ON THE CONSTRUCTION OF SETS OF MUTUALLY ORTHOGONAL LATIN SQUARES AND THE FALSITY OF A CONJECTURE OF EULER⁽¹⁾

BY

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1. Introduction. If $p_1^{n_1} p_2^{n_2} \cdots p_u^{n_u}$ is the prime power decomposition of an integer v , and we define the arithmetic function $n(v)$ by

$$n(v) = \min(p_1^{n_1}, p_2^{n_2}, \cdots, p_u^{n_u}) - 1$$

then it is known, MacNeish [16] and Mann [17] that there exists a set of at least $n(v)$ mutually orthogonal Latin squares (m.o.l.s.) of order v . It seemed plausible that $n(v)$ is also the maximum possible number of m.o.l.s. of order v . This would have implied the correctness of Euler's [13] conjecture about the nonexistence of two orthogonal Latin squares of order v when $v \equiv 2 \pmod{4}$, since $n(v) = 1$ in this case. If we denote by $N(v)$ the maximum possible number of mutually orthogonal Latin squares of order v , then Parker [18] showed that in certain cases $N(v) > n(v)$ by proving that if there exists a balanced incomplete block (BIB) design with v treatments, $\lambda = 1$, and block size k which is a prime power then $N(v) \geq k - 2$. This result cast a doubt on the validity of Euler's conjecture. In this paper we generalize Parker's method of constructing m.o.l.s., by showing that a very general class of designs, which we have called pairwise balanced designs of index unity, can be used for the construction of sets of m.o.l.s. Various applications of this method have been made. In particular we show that Euler's conjecture is false for an infinity of values of v , including all values of the form $36m + 22$. We also provide a table for all values of $v \leq 150$, for which our methods enable us to show that $N(v) > n(v)$. The smallest number for which we have been able to demonstrate the falsity of Euler's conjecture is 22. Two orthogonal squares of order 22 constructed by our method are given in the Appendix. Several attempts, necessarily erroneous, have been made in the past to prove Euler's conjecture, e.g., Peterson [19], Wernicke [20] and MacNeish [16]. Levi [14, p. 14] points

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out the inadmissibility of the argument used by Peterson and MacNeish. Wernicke's argument was shown to be fallacious by MacNeish [15].

2. Orthogonal Latin squares and orthogonal arrays. A Latin square of order v may be defined as an arrangement of v symbols, say $1, 2, \dots, v$ in a $v \times v$ square such that each symbol occurs once in every row and every column. Two Latin squares are said to be orthogonal, when if they are superposed each symbol of the first square occurs just once with each symbol of the second square. A set of mutually orthogonal Latin squares (m.o.l.s.) is a set of Latin squares any two of which are orthogonal.

Defining $N(v)$ and $n(v)$ as in the introduction, we have already referred to the MacNeish-Mann theorem viz. $N(v) \geq n(v)$. We shall now give a useful lemma which generalizes this result.

LEMMA 1. *If $v = v_1 v_2 \dots v_u$ then $N(v) \geq \min(N(v_1), N(v_2), \dots, N(v_u))$.*

We shall derive this lemma from a general result on orthogonal arrays due to Bush and one of the authors [5; 10]. Let us consider a matrix $A = (a_{ij})$, where each a_{ij} represents one of the symbols $1, 2, \dots, v$, with $m = \lambda v^t$ columns and r rows. Consider all t rowed submatrices which can be formed from A by choosing t rows, $t \leq r$. Each column of the submatrix can be regarded as an ordered t -plet. The matrix A is called an orthogonal array (m, r, v, t) of size m , r constraints, v levels, strength t and index λ ; if each t rowed submatrix which can be formed from A contains every one of the v^t possible ordered t -plets each repeated λ times. If $f(m, v, t)$ denotes the maximum number of constraints for given m, v and t , then the following theorem is known [5; 10].

THEOREM. *If m_i is divisible by v_i^t for $i = 1, 2, \dots, u$, then*

$$f(m_1 m_2 \dots m_u, v_1 v_2 \dots v_u, t) \geq \min(r_1, r_2, \dots, r_u)$$

where $r_i = f(m_i, v_i, t)$.

Now it is known [11] that an orthogonal array $(v^2, q, v, 2)$ of strength 2 and index unity is equivalent to the existence of a set of $q-2$ m.o.l.s. of order v , which may be obtained from this array by using two of the rows for coordinatization. Conversely the array can be obtained from the squares. Hence $f(v^2, v, 2) = N(v) + 2$. The required lemma follows by taking $m_i = v_i^2$, $t = 2$ in the theorem.

The MacNeish-Mann theorem follows from the lemma by taking $v_i = p_i^{n_i}$, $i = 1, 2, \dots, u$. If $N(v)$ had always been equal to $n(v)$ the lemma would add nothing new. However this is not the case. For example it will be shown in §5, that $N(21) \geq 4$. It follows from the lemma that $N(105) \geq \min(N(21), N(5)) \geq 4$ whereas $n(105) = \min(3, 5, 7) - 1 = 2$.

3. Pairwise balanced designs of index unity. An arrangement of v objects (called treatments) in b sets (called blocks) will be called a pairwise balanced

design of index unity and type $(v; k_1, k_2, \dots, k_m)$ if each block contains either $k_1, k_2, \dots$ or k_m treatments which are all distinct ($k_i \leq v, k_i \neq k_j$), and every pair of distinct treatments occurs in exactly one block of the design. If the number of blocks containing k_i treatments is b_i , then clearly

$$(3.1) \quad b = \sum_{i=1}^m b_i, \quad v(v-1) = \sum_{i=1}^m b_i k_i (k_i - 1).$$

Consider a pairwise balanced design (D) of index unity and type $(v; k_1, k_2, \dots, k_m)$. The subdesign (D_i) formed by the blocks of size k_i will be called the i th equiblock component of (D) , $i=1, 2, \dots, m$.

A subset of blocks belonging to any equiblock component (D_i) will be said to be of type I if every treatment occurs in the subset exactly k_i times. The number of blocks in such a subset is clearly v . As noted by Levi using König's theorem on the decomposition of even regular graphs [14, pp. 4-6], we can rearrange the treatments within the blocks of the subset in such a way that every treatment comes in each position exactly once. If the v blocks of the subset are written out as columns, each treatment occurs exactly once in every row. When so written out the blocks will be said to be in the standard form.

A subset of blocks belonging to any equiblock component (D_i) will be said to be of type II if every treatment occurs in the subset exactly once. The number of blocks in such a subset is v/k_i .

