf{g}}_1} \tilde{D}_{1,3}$, and $U_{\tilde{\mathbf{g}}_2} = \tilde{D}_{0,4}^* V_{\tilde{\mathbf{g}}_2} \tilde{D}_{1,4}$, where $\tilde{V}_{\tilde{\mathbf{f}}_1} = [\mathbf{f}_1^*, F_1^*]$, $\tilde{V}_{\tilde{\mathbf{f}}_2} = [\mathbf{f}_2^*, F_2^*]$, $\tilde{V}_{\tilde{\mathbf{g}}_1} = [\mathbf{g}_1^*, G_1^*]$, and $\tilde{V}_{\tilde{\mathbf{g}}_2} = [\mathbf{g}_2^*, G_2^*]$ are nonsingular matrices with all the entries in $\mathbb{F}$. Moreover, $U := \text{diag}(U_{\tilde{\mathbf{f}}_1}, U_{\tilde{\mathbf{f}}_2}, U_{\tilde{\mathbf{g}}_1}, U_{\tilde{\mathbf{g}}_2}) = \tilde{D}_0^* V \tilde{D}_1$ with $\tilde{D}_1 := \text{diag}(\tilde{D}_{1,1}, \tilde{D}_{1,2}, \tilde{D}_{1,3}, \tilde{D}_{1,4})$ normalizes the vector $[\mathbf{f}_1, -\mathbf{f}_2, \mathbf{g}_1, -\mathbf{g}_2] \tilde{D}_0$ to the following form:

$$[\mathbf{f}_1, -\mathbf{f}_2, \mathbf{g}_1, -\mathbf{g}_2] \tilde{D}_0 U \\ = [d_1 f_1, 0, \dots, 0, -d_1 f_1, 0, \dots, 0, d_3 g_1, 0, \dots, 0, -d_4 g_2, 0, \dots, 0],$$

where we can choose d_1, d_3, d_4 and f_1, g_1, g_2 to be $d_1 = \|\tilde{\mathbf{f}}_1\|$, $d_3 = \|\tilde{\mathbf{g}}_1\|$, $d_4 = \|\tilde{\mathbf{g}}_2\|$ and $f_1 = g_1 = g_2 = 1$ (other choices are possible as long as d_1, d_3, d_4 and f_1, g_1, g_2 satisfy conditions in Lemma 7.1). In other words, under the action of U , $\tilde{\mathbf{f}}_1, \tilde{\mathbf{f}}_2, \tilde{\mathbf{g}}_1$, and $\tilde{\mathbf{g}}_2$ become $[d_1 f_1, \mathbf{0}]$, $[d_1 f_1, \mathbf{0}]$, $[d_3 g_1, \mathbf{0}]$, and $[d_4 g_2, \mathbf{0}]$, respectively. Note that $\tilde{\mathbf{f}}_1 U_{\tilde{\mathbf{f}}_1} = \tilde{\mathbf{f}}_2 U_{\tilde{\mathbf{f}}_2} = [d_1 f_1, 0, \dots, 0]$ follows from the fact that $\|\tilde{\mathbf{f}}_1\| = \|\tilde{\mathbf{f}}_2\|$, and $d_3 g_1$ or $d_4 g_2$ could be an empty entry.

- (3) Applying Lemma 7.1 to the 1×4 subvector of $\mathbf{p}_0 U = \mathbf{p} \mathbf{D}_p U$ consisting its 1st, $(1+s_1)$ th, $(1-\delta(s_3))(1+s_1+s_2)$ th, and $(1-\delta(s_4))(1+s_1+s_2+s_3)$ th entries, we obtain a paraunitary matrix $\mathbf{U}_0$ as described in Lemma 7.1. Note that $\mathbf{U}_0$ could be a 2×2 or 3×3 matrix depending on whether or not $d_3 g_1$ or $d_4 g_2$ is an empty entry. Let $\mathcal{I} := \{1, 1+s_1, (1-\delta(s_3))(1+s_1+s_2), (1-\delta(s_4))(1+s_1+s_2+s_3)\} \setminus \{0\}$ be an index set and $n_0 := \#(\mathcal{I})$ be its cardinality.
- (4) Extend $\mathbf{U}_0$ to an $s \times s$ paraunitary matrix $\mathbf{U}$ as follows:

