

THE LOCATION OF THE ZEROS OF THE HIGHER ORDER DERIVATIVES OF A POLYNOMIAL

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ABSTRACT. Let $p(z)$ be a complex polynomial of degree n having k zeros in a disk D . We deal with the problem of finding the smallest concentric disk containing $k - l$ zeros of $p^{(l)}(z)$. We obtain some estimates on the radius of this disk in general as well as in the special case, where k zeros in D are isolated from the other zeros of $p(z)$. We indicate an application to the root-finding algorithms.

1. INTRODUCTION

Let us consider the following problem: If k zeros of a polynomial $p(z)$ of degree n ($2 \leq k \leq n$) lie in a disk D of radius r , what is the smallest concentric disk that contains $k - l$ zeros of $p^{(l)}(z)$ ($1 \leq l \leq k - 1$)? Since the problem is scaling and translation invariant, we can assume that the disk D is the closed unit disk $\overline{\Delta} := \{z \in \mathbb{C} : |z| \leq 1\}$. Let $\mathcal{P}_{n,k}$ denote the class of complex polynomials of degree n having exactly k zeros in $\overline{\Delta}$. We define the function $\rho(n, k, l)$, $n \geq k > l$, as follows:

$$(1.1) \quad \rho(n, k, l) = \sup_{p \in \mathcal{P}_{n,k}} \min \left\{ R > 0 : \overline{D}(0, R) \text{ contains at least } k - l \text{ zeros of } p^{(l)}(z) \right\}.$$

Because of scaling and translation invariance we can conclude that if $\overline{D}(c, r)$ contains k zeros of the polynomial $p(z)$, then $\overline{D}(c, r\rho(n, k, l))$ contains $k - l$ zeros of $p^{(l)}(z)$.

The problem of estimating $\rho(n, k, l)$ has a long history in the case $l = 1$. The results listed below can be found in Marden's book ([2]). The Gauss-Lucas Theorem states that $\rho(n, n, 1) = 1$. Result $\rho(n, 2, 1) = \cot(\pi/n)$ is due to Alexander, Kakeya and Szegő. Biernacki proved that

$$(1.2) \quad \rho(n, n - 1, 1) \leq (1 + 1/n)^{1/2} \quad \text{and} \quad \rho(n, k, 1) \leq \prod_{i=1}^{n-k} [(n+i)/(n-i)],$$

and Marden showed that

$$(1.3) \quad \rho(n, k, 1) \leq \csc \frac{\pi}{2(n-k+1)}.$$

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More recently, Coppersmith and Neff ([1]) proved that, under the condition that k zeros in $\overline{\Delta}$ are centered at 0,

$$(1.4) \quad \rho(n, k, l) \leq 10.477 \frac{(n-k+1)(k-l)}{\sqrt{k}}$$

and

$$(1.5) \quad \rho(n, k, k-1) \leq C \max \left\{ (n-k+1)^{1/2} k^{-1/4}, (n-k+1) k^{-2/3} \right\}.$$

One can also impose an additional condition on $p \in \mathcal{P}_{n,k}$ to have zeros in Δ isolated from the other zeros. Following Pan ([3]), we define *the isolation ratio of a polynomial p with respect to a disk D* , $I(p, D)$, as:

$$(1.6) \quad I(p, D(c, r)) := \sup \{ \mu > 0 : D(c, \mu r) \text{ contains exactly the same zeros of } p \text{ as } D(c, r) \}.$$

We say that a disk D is *f-isolated* if $I(p, D) \geq f$. The isolation ratio is also scaling and translation invariant. We define *the isolation ratio of p* as $I(p) := I(p, \overline{\Delta})$. Renegar ([5]) proved the following result:

$$(1.7) \quad \text{If } I(p) \geq 15n^3, \text{ then } \rho(n, k, k-1) \leq \frac{3n}{2} \text{ and } I(p^{(k-1)}, \overline{D}(0, 3n/2)) \geq \frac{I(p)}{10n^2}.$$

In other words, if k zeros of a polynomial p are well-isolated in some disk, then $p^{(k-1)}$ has an isolated single zero in a larger concentric disk. Using Walsh's Coincidence Theorem ([2]) we improve the bound in (1.7). We also generalize Biernacki's proof of (1.2) and obtain an upper bound for $\rho(n, k, l)$ which is smaller than (1.4) for some special choices of n, k and l .