The component (D_i) will be defined to be separable if the blocks can be divided into subsets of type I or type II (both types may occur in (D_i) at the same time). Let r_i be the number of subsets of type I and s_i the number of subsets of type II in (D_i) , then clearly

$$(3.2) \quad v(r_i k_i + s_i) = k_i b_i.$$

The design (D) is defined to be separable if each equiblock component (D_i) is separable. It follows from (3.1) that

$$(3.3) \quad v-1 = \sum_{i=1}^m (r_i k_i + s_i)(k_i - 1).$$

4. Main theorem. We shall first prove the following:

LEMMA 2. *Suppose there exists a set Σ of $q-1$ m.o.l.s. of order k , then we can construct a $q \times k(k-1)$ matrix P , whose elements are the symbols $1, 2, \dots, k$ and such that (i) any ordered pair (i, j) , $i \neq j$ occurs as a column exactly once in any two rowed submatrix of P , (ii) P can be subdivided into $k-1$ submatrices $P_1, P_2, \dots, P_{k-1}$ of order $q \times k$ such that in each row of P_c , $1 \leq c \leq k-1$, each of the symbols $1, 2, \dots, k$ occurs exactly once.*

Proof. It has been shown in [9] that matrix P with the property (i) can be obtained in the following manner. Without loss of generality we can take

the set Σ in a form in which the first row of each Latin square contains the symbols $1, 2, \dots, k$ in that order and prefix to the set Σ a square containing the symbol i in each position in the i th column. If we then write the elements of each square in a single row such that the symbol in the i th row and j th column occupies the n th position in the row where $n = k(i-1) + j$, and from the resulting $q \times k^2$ matrix we delete the first k columns, we get the matrix P . If P_c is the submatrix of P formed by the columns $k(c-1) + j, j = 1, 2, \dots, k$, then the property (ii) follows from the fact that each of the symbols $1, 2, \dots, k$ appears exactly once in any given row of a Latin square belonging to Σ .

Let γ be a column of k distinct treatments chosen from the set $t_1, t_2, \dots, t_v$ then following the notation used in [9], we shall denote by $P(\gamma)$, the $q \times k(k-1)$ matrix obtained from P on replacing the symbol i by the treatment occurring in the i th position in γ . A similar meaning will be assigned to $P_i(\gamma)$ and $\pi_{cj}(\gamma)$ where π_{cj} denotes the j th column of P_c .

Clearly every treatment of γ occurs once in every row of $P_c(\gamma)$, $c = 1, 2, \dots, k-1$, and if t_a and t_b are any two treatments of γ then the ordered pair (t_a, t_b) occurs as a column exactly once in any two rowed submatrix of $P(\gamma)$.

THEOREM 1. *Let there exist a pairwise balanced design (D) of index unity and type $(v; k_1, k_2, \dots, k_m)$ and suppose there exist $q_i - 1$ m.o.l.s. of order k_i . If*

$$q = \min(q_1, q_2, \dots, q_m)$$

then there exist at least $q-2$ m.o.l.s. of order v . If the design (D) is separable then the number of m.o.l.s. of order v is at least $q-1$.

Proof. The first part of the theorem was proved in [9]. However in order to prove the second part it will be necessary to reproduce this proof with slightly altered notation. Let the treatments of the design be $t_1, t_2, \dots, t_v$, and let the blocks of the design (written out as columns) belonging to the equiblock component (D_i) be $\delta_{i1}, \delta_{i2}, \dots, \delta_{ib_i}$ ($i = 1, 2, \dots, m$). We now define the $k_i \times b_i$ matrix D_i by

$$D_i = [\delta_{i1}, \delta_{i2}, \dots, \delta_{ib_i}].$$

Let P_i be the matrix of order $q_i \times k_i(k_i-1)$ defined in Lemma 2, the elements of P_i being the symbols $1, 2, \dots, k_i$. Let P_{ic} , $c = 1, 2, \dots, k_i-1$ be the submatrices of P_i , such that each row of P_{ic} contains the symbols $1, 2, \dots, k_i$ exactly once and let π_{icj} be the j th column of P_{ic} , $j = 1, 2, \dots, k_i$. Let $P_i(\delta_{iu})$, $u = 1, 2, \dots, b_i$, be the matrix obtained from P_i and δ_{iu} , and let

$$P_i(D_i) = [P_i(\delta_{i1}), P_i(\delta_{i2}), \dots, P_i(\delta_{ib_i})].$$

Then $P_i(D_i)$ is of order $q_i \times b_i k_i(k_i-1)$. If t_a and t_b are any two treatments occurring in the same block of (D_i) , then the ordered pair (t_a, t_b) occurs exactly once as a column in any two rowed submatrix of $P_i(D_i)$. Let Δ_i be the matrix obtained from $P_i(D_i)$ by retaining only the first q rows, and let

$$\Delta = [\Delta_1, \Delta_2, \dots, \Delta_m].$$

Then from (3.1), Δ is of order $q \times v(v-1)$, and it follows from the properties of the design that any two rowed submatrix of Δ contains as a column each ordered pair of two distinct treatments chosen from $t_1, t_2, \dots, t_v$ exactly once. This property of Δ will be referred to as the property τ_2 .

Let Δ_0 be a $q \times v$ matrix whose i th column contains t_i in every position ($i=1, 2, \dots, v$). Then from [11] the matrix $[\Delta_0, \Delta]$ is an orthogonal array $[v^2, q, v, 2]$ of order v^2 , q constraints, v levels and strength 2. Using two rows to coordinatize we get a set of $q-2$ m.o.l.s. which proves the first part of the theorem.

To prove the second part of our theorem we shall introduce the concept "column equivalence." Two matrices A and B will be called column equivalent if they can be obtained from each other by a rearrangement of columns. This relation will be denoted by $A \sim B$.

Suppose now that the design (D) is separable. Let $(D_{i1}), \dots, (D_{ir_i})$ be the subsets of type I, and $(D_{i1}^*), \dots, (D_{is_i}^*)$ be the subsets of type II into which the blocks of (D_i) are divisible, the blocks of subsets of type I being written in the standard form. We shall show that after a rearrangement of columns, Δ can be divided into submatrices of order $q \times v$, possessing the property τ_1 that each row contains each of the treatments $t_1, t_2, \dots, t_v$ exactly once. Now

$$P_i(D_{if}) \sim [\dots, \pi_{icj}(D_{if}), \dots], \quad c = 1, 2, \dots, k_i - 1, j = 1, 2, \dots, k_i.$$

Since $\pi_{icj}(D_{if}), f=1, 2, \dots, r_i$, clearly possesses the property τ_1 , $P_i(D_{if})$ is divisible (up to a rearrangement of columns) into $k_i(k_i-1)$ submatrices each having the property τ_1 .

Again

$$P_i(D_{ig}^*) \sim [\dots, P_{ic}(D_{ig}^*), \dots], \quad c = 1, 2, \dots, k_i - 1.$$

Since $P_{ic}(D_{ig}^*), g=1, 2, \dots, s_i$, possesses the property τ_1 , $P_i(D_{ig}^*)$ is divisible (up to a rearrangement of columns) into k_i-1 submatrices each with the property τ_1 . Since

$$P_i(D_i) = [P_i(D_{i1}), \dots, P_i(D_{ir_i}), P_i(D_{i1}^*), \dots, P_i(D_{is_i}^*)]$$

and since the property τ_1 remains unchanged by deletion of rows it follows from (3.3) that after a rearrangement of columns Δ can be divided into $v-1$ submatrices $\Delta^{(1)}, \Delta^{(2)}, \dots, \Delta^{(v-1)}$ each with the property τ_1 . It should also be noted that the property τ_2 of Δ remains unchanged by reordering the columns. It follows that if α'_i denotes a $1 \times v$ row vector each element of which is t_i , the matrix

$$\begin{bmatrix} \alpha'_v & \alpha'_1 & \dots & \alpha'_{v-1} \\ \Delta_0 & \Delta^{(1)} & \dots & \Delta^{(v-1)} \end{bmatrix}$$

is an orthogonal array $[v^2, q+1, v, 2]$ of order v^2 , $q+1$ constraints, v levels and strength 2, and can therefore be used to construct a set of $q-1$ m.o.l.s. of order v .