$$[\mathbf{U}]_{\mathcal{I}_j, \mathcal{I}_k} := [\mathbf{U}_0]_{j,k}, \quad j, k = 1, \dots, n_0; \quad [\mathbf{U}]_{j,j} := 1, \quad j \notin \mathcal{I},$$

and all other entries of $\mathbf{U}$ are zero. Then $\mathbf{U}$ is of the form $\mathbf{U} = \tilde{D}_2^* \mathbf{V}_2 \tilde{D}_3$ for some strongly invertible matrix $\mathbf{V}_2$ extended from (7.1), where

$$\tilde{D}_2 = \text{diag}(d_1, \mathbf{1}_{s_1-1}, d_1, \mathbf{1}_{s_2-1}, d_3, \mathbf{1}_{s_3-1}, d_4, \mathbf{1}_{s_4-1})$$

and

$$\tilde{D}_3 = \text{diag}(\frac{1}{c}, \mathbf{1}_{s_1-1}, \frac{1}{c}, \mathbf{1}_{s_2-1}, \frac{d_3}{d_1 c}, \mathbf{1}_{s_3-1}, \frac{d_4}{d_1 c}, \mathbf{1}_{s_4-1})$$

with c defined in (7.1).

- (5) Define $\mathbf{B}_{\mathbf{p}_0}(z) := \mathbf{U} \mathbf{U}(z) = (\tilde{D}_0^* V \tilde{D}_1)(\tilde{D}_2^* \mathbf{V}_1 \tilde{D}_3)$. Then $\mathbf{B}_{\mathbf{p}_0}$ satisfies item (i) – (iv) of Lemma 4.2 for $\mathbf{p}_0$, and by our construction, $\mathbf{B}_{\mathbf{p}_0}$ is the form $\mathbf{B}_{\mathbf{p}_0} = \tilde{D}_0^* \mathbf{V}_0 \tilde{D}_4$:

$$\mathbf{V}_0^*(z) := \begin{bmatrix} \begin{array}{c} \mathbf{f}_1(z + c_0 + \frac{1}{z}) \\ F_1 \end{array} & \begin{array}{c} \mathbf{f}_2(z - \frac{1}{z}) \\ \mathbf{0} \end{array} & \begin{array}{c} \mathbf{g}_1(1 + \frac{1}{z}) \\ \mathbf{0} \end{array} & \begin{array}{c} \mathbf{g}_2(1 - \frac{1}{z}) \\ \mathbf{0} \end{array} \\ \begin{array}{c} -\mathbf{f}_1(z - \frac{1}{z}) \\ \mathbf{0} \end{array} & \begin{array}{c} -\mathbf{f}_2(z - c_0 + \frac{1}{z}) \\ F_2 \end{array} & \begin{array}{c} -\mathbf{g}_1(1 - \frac{1}{z}) \\ \mathbf{0} \end{array} & \begin{array}{c} -\mathbf{g}_2(1 + \frac{1}{z}) \\ \mathbf{0} \end{array} \\ \begin{array}{c} \mathbf{f}_1(1 + z) \\ \mathbf{0} \end{array} & \begin{array}{c} -\mathbf{f}_2(1 - z) \\ \mathbf{0} \end{array} & \begin{array}{c} \mathbf{g}_1 \\ G_1 \end{array} & \begin{array}{c} \mathbf{0} \\ \mathbf{0} \end{array} \\ \begin{array}{c} \mathbf{f}_1(1 - z) \\ \mathbf{0} \end{array} & \begin{array}{c} \mathbf{f}_2(1 + z) \\ \mathbf{0} \end{array} & \begin{array}{c} \mathbf{0} \\ \mathbf{0} \end{array} & \begin{array}{c} \mathbf{g}_2 \\ G_2 \end{array} \end{bmatrix}$$

where $c_0 = \frac{1}{\|\tilde{\mathbf{f}}_1\|^2} \text{coeff}(\mathbf{p}_0, \ell_1 + 1) \text{coeff}(\mathbf{p}_0^*, -\ell_2) \in \mathbb{F}$ and $\tilde{D}_4$ is a diagonal matrix determined by $\tilde{D}_1$, $\tilde{D}_2$ and $\tilde{D}_3$.

- (6) Define $\mathbf{B}_p(z) := \mathbf{D}_p(z) \mathbf{B}_{p_0}(z) \mathbf{D}_p^*(z)$. Then $\mathbf{B}_p$ is a desired paraunitary matrix.