Results of this type have found applications in constructing low complexity algorithms for finding zeros of polynomials. Smale and Renegar ([7], [5]) derived quantitative criteria for a point to be in the domain of quadratic convergence of Newton's algorithm. The basic idea is that if a zero is isolated from other zeros, then we need only crude approximation to obtain fast convergence. However, if we have an isolated cluster of k zeros, then $p^{(k-1)}$ should have a single zero nearby. In this case, the result like (1.7) allows us to apply Newton's algorithm to $p^{(k-1)}$ with quadratic convergence. To establish initial isolation of clusters of zeros, various global search algorithms are used ([4], [5])

2. MAIN RESULTS

First we prove the following proposition:

Proposition 2.1.

$$(2.1) \quad \text{If } I(p) \geq 1 + 2 \frac{l(n-k)}{k-l+1}, \text{ then } \rho(n, k, l) \leq 1 \text{ and}$$

$$I(p^{(l)}) \geq \frac{(k-l+1)I(p) - l(n-k)}{k-l+1 + l(n-k)}.$$

Proof. Let $I(p) = R$ and let $p(z) = f(z)g(z)$ where $f(z)$ is a polynomial of degree k whose zeros lie in $\overline{\Delta}$ and $g(z)$ is a polynomial of degree $n-k$ whose zeros lie in

$\{z : |z| > R\}$. Let $m = \min(n - k, l)$. If z is a solution of the equation

$$(2.2) \quad p^{(l)}(z) = \sum_{i=0}^m \binom{l}{i} g^{(i)}(z) f^{(l-i)}(z) = 0,$$

then the Walsh Coincidence Theorem implies that there exist $x_1, |x_1| \leq 1$ and $x_2, |x_2| \geq R$ such that

$$(2.3) \quad \sum_{i=0}^m \binom{l}{i} ((z - x_2)^{n-k})^{(i)} ((z - x_1)^k)^{(l-i)} = 0,$$

i.e.

$$(2.4) \quad \sum_{i=0}^m \binom{l}{i} \frac{(n-k)!}{(n-k-i)!} (z - x_2)^{n-k-i} \frac{k!}{(k-l+i)!} (z - x_1)^{k-l+i} = 0,$$

i.e.

$$(2.5) \quad (z - x_1)^{k-l} (z - x_2)^{n-k-m} \sum_{i=0}^m \binom{l}{i} \binom{n-l}{n-k-i} (z - x_2)^{m-i} (z - x_1)^i = 0.$$

Also z must lie in the interval $[x_1, x_2]$ and either $z = x_1$ or $z = x_2$ or z is a zero of the polynomial

$$(2.6) \quad q(z) = \sum_{i=0}^m \binom{l}{i} \binom{n-l}{n-k-i} (z - x_2)^{m-i} (z - x_1)^i.$$

Let $w = \frac{z - x_1}{z - x_2}$ and let

$$(2.7) \quad h(w) = \sum_{i=0}^m \binom{l}{i} \binom{n-l}{n-k-i} w^i = \sum_{i=0}^m a_i w^i.$$

All zeros of $h(w)$ lie on the negative real axis and by the Eneström-Kakeya Theorem they satisfy the inequality

$$(2.8) \quad \min_{0 \leq i \leq m-1} \frac{a_i}{a_{i+1}} \leq |w| \leq \max_{0 \leq i \leq m-1} \frac{a_i}{a_{i+1}}.$$

We have that

$$(2.9) \quad \frac{a_i}{a_{i+1}} = \frac{(i+1)(k-l+i+1)}{(l-i)(n-k-i)};$$

therefore

$$(2.10) \quad \zeta := \min_{0 \leq i \leq m-1} \frac{a_i}{a_{i+1}} = \frac{k-l+1}{l(n-k)}$$

and

$$(2.11) \quad \xi := \max_{0 \leq i \leq m-1} \frac{a_i}{a_{i+1}} = \frac{m(m+k-l)}{(l-m+1)(n-k-m+1)}.$$

Since all zeros of $h(w)$ lie in the interval $[-\xi, -\zeta]$, all zeros of $q(z)$ lie in the interval $\left[\frac{\zeta x_2 + x_1}{\zeta + 1}, \frac{\xi x_2 + x_1}{\xi + 1} \right]$. By continuity, $p^{(l)}(z)$ has $k-l$ zeros in $\overline{\Delta}$ if $\left| \frac{\zeta x_2 + x_1}{\zeta + 1} \right| \geq 1$. Since

$\left| \frac{\zeta x_2 + x_1}{\zeta + 1} \right| \leq \frac{\zeta R - 1}{\zeta + 1}$, the above inequality holds if $\frac{\zeta R - 1}{\zeta + 1} \geq 1$, i.e. if $R \geq \left(1 + \frac{2}{\zeta}\right) = \left(1 + 2\frac{l(n-k)}{k-l+1}\right)$. These are the only zeros of $p^{(l)}(z)$ in the disk

$$(2.12) \quad D\left(0, \frac{\zeta R - 1}{\zeta + 1}\right) = D\left(0, \frac{(k-l+1)R - l(n-k)}{k-l+1+l(n-k)}\right).$$

□

We use Proposition 2.1 to prove the following result analogous to (1.7).