5. Use of balanced incomplete block designs for obtaining sets of mutually orthogonal Latin squares. A balanced incomplete block (BIB) design with parameters v, b, r, k, λ is an arrangement of v objects or treatments into b sets or blocks such that (i) each block contains $k < v$ different treatments, (ii) each treatment occurs in r different blocks, (iii) each pair of treatments occurs together in exactly λ blocks. The parameters satisfy the relations

$$\lambda(v-1) = r(k-1), \quad bk = vr, \quad b \geq v.$$

These conditions are necessary but not sufficient for the existence of a BIB design. BIB designs were first introduced into statistical studies by Yates [21], but occur in earlier literature in connection with various combinatorial problems. Subsequent to Yates many authors have dealt with the problem of constructing these designs. Without attempting a complete bibliography we shall only refer to [2]. A BIB design is said to be symmetrical if $v=b$, and in consequence $k=r$. A BIB design is said to be resolvable [3] if the blocks can be divided into sets, such that the blocks of any given set contain each treatment exactly once.

A BIB design with $\lambda=1$ is clearly a pairwise balanced design of index unity and type $(v; k)$. We shall denote such a design by BIB $(v; k)$. It is evident that a symmetric or a resolvable BIB design is separable, since in the first case the blocks constitute a single set of type I, and in the second case the blocks can be divided into r subsets of type II. From Theorem 1 we have

THEOREM 2. *If there exists a BIB design with v treatments, block size k and $\lambda=1$, then $N(v) \geq N(k)-1$. If the BIB design is separable $N(v) \geq N(k)$.*

When k is a prime power $N(k)=k-1$, hence we have Parker's theorem [18] mentioned in the introduction which asserts that $N(v) \geq k-2$, when k is a prime power. This result can be improved by unity when the design is separable.

THEOREM 3. *$N(v) \geq N(k)$ in the following cases:*

- (i) $v=s^2+s+1, k=s+1$, where s is a prime power,
- (ii) $v=s^3+1, k=s+1$, where s is a prime power,
- (iii) $v=s(s-1)/2, k=s/2=2^{m-1}, m \geq 2$.

Proof. Part (i) follows from Theorem 2, since there exists a symmetrical BIB design with parameters $v=b=s^2+s+1, r=k=s+1, \lambda=1$ when s is a prime power [2, p. 364]. The existence of the design is equivalent to the existence of a projective plane PG(2, s) of order s .

EXAMPLE (1). If $s=4, N(k)=N(5)=4$. Hence there exist at least 4 m.o.l.s. of order 21. The MacNeish-Mann result only assures the existence of 2 orthogonal Latin squares.

EXAMPLE (2). If s is a Mersenne prime > 3 then $k = s + 1 = 2^p$ and $N(k) = 2^p - 1$. Hence $N(s^2 + s + 1) \geq s$, whereas $n(s^2 + s + 1) = 2$, e.g., if we take $s = 31$ then $N(993) \geq 31$, $n(993) = 2$.

EXAMPLE (3). If $s = 49$, $k = 50$. Parker's theorem is inapplicable. It is shown in §6, Example (7), that $N(50) \geq 5$. Hence $N(49^2 + 49 + 1) \geq 5$ whereas $n(49^2 + 49 + 1) = 2$.

Part (ii) follows from the existence of the series of resolvable BIB designs

$$v = s^3 + 1, \quad b = s^2(s^2 - s + 1), \quad r = s^2, \quad k = s + 1, \quad \lambda = 1,$$

which has been proved in [6].

EXAMPLE (4). Let $v = 31^3 + 1 = 2^5 \cdot 7^2 \cdot 19 = 29792$, $k = 32 = 2^5$. Hence $N(v) \geq 31$ whereas $n(v) = 18$.

To prove part (iii) we shall establish the existence of a series of resolvable BIB designs with parameters

$$(5.1) \quad v = s(s - 1)/2, \quad b = s^2 - 1, \quad r = s + 1, \quad k = s/2 = 2^{m-1}, \quad \lambda = 1.$$

Consider the finite projective plane $PG(2, 2^m)$, and take a nondegenerate conic C on this plane, e.g., the conic

$$xz = y^2.$$

We know [4, p. 158] that there are $s + 1$ points $P_1, P_2, \dots, P_{s+1}$ on this conic, where $s = 2^m$. Through any point $P_i = (x', y', z')$ on the conic there will pass $s + 1$ lines s of which meet the conic in the other s points on the conic, and the remaining line $z'x + x'z = 0$, meets the conic in the single point P_i . This may be called a tangent to the conic at P_i . The $s + 1$ tangents to the conic all pass through the point $P_0 = (0, 1, 0)$ which may be called the pole of the conic. The $s^2 + s + 1$ lines of the plane may be divided into three classes: (a) The $s(s + 1)/2$ intersectors each of which meets the conic in 2 points but does not pass through P_0 ; (b) The $s + 1$ tangents each of which meets the conic in one point and passes through P_0 ; (c) The $s(s - 1)/2$ nonintersectors which do not meet the conic and hence do not pass through P_0 . Let the $s^2 - 1$ points of the plane other than P_0 and the points of the conic be called retained points. Consider the configuration of the retained points and nonintersectors. If the retained points are called blocks, and the nonintersectors are called treatments, and if treatment is considered to belong to a block if the corresponding line and point are incident then we get the design (5.1). This follows from the fact that on each nonintersector there lie $s + 1$ points. This gives $r = s + 1$. Also through each retained point P there passes one tangent, viz. PP_0 . If the point of contact of this tangent is P_i then the other s points on the conic lie two by two on $s/2$ intersectors through P . Hence there pass $s/2$ nonintersectors through P . This gives $k = s/2$. That $\lambda = 1$ follows from the fact that any two nonintersectors meet in one and only one point which is a retained point.

Next we have to prove that the design is resolvable, i.e., the blocks

can be separated in groups. We have to find sets of $s-1$ blocks which contain each treatment exactly once, i.e., sets of $s-1$ retained points such that each nonintersector passes through one and only one point of the set. The $s-1$ retained points on any intersector or tangent provide such a set. Finally we have to choose a set of $s+1$ intersectors or tangents which contain between them all the retained points. Such a set is provided by the $s+1$ tangents, or the $s+1$ lines through any point on the conic. This proves the resolvability of the design.