Properties (i) – (v) of $\mathbf{B}_p$ follow from our construction. $\square$

We complete the paper by proving Lemma 4.3.

Proof of Lemma 4.3. The proof is similar to that of Lemma 4.2. First, one can show that there exists a permutation matrix $E_{(\mathbf{q}_1, \mathbf{q}_2)}$ such that $\mathbf{q}_1 E_{(\mathbf{q}_1, \mathbf{q}_2)}$ and $\mathbf{q}_2 E_{(\mathbf{q}_1, \mathbf{q}_2)}$ take the following form:

$$\begin{bmatrix} \mathbf{p}_1 \\ \mathbf{p}_2 \end{bmatrix} := \begin{bmatrix} \mathbf{q}_1 \\ \mathbf{q}_2 \end{bmatrix} E_{(\mathbf{q}_1, \mathbf{q}_2)} = \begin{bmatrix} \mathbf{0} & \mathbf{0} & \tilde{\mathbf{g}}_3 & -\tilde{\mathbf{g}}_4 \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \end{bmatrix} z^{-k} + \begin{bmatrix} \tilde{\mathbf{f}}_5 & -\tilde{\mathbf{f}}_6 & \tilde{\mathbf{g}}_7 & -\tilde{\mathbf{g}}_8 \\ \tilde{\mathbf{g}}_1 & -\tilde{\mathbf{g}}_2 & \tilde{\mathbf{f}}_7 & -\tilde{\mathbf{f}}_8 \end{bmatrix} z^{-k+1} \\ + \cdots + \begin{bmatrix} \tilde{\mathbf{f}}_5 & \tilde{\mathbf{f}}_6 & \tilde{\mathbf{g}}_3 & \tilde{\mathbf{g}}_4 \\ \tilde{\mathbf{g}}_5 & \tilde{\mathbf{g}}_6 & \tilde{\mathbf{f}}_7 & \tilde{\mathbf{f}}_8 \end{bmatrix} z^{k-1} + \begin{bmatrix} \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \\ \tilde{\mathbf{g}}_1 & \tilde{\mathbf{g}}_2 & \mathbf{0} & \mathbf{0} \end{bmatrix} z^k,$$

where $\tilde{\mathbf{g}}_1, \tilde{\mathbf{g}}_2, \tilde{\mathbf{g}}_3, \tilde{\mathbf{g}}_4$ are all nonzero row vectors of size $1 \times s_1, 1 \times s_2, 1 \times s_3, 1 \times s_4$, respectively. And $[\tilde{\mathbf{g}}_1, \tilde{\mathbf{g}}_2, \tilde{\mathbf{g}}_3, \tilde{\mathbf{g}}_4] = [\mathbf{g}_1, \mathbf{g}_2, \mathbf{g}_3, \mathbf{g}_4] \tilde{D}_0$ for some vectors $\mathbf{g}_1, \dots, \mathbf{g}_4$ with all their entries in $\mathbb{F}$ and having same size as $\tilde{\mathbf{g}}_1, \dots, \tilde{\mathbf{g}}_4$, respectively, and for some diagonal matrix $\tilde{D}_0$ with all the entries of $\tilde{D}_0 \tilde{D}_0^*$ in $\mathbb{F}$. Note that $\|\tilde{\mathbf{g}}_1\| = \|\tilde{\mathbf{g}}_2\|$ and $\|\tilde{\mathbf{g}}_3\| = \|\tilde{\mathbf{g}}_4\|$.