Theorem 2.2.

$$(2.13) \quad \text{If } I(p) \geq \frac{n^2}{4}, \text{ then } \rho(n, k, k-1) \leq 1 \text{ and } I\left(p^{(k-1)}\right) \geq \frac{4I(p)}{n^2}.$$

Proof. If $l = k - 1$ and $I(p) \geq 1 + (k-1)(n-k)$, then, by Proposition 2.1, $\overline{\Delta}$ contains a zero of $p^{(k-1)}$. Since $1 + (k-1)(n-k) \leq k(n-k) \leq n^2/4$, the same is true if $I(p) \geq n^2/4$. Also,

$$(2.14) \quad I\left(p^{(k-1)}\right) \geq \frac{2I(p) - (k-1)(n-k)}{2 + (k-1)(n-k)} > \frac{I(p)}{1 + (k-1)(n-k)} \geq \frac{I(p)}{k(n-k)} \geq \frac{4I(p)}{n^2}.$$

□

Remark 2.3. We don't know of any estimates for $\rho(n, k, k-1)$ if the isolation ratio of p is smaller than $n^2/4$. In applications to zero-finding we can use the classical Graeffe process of root-squaring to control the isolation ratio of a polynomial ([3],[6]). If p_0 has zeros z_1, \dots, z_n and isolation ratio $C > 1$, then the polynomial

$$(2.15) \quad p_1(z) := p_0(\sqrt{z})p_0(-\sqrt{z})$$

has zeros z_1^2, \dots, z_n^2 and isolation ratio C^2 . By iterating this process we obtain polynomial p_m with zeros $z_1^{2^m}, \dots, z_n^{2^m}$ and isolation ratio C^{2^m} .

We can also apply Proposition 2.1 to obtain an upper bound for $\rho(n, k, l)$.

Theorem 2.4.

$$(2.16) \quad \rho(n, k, l) \leq \frac{n+l}{n-l} \prod_{i=k}^{n-2} \left(1 + \frac{2l(n-i)}{i-l+1}\right).$$

Proof. Let us assume that the zeros of $p(z)$ are ordered by their moduli, $|z_1| \leq |z_2| \leq \dots \leq |z_n|$. If $k = n - 1$, then $p(z) = f(z)(z - z_n)$, where $f(z)$ is a polynomial with zeros $z_i, |z_i| \leq 1, 1 \leq i \leq n - 1$. We have

$$(2.17) \quad p^{(l)}(z) = f^{(l)}(z)(z - z_n) + lf^{(l-1)}(z).$$

By the Walsh Coincidence Theorem, if z is a zero of $p^{(l)}(z)$ it is also a zero of the equation:

$$(2.18) \quad ((z - \gamma)^{n-1})^{(l)}(z)(z - z_n) + l((z - \gamma)^{n-1})^{(l-1)}(z) = 0$$

where $|\gamma| \leq 1$, i.e.,

$$(2.19) \quad (z - \gamma)^{n-1-l}((n-l)(z - z_n) + l(z - \gamma)) = 0.$$

Therefore $z = \gamma$ or $z = ((n-l)z_n + l\gamma)/n$. Also if $|((n-l)z_n + l\gamma)/n| > 1$, then there are exactly $k-1-l$ zeros of $p^{(l)}(z)$ in $\overline{\Delta}$. This is the case if $|z_n| > \frac{n+l}{n-l}$. Otherwise, by the Gauss-Lucas Theorem, all zeros of $p^{(l)}(z)$ are in the disk $D\left(0, \frac{n+l}{n-l}\right)$. Now let's fix K and suppose that (2.16) holds for $k = n-1 \dots K-1$. By Proposition 2.1, if $|z_{K+1}| > 1 + 2\frac{l(n-K)}{K-l+1}$, then there are exactly $K-l$ zeros of $p^{(l)}(z)$ in $\overline{\Delta}$. Otherwise there are $K+1$ zeros of $p(z)$ in the disk $D\left(0, 1 + 2\frac{l(n-K)}{K-l+1}\right)$ and by induction there are at least $K-l$ zeros of $p^{(l)}(z)$ in $D(0, R)$, where

$$(2.20) \quad R \leq \frac{n+l}{n-l} \prod_{i=K}^{n-2} \left(1 + \frac{2l(n-i)}{i-l+1}\right).$$

□

Inequality (2.16) can be better in some cases than (1.4), namely if k and l are close to n . In the most interesting case $l = k-1$, Coppersmith and Neff gave a lower bound on $\rho(n, k, k-1)$ and also conjectured that if $k > n/2$, then $\rho(n, k, k-1)$ is bounded by a constant.

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