EXAMPLE (5). Parker leaves undecided the question whether there are an infinity of values of v for which $N(v) > n(v)$. This question is decided in the affirmative by part (iii) of our theorem, since if we take $m=2n$, then $N(v) \geq N(2^{m-1}) = 2^{2n-1} - 1$. Also $v = 2^{2n-1}(2^n + 1)(2^n - 1)$. Hence $n(v) \leq 2^n - 2 < N(v)$. For example, let $m=4$, then $v=120$, $k=8$, $N(120) \geq N(k)=7$, $n(120) = n(8 \cdot 5 \cdot 3) = 2$.

If from a BIB design with v treatments, block size k and $\lambda=1$, we omit $x \leq k$ treatments belonging to any block, we obviously get a pairwise balanced design of index unity and type $(v-x; k, k-1, k-x)$. Applying Theorem 1, we get

THEOREM 4. *Existence of a BIB design with v treatments, block size k and $\lambda=1$, implies that*

$$N(v-x) \geq \min(N(k), N(k-1), N(k-x)) - 1$$

if $x \leq k$.

EXAMPLE (6). Consider the symmetrical BIB design $v=b=s^2+s+1$, $r=k=s+1$, $\lambda=1$, when $s=16$. If we take $x=6, 8$ and 9 respectively, we obtain $N(267) \geq 9$, $N(265) \geq 7$, $N(264) \geq 6$, whereas $n(267)=2$, $n(265)=4$, $n(264)=2$.

6. The falsity of Euler's conjecture regarding the nonexistence of two orthogonal Latin squares of order $v \equiv 2 \pmod{4}$, and the method of adjunction of treatments.

LEMMA 3. *There exists a resolvable series of BIB designs with parameters*

$$(6.1) \quad v^* = 24m + 15, \quad b^* = (8m + 5)(12m + 7), \quad r^* = 12m + 7, \quad k^* = 3, \quad \lambda^* = 1.$$

Proof. The existence of the design is equivalent to a solution of the Kirkman school girl problem for triplets with $n=24m+15$ girls. There is an extensive literature on this problem and a complete bibliography of all work up to 1911 is given in [12], and an excellent summary of the known results is contained in [1, Chapter X, pp. 267-298]. The problem is solvable for all values of n of the form $6t+3$. The special case required by us is $t=4m+2$. The girls can be considered as the treatments of (6.1), triplets as blocks, and each day's walk as a complete replication.

In the special case when $12m+7$ is a prime power, an analytical solution can be given as follows. We first prove

LEMMA 4. *If x is a primitive element of the Galois field $GF(12m+7)$, where $12m+7$ is a prime power, then there exist odd integers q and α such that*

$$x^q = (x^\alpha + 1)/(x^\alpha - 1).$$

Proof. The lemma is proved by considering the $(1, 1)$ correspondence

$$y = (z + 1)/(z - 1)$$

between the elements of $GF(12m+7)$, augmented by infinity. The correspondence is an involution and it is easily shown that there must exist a pair of corresponding odd powers of x .

To each element u of the field $GF(12m+7)$ let there correspond two treatments u_1 and u_2 . To the $24m+14$ treatments thus obtained adjoin the treatment ∞ . The existence and resolvability of the design can now be proved by using the second fundamental theorem of the method of differences developed in [2]. Consider the set of $8m+5$ initial blocks

$$\begin{aligned} & (x_1^{2i}, x_1^{2i+4m+2}, x_1^{2i+8m+4}), \quad i = 0, 1, \dots, 2m; \\ & (x_1^{\alpha+j}, x_2^{2j}, x_2^{6m+3+2j}), \quad j = 0, 1, \dots, 6m+2; \\ & (0_1, 0_2, \infty). \end{aligned}$$

Using methods on pages 374 and 384 of [2] it is easy to verify that all the nonzero pure differences and all the mixed differences are symmetrically repeated once. Hence the design (6.1) can be obtained by developing the initial blocks. Since the set of initial blocks gives a complete replication, the design is resolvable.

We can now proceed to prove the main theorem of this section.

THEOREM 5. *There exist at least two orthogonal Latin squares of order $v=36m+22$, where m is any non-negative integer.*

Proof. Let us start with the BIB design whose parameters are given by (6.1). Take $12m+7$ new treatments $\theta_1, \theta_2, \dots, \theta_{12m+7}$ and to each block of the i th replication add the treatment θ_i . Also take a new block consisting of the new treatments. The augmented design is now a pairwise balanced design of index unity of the type $(v; k_1, k_2)$ where $v=36m+22$, $k_1=4$, $k_2=12m+7$. Since $q_1-1=N(k_1)=3$, $q_2-1=N(k_2) \geq n(12m+7) \geq 4$, $q=\min(q_1, q_2)=4$, it follows from Theorem 1, that there exist at least two orthogonal Latin squares of order $v=36m+22$.

Two superposed orthogonal Latin squares of order 22 constructed by this method are exhibited in the Appendix, the capital letters corresponding to the first square and the small letters to the second.

We shall now show that there exist numbers $v \equiv 2 \pmod{4}$ for which there exist many more than two mutually orthogonal Latin squares of order v .

THEOREM 6. *If s and $s+1$ are both prime powers, then for any integer $m \geq 2$.*

$$N(s^m + 1) \geq s - 2.$$

Proof. Consider the Euclidean geometry $EG(m, s)$. Taking the points and lines of this geometry as the treatments and blocks respectively, we get a BIB design with parameters $v^* = s^m$, $b^* = r^* s^{m-1}$, $r^* = 1 + s + \dots + s^{m-1}$, $k^* = s$, $\lambda^* = 1$. Further the design is resolvable into r^* replications corresponding to the lines of the Euclidean geometry passing through a point of the $(m-1)$ -flat at infinity of the projective geometry $PG(m, s)$ in which the Euclidean geometry can be embedded. Add a new treatment θ to the blocks of the first replication. This is equivalent to the adjunction of a point at infinity to the lines of a parallel pencil. The augmented design is a pairwise balanced design of index unity and type $(s^m + 1; s + 1, s)$. Since $s+1$ and s are both prime powers Theorem 1 gives $N(s^m + 1) \geq s - 2$.

We may note that s and $s+1$ will be prime powers if (i) s is a Mersenne number, (ii) $s+1$ is a Fermat prime, (iii) $s = 8$.

REMARK. If s is a Mersenne number, $s^{2^l} + 1 \equiv 2 \pmod{4}$. Hence if $s > 3$, this provides another infinite set of values v of the form $4t+2$ for which at least $s-2$ m.o.l.s. of order v can be constructed, thus again disproving Euler's conjecture.

EXAMPLE (7). Taking $l=1$ and $s=7, 8$ or 31 , we get $N(50) \geq 5$, $N(65) \geq 6$, $N(962) \geq 29$.