Using Lemma 2.2, we can construct unitary matrices $U_{\tilde{\mathbf{g}}_1}, U_{\tilde{\mathbf{g}}_2}, U_{\tilde{\mathbf{g}}_3}, U_{\tilde{\mathbf{g}}_4}$ of the form $U_{\tilde{\mathbf{g}}_1} = \tilde{D}_{0,1}^* [\mathbf{g}_1^*, G_1^*] \tilde{D}_{1,1}^*$, $U_{\tilde{\mathbf{g}}_2} = \tilde{D}_{0,2}^* [\mathbf{g}_2^*, G_2^*] \tilde{D}_{1,2}^*$, $U_{\tilde{\mathbf{g}}_3} = \tilde{D}_{0,3}^* [\mathbf{g}_3^*, G_3^*] \tilde{D}_{1,3}^*$, and $U_{\tilde{\mathbf{g}}_4} = \tilde{D}_{0,4}^* [\mathbf{g}_4^*, G_4^*] \tilde{D}_{1,4}^*$, where $\tilde{D}_0 =: \text{diag}(\tilde{D}_{0,1}, \tilde{D}_{0,2}, \tilde{D}_{0,3}, \tilde{D}_{0,4})$, $G_1, \dots, G_4$ are nonsingular matrices with all the entries in $\mathbb{F}$, and $\tilde{D}_1 =: \text{diag}(\tilde{D}_{1,1}, \tilde{D}_{1,2}, \tilde{D}_{1,3}, \tilde{D}_{1,4})$ is some diagonal matrix with all the entries of $\tilde{D}_1 \tilde{D}_1^*$ in $\mathbb{F}$.

Define $U := \text{diag}(U_{\tilde{\mathbf{g}}_1}, U_{\tilde{\mathbf{g}}_2}, U_{\tilde{\mathbf{g}}_3}, U_{\tilde{\mathbf{g}}_4}) =: \tilde{D}_0 V \tilde{D}_1$. Applying U to the pair $(\mathbf{p}_1, \mathbf{p}_2)$, we normalize the pair to be supported on $[-k+1, k-1]$ except those entries on 1st, $(1+s_1)$ th, $(1+s_1+s_2)$ th, and $(1+s_1+s_2+s_3)$ th positions. Hence, we only need to consider a problem of constructing a 4×4 paraunitary matrix $\mathbf{U}_0$, which can be constructed by Lemma 7.1 similarly as in the proof Lemma 4.2, corresponding to the 1×4 vector of Laurent polynomials with symmetry. Finally, we can extend $\mathbf{U}_0$ into a full $s \times s$ paraunitary matrix $\mathbf{U}$ of the form $\mathbf{U} = \tilde{D}_2 \mathbf{V}_1 \tilde{D}_3$. Consequently, the paraunitary matrix for the pair $(\mathbf{p}_1, \mathbf{p}_2)$ is given by $\mathbf{B}_{(\mathbf{p}_1, \mathbf{p}_2)} = \mathbf{U} \mathbf{U}^* = (\tilde{D}_0^* V \tilde{D}_1^*) (\tilde{D}_2^* \mathbf{V}_1 \tilde{D}_3^*)$, and by our construction, $\mathbf{B}_{(\mathbf{p}_1, \mathbf{p}_2)}$ can be written as $\mathbf{B}_{(\mathbf{p}_1, \mathbf{p}_2)} = \tilde{D}_0 \mathbf{V}_0 \tilde{D}_4$ with $\mathbf{V}_0$ being given as follows:

$$\mathbf{V}_0^*(z) := \begin{bmatrix} \begin{array}{c|c|c|c} \mathbf{g}_1 & \mathbf{0} & \mathbf{g}_3(1 + \frac{1}{z}) & \mathbf{g}_4(1 - \frac{1}{z}) \\ \hline G_1 & \mathbf{0} & \mathbf{0} & \mathbf{0} \\ \hline \mathbf{0} & \mathbf{g}_2 & -\mathbf{g}_3(1 - \frac{1}{z}) & -\mathbf{g}_4(1 + \frac{1}{z}) \\ \hline \mathbf{0} & G_2 & \mathbf{0} & \mathbf{0} \\ \hline \mathbf{g}_1(1+z) & -\mathbf{g}_2(1-z) & -\mathbf{g}_3 & \mathbf{0} \\ \hline \mathbf{0} & \mathbf{0} & G_3 & \mathbf{0} \\ \hline \mathbf{g}_1(1-z) & -\mathbf{g}_2(1+z) & \mathbf{0} & -\mathbf{g}_4 \\ \hline \mathbf{0} & \mathbf{0} & \mathbf{0} & G_4 \end{array} \end{bmatrix},$$

and $\tilde{D}_4$ being determined by $\tilde{D}_1$, $\tilde{D}_2$, and $\tilde{D}_3$. Let $B_{(q_1, q_2)} := E_{(q_1, q_2)} B_{(p_1, p_2)} E_{(q_1, q_2)}^T$. Then $B_{(q_1, q_2)}$ is a desired matrix and properties (i) – (v) follow from our construction. $\square$

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