Instead of one we may add x new treatments, $x \leq r^*$, one to each block of a complete replication, and add a new block containing the new treatments. We get a pairwise balanced design of index unity and type $(s^m + x; s + 1, s, x)$. Hence the following:

COROLLARY 1. *If s and $s+1$ are prime powers and $1 < x \leq 1 + s + \dots + s^{m-1}$, $m \geq 2$, $N(s^m + x) \geq \min(s-1, N(x)) - 1$.*

EXAMPLE (8). Taking $s=7$, $m=2$, $x=5$, we get $N(54) \geq 3$.

If we do not insist that s and $s+1$ are both prime powers in the above theorem, then we get

COROLLARY 2. *If s is a prime power, and $m \geq 2$, then*

$$N(s^m + 1) \geq \min(s - 1, N(s + 1)) - 1,$$

and if $1 < x \leq 1 + s + \dots + s^{m-1}$

$$N(s^m + x) \geq \min(s - 1, N(s + 1), N(x)) - 1.$$

EXAMPLE (9). Let $s=19$, $m=2$. Then $N(s+1) = N(20) \geq n(20) = 3$. Hence $N(362) \geq 2$.

THEOREM 7. *If there exist at least two orthogonal Latin squares of order $4t+2$, then there exist at least two orthogonal Latin squares of any order v which is an odd multiple of $4t+2$.*

Proof. Let $v_1 = 4t+2$, $v_2 = 2u+1$ and $v = v_1v_2$. Then from Lemma 1

$$N(v) \geq \min(N(v_1), N(v_2)).$$

Now $N(v_1) \geq 2$ by hypothesis, and since v_2 is odd $N(v_2) \geq n(v_2) \geq 2$. This proves the theorem.

EXAMPLE (10). Let $v = 110 = 22 \cdot 5$. From Theorem 5, $N(22) \geq 2$. Hence $N(110) \geq 2$.

The method of adjunction can obviously be applied to any resolvable BIB designs. In particular taking the design given by parameters (5.1), we get

THEOREM 8. *If $s = 2^m$, $m \geq 3$, then*

$$N\left(\frac{s(s-1)}{2} + 1\right) \geq \min\left(\frac{s}{2} - 1, N\left(\frac{s}{2} + 1\right)\right) - 1,$$

and if $1 < x \leq s+1$,

$$N\left(\frac{s(s-1)}{2} + x\right) \geq \min\left(\frac{s}{2} - 1, N\left(\frac{s}{2} + 1\right), N(x)\right) - 1.$$

EXAMPLE (11). Taking $s = 16$, $x = 9$, we get $N(129) \geq 6$, whereas $n(129) = 2$.

7. Certain pairwise balanced designs of index unity associated with the finite planes $\text{PG}(2, 2^3)$ and $\text{EG}(2, 3^2)$, and the corresponding sets of mutually orthogonal Latin squares. We shall now consider certain designs based on finite geometry which take advantage of the fact that 7, 8 and 9 are three consecutive integers which are prime powers.

Let us consider the finite projective plane $\text{PG}(2, 2^3)$. It has been shown in proving part (iii) of Theorem 3, that there exists on this plane a set Ω of 10 points, viz., the set of points $P_1, P_2, \dots, P_9$ on a nondegenerate conic C , and its pole P_0 , such that every line on the plane intersects Ω in either 2 points or no points. Let Ω^* be a subset of Ω . Then the number of points in Ω^* is $x \leq 10$. Consider the set of retained points obtained by deleting from the plane the points of Ω^* . If we take the set of retained points as treatments, and lines as blocks, then clearly each block contains either 7, 8 or 9 points. Also any pair of retained points occurs exactly in one block. Hence we obtain a pairwise balanced design of index unity and type $(73-x; 7, 8, 9)$. Applying Theorem 1, we have

THEOREM 9. $N(73-x) \geq 5$ if $x \leq 10$.

EXAMPLE (12). Taking $x = 3$ or $x = 7$ we find that there exist at least 5 m.o.l.s. of orders 70 or 66.

Again consider the Euclidean plane $EG(2, 3^2)$. If $ax^2 + 2hxy + by^2$ is irreducible in $GF(3^2)$, where a, b, c belong to this field, and if f, g, h are so chosen that

$$\Delta = abc + 2fgh - af^2 - bg^2 - ch^2 \neq 0$$

then the conic $ax^2 + 2hxy + by^2 + 2gx + 2fy + c = 0$ is nondegenerate and does not meet the line at infinity. If Ω is the set of 10 points on this conic, then no line meets Ω in more than two points. If we take Ω^* to be a subset of Ω , then the set of retained points obtained by deleting the points of Ω^* from the plane, provide a pairwise balanced design of the type $(81-x; 7, 8, 9)$ where $x \leq 10$. Hence using Theorem 1 we have

THEOREM 10. $N(81-x) \geq 5$, if $x \leq 10$.

EXAMPLE (13). Taking $x = 3$ or 7 we find that there exist at least 5 m.o.l.s. of orders 78 and 74.

8. Use of group divisible designs in constructing sets of mutually orthogonal Latin squares. An arrangement of v objects (treatments) in b sets (blocks) each containing k distinct treatments is said to be a group divisible (GD) design if the treatments can be divided into l groups of m treatments each, so that any two treatments belonging to the same group occur together in λ_1 blocks, and any two treatments from different groups occur together in λ_2 blocks. We will denote such a design by the notation $GD(v; k, m; \lambda_1, \lambda_2)$. The combinatorial properties of these designs have been studied in [7] where it has been shown that

$$(8.1) \quad v = lm, \quad bk = vr, \quad \lambda_1(m-1) + \lambda_2m(l-1) = r(k-1),$$

r being the number of replications, i.e., the number of times each treatment occurs in the design. It has also been shown that

$$(8.2) \quad P = rk - \lambda_2v \geq 0, \quad Q = r - \lambda_1 \geq 0.$$

The GD designs can be divided into three classes.

- (i) Regular (R) characterized by $P > 0, Q > 0$.
- (ii) Semiregular (SR) characterized by $P = 0, Q > 0$.
- (iii) Singular (S) characterized by $Q = 0$.

Methods of constructing these designs have been given in [8]. So far as the construction of m.o.l.s. is concerned a special role is played by GD designs with $\lambda_1 = 0, \lambda_2 = 1$, which in our notation can be denoted by $GD(v; k, m; 0, 1)$. If further this design is regular we shall denote it by $RGD(v; k, m; 0, 1)$ and if it is semiregular we shall denote it by $SRGD(v; k, m; 0, 1)$.

LEMMA 5. *Existence of a $GD(v; k, m; 0, 1)$ implies the existence of pairwise balanced design (D) of index unity and type $(v; k, m)$.*

If to the b blocks of the GD design, we add l new blocks corresponding to

the groups we get a pairwise balanced design of index unity and type $(v; k, m)$. It follows from Theorem 1:

THEOREM 11. *If there exists a $\text{GD}(v; k, m; 0, 1)$, then $N(v) \geq \min(N(k), N(m)) - 1$. Further if the GD design is separable $N(v) \geq \min(N(k), N(m))$.*

COROLLARY. *If s is a prime power $N(s^2 - 1) \geq N(s - 1)$.*

Proof. Since s is a prime power $N(s) = s - 1$, and we know from [8, p. 173] that there exists a symmetric and hence separable $\text{RGD}(s^2 - 1; s, s - 1; 0, 1)$.

EXAMPLE (14). Taking $s = 9$, $N(80) \geq N(8) = 7$, whereas $n(80) = 4$. We note that Parker's result is inapplicable in this case since there does not exist a BIB design with $v = 80$, and $k > 2$, $\lambda = 1$. Taking $s = 32$, we get $N(1023) \geq N(31) = 30$ whereas $n(1023) = 2$.

THEOREM 12. *Suppose there exists a $\text{GD}(v; k, m; 0, 1)$. Then $N(v - 1) \geq q - 2$ where*

$$q - 1 = \min(N(k), N(k - 1), N(m), N(m - 1)).$$

If further the GD design is resolvable then $N(v - 1) \geq q - 1$.

Proof. From Lemma 5, we get a pairwise balanced design (D) of index unity and type $(v; k, m)$. Omitting a particular treatment say α from the blocks in which it occurs, we get a pairwise balanced design (D_α) of the type $(v - 1; k, k - 1, m, m - 1)$. Theorem 1 now gives $N(v - 1) \geq q - 2$.

We shall now show that we can get $q - 1$ m.o.l.s. when the GD design is resolvable. There are l blocks in (D) , each of which contains m treatments, viz., the blocks corresponding to the groups. Consider the corresponding set of blocks of (D_α) . Since only one group contains the treatment α , this set of blocks consists of one block $\delta_{g, m-1}$ containing $m - 1$ treatments, and a set $(S_{g, m})$ of $l - 1$ blocks each of which contains m treatments. Since the initial GD design is resolvable, the other blocks of (D) are divisible into r sets, such that each set of blocks contains every treatment exactly once. From the i th of these sets, we obtain by dropping α the corresponding set of blocks of (D_α) , $i = 1, 2, \dots, r$. This set consists of v/k blocks, one of which $\delta_{i, k-1}$ contains $k - 1$ treatments and a set $(S_{i, k})$ of $(v/k) - 1$ blocks each of which contains k treatments.

Put $q_3 = N(m - 1) + 1$ and $q_4 = N(m) + 1$. Let P_3 and P_{3c} , $c = 1, 2, \dots, m - 2$ be the matrices of order $q_3 \times (m - 1)(m - 2)$ and $q_3 \times (m - 1)$ respectively of Lemma 2. Similarly let P_4 and P_{4d} , $d = 1, 2, \dots, m - 1$ be matrices of order $q_4 \times m(m - 1)$ and $q_4 \times m$ respectively.

Let Δ_g be the matrix of order $q \times ((m - 1)(m - 2) + m(m - 1)(l - 1))$ obtained from

$$[P_3(\delta_{g, m-1}), P_4(S_{g, m})]$$

by retaining the first q rows. Defining column equivalence as in the proof of Theorem 1 and remembering that (i) the number of submatrices P_{4d} is one

greater than the number of submatrices P_{3c} , (ii) the block $\delta_{g,m-1}$ and the set of blocks $(S_{g,m})$ contain between them all the $v-1$ treatments in the design (D_α) we can write

$$\Delta_g \sim [\Delta_g^{(1)}, \Delta_g^{(2)}, \dots, \Delta_g^{(m-2)}, \Delta_g^*],$$

where each row of $\Delta_g^{(j)}$ contains each of the $v-1$ retained treatments exactly once ($j=1, 2, \dots, m-2$). Thus $\Delta_g^{(j)}$ possesses the property τ_1 relative to the $v-1$ retained treatments. Each row of Δ_g^* contains all the $v-m$ retained treatments other than those contained in $\delta_{g,m-1}$. Also any pair of treatments occurring in the block $\delta_{g,m-1}$ or blocks $(S_{g,m})$ of (D_α) occurs as a column exactly once in any two rowed submatrix of Δ_g .

Likewise from $\delta_{i,k-1}$ and $(S_{i,k})$ we get a $q \times ((k-1)(k-2) + (k-1)(v-k))$ matrix Δ_i

$$\Delta_i \sim [\Delta_i^{(1)}, \Delta_i^{(2)}, \dots, \Delta_i^{(k-2)}, \Delta_i^*], \quad i = 1, 2, \dots, r,$$

possessing similar properties in relation to the blocks $\delta_{i,k-1}$ and $(S_{i,k})$.

Let Δ_0 be a $q \times (v-1)$ matrix whose j th column contains the j th treatment in (D_α) in every position. It follows from the property of pairwise balance possessed by (D) , that the set of blocks $\delta_{g,m-1}, \delta_{1,k-1}, \dots, \delta_{r,k-1}$ taken together contain each of the $v-1$ retained treatments exactly once. Let Δ_g^{**} be a $q \times (m-1)$ matrix each of whose rows is $\delta'_{g,m-1}$. Similarly let Δ_i^{**} be a $q \times (k-1)$ matrix each of whose rows is $\delta'_{i,k-1}$ ($i=1, 2, \dots, r$). Then

$$\Delta_0 \sim [\Delta_g^{**}, \Delta_1^{**}, \dots, \Delta_r^{**}],$$

clearly

$$\Delta = [\Delta_0, \Delta_1, \dots, \Delta_r],$$

has the property τ_2 in relation to the $v-1$ treatments of (D_α) . It follows from this or can be directly verified using (8.1) that Δ is of order $q \times (v-1)(v-2)$. Hence $[\Delta_0, \Delta]$ is an orthogonal array $[(v-1)^2, q, v-1, 2]$ of order $(v-1)^2$, q constraints, $v-1$ levels and strength 2.

We now note that the matrices $[\Delta_g^*, \Delta_g^{**}]$ and $[\Delta_i^*, \Delta_i^{**}]$, $i=1, 2, \dots, r$ have the property τ_1 , that each row contains every treatment of (D_α) exactly once. Hence by a reordering of columns, $[\Delta_0, \Delta]$ can be divided into $(v-1)$ submatrices of order $q \times (v-1)$ such that each submatrix has the property τ_1 . After this reordering has been made $[\Delta_0, \Delta]$ can be converted into an orthogonal array $[(v-1)^2, q+1, v-1, 2]$ of $q+1$ constraints, by adding another row such that the part of this row corresponding to the c th submatrix with the property τ_1 consists of the c th treatment repeated $v-1$ times. It follows that there exist $q-1$ m.o.l.s. of order $v-1$. This proves the second part of Theorem 12.

In the same manner we can prove:

COROLLARY. *If there exists a resolvable BIB $(v; k)$ then*

$$N(v-1) \geq \min(N(k), N(k-1)).$$

Let there exist a $\text{GD}(v; k, m; 0, 1)$. Then by taking a new treatment θ , and adding it to each of the l blocks of the design (D) of Lemma 5, we get a pairwise balanced design of index unity and type $(v+1; k, m+1)$. Using Theorem 1, we have

THEOREM 13. *If there exists a $\text{GD}(v; k, m; 0, 1)$ then*

$$N(v+1) \geq \min(N(k), N(m+1)) - 1.$$

Suppose the design $\text{GD}(v; k, m; 0, 1)$ is resolvable and r is the number of replications. We can take new treatments $\theta_1, \theta_2, \dots, \theta_x, x \leq r$, and add θ_i to each block of the i th replication. Also we take blocks corresponding to the groups of treatments and a new block consisting of the treatments $\theta_1, \theta_2, \dots, \theta_x$ when $x > 1$. We then get a pairwise balanced design of index unity and type $(v+x; k+1, k, m)$ if $x=1$, type $(v+x; k+1, k, m, x)$ if $1 < x < r$ and type $(v+x; k+1, m, x)$ if $x=r$. Using Theorem 1, we get

THEOREM 14. *Suppose there exists a resolvable $\text{GD}(v; k, m; 0, 1)$, then $N(v+x) \geq \min(N(k+1), N(k), N(m)) - 1$ if $x=1$, $N(v+x) \geq \min(N(k+1), N(k), N(m), N(x)) - 1$ if $1 < x < r$ and $N(v+x) \geq \min(N(k+1), N(m), N(x)) - 1$, if $x=r$.*

Theorems 12, 13 and 14 can be applied to semiregular GD designs, which provide a large class of resolvable GD designs. The following result is taken from [8, p. 181].

LEMMA 6. *The existence of an $\text{SRGD}(lm; l, m; 0, 1)$ implies the existence of an orthogonal array $A = (m^2, l, m, 2)$ of strength 2 (or equivalently the existence of a set of $l-2$ m.o.l.s. of order m), and conversely.*

Suppose there exists an $\text{SRGD}(lm; l, m; 0, 1)$. Then from Lemma 6, there exists an orthogonal array $A = (m^2, l, m, 2)$. Without loss of generality we can take the first row of A as $(\epsilon'_1, \epsilon'_2, \dots, \epsilon'_m)$ where ϵ'_i is a $1 \times m$ row vector each of whose elements is i . If we now delete the first row, there remains an orthogonal array $A_1 = (m^2, l-1, m, 2)$ of $l-1$ constraints which can be subdivided into m submatrices each of order $(l-1) \times m$ possessing the property τ_1 that every row of each submatrix contains the symbols $1, 2, \dots, m$ exactly once. This property remains unchanged if we get $A_2 = (m^2, l_1, m, 2)$ by deleting some rows of A_1 , $l_1 \leq l-1$. Replace the symbol j in the i th row of A_2 by the treatment θ_{ij} ($i=1, 2, \dots, l_1; j=1, 2, \dots, m$), and take the columns as blocks. We then get an $\text{SRGD}(l_1 m; l_1, m; 0, 1)$ for which the groups are given by the m treatments occurring in any given row. This design is resolvable, the blocks corresponding to the submatrices with the prop-

TABLE I

v	$n(v)$	l.b. for $N(v)$	Remarks
21	2	4	Th. 3, Ex. (1)
22*	1	2	Th. 5, $m=0$
24	2	3	Th. 11, Cor. $s=5$
39	2	3	Th. 12, Ex. (15)
50*	1	5	Th. 6, Ex. (7)
54*	1	3	Th. 6, Ex. (8)
57	2	7	Th. 2, BIB (57; 8)
58*	1	2	Th. 5, $m=1$
65	4	6	Th. 6, Ex. (7)
66*	1	5	Th. 9, $x=7$
68	3	5	Th. 9, $x=5$
69	2	5	Th. 9, $x=4$
70*	1	5	Th. 9, $x=3$
74*	1	5	Th. 10, $x=7$
75	2	5	Th. 10, $x=6$
76	3	5	Th. 10, $x=5$
78*	1	5	Th. 10, $x=3$
80	4	7	Th. 11, Cor. $s=9$
82*	1	3	Th. 14, Ex. (17)
84	2	5	SRGD (77; 7, 11; 0, 1), Th. 14, $x=7$
85	4	5	SRGD (77; 7, 11; 0, 1), Th. 14, $x=8$
86*	1	5	Th. 14, Ex. (17)
92	3	5	SRGD (91; 7, 13; 0, 1), Th. 14, $x=1$
94*	1	2	Th. 5, $m=2$
95	4	5	SRGD (88; 8, 11; 0, 1), Th. 14, $x=7$
96	2	6	SRGD (88; 8, 11; 0, 1), Th. 14, $x=8$
98*	1	5	SRGD (91; 7, 13; 0, 1), Th. 14, $x=7$
100	3	5	SRGD (91; 7, 13; 0, 1), Th. 14, $x=9$
102*	1	5	SRGD (91; 7, 13; 0, 1), Th. 14, $x=11$
105	2	6	SRGD (104; 8, 13; 0, 1), Th. 14, $x=1$
110*	1	2	Th. 7, Ex. (10)
111	2	5	SRGD (104; 8, 13; 0, 1), Th. 14, $x=7$
115	4	6	SRGD (104; 8, 13; 0, 1), Th. 14, $x=11$
120	2	7	Th. 3, pt. (iii), $s=16$
123	2	5	SRGD (112; 7, 16; 0, 1), Th. 14, $x=11$
126*	1	5	SRGD (119; 7, 17; 0, 1), Th. 14, $x=7$
129	2	6	Th. 8, Ex. (11)
130*	1	5	SRGD (119; 7, 17; 0, 1), Th. 14, $x=11$
132	2	5	SRGD (119; 7, 17; 0, 1), Th. 14, $x=13$
134*	1	5	SRGD (133; 7, 19; 0, 1), Th. 14, $x=1$
135	4	6	Th. 12, Ex. (15)
138*	1	3	SRGD (133; 7, 19; 0, 1), Th. 14, $x=5$
140	3	5	SRGD (133; 7, 19; 0, 1), Th. 14, $x=7$
141	2	6	SRGD (128; 8, 16; 0, 1), Th. 14, $x=13$
142*	1	5	SRGD (133; 7, 19; 0, 1), Th. 14, $x=9$
145	4	7	Th. 13, Ex. (16)
146*	1	5	SRGD (133; 7, 19; 0, 1), Th. 14, $x=13$
147	2	6	SRGD (136; 8, 17; 0, 1), Th. 14, $x=11$
150*	1	5	SRGD (133; 7, 19; 0, 1), Th. 14, $x=17$

erty τ_1 give a complete replication. Clearly $r = m$ for this design. We thus have

LEMMA 7. *The existence of an SRGD($lm; l, m; 0, 1$) implies the existence of a resolvable SRGD($l_1m; l_1, m; 0, 1$), and an orthogonal array $A_2 = (m^2, l_1, m, 2)$ divisible into submatrices with property τ_1 , for any $l_1 < l$.*

COROLLARY. *There exists a resolvable SRGD($l_1m; l_1, m; 0, 1$) if $l_1 \leq N(m) + 1$.*

Proof. Let $l = N(m) + 2$. Then there exists an orthogonal array $(m^2, l, m, 2)$. From Lemma 6, there exists an SRGD($lm; l, m; 0, 1$). The required result now follows from Lemma 7.

EXAMPLE (15). From SRGD(40; 5, 8; 0, 1) which is resolvable, using Theorem 12, we get $N(39) \geq 3$, whereas $n(39) = 2$. Similarly from SRGD(136; 8, 17; 0, 1) we get $N(135) \geq 6$, whereas $n(135) = 4$.

EXAMPLE (16). From SRGD(144; 9, 16; 0, 1) and Theorem 13, we get $N(145) \geq 7$, whereas $n(145) = 4$.

EXAMPLE (17). From SRGD(77; 7, 11; 0, 1) and Theorem 14 with $x = 5$ and 9 respectively, we get $N(82) \geq 3$, $N(86) \geq 5$, whereas $n(82) = n(86) = 1$.

9. Table of lower bounds of $N(v)$ for $v \leq 150$. We provide below a table of lower bounds (l.b.) of $N(v)$ for those values of $v \leq 150$, where the lower bound exceeds $n(v)$. In the last column under remarks we indicate how the corresponding lower bound was obtained. The numbers for which $v \equiv 2 \pmod{4}$ are starred.

10. Concluding remarks. We are conscious of the fact that the table given in the last section may be incomplete. There may be other values of v for which $N(v) > n(v)$, as the pairwise balanced designs of index unity have not been exhaustively studied. Indeed they were improvised for the purposes of the present paper, though we hope that they may be useful in other combinatorial problems. Only a lower bound for $N(v)$ has been given and it may be possible to improve it in many cases.

Many new interesting problems arise as a result of the present work. A positive integer n may be called an Eulerian number if there do not exist more than one orthogonal Latin square of order n . The density and distribution of Eulerian numbers is a question of great interest. The table of the last section shows that there exists a run of 24 consecutive non-Eulerian numbers from 63 to 86. We shall show in a later communication by extending the methods of §7, that there exist arbitrarily large runs of consecutive non-Eulerian numbers. In other words there exist arbitrarily large gaps between consecutive Eulerian numbers.

It may also be observed that analogous methods for constructing orthogonal arrays of strength two and index λ can be made to depend on pairwise balanced designs of index λ , which are defined as in §3, except that every pair of distinct treatments occurs in λ different blocks. It is hoped to consider these and allied problems in a subsequent communication.

APPENDIX
22×22 ORTHOGONAL LATIN SQUARES

1	Aa	Dp	Gv	Pb	Ft	Te	Vc	Oq	Sn	Um	Lr	Rk	Ju	Is	Qh	Bd	Ho	Kl	Ni	Ef	Mj	Cg
2	Pd	Bb	Eq	Ap	Qc	Gu	Uf	Jt	Or	Th	Vn	Ms	Sl	Kv	Ri	Da	Ce	Io	Lm	Hj	Fg	Nk
3	Vg	Qe	Cc	Fr	Bq	Rd	Av	Lp	Ku	Os	Ui	Ph	Nt	Tm	Sj	Hl	Eb	Df	Jo	Mn	Ik	Ga
4	Bp	Pa	Rf	Dd	Gs	Cr	Se	Un	Ds	Nr	Ot	Vj	Qi	Hu	Tk	Im	Fc	Jn	Eg	Ko	Nh	Jl
5	Tf	Cq	Qb	Sg	Ee	At	Ds	Iv	Vh	Lv	Mp	Ou	Pk	Rj	Ul	Bc	Bc	Jn	Gd	Fa	Lo	Hi
6	Et	Ug	Dr	Rc	Ta	Ff	Bu	Sk	Jp	Pi	Hs	Nq	Ov	Ql	Vm	Ij	Ln	Cd	Kh	Ae	Gb	Mo
7	Cv	Fu	Va	Es	Sd	Ub	Gg	Rm	Tl	Kq	Qj	It	Hr	Op	Pn	No	Jk	Mh	De	Li	Bf	Ac
8	Qo	Tj	Pl	Nu	Vi	Ks	Mr	Hh	Ev	Bt	Sf	Cp	Rg	Ud	Aq	Lc	Oa	Gm	Fk	Bj	Dn	Ie
9	Ns	Ro	Uk	Qm	Hv	Pj	Lt	Ve	Ii	Fp	Cu	Tg	Dq	Sa	Br	Jf	Md	Ob	An	Gl	Kc	Eh
10	Mu	Ht	So	Vl	Rn	Ip	Qk	Tb	Pf	Jj	Gq	Dv	Ua	Er	Cs	Fi	Kg	Ne	Oc	Bh	Am	Ld
11	Rl	Nv	Iu	To	Pm	Sh	Jq	Fs	Uc	Qg	Kk	Ar	Ep	Vb	Dt	Me	Gj	La	Hf	Od	Ci	Bn
12	Kr	Sm	Hp	Jv	Uo	Qn	Ti	Pc	Gt	Vd	Ra	Ll	Bs	Fq	Eu	Ch	Nf	Ac	Mb	Ig	Oe	Dj
13	Uj	Js	Tn	Iq	Kp	Vo	Rh	Gr	Qd	Au	Pe	Sb	Mm	Ct	Fv	Ek	Di	Hg	Bb	Nc	Ja	Of
14	Si	Vk	Mt	Uh	Jr	Lq	Po	Du	As	Re	Bv	Qf	Tc	Nn	Gp	Og	Fl	Ej	Ia	Cm	Hd	Kb
15	Hq	Ir	Js	Kt	Lu	Mv	Np	Qa	Rb	Sc	Td	Ue	Vf	Pg	Oo	Gn	Ah	Bi	Cj	Dk	El	Fm
16	Db	Ad	Lh	Ba	Mk	Ji	On	Cl	Fj	If	Em	Hc	Ke	Go	Ng	Pp	Rv	Tu	Vt	Qs	Sr	Uq
17	Oh	Ec	Be	Mi	Cb	Nl	Kj	Ao	Dm	Gk	Jg	Fn	Id	Lf	Ha	Vr	Qq	Sp	Uv	Pu	Rt	Ts
18	Lk	Oi	Fd	Cf	Nj	Dc	Hm	Mg	Bo	En	Al	Ka	Gh	Je	Ib	Ut	Ps	Rr	Tq	Vp	Qv	Su
19	In	Ml	Oj	Ge	Dg	Hk	Ed	Kf	Na	Co	Fh	Bm	Lb	Ai	Jc	Tv	Vu	Qt	Ss	Ur	Pq	Rp
20	Fe	Jh	Nm	Ok	Af	Ea	Il	Bj	Lg	Hb	Do	Gi	Cn	Mc	Kd	Sq	Up	Pv	Ru	Tt	Vs	Qr
21	Jm	Gf	Ki	Hn	Ol	Bg	Fb	Nd	Ck	Ma	Ic	Eo	Aj	Dh	Le	Rs	Tr	Vq	Pr	Sv	Uu	Pt
22	Gc	Kn	Ag	Lj	Ih	Om	Ca	Ei	He	Di	Nb	Jd	Fo	Bk	Mf	Qu	St	Us	Qp	Rq	Tp	Vv

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